## Machine Learning Exercise 4

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#### 1. Give a proof of $VC(\mathcal{H}) = VC(\Phi)$

Proof:

Given the condition  $VC(\Phi) = d$ , we know  $\exists z_1, z_2, \dots, z_d$  such that

$$|\{(\phi(z_1), \cdots, \phi(z_d))|\phi \in \Phi\}| = 2^d$$

and there are no  $z_1, \dots, z_d, z_{d+1}$  such that

$$|\{(\phi(z_1),\cdots,\phi(z_{d+1})|\phi\in\Phi\}|=2^{d+1}$$

Define set  $\mathcal{S}_Z = \{(f(x_1), \dots, f(x_d)) | f \in \mathcal{H}\}$  and  $\mathcal{T}_Z = \{(\phi(z_1), \dots, \phi(z_d)) | \phi \in \Phi\}$  for a fixed sequence  $Z = (z_1, \dots, z_d), z_i = (x_i, y_i).$ 

Let  $g_Z(S) = (I(y_1 \neq s_1), \dots, I(y_d \neq s_d)) = (y_1, \dots, y_d) \oplus (s_1, \dots, s_d)$ ( $\oplus$  stands for XOR). Since for all binary sequence  $S_1, S_2, S_1 \neq S_2, g_Z(S_1) \neq$   $g_Z(S_2)$  holds true, and for all binary sequence  $T, \exists S_T$  such that  $g_Z(S_T) =$ T, function  $g_Z(\cdot) : \mathcal{S}_Z \to \mathcal{T}_Z$  is **bijection**, and two sets  $\mathcal{S}_Z$  and  $\mathcal{T}_Z$  are equipotent.

If  $VC(\Phi) = d$ , then  $\exists z_1, z_2, \cdots, z_d$  such that

$$|\{(f(x_1),\cdots,f(x_d))|f\in\mathcal{H}\}|=|\{(\phi(z_1),\cdots,\phi(z_d))|\phi\in\Phi\}|=2^d$$

and for all  $z_1, \dots, z_d, z_{d+1}$ 

$$|\{(f(x_1), \cdots, f(x_{d+1}))| f \in \mathcal{H}\}| = |\{(\phi(z_1), \cdots, \phi(z_{d+1})| \phi \in \Phi\}| \neq 2^{d+1}$$

Therefore  $VC(\mathcal{H}) = d$ , and vice versa.

#### 2. Prove the equivalence of the two optimization problems

$$\max_{w,b,t} t$$

$$s.t. \quad y_i(w^T x_i + b) \ge t, \forall i \in [n]$$

$$||w|| = 1$$

and

$$\min_{w,b} \quad \frac{1}{2} ||w||^2$$
s.t. 
$$y_i(w^T x_i + b) \ge 1, \forall i \in [n]$$

Proof:

The primal problem  $\Leftrightarrow$ 

$$\max_{w,b,t} \quad \frac{t}{||w||}$$

$$s.t. \quad y_i((\frac{w}{t})^T x_i + \frac{b}{t}) \ge 1, \forall i \in [n]$$

$$||w|| = 1$$

Let  $v = \frac{w}{t}$  and  $c = \frac{b}{t}$ , then the primal  $\Leftrightarrow$ 

$$\max_{v,c} \quad \frac{1}{||v||}$$

$$s.t. \quad y_i(v^T x_i + c) \ge 1, \forall i \in [n]$$

It is obvious that maximizing  $\frac{1}{||v||}$  is equivalent to minimizing  $\frac{1}{2}||v||^2,$  and the primal  $\Leftrightarrow$ 

$$\begin{aligned} & \min_{v,c} & \frac{1}{2} ||v||^2 \\ & s.t. & y_i(v^T x_i + c) \ge 1, \forall i \in [n] \end{aligned}$$

# 3. Give the dual problem of the latter one of Ex 2, namely the dual problem of

$$\min_{w,b} \quad \frac{1}{2} ||w||^2$$
s.t. 
$$y_i(w^T x_i + b) \ge 1, \forall i \in [n]$$

#### Solution:

Define the Langrangian function

$$\mathcal{L}(w, b, \alpha) = \frac{1}{2} ||w||^2 - \sum_{i=1}^{n} \alpha_i (y_i(w^T x_i + b) - 1), \forall i \in [n], \alpha_i \ge 0$$

Set the partial derivatives to zero, w.r.t.  $\boldsymbol{w}$  and  $\boldsymbol{b}$  respectively, and we get

$$\frac{\partial \mathcal{L}}{\partial w} = 0 \quad \Rightarrow \quad w = \sum_{i=1}^{n} \alpha_i y_i x_i$$
$$\frac{\partial \mathcal{L}}{\partial b} = 0 \quad \Rightarrow \quad \sum_{i=1}^{n} \alpha_i y_i = 0$$

Therefore

$$\mathcal{L}(w, b, \alpha) = \frac{1}{2} ||w||^2 - \sum_{i=1}^n \alpha_i (y_i (w^T x_i + b) - 1)$$

$$= \frac{1}{2} w^T w - \sum_{i=1}^n \alpha_i y_i w^T x_i - \sum_{i=1}^n \alpha_i y_i b + \sum_{i=1}^n \alpha_i$$

$$= \frac{1}{2} w^T \sum_{i=1}^n \alpha_i y_i x_i - w^T \sum_{i=1}^n \alpha_i y_i x_i - b \sum_{i=1}^n \alpha_i y_i + \sum_{i=1}^n$$

$$= -\frac{1}{2} w^T \sum_{i=1}^n \alpha_i y_i x_i + \sum_{i=1}^n \alpha_i$$

$$= -\frac{1}{2} (\sum_{i=1}^n \alpha_i y_i x_i)^T \sum_{i=1}^n \alpha_i y_i x_i + \sum_{i=1}^n \alpha_i$$

$$= \sum_{i=1}^n \alpha_i - \frac{1}{2} \sum_{i=1}^n \alpha_i \alpha_j y_i y_j x_i^T x_j$$

The dual problem can be described as follow

$$\begin{aligned} \max_{\alpha} & \sum_{i=1}^{n} -\frac{1}{2} \sum_{i,j=1}^{n} \alpha_{i} \alpha_{j} y_{i} y_{j} x_{i}^{T} x_{j} \\ s.t. & \alpha_{i} \geq 0, \forall i \in [n] \\ & \sum_{i=1}^{n} \alpha_{i} y_{i} = 0 \end{aligned}$$