

DUAL SELECTION GAMES

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ABSTRACT. (an investigation of dual selection games)

1. INTRODUCTION

Definition 1. The *selection game* $G_1(\mathcal{A}, \mathcal{B})$ is an ω -length game involving Players I and II. During round n , I chooses $A_n \in \mathcal{A}$, followed by II choosing $B_n \in A_n$. Player II wins in the case that $\{B_n : n < \omega\} \in \mathcal{B}$, and Player I wins otherwise.

For brevity, let

$$G_1(\mathcal{A}, \neg\mathcal{B}) = G_1(\mathcal{A}, \mathcal{P}\left(\bigcup \mathcal{A}\right) \setminus \mathcal{B}).$$

That is, II wins in the case that $\{B_n : n < \omega\} \notin \mathcal{B}$, and I wins otherwise.

Definition 2. For a set X , let $\mathbf{C}(X)$ be the collection of all choice functions on X , functions $f : X \rightarrow \bigcup X$ such that $f(x) \in x$ for all $x \in X$.

Definition 3. The set \mathcal{R} is said to be a *reflection* of the set \mathcal{A} if

$$\mathcal{A} = \{\text{range}(f) : f \in \mathbf{C}(\mathcal{R})\}.$$

For example, a reflection of the collection \mathcal{O}_X of basic open covers of X would be $\mathcal{P}_X = \{\mathcal{T}_{X,x} : x \in X\}$, where $\mathcal{T}_{X,x}$ is the corresponding point-base at $x \in X$. Likewise for the collection $\Omega_{X,x}$ of sets with $x \in X$ as a limit point, $\mathcal{T}_{X,x}$ is itself a reflection.

Lemma 4. Let \mathcal{R} be a reflection of \mathcal{A} . Then $\bigcup \mathcal{R} = \bigcup \mathcal{A}$.

Proof. If $x \in \bigcup \mathcal{A}$, then $x \in \text{range}(f)$ for some $f \in \mathbf{C}(\mathcal{R})$. Thus $x = f(R) \in R$ for some $R \in \mathcal{R}$, showing $x \in \bigcup \mathcal{R}$.

Likewise if $x \in \bigcup \mathcal{R}$, so $x \in R$ for some $R \in \mathcal{R}$. Let $f \in \mathbf{C}(\mathcal{R})$ satisfy $f(R) = x$, so $x \in \text{range}(f)$, showing $x \in \bigcup \mathcal{A}$. \square

Theorem 5. Let \mathcal{R} be a reflection of \mathcal{A} .

Then I $\uparrow_{\text{pre}} G_1(\mathcal{A}, \mathcal{B})$ if and only if II $\uparrow_{\text{mark}} G_1(\mathcal{R}, \neg\mathcal{B})$.

Proof. Let σ witness I $\uparrow_{\text{pre}} G_1(\mathcal{A}, \mathcal{B})$. Since $\sigma(n) \in \mathcal{A} = \{\text{range}(f) : f \in \mathbf{C}(\mathcal{R})\}$, $\sigma(n) = \text{range}(f_n)$ for some $f_n \in \mathbf{C}(\mathcal{R})$. So let $\tau(R, n) = f_n(R)$ for all $R \in \mathcal{R}$ and $n < \omega$. Suppose $R_n \in \mathcal{R}$ for all $n < \omega$. Note that since σ is winning and $\tau(R_n, n) = f_n(R_n) \in \text{range}(f_n) = \sigma(n)$, $\{\tau(R_n, n) : n < \omega\} \notin \mathcal{B}$. Thus τ witnesses II $\uparrow_{\text{mark}} G_1(\mathcal{R}, \neg\mathcal{B})$.

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Now let σ witness $\text{II} \uparrow_{\text{mark}} G_1(\mathcal{R}, \neg\mathcal{B})$. Let $f_n \in \mathbf{C}(\mathcal{R})$ be defined by $f_n(R) = \sigma(R, n)$. Since $\tau(n) \in \mathcal{A} = \{\text{range}(f) : f \in \mathbf{C}(\mathcal{R})\}$, let $\tau(n) = \text{range}(f_n)$. Suppose that $B_n \in \tau(n) = \text{range}(f_n)$ for all $n < \omega$. Choose $R_n \in \mathcal{R}$ such that $B_n = f_n(R_n) = \sigma(R_n, n)$. Since σ is winning, $\{B_n : n < \omega\} \notin \mathcal{B}$. Thus τ witnesses $\text{I} \uparrow_{\text{pre}} G_1(\mathcal{A}, \mathcal{B})$. \square

Theorem 6. *Let \mathcal{R} be a reflection of \mathcal{A} .*

Then $\text{II} \uparrow_{\text{mark}} G_1(\mathcal{A}, \mathcal{B})$ if and only if $\text{I} \uparrow_{\text{pre}} G_1(\mathcal{R}, \neg\mathcal{B})$.

Proof. Let σ witness $\text{II} \uparrow_{\text{mark}} G_1(\mathcal{A}, \mathcal{B})$. Let $n < \omega$. Suppose that for each $R \in \mathcal{R}$, there was $g(R) \in \mathcal{A}$ such that for all $A \in \mathcal{A}$, $\sigma(A, n) \neq g(R)$. Then $g \in \mathbf{C}(\mathcal{R})$, and $\sigma(\text{range}(g), n) \neq g(R)$ for all $R \in \mathcal{R}$, a contradiction.

So choose $\tau(n) \in \mathcal{R}$ such that for all $r \in \tau(n)$ there exists $A_{r,n} \in \mathcal{A}$ such that $\sigma(A_{r,n}, n) = r$. It follows that when $r_n \in \tau(n)$ for $n < \omega$, $\{r_n : n < \omega\} = \{\sigma(A_{r_n,n}, n) : n < \omega\} \in \mathcal{B}$, so τ witnesses $\text{I} \uparrow_{\text{pre}} G_1(\mathcal{R}, \neg\mathcal{B})$.

Now let σ witness $\text{I} \uparrow_{\text{pre}} G_1(\mathcal{R}, \neg\mathcal{B})$. Then $\sigma(n) \in \mathcal{R}$, so for $A \in \mathcal{A}$, let $f_A \in \mathbf{C}(\mathcal{R})$ satisfy $A = \text{range}(f_A)$, and let $\tau(A, n) = f_A(\sigma(n))$. Then if $A_n \in \mathcal{A}$ for $n < \omega$, $\tau(A_n, n) \in \sigma(n)$, so $\{\tau(A_n, n) : n < \omega\} \in \mathcal{B}$. Thus τ witnesses $\text{II} \uparrow_{\text{mark}} G_1(\mathcal{A}, \mathcal{B})$. \square

Theorem 7. *Let \mathcal{R} be a reflection of \mathcal{A} .*

Then $\text{I} \uparrow G_1(\mathcal{A}, \mathcal{B})$ if and only if $\text{II} \uparrow G_1(\mathcal{R}, \neg\mathcal{B})$.

Proof. Let σ witness $\text{I} \uparrow G_1(\mathcal{A}, \mathcal{B})$. Let $c(\emptyset) = \emptyset$. Suppose $c(s) \in (\bigcup A)^{<\omega} = (\bigcup R)^{<\omega}$ is defined for $s \in \mathcal{R}^{<\omega}$. Since $\sigma(c(s)) \in \mathcal{A}$, let $f_s \in \mathbf{C}(\mathcal{R})$ satisfy $\sigma(c(s)) = \text{range}(f_s)$, and let $c(s \smallfrown \langle R \rangle) = c(s) \smallfrown \langle f_s(R) \rangle$. Then let $c(\alpha) = \bigcup \{c(\alpha \upharpoonright n) : n < \omega\}$ for $\alpha \in \mathcal{R}^\omega$, so

$$c(\alpha)(n) = f_{\alpha \upharpoonright n}(\alpha(n)) \in \text{range}(f_{\alpha \upharpoonright n}) = \sigma(c(\alpha \upharpoonright n))$$

demonstrating that $c(\alpha)$ is a legal attack against σ .

Let $\tau(s \smallfrown \langle R \rangle) = f_s(R)$. Consider the attack $\alpha \in \mathcal{R}^\omega$ against τ . Then since σ is winning and $\tau(\alpha \upharpoonright n+1) = f_{\alpha \upharpoonright n}(\alpha(n)) \in \text{range}(f_{\alpha \upharpoonright n}) = \sigma(c(\alpha \upharpoonright n))$, it follows that $\{\tau(\alpha \upharpoonright n+1) : n < \omega\} \notin \mathcal{B}$. Thus τ witnesses $\text{II} \uparrow G_1(\mathcal{R}, \neg\mathcal{B})$.

Now let σ witness $\text{II} \uparrow G_1(\mathcal{R}, \neg\mathcal{B})$. For $s \in \mathcal{R}^{<\omega}$, define $f_s \in \mathbf{C}(\mathcal{R})$ by $f_s(R) = \sigma(s \smallfrown \langle R \rangle)$. Let $\tau(\emptyset) = \text{range}(f_\emptyset)$, and for $x \in \tau(\emptyset)$, choose $R_{\langle x \rangle} \in \mathcal{R}$ such that $x = f_\emptyset(R_{\langle x \rangle})$ (for other $x \in \bigcup A$, choose $R_{\langle x \rangle}$ arbitrarily as it won't be used). Now let $s \in (\bigcup A)^{<\omega} \setminus \emptyset$, and suppose $\tau(s \upharpoonright n) \in \mathcal{A}$ and $R_{s \upharpoonright n+1} \in \mathcal{R}$ have been defined for $n < |s|$. Then let $\tau(s) = \text{range}(f_{\langle R_{s \upharpoonright 0}, \dots, R_{s \upharpoonright n} \rangle})$ and for $x \in \tau(s)$ choose $R_{s \smallfrown \langle x \rangle}$ such that $x = f_{\langle R_{s \upharpoonright 0}, \dots, R_{s \upharpoonright n} \rangle}(R_{s \smallfrown \langle x \rangle})$ (and again, choose $R_{s \smallfrown \langle x \rangle}$ arbitrarily for other $x \in \bigcup A$ as it won't be used).

Then let α attack τ , so $\alpha(n) \in \tau(\alpha \upharpoonright n)$ and thus $\alpha(n) = f_{\langle R_{\alpha \upharpoonright 0}, \dots, R_{\alpha \upharpoonright n} \rangle}(R_{\alpha \upharpoonright n+1}) = \sigma(\langle R_{\alpha \upharpoonright 0}, \dots, R_{\alpha \upharpoonright n+1} \rangle)$. Since σ is winning, $\{\sigma(\langle R_{\alpha \upharpoonright 0}, \dots, R_{\alpha \upharpoonright n+1} \rangle) : n < \omega\} = \{\alpha(n) : n < \omega\} \notin \mathcal{B}$. Thus τ witnesses $\text{I} \uparrow G_1(\mathcal{A}, \mathcal{B})$. \square

REFERENCES

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