K-Dimensional Trees

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Introduction

- Invented by Jon Louis Bentley in 1975
- Index structure for managing multidimensional objects
- k is the dimensionality of the search space
- Space partitioning is performed via building a binary search tree
- Each node is equal to a k-dimensional point
- Performs only axis aligned splits

Application Domain

- Mostly theoretical
- Can be used for
 - Nearest Neighbor Search
 - Search k-dimensional point in tree nearest to a given point
 - Suffers from curse of dimensionality
 - Range searches
 - Search for ranges of attributes
 - For example query income-ranges for persons
 - Worst complexity O(k * N^(1-1/k))

Complexity

Action	Average Complexity	Worst Complexity
Space	O(n)	O(n)
Insertion	O(log n)	O(n)
Query	O(log n)	O(n)
Removal	O(log n)	O(n)

Examples

- Construction
- Insertion
- Deletion
- Query/Lookup
- NN-Search

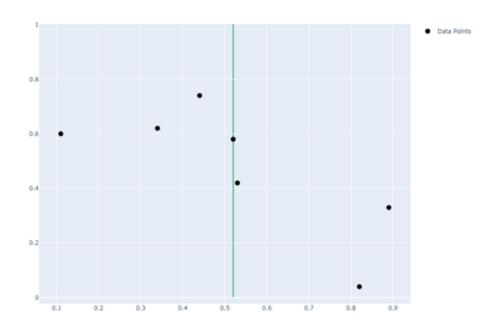
Construction

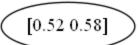
Construction

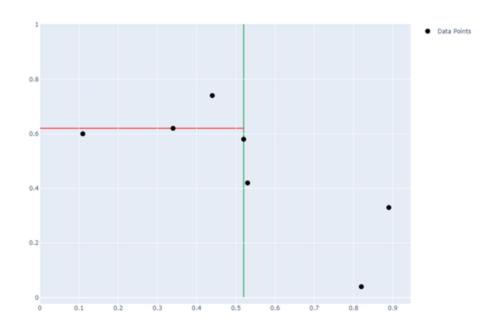
1. as long as >= 1 element remaining:

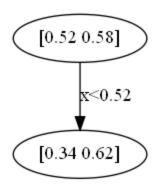
Pick a split dimension (f.e. x or y)

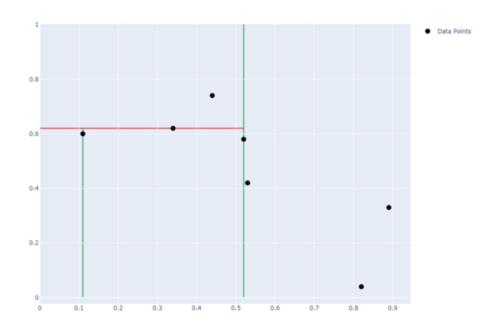
- canonical (increasing)
- o random
- 2. Compute the median
- 3. Axis aligned split into left and right subtree at median
- 4. repeat 1

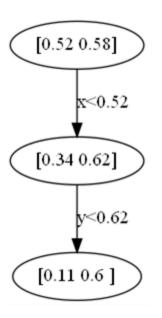


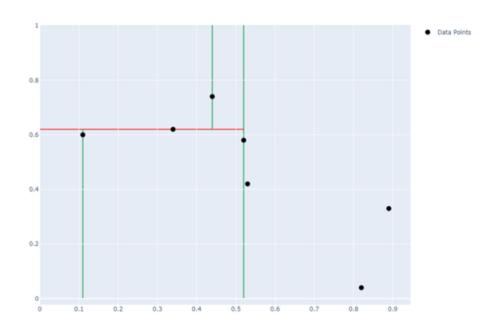


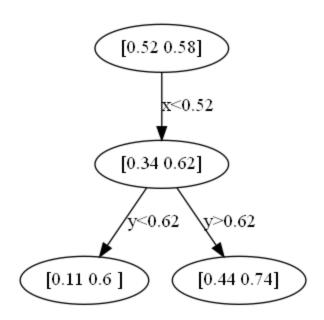


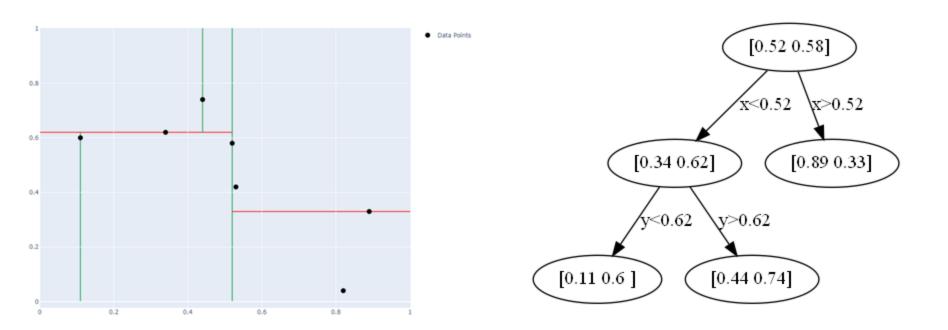


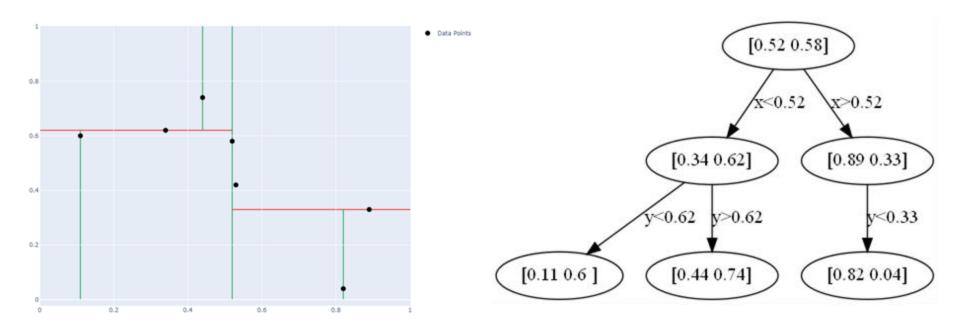


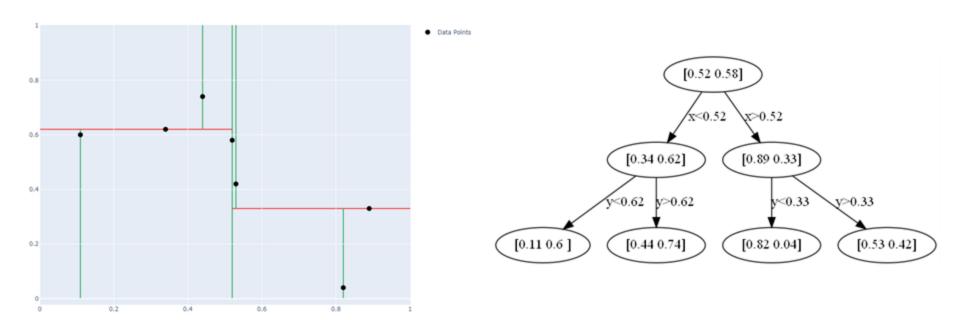








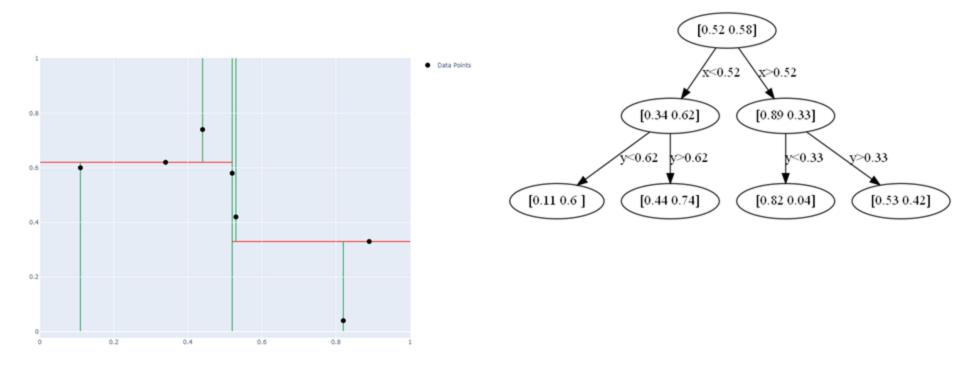


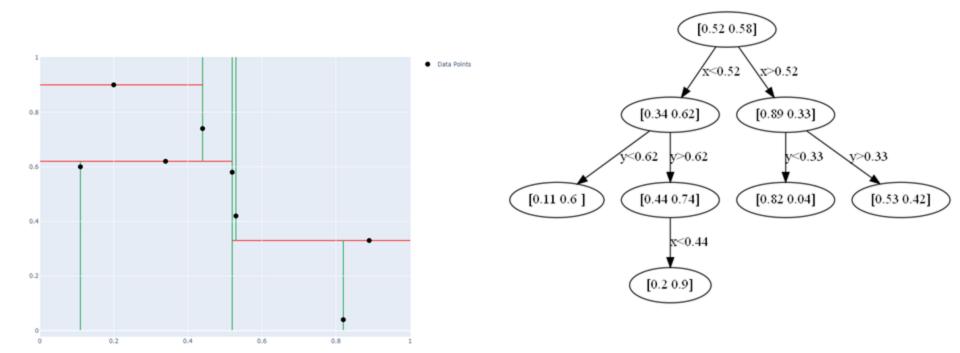


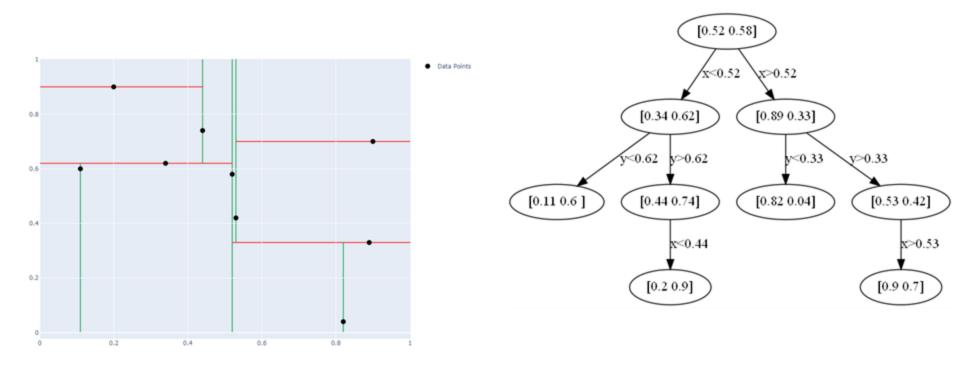
Insertion

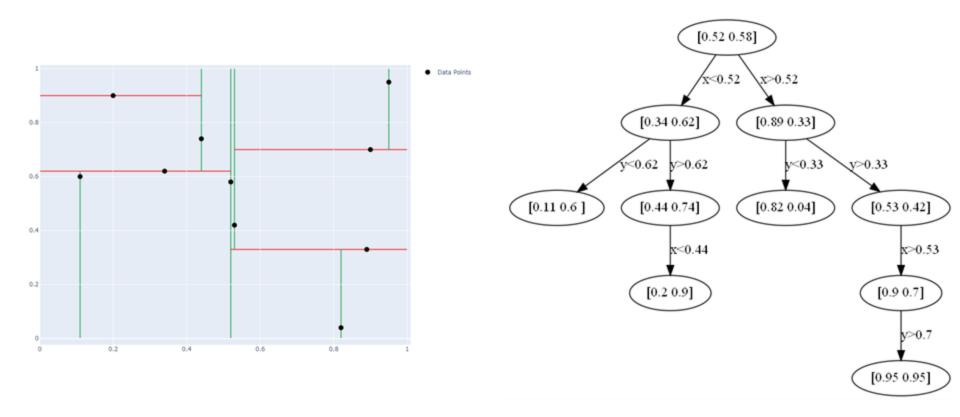
Insertion

- 1. Start traversing the tree from the root
- 2. Compare point to insert with current node
- 3. Choose left or right subtree depending on comparison
- 4. Traverse until leaf is reached
- 5. Insert node below the leaf depending on its relative position
- Warning:
 - Tree may become unbalanced during insertions (NN-Search performance suffers)
 - Rebalancing steps possible





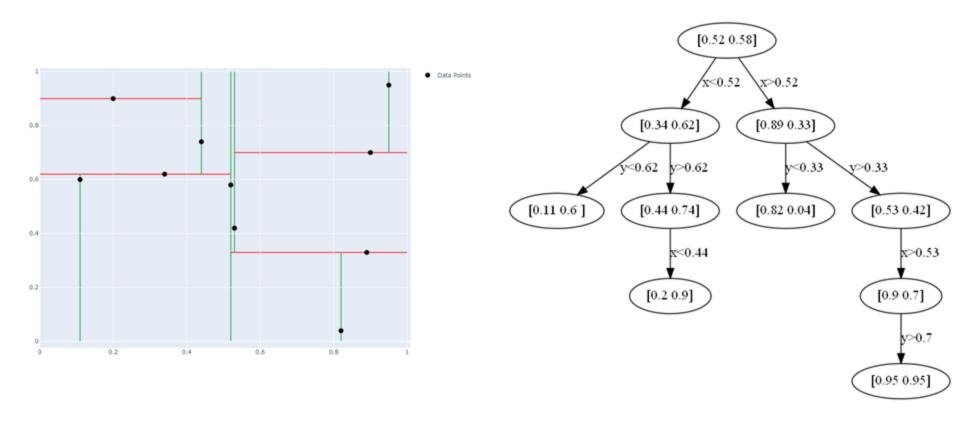


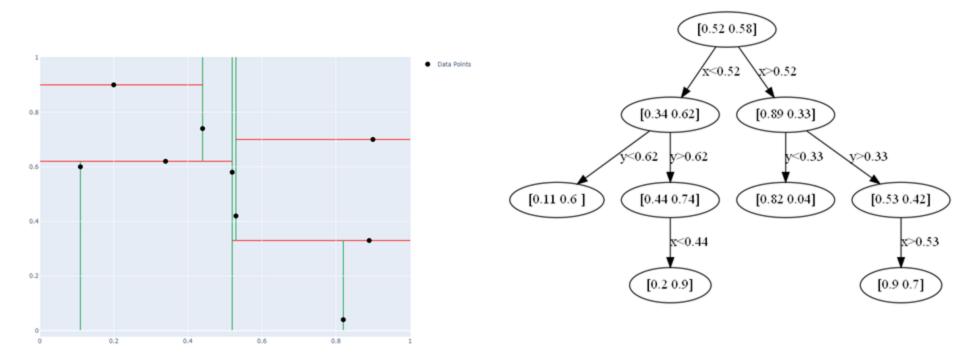


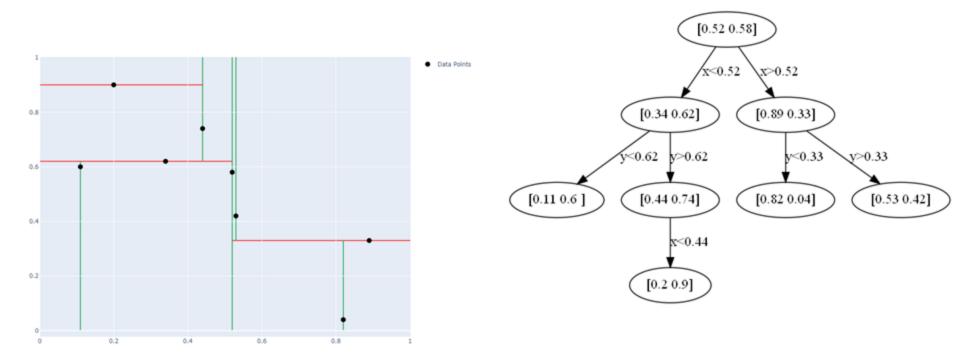
Deletion

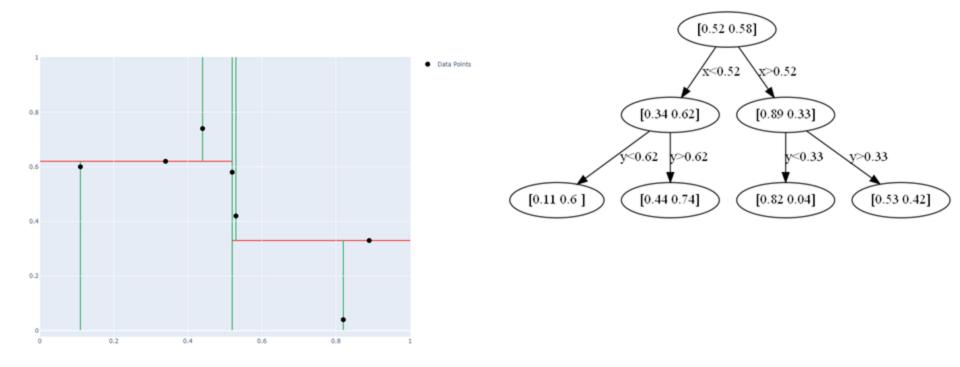
Deletion

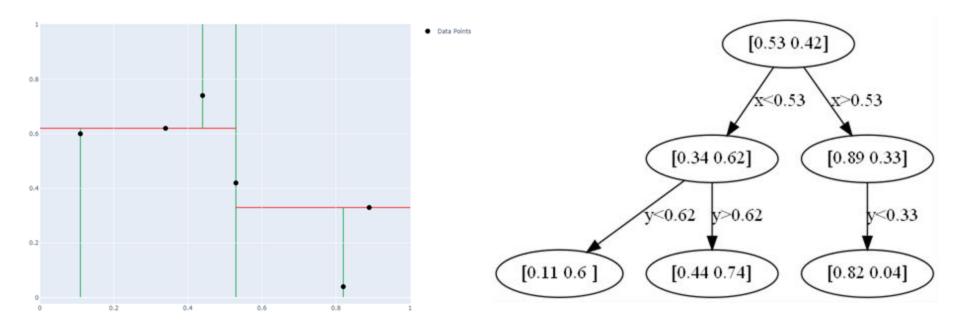
- 1. Start traversing the tree from the root
- 2. if current node == target node
 - a. find "in-order successor" in right subtree
 - b. replace current node with "in-order successor"
 - c. recursively start deletion from "in-order successor"
- 3. else
 - a. if current.left.contains(target): repeat 1. from current.left
 - b. else: repeat 1. from current.right
- Warning:
 - Tree may become unbalanced during insertions (NN-Search performance suffers)
 - Rebalancing steps possible









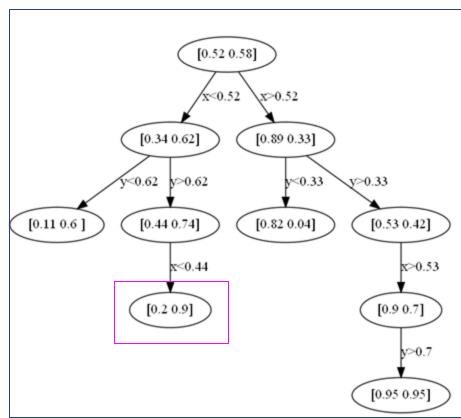


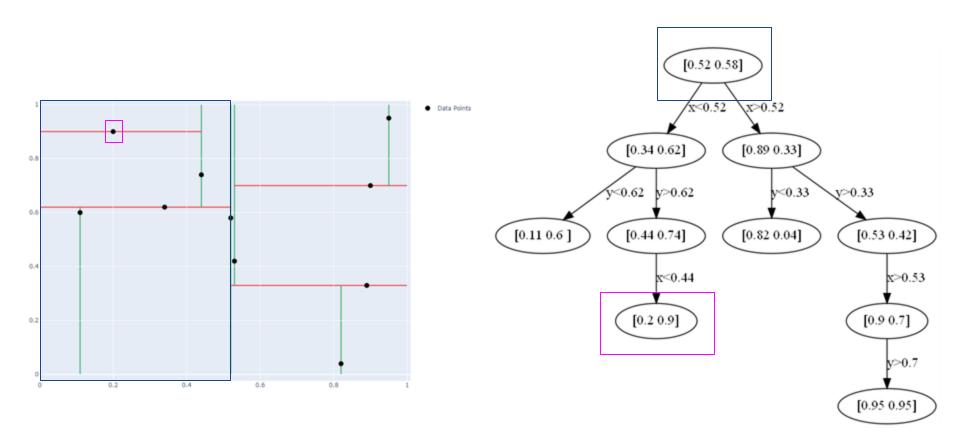
Query / Lookup

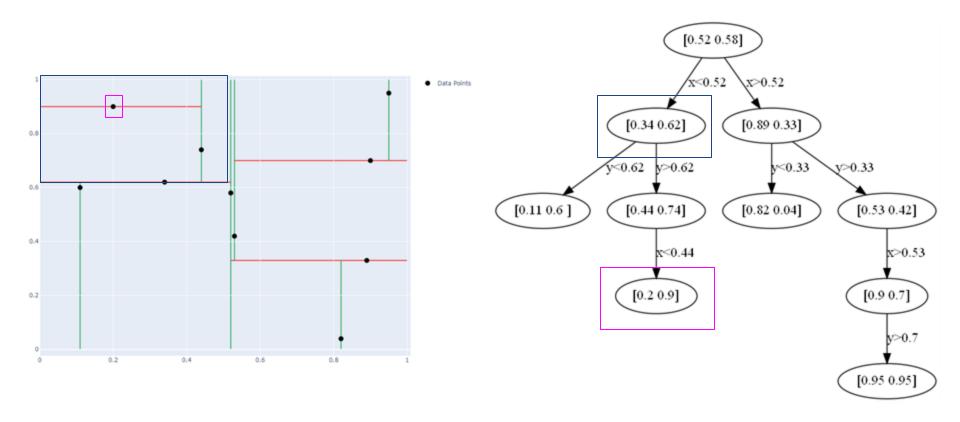
Query

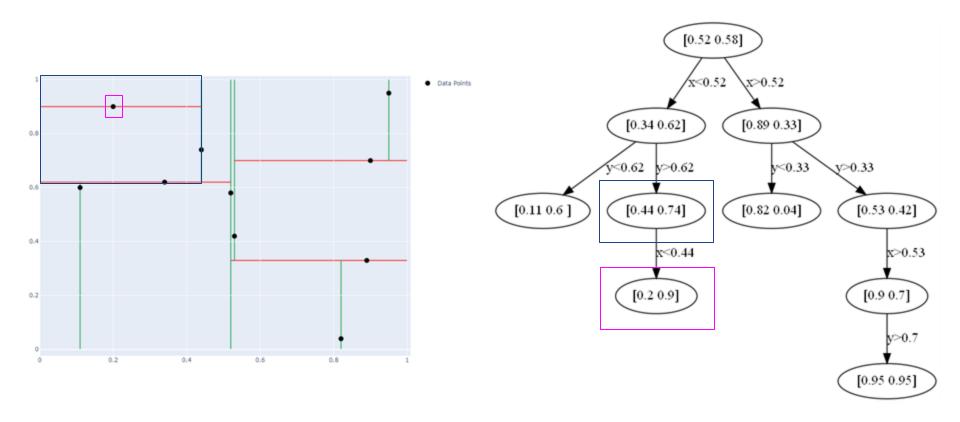
- 1. Start at root
- 2. Traverse tree down till search element is found or a not matching leaf is reached
- 3. At every node compare value to given condition
- 4. Traverse the subtree depending where the condition is true

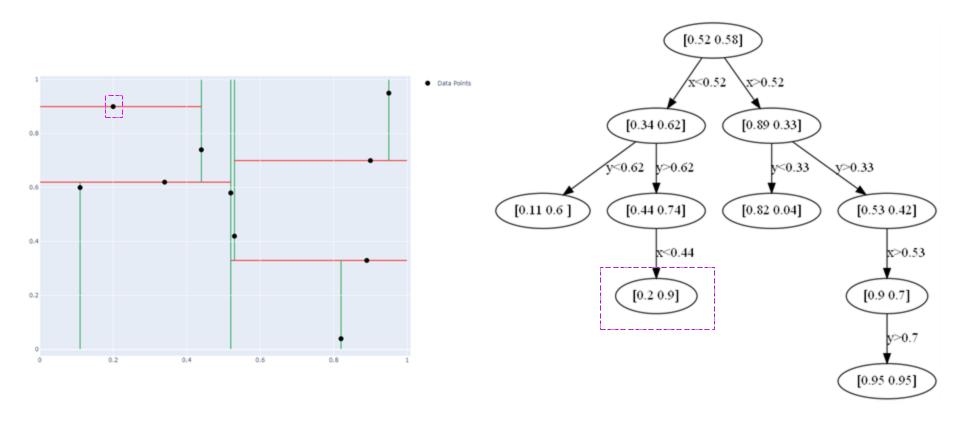












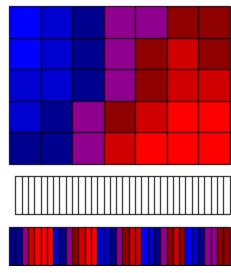
Nearest Neighbor Search

- start at root
- move recursively down to leave
- hypersphere around search point (radius = nearest distance)
 - o crossed hyperplane to other branch -> maybe nearer points on the other side
 - o if so search point moves down on the other side of the plane
 - o if not other side is eliminated, search point moves up
- ends at root

Variations

Variations and Related Versions

- Adaptive kd-Tree
 - successive levels split along different dimensions
- Implicit kd-Tree
 - defined implicitly above a rectilinear grid
- Min/Max kd-Tree
 - min/max values of each inner node are equal to min/max values of a child note -> save memory
- Relaxed kd-Tree
 - leaves of node do not have to have the same split axis
 - this relaxation allows insertions at arbitrary dimensions without reordering





Implicit kd-tree example

https://upload.wikimedia.org/wikipedia/commo ns/b/b7/Implicitmaxkdtree.gif Last accessed 25.5.20