

**Definition 1.** A **uniform space**  $\langle X, \mathcal{D} \rangle$  is a set  $X$  paired with a filter  $\mathcal{D}$  (called its **uniformity**) of relations (called **entourages**) on  $X$  such that for each entourage  $D \in \mathcal{D}$ :

- $D$  is reflexive, i.e., the diagonal  $\Delta \subseteq D$ .
- Its inverse  $D^{-1} = \{\langle y, x \rangle : \langle x, y \rangle \in D\} \in \mathcal{D}$ .
- There exists  $\frac{1}{2}D \in \mathcal{D}$  such that

$$2(\frac{1}{2}D) = \frac{1}{2}D \circ \frac{1}{2}D = \{\langle x, z \rangle : \exists y(\langle x, y \rangle, \langle y, z \rangle \in \frac{1}{2}D)\} \subseteq D$$

Note that since  $\mathcal{D}$  is a filter, for each  $D \in \mathcal{D}$ , the symmetric relation  $D \cap D^{-1} \in \mathcal{D}$ .

**Proposition 2.** For each  $D \in \mathcal{D}$  and  $n < \omega$  there exists  $\frac{1}{2^{n+1}}D \in \mathcal{D}$  such that

$$2(\frac{1}{2^{n+1}}D) = \frac{1}{2^{n+1}}D \circ \frac{1}{2^{n+1}}D \subseteq \frac{1}{2^n}D$$

and if  $2E \subseteq \frac{1}{2^n}D$ , then  $E \subseteq \frac{1}{2^{n+1}}D$ .

**Definition 3.** For an entourage  $D \in \mathcal{D}$ , let  $D[x] = \{y : \langle x, y \rangle \in D\}$  be the  $D$ -**neighborhood** of  $x$ . The uniform topology for a uniform space  $\langle X, \mathcal{D} \rangle$  is generated by the base  $\{D[x] : x \in X, D \in \mathcal{D}\}$ .

**Theorem 4.** A space  $X$  is uniformizable (its topology is the uniform topology for some uniformity) if and only if  $X$  is completely regular ( $T_{3\frac{1}{2}}$ ).

**Proposition 5.** If  $X$  is a uniform space, then for all  $x \in X$  and symmetric entourages  $D$ :

$$x \in \frac{1}{2}D[y] \text{ and } y \in \frac{1}{2}D[z] \Rightarrow x \in D[z]$$

and

$$\frac{1}{2}D[x] \subseteq \overline{\frac{1}{2}D[x]} \subseteq D[x]$$

*Proof.* The first is by definition of  $\frac{1}{2}D$ .

If  $z \in \overline{\frac{1}{2}D[x]}$ , it follows that there is  $y \in \frac{1}{2}D[x] \cap \frac{1}{2}D[z]$  since  $\frac{1}{2}D[z]$  is an open neighborhood of  $z$ . Thus  $(x, z) \in D \Rightarrow z \in D[x] \Rightarrow \overline{\frac{1}{2}D[x]} \subseteq D[x]$ .  $\square$

**Definition 6.** For a uniform space  $X$ , Bell's proximity game proceeds as follows.

In round 0,  $\mathcal{D}$  chooses an entourage  $D_0$ , followed by  $\mathcal{P}$  choosing a point  $p_0 \in X$ .

In round  $n + 1$ ,  $\mathcal{D}$  chooses an entourage  $D_{n+1} \subseteq D_n$ , followed by  $\mathcal{P}$  choosing a point  $p_{n+1} \in 4D_n[p_n]$ .

Player  $\mathcal{D}$  wins if either  $\bigcap_{n < \omega} 4D_n[p_n] = \emptyset$  or  $\langle p_0, p_1, \dots \rangle$  converges.

**Definition 7.** For a uniform space  $X$ , the simplified proximal game  $Prox_{D,P}(X)$  can be defined as follows:

In round 0,  $\mathcal{D}$  chooses a symmetric entourage  $D_0$ , followed by  $\mathcal{P}$  choosing a point  $p_0 \in X$ .

In round  $n+1$ ,  $\mathcal{D}$  chooses a symmetric entourage  $D_{n+1}$ , followed by  $\mathcal{P}$  choosing a point  $p_{n+1} \in \left(\bigcap_{m \leq n} D_m\right)[p_n]$ .

Player  $\mathcal{D}$  wins if either  $\bigcap_{n < \omega} \left(\bigcap_{m \leq n} D_m\right)[p_n] = \emptyset$  or  $\langle p_0, p_1, \dots \rangle$  converges.

**Theorem 8.**  $\mathcal{D}$  has a winning perfect-information strategy in Bell's game if and only if  $\mathcal{D} \uparrow Prox_{D,P}(X)$ .

*Proof.* Let  $\sigma$  be a winning perfect information strategy for  $\mathcal{D}$  in Bell's game. We define a perfect information strategy  $\tau$  in the simplified game to yield symmetric entourages  $\tau(p \upharpoonright n) = \sigma(p \upharpoonright n) \cap (\sigma(p \upharpoonright n))^{-1}$  for all partial attacks  $p \upharpoonright n$ . Note that  $\tau(p \upharpoonright n) = \bigcap_{m \leq n} \tau(p \upharpoonright m)$ .

If  $p$  attacks  $\tau$  in the simplified game,  $p(n+1) \in \left(\bigcap_{m \leq n} \tau(p \upharpoonright m)\right)[p(n)] = \tau(p \upharpoonright n)[p(n)] \subseteq \sigma(p \upharpoonright n)[p(n)] \subseteq 4\sigma(p \upharpoonright n)[p(n)]$ , so  $p$  attacks  $\sigma$  in Bell's game. Thus either  $p$  converges, or

$$\emptyset = \bigcap_{n < \omega} 4\sigma(p \upharpoonright n)[p(n)] \supseteq \bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)] = \bigcap_{n < \omega} \left( \bigcap_{m \leq n} \tau(p \upharpoonright m) \right)[p(n)]$$

For the other direction, let  $\sigma$  be a winning perfect information strategy for  $\mathcal{D}$  in the simplified game such that  $\sigma(p \upharpoonright n) = \bigcap_{m \leq n} \sigma(p \upharpoonright m)$ . Define the perfect information strategy  $\tau$  in Bell's Game such that  $4\tau(p \upharpoonright n) \subseteq \sigma(p \upharpoonright n)$  and  $\tau(p \upharpoonright n) = \bigcap_{m \leq n} \tau(p \upharpoonright m)$  for all partial attacks  $p \upharpoonright n$ .

If  $p$  attacks  $\tau$  in Bell's game,  $p(n) \in 4\tau(p \upharpoonright n) \subseteq \sigma(p \upharpoonright n) = \bigcap_{m \leq n} \sigma(p \upharpoonright m)$ , so  $p$  attacks  $\sigma$  in the simplified game. Thus either  $p$  converges, or

$$\emptyset = \bigcap_{n < \omega} \left( \bigcap_{m \leq n} \sigma(p \upharpoonright m) \right)[p(n)] = \bigcap_{n < \omega} \sigma(p \upharpoonright n)[p(n)] \supseteq \bigcap_{n < \omega} 4\tau(p \upharpoonright n)[p(n)] \supseteq \bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)]$$

□

**Proposition 9.**  $\mathcal{P}$  has a winning perfect-information strategy in Bell's game if and only if  $\mathcal{P} \uparrow Prox_{D,P}(X)$ .

*Proof.* Similar to the previous.

□

**Definition 10.** A uniform space is **proximal** if  $\mathcal{D} \uparrow \text{Prox}_{D,P}(X)$ .

**Definition 11.** For a space  $X$  and a point  $x \in X$ , the  **$W$ -convergence-game**  $\text{Con}_{O,P}(X, x)$  proceeds as follows.

In round 0,  $\mathcal{O}$  chooses a neighborhood  $U_n$  of  $x$ , followed by  $\mathcal{P}$  choosing a point  $p_n \in \bigcap_{m \leq n} U_m$ .

Player  $\mathcal{O}$  wins if  $\langle p_0, p_1, \dots \rangle$  converges.

**Definition 12.** A space is  **$W$**  if  $\mathcal{O} \uparrow \text{Con}_{O,P}(X, x)$  for all  $x \in X$ .

**Definition 13.** For each finite tuple  $(m_0, \dots, m_{n-1})$ , we define the  **$k$ -tactical fog-of-war**

$$T_k(\langle m_0, \dots, m_{n-1} \rangle) = \langle m_{n-k}, \dots, m_{n-1} \rangle$$

and the  **$k$ -Marköv fog-of-war**

$$M_k(\langle m_0, \dots, m_{n-1} \rangle) = \langle \langle m_{n-k}, \dots, m_{n-1} \rangle, n \rangle$$

So  $P \uparrow G$  if and only if there exists a winning strategy for  $P$  of the form  $\sigma \circ T_k$ , and  $P \uparrow_{k\text{-tact}} G$  if and only if there exists a winning strategy of the form  $\sigma \circ M_k$ .

**Theorem 14.** For all  $x \in X$ :

- $\mathcal{D} \uparrow \text{Prox}_{D,P}(X) \Rightarrow \mathcal{O} \uparrow \text{Con}_{O,P}(X, x)$
- $\mathcal{D} \uparrow_{2k\text{-tact}} \text{Prox}_{D,P}(X) \Rightarrow \mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X, x)$
- $\mathcal{D} \uparrow_{2k\text{-mark}} \text{Prox}_{D,P}(X) \Rightarrow \mathcal{O} \uparrow_{k\text{-mark}} \text{Con}_{O,P}(X, x)$

*Proof.* Let  $\sigma$  witness  $\mathcal{D} \uparrow_{2k\text{-tact}} \text{Prox}_{D,P}(X)$  (resp.  $\mathcal{D} \uparrow_{2k\text{-mark}} \text{Prox}_{D,P}(X)$ ,  $\mathcal{D} \uparrow \text{Prox}_{D,P}(X)$ ).

We define the  $k$ -tactical (resp.  $k$ -Marköv, perfect info) strategy  $\tau$  such that

$$\tau \circ L_k(p) = \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(|p| - 1) \rangle)[x] \cap \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(|p| - 1), x \rangle)[x]$$

where  $L_{2k}$  is the  $2k$ -tactical fog-of-war (resp.  $2k$ -Marköv fog-of-war, identity) and  $L_k$  is the  $k$ -tactical fog-of-war (resp.  $k$ -Marköv fog-of-war, identity).

Let  $p$  attack  $\tau$ . Consider the attack  $q$  against the winning strategy  $\sigma$  such that  $q(2n) = x$  and  $q(2n + 1) = p(n)$ , and let  $D_n = \sigma \circ L_{2k}(q)$  and  $E_n = \bigcap_{m \leq n} D_m$ .

Certainly,  $x \in E_{2n}[x] = E_{2n}[q(2n)]$  for any  $n < \omega$ . Note also for any  $n < \omega$  that

$$p(n) \in \bigcap_{m \leq n} \tau \circ L_k(p \upharpoonright m)$$

$$\begin{aligned}
 &= \bigcap_{m \leq n} (\sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(m-1) \rangle)[x] \cap \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(m-1), x \rangle)[x]) \\
 &= \bigcap_{m \leq n} (D_{2m}[x] \cap D_{2m+1}[x]) = \bigcap_{m \leq 2n+1} D_m[x] = E_{2n+1}[x]
 \end{aligned}$$

so by the symmetry of  $E_{2n+1}$ ,  $x \in E_{2n+1}[p(n)] = E_{2n+1}[q(2n+1)]$ . Thus  $x \in \bigcap_{n < \omega} E_n[q(n)] \neq \emptyset$ , and since  $\sigma$  is a winning strategy, the attack  $q$  converges. Since  $q(2n) = x$ ,  $q$  must converge to  $x$ . Thus its subsequence  $p$  converges to  $x$ , and  $\tau$  is a winning strategy in  $Con_{O,P}(X, x)$ .  $\square$

**Corollary 15.** *For all  $x \in X$ :*

- $\mathcal{D} \uparrow_{k\text{-tact}} Prox_{D,P}(X) \Rightarrow \mathcal{O} \uparrow_{k\text{-tact}} Con_{O,P}(X, x)$
- $\mathcal{D} \uparrow_{k\text{-mark}} Prox_{D,P}(X) \Rightarrow \mathcal{O} \uparrow_{k\text{-mark}} Con_{O,P}(X, x)$

**Corollary 16.** *All proximal spaces are  $W$ -spaces.*

**Theorem 17.** *Let  $X \cup \{\infty\}$  be a uniformizable space such that  $X$  is discrete. Then*

- $\mathcal{O} \uparrow Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathcal{D} \uparrow Prox_{D,P}(X \cup \{\infty\})$
- $\mathcal{O} \uparrow_{k\text{-tact}} Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathcal{D} \uparrow_{k\text{-tact}} Prox_{D,P}(X \cup \{\infty\})$
- $\mathcal{O} \uparrow_{k\text{-mark}} Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathcal{D} \uparrow_{k\text{-mark}} Prox_{D,P}(X \cup \{\infty\})$

*Proof.* Note that the topology on  $X \cup \{\infty\}$  is induced by the uniformity with equivalence relation entourage  $D(U) = \Delta \cup U^2$  for each open neighborhood  $U$  of  $\infty$ .

Let  $\sigma$  witness  $\mathcal{D} \uparrow_{k\text{-tact}} Con_{O,P}(X \cap \{\infty\}, \infty)$  (resp.  $\mathcal{D} \uparrow_{k\text{-mark}} Con_{O,P}(X \cap \{\infty\}, \infty)$ ,  $\mathcal{D} \uparrow Con_{O,P}(X \cap \{\infty\}, \infty)$ ). We define the  $k$ -tactical (resp.  $k$ -Marköv, perfect info) strategy  $\tau$  such that

$$\tau \circ L(p) = D(\sigma \circ L(p))$$

where  $L$  is the  $k$ -tactical fog-of-war (resp.  $k$ -Marköv fog-of-war, identity).

Let  $p \in (X \cup \{\infty\})^\omega$  attack  $\tau$  such that  $\bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)] \neq \emptyset$ .

If  $\infty \in \bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)]$ , it follows that  $p$  is an attack on  $\sigma$ . Since  $\sigma$  is a winning strategy, it follows that  $q$  and its subsequence  $p$  must converge to  $\infty$ .

Otherwise,  $\infty \notin \tau(p \upharpoonright N)[p(N)]$  for some  $N < \omega$ , and then  $\tau(p \upharpoonright N)[p(N)] = \{p(N)\}$  implies  $p \rightarrow p(N)$ .

Thus  $\tau \circ L$  is a winning strategy.  $\square$

**Corollary 18.** *Let  $X \cup \{\infty\}$  be a uniformizable space such that  $X$  is discrete. Then*

- $\mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathcal{D} \uparrow_{k\text{-tact}} \text{Prox}_{D,P}(X \cup \{\infty\})$
- $\mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathcal{D} \uparrow_{k\text{-tact}} \text{Prox}_{D,P}(X \cup \{\infty\})$
- $\mathcal{O} \uparrow_{k\text{-mark}} \text{Con}_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathcal{D} \uparrow_{k\text{-mark}} \text{Prox}_{D,P}(X \cup \{\infty\})$

**Proposition 19.** *For any  $x \in X$  and  $k \geq 1$ ,*

- $\mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X, x) \Leftrightarrow \mathcal{O} \uparrow_{\text{tact}} \text{Con}_{O,P}(X, x)$
- $\mathcal{O} \uparrow_{k\text{-mark}} \text{Con}_{O,P}(X, x) \Leftrightarrow \mathcal{O} \uparrow_{\text{mark}} \text{Con}_{O,P}(X, x)$

*Proof.* If  $\sigma$  witnesses  $\mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X, x)$ , let  $\tau(\emptyset) = \sigma(\emptyset)$  and

$$\tau(\langle q \rangle) = \bigcap_{i < k} \sigma(\langle \underbrace{x, \dots, x}_{k-i-1}, \underbrace{q, x, \dots, x}_i \rangle)$$

This is easily verified to be a winning strategy. The proof for  $\mathcal{O} \uparrow_{k\text{-mark}} \text{Con}_{O,P}(X, x)$  is analogous.  $\square$

**Corollary 20.** *Let  $X \cup \{\infty\}$  be a uniformizable space such that  $X$  is discrete, and  $k \geq 1$ . Then*

- $\mathcal{D} \uparrow_{k\text{-tact}} \text{Prox}_{D,P}(X \cup \{\infty\}) \Leftrightarrow \mathcal{O} \uparrow_{\text{tact}} \text{Prox}_{D,P}(X \cup \{\infty\})$
- $\mathcal{D} \uparrow_{k\text{-mark}} \text{Prox}_{D,P}(X \cup \{\infty\}) \Leftrightarrow \mathcal{O} \uparrow_{\text{mark}} \text{Prox}_{D,P}(X \cup \{\infty\})$

**Proposition 21.** *For any uniform space  $X$ ,*

- $\mathcal{O} \uparrow_{k\text{-tact}} \text{Prox}_{D,P}(X) \Leftrightarrow \mathcal{O} \uparrow_{2\text{-tact}} \text{Prox}_{D,P}(X)$
- $\mathcal{O} \uparrow_{k\text{-mark}} \text{Prox}_{D,P}(X) \Leftrightarrow \mathcal{O} \uparrow_{2\text{-mark}} \text{Prox}_{D,P}(X)$

*Proof.* If  $\sigma$  witnesses  $\mathcal{O} \uparrow_{k\text{-tact}} \text{Con}_{O,P}(X, x)$ , let  $\tau(\emptyset) = \sigma(\emptyset)$  and

$$\tau(\langle q \rangle) = \bigcap_{i < k} \sigma(\langle \underbrace{q, \dots, q}_i \rangle)$$

$$\tau(\langle q, q' \rangle) = \bigcap_{i < k} \sigma(\underbrace{\langle q, \dots, q \rangle}_{k-i}, \underbrace{\langle q', \dots, q' \rangle}_i)$$

This is easily verified to be a winning strategy. The proof for  $\mathcal{O} \uparrow_{k\text{-mark}} \text{Con}_{O,P}(X, x)$  is analogous. □

**Definition 22.** The absolute proximal game  $aProx_{D,P}(X)$  is analogous to  $Prox_{D,P}(X)$ , except  $\mathcal{D}$  may only win if  $p$  converges.

**Definition 23.** A **uniformly locally compact** space is a uniformizable space with a **uniformly compact entourage**  $M$  where  $\overline{M[x]}$  is compact for all  $x$ .

**Theorem 24.** For any uniformly locally compact space  $X$ ,  $\mathcal{D} \uparrow Prox_{D,P}(X) \Leftrightarrow \mathcal{D} \uparrow aProx_{D,P}(X)$

*Proof.* Let  $M$  be a uniformly locally compact entourage. Let  $\sigma$  witness  $\mathcal{D} \uparrow Prox_{D,P}(X)$  such that  $\sigma(a) \subseteq M$  always (so  $\overline{\sigma(a)[x]} \subseteq \overline{M[x]}$  is compact), and  $a \supseteq b$  implies  $\sigma(a) \subseteq \frac{1}{4}\sigma(b)$ .

Let  $\tau(p \upharpoonright n) = \frac{1}{2}\sigma(p \upharpoonright n)$ . If  $p$  attacks  $\tau$  in  $aProx_{D,P}(X)$ , then

$$p(n+1) \in \tau(p \upharpoonright n)[p(n)] = \frac{1}{2}\sigma(p \upharpoonright n)[p(n)]$$

and for

$$x \in \overline{\sigma(p \upharpoonright (n+1))[p(n+1)]} \subseteq \overline{\frac{1}{4}\sigma(p \upharpoonright n)[p(n+1)]} \subseteq \frac{1}{2}\sigma(p \upharpoonright n)[p(n+1)]$$

we can conclude  $x \in \sigma(p \upharpoonright n)[p(n)]$ . Thus

$$\sigma(p \upharpoonright (n+1))[p(n+1)] \subseteq \overline{\sigma(p \upharpoonright (n+1))[p(n+1)]} \subseteq \sigma(p \upharpoonright n)[p(n)]$$

Finally, note that  $p$  attacks the winning strategy  $\sigma$  in  $Prox_{D,P}(X)$ , but since the intersection of a chain of nonempty compact sets is nonempty:

$$\bigcap_{n < \omega} \sigma(p \upharpoonright n)[p(n)] = \bigcap_{n < \omega} \overline{\sigma(p \upharpoonright n)[p(n)]} \neq \emptyset$$

We conclude that  $p$  converges. □

**Corollary 25.** A uniformly locally compact space  $X$  is proximal if and only if  $\mathcal{D} \uparrow aProx_{D,P}(X)$ .

**Theorem 26.** For any uniformly locally compact proximal space  $X$ ,  $\mathcal{O} \uparrow Clus_{O,P}(X, H)$  for all compact  $H \subseteq X$ .

*Proof.* Let  $\sigma$  witness  $\mathcal{D} \uparrow aProx_{D,P}(X)$  such that  $p \supseteq q$  implies  $\sigma(p) \subseteq \frac{1}{4}\sigma(q)$ .

Let  $o(t)$  be the subsequence of  $t$  consisting of its odd-indexed terms.

We define  $T(\emptyset)$ , etc. as follows:

- Let  $\emptyset \in T(\emptyset)$ .
- Choose  $m_\emptyset < \omega$ ,  $h_{\emptyset,i} \in H$  for  $i < m_\emptyset$ , and  $h_{\emptyset,i,j} \in H \cap \overline{\frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}]}$  for  $i, j < m_\emptyset$  such that

$$\{\frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}] : i < m_\emptyset\}$$

is a cover for  $H$  and such that for each  $i < m_\emptyset$

$$\{\frac{1}{4}\sigma(\langle h_{\emptyset,i} \rangle)[h_{\emptyset,i,j}] : j < m_\emptyset\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}]}$ .

- Let  $\langle i \rangle \in T(\emptyset)$ ,  $\langle i, h_{\emptyset,i} \rangle \in T(\emptyset)$ , and  $\langle i, h_{\emptyset,i}, j \rangle \in T(\emptyset)$  for  $i, j < m_\emptyset$ .

Suppose  $T(a)$ , etc. are defined. We then define  $T(a \smallfrown \langle x \rangle)$ , etc. for

$$x \in \bigcup_{s \smallfrown \langle i, h_{s,i}, j \rangle \in \max(T(a))} \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$$

as follows:

- Let  $T(a) \subseteq T(a \smallfrown \langle x \rangle)$ .
- Choose  $t = s \smallfrown \langle i, h_{s,i}, j, x \rangle$  such that  $s \smallfrown \langle i, h_{s,i}, j \rangle \in \max(T(a))$  and  $x \in \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$ .
- Note that, assuming  $o(s) \smallfrown \langle h_{s,i} \rangle$  is a legal partial attack against  $\sigma$ , then

$$x \in \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}] \subseteq \frac{1}{4}\sigma(o(s))[h_{s,i,j}]$$

and

$$h_{s,i,j} \in \overline{\frac{1}{4}\sigma(o(s))[h_{s,i}]} \subseteq \frac{1}{2}\sigma(o(s))[h_{s,i}]$$

implies

$$x \in \sigma(o(s))[h_{s,i}]$$

and thus  $o(s) \smallfrown \langle h_{s,i}, x \rangle = o(t)$  is a legal partial attack against  $\sigma$ .

- Choose  $m_t < \omega$ ,  $h_{t,k} \in H \cap \overline{\frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]}$  for  $k < m_t$ , and  $h_{t,k,l} \in H \cap \overline{\frac{1}{4}\sigma(t)[h_{t,k}]}$  for  $k, l < m_t$  such that

$$\{\frac{1}{4}\sigma(o(t))[h_{t,k}] : k < m_t\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]}$  and such that for each  $k < m_t$

$$\{\frac{1}{4}\sigma(o(t) \smallfrown \langle h_{t,k} \rangle)[h_{t,i,j}] : l < m_t\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(o(t))[h_{t,k}]}$ .



- Note that, assuming  $o(t)$  is a legal partial attack against  $\sigma$ , then

$$h_{t,k} \in \overline{\frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]} \subseteq \frac{1}{2}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$$

and

$$x \in \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$$

implies

$$h_{t,k} \in \sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[x]$$

and thus  $o(t) \smallfrown \langle h_{t,k} \rangle$  is a legal partial attack against  $\sigma$ .

- Let  $t \in T(a \smallfrown \langle x \rangle)$ ,  $t \smallfrown \langle k \rangle \in T(a \smallfrown \langle x \rangle)$ ,  $t \smallfrown \langle k, h_{t,k} \rangle \in T(a \smallfrown \langle x \rangle)$ , and  $t \smallfrown \langle k, h_{t,k}, l \rangle \in T(a \smallfrown \langle x \rangle)$  for  $k, l < m_t$ .
- Note that assuming

$$\{\frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}] : s \smallfrown \langle i, h_{s,i}, j \rangle \in \max(T(a))\}$$

covers  $H$ , then since

$$\{\frac{1}{4}\sigma(o(t) \smallfrown \langle h_{t,k} \rangle)[h_{t,k,l}] : s \smallfrown \langle i, h_{s,i}, j, x, k, h_{t,k}, l \rangle \in \max(T(a \smallfrown \langle x \rangle)) \setminus \max(T(a))\}$$

covers  $H \cap \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$ , we have that

$$\{\frac{1}{4}\sigma(o(t) \smallfrown \langle h_{t,k} \rangle)[h_{t,k,l}] : t \smallfrown \langle k, h_{t,k}, l \rangle \in \max(T(a \smallfrown \langle x \rangle))\}$$

covers  $H$ .

With this we may define the perfect information strategy  $\tau$  for  $\mathcal{O}$  in  $Con_{O,P}(X, H)$  such that:

$$\tau(p \upharpoonright n) = \bigcup_{s \smallfrown \langle i, h_{s,i}, j \rangle \in \max(T(p \upharpoonright n))} \frac{1}{4}\sigma(o(s) \smallfrown \langle h_{s,i} \rangle)[h_{s,i,j}]$$

If  $p$  attacks  $\tau$ , then it follows that  $T(p \upharpoonright n)$  is defined for all  $n < \omega$ , so let  $T(p) = \bigcup_{n < \omega} T(p \upharpoonright n)$ . We note  $T(p)$  is an infinite tree with finite levels:

- $\emptyset$  has exactly  $m_\emptyset$  successors  $\langle i \rangle$ .
- $s \smallfrown \langle i \rangle$  has exactly one successor  $t \smallfrown \langle i, h_{s,i} \rangle$
- $s \smallfrown \langle i, h_{s,i} \rangle$  has exactly  $m_s$  successors  $t \smallfrown \langle i, h_{s,i}, j \rangle$
- $s \smallfrown \langle i, h_{s,i}, j \rangle$  has either no successors or exactly one successor  $t \smallfrown \langle i, h_{s,i}, j, x \rangle$

- $t = s^\frown \langle i, h_{s,i}, j, x \rangle$  has exactly  $m_t$  successors  $t^\frown \langle k \rangle$

Let  $q' = \langle i_0, h_0, j_0, x_0, i_1, h_1, j_1, x_1, \dots \rangle$  correspond to this infinite branch in  $T(p)$ , and let  $q = o(q') = \langle h_0, x_0, h_1, x_1, \dots \rangle$ . Note that by the construction of  $T(p)$ ,  $q$  is an attack on the winning strategy  $\sigma$  in  $aProx_{D,P}(X)$ , so it must converge. Since every other term of  $q$  is in  $H$ , it must converge to  $H$ . Then since  $q$  is a subsequence of  $p$ ,  $p$  must cluster at  $H$ .  $\square$

**Corollary 27.** *For any uniformly locally compact proximal space,  $\mathcal{O} \uparrow Con_{O,P}(X, H)$  for all compact  $H \subseteq X$ .*

*Proof.*  $\mathcal{O} \uparrow Con_{O,P}(X, H)$  if and only if  $\mathcal{O} \uparrow Clus_{O,P}(X, H)$ .  $\square$

**Corollary 28.** *A compact uniform space  $X$  is Corson compact if and only if it is proximal.*

*Proof.* A characterization of Corson compact is having a  $W$ -set diagonal. If  $X$  is proximal compact, then  $X^2$  is proximal compact, and its compact diagonal is a  $W$ -set.  $\square$

**Theorem 29.**  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, H)$  if and only if there exists a countable base around  $H$ .

*Proof.* Let  $\{U_n : n < \omega\}$  be a countable base around  $H$ . We define the predetermined strategy  $\sigma(n) = \bigcap_{m \leq n} U_m$ . Let  $p$  attack  $\sigma(n)$  - then if  $U$  is any neighborhood of  $H$ , we may choose  $H \subseteq U_m \subseteq U$ , and note that  $\sigma(n) \subseteq U_m$  for  $n \geq m$ , and thus  $p(n) \in U_m \subseteq U$  for all  $n \geq m$ . Thus  $\sigma$  is a winning strategy.

For the other direction, suppose there does not exist a countable base around  $H$ , and let  $\sigma(n)$  be an arbitrary predetermined strategy. Since  $\{\bigcap_{m \leq n} \sigma(m) : n < \omega\}$  is not a countable base around  $H$ , we may choose an open set  $U$  around  $H$  such that  $\bigcap_{m \leq n} \sigma(m) \not\subseteq U$  for all  $n < \omega$ . We may easily verify that if  $p(n) \in \bigcap_{m \leq n} \sigma(m) \setminus U$  for all  $n < \omega$ , then  $p$  is a successful counterattack to  $\sigma$ .  $\square$

**Corollary 30.**  $X$  is first countable if and only if  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, x)$  for all  $x \in X$

**Corollary 31.**  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$  implies  $X$  is first countable.

**Definition 32.** Scattered Eberlein compact spaces are known as **strong Eberlein compact** spaces.

**Theorem 33** (folklore). *Scattered compact first-countable spaces are metrizable.*

**Corollary 34.** If  $X$  is scattered compact and  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, x)$  for all  $x \in X$  (or  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$ ), then  $X$  is metrizable.

**Example 35.**  $\mathcal{D} \nuparrow_{pre} Prox_{D,P}(\omega_1^*)$

*Proof.* There does not exist a countable base around  $\infty$ , so  $\mathcal{O} \nuparrow_{pre} Con_{O,P}(X, \omega_1)$ .  $\square$

**Example 36.**  $\mathcal{O} \uparrow_{tact} Con_{O,P}(\kappa^*, \infty)$  and  $\mathcal{D} \uparrow_{tact} Prox_{D,P}(\kappa^*)$  for all cardinals  $\kappa$

*Proof.* For  $Con_{O,P}(\kappa^*, \infty)$ , let  $\sigma() = \sigma(\infty) = \kappa^*$  and  $\sigma(x) = \kappa^* \setminus \{x\}$  otherwise.  $\square$

**Theorem 37.** If  $H$  is a closed subset of  $X$ , then  $\mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(X) \Rightarrow \mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(H)$  where  $\uparrow_{\text{limit}}$  is any of  $\uparrow$ ,  $\uparrow_{k\text{-tact}}$ , or  $\uparrow_{k\text{-mark}}$ .

*Proof.* Let  $\sigma \circ L$  witness  $\mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(X)$ . We define  $\tau \circ L$  for  $\mathcal{D}$  in  $\text{Prox}_{D,P}(H)$  as follows:

$$\tau \circ L(p \upharpoonright n) = \sigma \circ L(p \upharpoonright n) \cap H^2$$

Let  $p$  attack  $\tau \circ L$ .  $p$  also attacks the winning strategy  $\sigma \circ L$ , so either

$$\bigcap_{n < \omega} \left( \bigcap_{m \leq n} \tau \circ L(p \upharpoonright m) \right) [p(n)] \subseteq \bigcap_{n < \omega} \left( \bigcap_{m \leq n} \sigma \circ L(p \upharpoonright m) \right) [p(n)] = \emptyset$$

or  $p$  converges in  $X$ , and thus converges in  $H$ .  $\square$

**Theorem 38.** If  $\mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(X_i)$  for  $i < \omega$ , then  $\mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(\prod_{i < \omega} X_i)$ , where  $\uparrow_{\text{limit}}$  is either  $\uparrow$  or  $\uparrow_{k\text{-mark}}$ .

*Proof.* A subbase for  $\prod_{i < \omega} X_i$  is

$$\{\pi_i^{-1}(D) : i < \omega, D \in \mathcal{D}_i\}$$

where  $\pi_i$  is the natural projection from  $(\prod_{i < \omega} X_i)^2$  onto  $X_i^2$ . (See Bell.)

For  $p \in (\prod_{i < \omega} X_i)^\omega$ , let  $p_i \in X_i^\omega$  such that  $p_i(n) = p(n)(i)$ .

Let  $\sigma_i \circ L$  witness  $\mathcal{D} \uparrow_{\text{limit}} \text{Prox}_{D,P}(X_i)$  for  $i < \omega$ , and assume without loss of generality that  $\sigma_i \circ L$  always yields  $X_i^2$  before round  $i$ .

Then we define the strategy  $\tau \circ L$  for  $\mathcal{D}$  in  $\text{Prox}_{D,P}(\prod_{i < \omega} X_i)$  as follows:

$$\tau \circ L(p \upharpoonright n) = \bigcap_{i \leq n} \pi_i^{-1}(\sigma_i \circ L(p_i \upharpoonright n))$$

Let  $p$  attack  $\tau \circ L$ . If  $\bigcap_{n < \omega} \left( \bigcap_{m \leq n} \sigma_i(p_i \upharpoonright m) \right) [p_i(n)] = \emptyset$  for any  $i < \omega$ , it easily follows that  $\bigcap_{n < \omega} \left( \bigcap_{m \leq n} \tau(p \upharpoonright m) \right) [p(n)] = \emptyset$ .

Otherwise, we assume that for each  $i < \omega$ ,  $p_i$  converges to some  $x_i \in X_i$ . Thus  $p$  converges to  $x = \langle x_0, x_1, \dots \rangle$ .  $\square$

Note: I expect I should be able to do some clever things assuming  $S(\kappa, \omega, \omega)$  to get a similar result for sigma products of dimension  $\kappa$ .

**Example 39.**  $\mathcal{D} \uparrow_{\text{mark}} \text{Prox}_{D,P}((\kappa^*)^\omega)$

*Proof.*  $\mathcal{D} \uparrow_{\text{tact}} \text{Prox}_{D,P}(\kappa^*) + \text{previous result}$

□

**Lemma 40.**  $\mathcal{O} \uparrow_{pre} Clus_{O,P}(X, S)$  if and only if  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, S)$ .

*Proof.* Suppose that  $\sigma$  is a predetermined winning strategy for  $Clus_{O,P}(X, S)$ . Let  $p$  attack  $\sigma$ , and  $q$  be a subsequence of  $p$ . It follows that  $q$  also attacks  $\sigma$ , so  $q$  clusters at  $S$ . Thus  $p$  conveys to  $S$ , and  $\sigma$  is a predetermined winning strategy for  $Con_{O,P}(X, S)$ .  $\square$

**Theorem 41.** For any predetermined absolutely proximal space  $X$ ,  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, H)$  for all compact  $H \subseteq X$ .

*Proof.* Let  $\sigma(n)$  be a winning predetermined strategy for  $\mathcal{D}$  in the absolutely proximal game such that  $\sigma(n+1) \subseteq \sigma(n)$ . For a given tree  $T$ , let  $\max(T)$  denote its maximal nodes.

First we define  $T(0) \subseteq \omega^{\leq 2}$ .

- Let  $\emptyset \in T(0)$ .
- Choose

$$m_\emptyset < \omega$$

and for  $i < m_\emptyset$  choose

$$h_{\langle i \rangle} \in H$$

and for  $i, j < m_\emptyset$  choose

$$h_{\langle i, j \rangle} \in H \cap \overline{\frac{1}{4}\sigma(0)[h_{\langle i \rangle}]}$$

such that

$$\left\{ \frac{1}{4}\sigma(0)[h_{\langle i \rangle}] : i < m_\emptyset \right\}$$

is a cover for  $H$  and such that for each  $i < m_\emptyset$

$$\left\{ \frac{1}{4}\sigma(1)[h_{\langle i, j \rangle}] : j < m_\emptyset \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(0)[h_{\langle i \rangle}]}$ .

- Let  $\langle i \rangle$  and  $\langle i, j \rangle$  be in  $T(0)$  for  $i, j < m_\emptyset$ .

Now suppose  $T(n) \subseteq \omega^{\leq 2n+2}$  is defined. We then define  $T(n+1) \subseteq \omega^{\leq 2n+4}$  as follows:

- Let  $T(n) \subseteq T(n+1)$ .

- For each  $t \smallfrown \langle i, j \rangle \in \max(T(n))$ , choose

$$m_{t \smallfrown \langle i, j \rangle} < \omega$$

and for  $k < m_{t \smallfrown \langle i, j \rangle}$  choose

$$h_{t \smallfrown \langle i, j, k \rangle} \in H \cap \overline{\frac{1}{4}\sigma(2n+2)[h_{t \smallfrown \langle i, j \rangle}]}$$

and for  $k, l < m_{t \smallfrown \langle i, j \rangle}$  choose

$$h_{t \smallfrown \langle i, j, k, l \rangle} \in H \cap \overline{\frac{1}{4}\sigma(2n+3)[h_{t \smallfrown \langle i, j, k \rangle}]}$$

such that

$$\left\{ \frac{1}{4}\sigma(2n+2)[h_{t \smallfrown \langle i, j, k \rangle}] : k < m_{t \smallfrown \langle i, j \rangle} \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(2n+1)[h_{t \smallfrown \langle i, j \rangle}]}$ , and such that for each  $k < m_{t \smallfrown \langle i, j \rangle}$

$$\left\{ \frac{1}{4}\sigma(2n+3)[h_{t \smallfrown \langle i, j, k \rangle}] : l < m_t \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(2n+2)[h_{t \smallfrown \langle i, j, k \rangle}]}$ .

- For each  $t \in \max(T(n))$  and each  $i, j < m_t$ , put  $t \smallfrown \langle i \rangle$  and  $t \smallfrown \langle i, j \rangle$  in  $T(n+1)$ .

We now define the predetermined strategy  $\tau$  for  $\mathcal{O}$  in  $Clus_{O,P}(X, H)$  such that:

$$\tau(n) = \bigcup_{t \in \max(T(n))} \bigcap_{m < 2n+2} \frac{1}{4}\sigma(m)[h_{t \smallfrown m+1}]$$

noting that  $\tau(n) = \bigcap_{m \leq n} \tau(m)$  by definition.

Since  $\{\frac{1}{4}\sigma(2n+2)[h_{t \smallfrown \langle i \rangle}] : i < m_t\}$  is a cover for  $H \cap \frac{1}{4}\sigma(2n+1)[h_t]$ , since  $\{\frac{1}{4}\sigma(2n+1)[h_{t \smallfrown \langle i, j \rangle}] : j < m_t\}$  is a cover for  $H \cap \frac{1}{4}\sigma(2n)[h_{t \smallfrown \langle i \rangle}]$ , and since  $\{\frac{1}{4}\sigma(0)[h_{\langle i \rangle}] : i < m_\emptyset\}$  is a cover for  $H$ , it follows that  $\tau(n)$  contains  $H$  and is  $\tau$  is a legal strategy.

Let  $p$  be an attack against  $\tau$  such that  $p(n) \in \tau(n)$ . If we can construct an attack  $q$  against  $\sigma$  which shares a subsequence of  $p$ , then  $p$  must cluster since  $q$  must converge. To find such a  $q$ , we construct a subtree  $T' \subseteq T$ .

We begin by setting  $T'(0) = T(0)$ .

For  $n < \omega$ , suppose  $T'(n) \subseteq T(n)$  is defined such that:

- If  $t' \smallfrown \langle i, j \rangle \in T'(n)$ , then  $|t'| \leq 2n$

- If  $s' \leq t' \in T'(n)$ , then  $s \in T'(n)$ .
- If  $t' \in T'(n) \setminus \max(T'(n))$  and  $|t'|$  is even, then  $t' \frown \langle i, j \rangle \in T'(n)$  for  $i, j < m_{t'}$ .

Since  $p(n) \in \tau(n)$ , there exists some  $t_n \in \max(T(n))$  such that

$$p(n) \in \bigcap_{m < 2n+2} \frac{1}{4} \sigma(m) [h_{t_n \upharpoonright m+1}]$$

and in turn, there exists  $t'_n \frown \langle i, j \rangle \in \max(T'(n))$  with  $t'_n \frown \langle i, j \rangle \leq t_n$  and

$$p(n) \in \bigcap_{m < |t'_n|+2} \frac{1}{4} \sigma(m) [h_{t'_n \frown \langle i, j \rangle \upharpoonright m+1}]$$

since  $|t'_n| \leq 2n$ .

Let  $p_{t'_n \frown \langle i, j \rangle} = p(n)$

Take note that, in particular,

$$p_{t'_n \frown \langle i, j \rangle} \in \sigma(|t'_n|) [h_{t'_n \frown \langle i \rangle}]$$

and

$$p_{t'_n \frown \langle i, j \rangle} \in \frac{1}{4} \sigma(|t'_n| + 1) [h_{t'_n \frown \langle i, j \rangle}]$$

We then define

$$T'(n+1) = T'(n) \cup \{t'_n \frown \langle i, j, k \rangle : k \leq m_{t'_n \frown \langle i, j \rangle}\} \cup \{t'_n \frown \langle i, j, k, l \rangle : k, l \leq m_{t'_n \frown \langle i, j \rangle}\}$$

while noting that for all  $k \leq m_{t'_n \frown \langle i, j \rangle}$ ,

$$h_{t'_n \frown \langle i, j, k \rangle} \in H \cap \overline{\frac{1}{4} \sigma(|t'_n| + 2) [h_{t'_n \frown \langle i, j \rangle}]} \subseteq \frac{1}{2} \sigma(|t'_n| + 1) [h_{t'_n \frown \langle i, j \rangle}]$$

and thus

$$h_{t'_n \frown \langle i, j, k \rangle} \in \sigma(|t'_n| + 1) [p_{t'_n \frown \langle i, j \rangle}]$$

Finally, we let  $T' = \bigcup_{n < \omega} T'(n)$ . Since  $T'$  is an infinite tree with finite levels, we may pick an infinite branch  $b$ . From  $b$ , we construct the sequence

$$q = \langle h_{b \upharpoonright 1}, p_{b \upharpoonright 2}, h_{b \upharpoonright 3}, p_{b \upharpoonright 4}, \dots \rangle$$

and claim it attacks  $\sigma$  and thus must converge. If so, since  $\langle p_{b \upharpoonright 2}, p_{b \upharpoonright 4}, \dots \rangle$  is a subsequence of  $p$ ,  $p$  must cluster. To see this, recall that for some  $t'_n$ :

$$p_{b \upharpoonright 2n+2} = p_{t'_n \frown \langle i, j \rangle} \in \sigma(|t'_n|) [h_{t'_n \frown \langle i \rangle}]$$

and

$$h_{b \upharpoonright 2n+3} = h_{t'_n \frown \langle i, j, k \rangle} \in \sigma(|t'_n| + 1) [p_{t'_n \frown \langle i, j \rangle}]$$

We have thus proven  $\mathcal{O} \upharpoonright_{\text{pre}} \text{Clus}_{O,P}(X, H)$ , and thus  $\mathcal{O} \upharpoonright_{\text{pre}} \text{Con}_{O,P}(X, H)$ . □



**Example 42.** Let  $X = I \times 2$  be the Alexandrov double interval. Then  $\mathcal{D} \uparrow_{\text{pre}} \text{Prox}_{D,P}(X)$ ,  
but  $\mathcal{D} \uparrow_{\text{mark}} \text{Prox}_{D,P}(X)$ .

*Proof.* We assume that the uniformity on  $X$  is given by entourages

$$\begin{aligned} D(\epsilon, F) = & \{ \langle x, 0 \rangle, \langle y, 0 \rangle : |x - y| < \epsilon \} \cup \{ \langle x, 1 \rangle, \langle y, 0 \rangle : |x - y| < \epsilon \vee x \notin F \} \\ & \cup \{ \langle x, 0 \rangle, \langle y, 1 \rangle : |x - y| < \epsilon \vee y \notin F \} \cup \{ \langle x, 1 \rangle, \langle y, 1 \rangle : x = y \} \end{aligned}$$

That is, points are  $D(\epsilon, F)$ -close if they are the same point, or the first coordinates are within  $\epsilon$  of each other while neither second coordinate is in  $F$ .

Suppose  $\mathcal{D}$  had a predetermined winning strategy  $\sigma(n) = D(\epsilon_n, F_n)$ . Then  $\mathcal{P}$  can choose  $x \notin \bigcup_{n < \omega} F_n$ , and play  $\langle x, 1 \rangle$  during even rounds, and  $\langle x_{2n+1}, 0 \rangle$  where  $|x - x_{2n+1}| < \epsilon_{2n}$  during odd rounds, preventing convergence.

However, assume  $\mathcal{D}$  uses the Marköv strategy  $\sigma(x, n) = D(2^{-n}, \{x\})$ . If  $\mathcal{P}$  repeats a point of the form  $\langle x, 1 \rangle$ , then since  $D(2^{-n}, \{x\})[\langle x, 1 \rangle] = \{\langle x, 1 \rangle\}$ ,  $\mathcal{P}$  must repeat  $\langle x, 1 \rangle$  for the rest of the game, and  $\mathcal{D}$  wins. Otherwise,  $\mathcal{P}$  cannot repeat points played in  $I \times \{1\}$ , and as the first coordinates form a Cauchy sequence and converge to some  $z$ , any open set about  $\langle z, 0 \rangle$  contains all but finitely many points of  $\mathcal{P}$ 's sequence, and  $\mathcal{D}$  wins.  $\square$

**Theorem 43.** For any uniformly locally compact space  $X$ ,  $\mathcal{D} \uparrow_{\text{pre}} \text{Prox}_{D,P}(X) \Leftrightarrow \mathcal{D} \uparrow_{\text{pre}} a\text{Prox}_{D,P}(X)$

*Proof.* Let  $M$  be a uniformly locally compact entourage. Let  $\sigma$  witness  $\mathcal{D} \uparrow_{\text{pre}} \text{Prox}_{D,P}(X)$  such that  $\sigma(n) \subseteq M$  always (so  $\overline{\sigma(a)[x]} \subseteq \overline{M[x]}$  is compact),  $\sigma(n+1) \subseteq \frac{1}{4}\sigma(n)$ .

Let  $\tau(n) = \frac{1}{2}\sigma(n)$ . If  $p$  attacks  $\tau$  in  $a\text{Prox}_{D,P}(X)$ , then

$$p(n+1) \in \tau(n)[p(n)] = \frac{1}{2}\sigma(n)[p(n)]$$

and for

$$x \in \overline{\sigma(n+1)[p(n+1)]} \subseteq \overline{\frac{1}{4}\sigma(n)[p(n+1)]} \subseteq \frac{1}{2}\sigma(n)[p(n+1)]$$

we can conclude  $x \in \sigma(n)[p(n)]$ . Thus

$$\sigma(n+1)[p(n+1)] \subseteq \overline{\sigma(n+1)[p(n+1)]} \subseteq \sigma(n)[p(n)]$$

Finally, note that  $p$  attacks the winning strategy  $\sigma$  in  $Prox_{D,P}(X)$ , but since the intersection of a chain of nonempty compact sets is nonempty:

$$\bigcap_{n < \omega} \sigma(n)[p(n)] = \bigcap_{n < \omega} \overline{\sigma(n)[p(n)]} \neq \emptyset$$

We conclude that  $p$  converges. □

**Proposition 44.** *If  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$ , then  $X$  has a  $G_\delta$  diagonal.*

*Proof.* If  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$  with strategy  $\sigma$ , then consider  $\langle x, y \rangle \in \bigcap_{n < \omega} \sigma(n)$ . It follows that  $\langle x, y, x, y, \dots \rangle$  attacks  $\sigma$ , and  $\{x, y\} \subseteq \bigcap_{n < \omega} \sigma(n)[x] \cap \bigcap_{n < \omega} \sigma(n)[y] \neq \emptyset$  so it must converge, and  $x = y$ . Thus  $\bigcap_{n < \omega} \sigma(n) = \Delta$  is  $G_\delta$ . □

**Example 45.** The Sorgenfrey line  $S$  has a  $G_\delta$  diagonal but  $\mathcal{D} \uparrow Prox_{D,P}(S)$ .

**Corollary 46.** *For  $X$  with uniformity  $\mathbb{D}$  inducing the compact Hausdorff topology  $\tau$ , the following are equivalent:*

- (a)  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$
- (b)  $\mathcal{D} \uparrow_{pre} aProx_{D,P}(X)$
- (c)  $X$  has a  $G_\delta$  diagonal
- (d)  $\mathbb{D}$  is metrizable
- (e)  $\tau$  is metrizable

*Proof.* For compact Hausdorff spaces, it is well known that there is exactly one uniformity inducing the topology. Thus (d)  $\Leftrightarrow$  (e). Since  $X$  is uniformly locally compact, (a)  $\Leftrightarrow$  (b). Also, compact spaces with a  $G_\delta$  diagonal are metrizable, so (c)  $\Rightarrow$  (e). Bell noted (d)  $\Rightarrow$  (a) for arbitrary uniform spaces, and the previous proposition shows (a)  $\Rightarrow$  (c). □

**Theorem 47.** *A uniformly locally compact space with a  $G_\delta$  diagonal is metrizable.*

*Proof.* Based on several folklore results.

Uniformly locally compact implies the topological sum of  $\sigma$ -compact spaces implies paracompact. Locally compact plus  $G_\delta$  diagonal implies locally metrizable. Locally metrizable plus paracompact characterizes metrizable. □

**Corollary 48.** *If  $X$  is uniformly locally compact, then  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$  implies  $X$ 's topology is metrizable.*

**Example 49.** Let  $R$  be the Michael Line. Then  $\mathcal{P} \uparrow \text{Prox}_{D,P}(X)$ .

*Proof.* During round 0,  $\mathcal{P}$  may choose  $m(0) = 0$  and  $p(0) = 1$ , and during round  $n + 1$ ,  $\mathcal{P}$  may choose  $m(n + 1) > m(n)$  and  $p(n + 1) = p(n) + \frac{1}{10^{m(n+1)}}$  such that  $p$  is a legal attack.

It follows that  $p$  “converges” to  $x = \sum_{n < \omega} \frac{1}{10^{m(n)}}$ , except  $x$  is an irrational number composed of 1s separated by strings of 0s of strictly increasing size.  $\square$

**Example 50.** Let  $\kappa$  be an uncountable regular cardinal with a ladder topology:

- All successor ordinals are isolated.
- Strictly increasing sequences (ladders)  $L_\alpha : \omega \rightarrow \alpha$  are defined for each limit ordinal  $\alpha$  such that  $L_\alpha$  converges to  $\alpha$  in the order topology, and each limit  $\alpha$  is given neighborhoods of the form  $\{\alpha\} \cup \{L_\alpha(n) : n \geq m\}$ . We assume that all successor ordinals are a part of some ladder.

Then  $\mathcal{P} \uparrow \text{Prox}_{D,P}(\kappa^*)$  where  $\kappa^*$  is its one-point compactification.

*Proof.* Entrouages of  $\kappa^*$  are then of the form  $D(F, n)$ , where  $F \in [\kappa^L]^{<\omega}$  and  $n < \omega$ .  $D(F, n)$  partitions  $\kappa^*$  such that  $\infty$ ’s part is the complement of the ladders leading to points in  $F$ . Each of those ladders is then partitioned by isolating the first  $n$  rungs of the ladder, and leaving the top of the ladder leading to a point in  $F$  as a whole part. (It’s possible that the tops of some ladders might overlap, so they must be considered the same part, but this could be prevented by  $\mathcal{D}$  by increasing  $n$  a sufficient amount to separate all the finite limits in  $F$  if desired.)

$\mathcal{P}$ ’s strategy involves first choosing two disjoint stationary subsets  $S_0, T_0$  of  $\kappa^L$ . During round 0,  $\mathcal{D}$ ’s move partitions ladders leading to limit ordinals in  $F_0 \in [\kappa^L]^{<\omega}$ . Let  $S'_0 = S_0 \setminus F_0$  and  $T'_0 = T_0 \setminus F_0$ , and observe that both are still stationary sets as only finitely many ordinals were removed.

For  $\mathcal{P}$ ’s initial move, she may apply the pressing down lemma to the sets  $S'_0, T'_0$  and the function  $f_i(\alpha) = L_\alpha(i)$  for  $i < \omega$  sufficiently large to identify stationary subsets  $S_1, T_1$  of  $S'_0, T'_0$  such that  $f_i(\alpha) = s_0$  for  $\alpha \in S_1$ ,  $f_i(\alpha) = t_0$  for  $\alpha \in T_1$ , and  $s_0, t_0$  are not in the range of  $L_\alpha$  for  $\alpha \in F_0$ .

$\mathcal{P}$  chooses  $s_0$  as her initial move.

During round  $n + 1$ , we assume that the disjoint stationary sets  $S_{n+1}, T_{n+1}$  were defined in the previous round.  $\mathcal{D}$ ’s move in this round again partitions ladders leading to limit ordinals in  $F_{n+1} \in [\kappa^L]^{<\omega}$ . Let  $S'_{n+1} = S_{n+1} \setminus F_{n+1}$  and  $T'_{n+1} = T_{n+1} \setminus F_{n+1}$ .

$\mathcal{P}$  then applies the pressing down lemma to the sets  $S'_{n+1}, T'_{n+1}$  and the function  $f_i(\alpha) = L_\alpha(i)$  for  $n < i < \omega$  sufficiently large to identify stationary subsets  $S_{n+2}, T_{n+2}$  of  $S'_{n+1}, T'_{n+1}$

such that  $f_i(\alpha) = s_{n+1}$  for  $\alpha \in S_{n+2}$ ,  $f_i(\alpha) = t_{n+1}$  for  $\alpha \in T_{n+2}$ , and  $s_{n+1}, t_{n+1}$  are not in the range of  $L_\alpha$  for  $\alpha \in F_{n+1}$ .

If  $n + 1$  is even,  $\mathcal{P}$  chooses  $s_{n+1}$  as her move; otherwise, she chooses  $t_{n+1}$ .

All choices of  $s_n, t_n$  by  $\mathcal{P}$  were within the partition containing  $\infty$ , and no choice was repeated infinitely often, so  $s_n$  and  $t_n$  must converge. (Need to disprove that either could converge to  $\infty$ , or could they? That would happen if  $\bigcap_{n < \omega} S_n = \emptyset$  or  $\bigcap_{n < \omega} T_n = \emptyset$ .)

□