Interfacing The Am9511 Arithmetic Processing Unit

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Introduction

If you are interested in a hardware solution to the problem of addition, subtraction, multiplication, division, and functions such as sine, cosine, tangent, square root, exponential, logarithm and their inverse functions, then the Am9511 integrated circuit will be of interest to you. The Am9511 Arithmetic Processing Unit is a product of Advanced Micro Devices Inc., 901 Thompson Place, Sunnyvale, CA 94086. It performs signed multiplication, addition, subtraction and division with either 16-bit integers or 32-bit integers, in twos complement form. It also does these operations and evaluates a variety of functions (mentioned above) in a 32-bit floating point form. In the floating point form, the mantissa of the number is represented by 24 bits (equivalent to approximately seven significant decimal digits). The exponent is represented by six bits and a sign bit, giving a range of numbers that can be represented from roughly 10-19 to 10 + 19. The one bit not accounted for so far is the sign of the mantissa. Thus, the Am9511 should satisfy most of the calculating needs of microcomputer users. It is important to point out that the Am9511 is a binary device as opposed to a BCD device. If you intend to use it like a calculator, then appropriate BCD-to-binary and binary-to-BCD routines will be needed to input and output numbers.

Timing of the various control pins on the Am9511 is one of the most important considerations in constructing an interface between it at the microprocessor. The timing requirements seem to be more relaxed in the most recent specification sheets, but my original specifications were quite complex. Perhaps it would be easy to interface the Am9511 somewhere in the address space, using address lines and control lines to operate it. However, given the complexities of the original timing diagrams, we used

an interface adapter (the 6522, although any of the other popular interface adapters such as the 6530 can also be used with our programs). One port is used for data transfers, while several pins of the other port on the interface adapter is used to control the Am9511. These techniques produce an extremely simple interface at the expense of some overhead in software.

Before proceeding to the details of the circuit and the driver programs it should be pointed out that if you are interested in building and using this or some other circuit that uses the Am9511, you will want to get complete specification sheets, a publication called "Algorithm Details for the Am9511 Arithmetic Processing Unit," and a card-type Am9511 reference card. All three of these publications are available from Advanced Micro Devices. The Am9511 itself costs about \$200, a number which may cause you to turn to the next article. A few mail order houses such as Advanced Computer Products are beginning to list the chip in their advertisements. Be sure to request all the literature mentioned above because you will need it to know how to use the chip. Space does not permit us to write a complete description of all the features of the chip.

The Am9511 Interface Circuit

The interface circuit is given in Figure 1. It is very simple because the complexity is absorbed in the software that must accompany this circuit. As noted, any 6502 system such as the SUPERKIM, KIM-1, AIM 65, etc., may be used, and any two-port interface adapter can be used. Be sure to include the 0.01 microfarad bypass capacitors, keep the leads between the Am9511 and the microcomputer short, and tie the unused control inputs (EACK and SVACK) to logic one as shown in Figure 1. I will not reveal how many hours of grief the failure to follow these standard procedures cost me. Keep it simple, neat, and don't try any shortcuts. Also follow the usual procedures in handling integrated circuits that are susceptible to damage by static discharge. This is not your typical El Cheapo IC: \$200 makes it irreplaceable. Avoid any Benjamin Franklin type experiments.

The Driver Subroutines

Listing 1 gives five subroutines that work with the interface circuit in Figure 1 to operate the Am9511. The subroutines are:

- RESET A subroutine that is used to reset the Am9511 either after power is applied or to clear the Am9511 to a known condition. This subroutine must be called after power-up and before using the Am9511.
- 2. WRITE This subroutine transfers a byte of data in the accumulator of the 6502 to the stack of the Am9511.
- COMMAND A subroutine that transfers an eight-bit command word from the accumulator



of the 6502 to the command register of the Am9511.

- READ Subroutine READ takes one byte of data (part of the answer) from the stack of the Am9511 and returns it to the X - register in the 6502.
- STATUS This subroutine reads the status register of the Am9511 and transfers its contents to the X - register in the 6502.

The comments in the various subroutines should be studied in connection with the Am9511 specification sheets to understand the functions of the various instructions. We only note here that each of the access subroutines, WRITE, COMMAND, READ, and STATUS, wait for the Am9511 to signal that an operation is complete when its PAUSE pin returns to logic one.

We will describe a few operations with the Am9511 to illustrate how the subroutines work. Refer to the literature mentioned previously for more details on the stack operation. The Am9511 stack may be regarded either as an eight-level, 16-bit wide stack, or as a four-level, 32-bit wide stack. Writing once to the Am9511 places an 8-bit word on the stack. However, since all of the "words" operated on by the Am9511 are either 16 bits or 32 bits wide, you must write at least 16 data bits (two bytes) to fill a 16-bit stack location. You must write four bytes to fill a 32-bit stack location. The last level filled (either 16 bits or 32 bits wide) is called TOS (acronym for top of stack). The level filled previously is referred to as NOS (next on stack).

An example will clarify the operation of the stack. Suppose we wish to add two 16-bit integers (they must be in twos complement form). Using the WRITE subroutine, we write the least-significant byte of one of the numbers to the Am9511 stack. Call this byte B1. Next we write B2, the mostsignificant byte of the same integer, to the Am9511. This puts a 16-bit integer onto TOS, the top level of the stack. The other addend, call it A1 and A2 for the least-significant and most-significant bytes respectively, is placed on the TOS by calling subroutine WRITE two more times. Now number B (B1 and B2) is in NOS and A (A1 and A2) is in TOS. The command code for a 16-bit addition, \$6C, is now placed in the 6502 accumulator and subroutine COMMAND is called. The Am9511 adds TOS to NOS and puts the result into TOS. The result R, consisting of the most-significant byte R1 and the least-significant byte R2 of the 16-bit answer, is obtained by calling subroutine READ. The first call of READ retrieves the most-significant byte R2, and the second call of READ retrieves the leastsignificant byte of the result R. The status register can be read to see if the addition produced a carry or an overflow.

Subtraction follows exactly the same pattern. The minuend M is loaded on the stack, followed by the subtrahend S to obtain the difference D where D = M - S. After M and S are loaded on the stack, the subtraction command (\$2D for a 32-bit word) will result in the difference D in TOS. Calling subroutine READ (twice for a 16-bit integer, four times for a 32-bit integer) gives the answer in the order from most-significant byte to least-significant byte. In division, the dividend is loaded on the stack followed by the divisor, and the quotient is read after the operation is completed. Some of you will recognize that the Am9511 uses RPN.

A program to illustrate these 16-bit operations is given in Listing 2. Suppose we wish to subtract \$32FC from \$FF5B. We would load \$5B into location \$0004, \$FF into location \$0003, \$FC into location \$0002, and \$32 would be loaded into location \$0001. The 16-bit subtraction command for the Am9511, \$6D, would be loaded into location \$0000. The program in Listing 2 will call the appropriate subroutines and place the answer in locations \$00FF (most-significant byte) and \$00FE (least-significant byte). This program can be used to test many of the operations of the Am9511, including sine, cosine, etc., by loading a 32-bit number (fixed or floatingpoint representation) on the stack, and then placing a command on the stack. It is a nice simple test program, but remember that many of the Am9511 functions require that the argument is in floating point form, so to find the square root of four requires that you convert four to a floating-point number. The Am9511 will do this if you either cannot or will not.

A word about execution time may be useful at this point. Instructions take from 16 clock cycles for a 16-bit integer addition to several thousand clock cycles for functions like sine, cosine, etc. We operated our Am9511 at 1MHz, but it can be operated at 2MHz and other versions go as high as 4MHz. Clearly the subroutines in Listing 1 require a significant amount of overhead for the simple integer operations, but become insignificant in terms of time overhead when the complex functions are called. Perhaps some reader will design an interface where instructions like STA DATA, STA COMMAND, LDA DATA, and LDA STATUS can be used instead of the subroutines. The difficulty is in working out the necessary timing requirements for the READ and WRITE operations of the 6502. The Am9511 timing seems to be more closely related to 8080A systems than either 6502 systems or 6800 systems.

Our final illustrative program is one that was designed to generate a sine table consisting of one cycle of a sine wave residing in one page of memory. The amplitude of the sine wave is \$7F00, in other words, we found \$7F00*Sin[Y*(Pi/128] where Y is a number that varied from \$00 to \$FF (0 to 255). This result was converted to a 16-bit fixed point format, and the most-significant byte was stored in a table in page \$0E, while the least-significant byte was stored in a table in page \$0F. Note that the result will be in twos complement form so at location \$0E80 in the

table when we are exactly half-way through the sine wave, you will find \$00, but at location \$0E81 you will find the first negative value of the sine wave and it is \$FC, the one in the most-significant bit of the 16-bit result indicating a minus number.

What do you do with a sine wave table? You could read it out to a D/A converter at various rates and play a tune, or you could add a series of sine waves to make a more complex sound. My purpose was to test the AM9511 and in the future I will use the sine wave table as part of a fast-Fourier transform program (I hope). Instead of synthesizing music I would really like to synthesize \$20 bills. Let me know if you succeed.

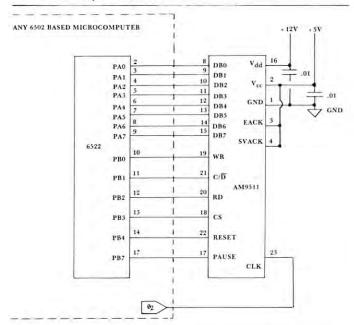


Figure 1. Interfacing the AM9511 Arithmetic Processing Unit to a 6522 VIA Chip. Other interface adapters that may be used include the 6520, the 6530 and the 6532. No special handshaking pins are used.

Listing 1	Subroutines to driv	e the AM9511	
0300 A9 1F RESE		Make PB0 - PB4	
0302 8D 02 A0	STA PBDD	output pins to con- trol the AM9511.	
0305 A9 0F	LDA \$0F	RESET pin to	
0307 8D 00 A0	STA PBD	logic zero.	
030A A9 1F	LDA \$1F	Hold RESET high	
030C 8D 00 A0	STA PBD	for at least five	
030F EA	NOP	clock cycles.	
0310 EA	NOP		
0311 A9 0F	LDA \$0F	Bring RESET pin	
0313 8D 00 A0	STA PBD	to logic zero to run the AM9511.	
0316 60	RTS	Return to the call- ing program.	
	*********	0.	
0320 8D 01 A0 WRIT	TE STA PAD	A contains the	
0323 A9 04	LDA \$04	byte to be written	
0325 8D 00 A0	STA PBD	to the AM9511 (A = accumulator) CS low, C/D	
		low, WR low.	
0328 AD 00 A0WAIT	r LDA PBD	Read PBD to see if PAUSE pin is at	

LDA \$FF STA PADD INC PBD LDA \$0F	transfer allowed). If PAUSE is high make PAD an output port to transfer data to the AM9511. Bring WR high to
STA PADD	make PAD an output port to transfer data to the AM9511. Bring WR high to
INC PBD	output port to transfer data to the AM9511. Bring WR high to
	transfer data to the AM9511. Bring WR high to
	the AM9511. Bring WR high to
	Bring WR high to
LDA \$0F	complete data
LDA \$0F	transfer.
A STATE OF THE PARTY OF THE PAR	Next bring CS,
	C/D high.
STA PBD	9
LDA \$00	Now make Port A
STA PADD	(PAD) an input
	port again.
RTS	Return to the
	calling program.

STA PAD	A contains the
	command for the
040-120	AM9511.
	CS low, C/D
	high, WR low.
	Is PAUSE low?
BPL LOAF	Yes, then wait
* 5 . 455	until it goes high.
	Make Port A an
	output port.
	Bring WR high.
	Bring other con-
	trol pins high.
	Return Port A to
	input status.
	CS low, C/D low,
	RD low.
	Read PBD to see
20.1.20	if PAUSE is low.
BPL LOITE	
LDX PAD	until it goes high.
	Am9511 output
	to X register.
LDA \$0F	Bring control pins
STA PBD	high.
RTS	Return to calling
	program with out-
	put in X.

LDA \$03	CS low, C/D
	high, RD low.
	Is PAUSE low?
BPL DELAY	Yes, then wait
	until it goes high,
	Read status regis-
LDA \$0F	ter of AM9511
	and keep it in the
CTA DDD	X register.
STA PBD	Bring control pins
DTC	high.
KIS	Status is in X
	upon return.
	STA PADD RTS STA PAD LDA \$06 STA PBD LDA PBD BPL LOAF LDA \$FF STA PADD INC PBD LDA \$00 STA PBD LDA \$00 STA PADD RTS LDA \$01 STA PBD LDA PBD LDA \$01 STA PBD LDA PBD LDA \$01 STA PBD LDA PBD RTS LDA \$01 STA PBD RTS LDA \$01 STA PBD RTS LDA \$01 STA PBD RTS

Listing 2 Program that loads four bytes (32 bits) and a command into the Am9511		0527 A9 00 0529 20 20 03 052C A9 1D	LDA \$00 JSR WRITE LDA \$1D	stack, Y into TOS. Change Y into	
			052E 20 40 03	JSR	floating point
0400 20 00 03 START	JSR RESET	Reset the AM9511		COMMAND	form.
		to start using it.	0531 A9 12	LDA \$12	Multiply to get
0403 A2 03	LDX #03	Initialize X to	0533 20 40 03	JSR COMMAND	Y*(Pi/128). Result to NOS.
0405 PF 04 4 00P	* D. I. D. I. M	count four bytes.		COMMAND	Pop stack up.
0405 B5 01 LOOP	LDA DATA,	KGet byte from the data table.	0536 A9 02	LDA \$02	Take SIN[Y*
0407 20 20 03	JSR WRITE	Write the byte in-	0538 20 40 03	JSR	(Pi/128)], result
0107 20 20 03	Jon While	to		COMMAND	to TOS.
		the Am9511.	053B A9 00	LDA \$00	Push \$7F00 on
040A CA	DEX	Decrement byte	053D 20 20 03	JSR WRITE	stack.
TALL TO LET	200 0.500	counter.	0540 A9 7F 0542 20 20 03	LDA \$7F JSR WRITE	
040B 10 F8	BPL LOOP	Loop until four	0545 A9 1D	LDA \$1D	Convert \$7F00
040D A5 00	LDA CMND	bytes are written. Get command	0547 20 40 03	JSR	= 32512 to
040D A3 00	LDA GMIND	byte from location		COMMAND	floating point
		\$0000.			form.
040F 20 40 03	JSR	Write command	054A A9 12	LDA \$12	Find 32512*
	COMMAND	to the AM9511.	054C 20 40 03	JSR	SIN[Y*(Pi/
0412 20 60 03	JSR READ	Get MSB of 16-		COMMAND	128)], result to NOS, pop
0.145.00.00		bit answer.			stack up.
0415 86 FF	STX MSB	Put most-signifi-	054F A9 1F	LDA \$1F	Convert that
0417 20 60 03	JSR READ	cant byte here. Get LSB of 16-	0551 20 40 03	JSR	number to
0117 20 00 03	Jon READ	bit answer.		COMMAND	fixed point
041A 86 FE	STX LSB	Put least-signifi-		****	format.
		cant byte in	0554 20 60 03	JSR READ TXA	Get MSB of
Contracts.	15 T. T.	\$00FE.	0557 8A	IAA	16-bit result in X register.
041C 00	BRK	End sample pro-	0558 99 00 0E	STA MSB,Y	Store it in a
		gram here.	3000 00 00 00	A - 1 - 1 - 1 - 1 - 1 - 1 - 1 - 1 - 1 -	table in page
Tisting ? Since table :					\$0E.
Listing 3. Sine table g	generator.		055B 20 60 03	JSR READ	Get LSB of 16-
			055E 8A	TXA	bit result.
0500 20 00 03 SINE	JSR RESET	Reset the	055F 99 00 0F	STA LSB,Y	Store it in a
	2.000.000.000	Am9511.			table in page \$0F.
0503 A9 1A	LDA \$1A	Push Pi	0562 C8	INY	Increment Y
0505 20 40 03	JSR	(3.14159) on		770.7	counter.
	COMMAND	TOS by writing	0563 D0 B9	BNE REPEAT	TRepeat until
		\$1A to Am9511.	036.00		table is filled.
0508 A9 80	LDA \$80	Load 128 =	0565 00	BRK	Break to the
050A 20 20 03	JSR WRITE	\$0080 on TOS,			monitor.
050D A9 00	LDA \$00	Pi is pushed			
050F 20 20 03	JSR WRITE	down to NOS.			
0512 A9 1D	LDA \$1D	Convert 128 =			
0514 20 40 03	JSR	\$0080 from			
	COMMAND	fixed point to to floating			
		point form.			
0517 A9 13	LDA \$13	Divide NOS by			
0519 20 40 03	JSR	TOS (Pi/128),			
	COMMAND	result onto			
0510 10 00	T DV 600	TOS.			
051C A0 00	LDY \$00	Y serves as			
		counter for 256 points.			
051E A9 37 REPEAT	LDA \$37	Duplicate NOS			
0520 20 40 03	JSR	with TOS.			
	COMMAND	Pi/128 is now			
		in TOS and			
0502 00	TVA	NOS.			
0523 98	TYA	Duplicate Y in			

0524 20 20 03

JSR WRITE

Duplicate Y in accumulator.

Push down TOS.

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