

Limited Information Strategies for Topological Games

by

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Abstract

I talk a lot about topological games.

TODO: Write this.

Acknowledgments

TODO: Thank people.

Table of Contents

| | |
|--|-----|
| Abstract | ii |
| Acknowledgments | iii |
| List of Figures | v |
| List of Tables | vi |
| 1 W convergence and clustering games | 1 |
| 1.1 Fort spaces | 4 |
| 1.2 Sigma-products | 8 |
| Bibliography | 12 |

List of Figures

List of Tables

Chapter 1

W convergence and clustering games

We begin by investigating a game due to Gary Gruenhage.

Game 1.0.1. Let $Con_{O,P}(X, S)$ denote the W -convergence game with players \mathcal{O} , \mathcal{P} , for a topological space X and $S \subseteq X$.

In round n , \mathcal{O} chooses an open neighborhood $O_n \supseteq S$, followed by \mathcal{P} choosing a point $x_n \in \bigcap_{m \leq n} O_m$.

\mathcal{O} wins the game if the points x_n converge to the set S ; that is, for every open neighborhood $U \supseteq S$, $x_n \in U$ for all but finite $n < \omega$.

If $S = \{x\}$ then we write $Con_{O,P}(X, x)$ for short. \diamond

The “W” in the name merely refers to \mathcal{O} ’s goal: to “win” the game. Gruenhage defined this game in his doctoral dissertation to define a class of spaces generalizing first-countability.

[1]

Definition 1.0.2. The spaces X for which $\mathcal{O} \uparrow Con_{O,P}(X, x)$ for all $x \in X$ are called W -spaces. \diamond

In fact, using limited information strategies, one may characterize the first-countable spaces using this game.

Proposition 1.0.3. X is first countable if and only if $\mathcal{O} \uparrow_{pre} Con_{O,P}(X, x)$ for all $x \in X$. \diamond

Proof. The forward implication shows that all W spaces are first-countable spaces, and was proven in [1]: if $\{U_n : n < \omega\}$ is a countable base at x , let $\sigma(n) = \bigcap_{m \leq n} U_m$. σ is easily seen to be a winning predetermined strategy.

If X is not first countable at some x , let σ be a predetermined strategy for \mathcal{O} in $Con_{\mathcal{O},P}^*(X, x)$. There exists an open neighborhood U of x which does not contain any $\bigcap_{m \leq n} \sigma(m)$ (otherwise $\{\bigcap_{m \leq n} \sigma(m) : n < \omega\}$ would be a countable base at x). Let x_n be an element of $\bigcap_{m \leq n} \sigma(m) \setminus U$ for all $n < \omega$. Then $\langle x_0, x_1, \dots \rangle$ is a winning counter-attack to σ for \mathcal{P} , so \mathcal{O} lacks a winning predetermined strategy. \square

At first glance, the difficulty of $Con_{\mathcal{O},P}(X, S)$ could be increased for \mathcal{O} by only restricting the choices for \mathcal{P} to be within the most recent open set played by \mathcal{O} , rather than all the previously played open sets.

Definition 1.0.4. Let $Con_{\mathcal{O},P}^*(X, S)$ denote the *hard W -convergence game* which proceeds as $Con_{\mathcal{O},P}(X, S)$, except that \mathcal{P} need only choose $x_n \in O_n$ rather than $x_n \in \bigcap_{m \leq n} O_m$ during each round. \diamond

This seemingly more difficult game for \mathcal{O} is Gruenhage's original formulation. But with perfect information, there is no real difference for \mathcal{O} .

Proposition 1.0.5. $\mathcal{O} \uparrow_{limit} Con_{\mathcal{O},P}(X, S)$ if and only if $\mathcal{O} \uparrow_{limit} Con_{\mathcal{O},P}^*(X, S)$, where \uparrow_{limit} is either \uparrow or \uparrow_{pre} . \diamond

Proof. The backwards implication is immediate.

For the forward implication, let σ be a winning predetermined (perfect information) strategy, and λ be the 0-Marköv fog-of-war μ_0 (the identity).

We define a new predetermined (perfect information) strategy τ by

$$\tau \circ \lambda(\langle x_0, \dots, x_{n-1} \rangle) = \bigcap_{m \leq n} \sigma \circ \lambda(\langle x_0, \dots, x_{m-1} \rangle)$$

so that each move by \mathcal{O} according to $\tau \circ \lambda$ is the intersection of \mathcal{O} 's previous moves. Then any attack against $\tau \circ \lambda$ is an attack against $\sigma \circ \lambda$, and since $\sigma \circ \lambda$ is a winning strategy, so is $\tau \circ \lambda$. \square

Put more simply, $\tau(n) = \bigcap_{m \leq n} \sigma(m)$ in the predetermined case, and $\tau(\langle x_0, \dots, x_{n-1} \rangle) = \bigcap_{m \leq n} \sigma(\langle x_0, \dots, x_{m-1} \rangle)$ in the perfect information case. The original proof would have been invalid if λ was required to be, say, the tactical fog-of-war ν_1 , since the value of \mathcal{O} 's own round 1 move $\sigma \circ \nu_1(\langle x_0 \rangle) = \sigma(\langle x_0 \rangle)$ could not be determined from the information she has during round 2: $\nu_1(\langle x_0, x_1 \rangle) = \langle x_1 \rangle$.

Due to the equivalency of the “hard” and “normal” variations of the convergence game in the perfect information case, many authors use them interchangeably. However, it is possible to find spaces for which the games are not equivalent when considering $k + 1$ -tactics and $k + 1$ -marks, as we will soon see.

In addition to the W -convergence games, we will also investigate “clustering” analogs to both variations.

Game 1.0.6. Let $Clus_{O,P}(X, S)$ ($Clus_{O,P}^*(X, S)$) be a variation of $Con_{O,P}(X, S)$ ($Con_{O,P}^*(X, S)$) such that x_n need only cluster at S , that is, for every open neighborhood U of S , $x_n \in U$ for infinitely many $n < \omega$. \diamond

This variation seems to make \mathcal{O} 's job easier, but Gruenhage noted that the clustering game is perfect-information equivalent to the convergence game for \mathcal{O} . This can easily be extended for some limited information cases as well.

Proposition 1.0.7. $\mathcal{O} \xrightarrow{\text{limit}} Con_{O,P}(X, S)$ if and only if $\mathcal{O} \xrightarrow{\text{limit}} Clus_{O,P}(X, S)$ where $\xrightarrow{\text{limit}}$ is any of $\xrightarrow{\text{pre}}$, $\xrightarrow{\text{tact}}$, $\xrightarrow{\text{mark}}$, or $\xrightarrow{\text{mark}}$. \diamond

Proof. For the perfect information case we refer to [1].

In the predetermined (resp. tactical) case, suppose that σ is a winning predetermined (resp. tactical) strategy for \mathcal{O} in $Clus_{O,P}(X, S)$. Let p be a legal attack against σ , and q be a subsequence of p . It's easily seen that q is also a legal attack against σ , so q clusters at S . Since every subsequence of p clusters at S , p converges to S , and σ is a winning predetermined (resp. tactical) strategy for \mathcal{O} in $Con_{O,P}(X, S)$ as well.

In the final case, note that any Marköv strategy σ' for \mathcal{O} may be strengthened to σ defined by $\sigma(x, n) = \bigcap_{m \leq n} \sigma'(x, m)$. So, suppose that σ is a winning Marköv strategy for \mathcal{O} in $Clus_{O,P}(X, S)$ such that $\sigma(x, m) \supseteq \sigma(x, n)$ for all $m \leq n$.

Let p be a legal attack against σ , and q be a subsequence of p . For $m < \omega$, there exists $f(m) \geq m$ such that $q(m) = p(f(m))$. It follows that $q(0) = p(f(0)) \in \sigma(\emptyset, 0) \cap \bigcap_{m \leq f(0)} \sigma(\langle p(m) \rangle, m) \subseteq \sigma(\emptyset, 0)$ and

$$\begin{aligned} q(n+1) = p(f(n+1)) &\in \sigma(\emptyset, 0) \cap \bigcap_{m < f(n+1)} \sigma(\langle p(m) \rangle, m+1) \\ &\subseteq \sigma(\emptyset, 0) \cap \bigcap_{m < n+1} \sigma(\langle p(f(m)) \rangle, f(m)+1) \\ &= \sigma(\emptyset, 0) \cap \bigcap_{m < n+1} \sigma(\langle q(m) \rangle, f(m)+1) \\ &\subseteq \sigma(\emptyset, 0) \cap \bigcap_{m < n+1} \sigma(\langle q(m) \rangle, m+1) \end{aligned}$$

so q is also a legal attack against σ . Since σ is a winning strategy, q clusters at S , and since every subsequence of p clusters at S , p must converge to S . Thus σ is also a winning Marköv strategy for \mathcal{O} in $Con_{O,P}(X, S)$ as well. \square

Two types of questions emerge from these results.

Question 1.0.8. Does $\mathcal{O} \uparrow_{2\text{-tact}} Clus_{O,P}(X, S)$ imply $\mathcal{O} \uparrow_{2\text{-tact}} Con_{O,P}(X, S)$? What about for $\uparrow_{2\text{-mark}} ?$ \diamond

Question 1.0.9. Could $\mathcal{O} \uparrow_{k+1\text{-tact}} Con_{O,P}(X, S)$ actually imply $\mathcal{O} \uparrow_{\text{tact}} Con_{O,P}(X, S)$? What about for $Clus_{O,P}(X, S)$? \diamond

1.1 Fort spaces

In his original paper, Gruenhage suggested the one-point-compactification of a discrete space as an example of a W -space which is not first-countable.

Definition 1.1.1. A *Fort space* $\kappa^* = \kappa \cup \{\infty\}$ is defined for each cardinal κ . Its subspace κ is discrete, and the neighborhoods of ∞ are of the form $\kappa^* \setminus F$ for each $F \in [\kappa]^{<\omega}$. \diamond

Proposition 1.1.2. $\mathcal{O} \uparrow_{tact} Con_{O,P}(\kappa^*, \infty)$ for all cardinals κ \diamond

Proof. Let $\sigma(\emptyset) = \sigma(\langle \infty \rangle) = \kappa^*$ and $\sigma(\langle \alpha \rangle) = \kappa^* \setminus \{\alpha\}$. Any legal attack against the tactic σ could not repeat non- ∞ points, so it must converge to ∞ . \square

Corollary 1.1.3. $\mathcal{O} \uparrow Con_{O,P}^*(\kappa^*, \infty)$ for all cardinals κ \diamond

Proof. Propositions 1.0.5 and 1.1.2. \square

Since it's trivial to show that $\mathcal{O} \uparrow_{pre} Con_{O,P}(\kappa^*, \infty)$ if and only if $\kappa \leq \omega$, this closes the question on limited information strategies for $Con_{O,P}(\kappa^*, \infty)$. However, limited information analysis of the harder $Con_{O,P}^*(\kappa^*, \infty)$ is more interesting.

Peter Nyikos noted Proposition 1.1.2 and the following in [2].

Theorem 1.1.4. $\mathcal{O} \not\uparrow_{mark} Con_{O,P}^*(\omega_1^*, \infty)$. \diamond

This actually can be generalized to any k -Marköv strategy with just a little more book-keeping.

Theorem 1.1.5. $\mathcal{O} \not\uparrow_{k-mark} Con_{O,P}^*(\omega_1^*, \infty)$. \diamond

Proof. Let σ be a k -mark for \mathcal{O} . Since the set

$$D_\sigma = \bigcap_{n < \omega, s \in \omega^{\leq k}} \sigma(s, n)$$

is co-countable, we may choose $\alpha_\sigma \in D_\sigma \cap \omega_1$. Thus, we may choose $n_0 < n_1 < \dots < \omega$ such that

$$\langle n_0, \dots, n_{k-1}, \alpha_\sigma, n_k, \dots, n_{2k-1}, \alpha_\sigma, \dots \rangle$$

is a legal counterattack, which fails to converge to ∞ since α_σ is repeated infinitely often. \square

However, while the clustering and convergence variants are equivalent for Marköv strategies in the “normal” version of the W game, they are *not* equivalent in the “hard” version.

Theorem 1.1.6. $\mathcal{O} \uparrow_{\text{mark}} \text{Clus}_{O,P}^*(\omega_1^*, \infty)$. \diamond

Proof. For each $\alpha < \omega_1$ let $A_\alpha = \langle A_\alpha(0), A_\alpha(1), \dots \rangle$ be a countable sequence of finite sets such that $A_\alpha(n) \subset A_\alpha(n+1)$ and $\bigcup_{n < \omega} A_\alpha(n) = \alpha + 1$.

We define the Marköv strategy σ by setting

$$\sigma(\emptyset, 0) = \sigma(\langle \infty \rangle, n) = \omega_1^*$$

and for all $\alpha < \omega_1$ setting

$$\sigma(\langle \alpha \rangle, n) = \omega_1^* \setminus A_\alpha(n)$$

Note that for any $\alpha_0 < \dots < \alpha_{k-1}$, there is some $n < \omega$ such that $\{\alpha_0, \dots, \alpha_{k-1}\} \subseteq A_{\alpha_i}(n)$ for all $i < k$. Thus for any legal attack p against σ , the range of p cannot be finite. Since the range of p is infinite, every open neighborhood of ∞ contains infinitely many points of p , so p clusters at ∞ . \square

However, knowledge of the round number is critical.

Theorem 1.1.7. $\mathcal{O} \nearrow_{k\text{-tact}} \text{Clus}_{O,P}^*(\omega_1^*, \infty)$. \diamond

Proof. Let σ be a k -tactic for \mathcal{O} in $\text{Clus}_{O,P}(\omega_1^*, \infty)$. By the closing-up lemma, the set

$$C_\sigma = \{\alpha < \omega_1 : s \in \alpha^{\leq k} \Rightarrow \omega_1^* \setminus \sigma(s) \subset \alpha\}$$

is closed and unbounded. Let $a_\sigma : \omega_1 \rightarrow C_\sigma$ be an order isomorphism.

Choose $n_0 < \dots < n_{k-1} < \omega$ such that for each $i < k$:

$$a_\sigma(n_i) \in \sigma(\langle a_\sigma(n_0), \dots, a_\sigma(n_{i-1}), a_\sigma(\omega + i), \dots, a_\sigma(\omega + k - 1) \rangle)$$

Finally, observe that the legal counterattack

$$\langle a_\sigma(n_0), \dots, a_\sigma(n_{k-1}), a_\sigma(\omega), \dots, a_\sigma(\omega+k-1), a_\sigma(n_0), \dots, a_\sigma(n_{k-1}), a_\sigma(\omega), \dots, a_\sigma(\omega+k-1), \dots \rangle$$

has a range outside the open neighborhood

$$\omega_1^* \setminus \{a_\sigma(n_0), \dots, a_\sigma(n_{k-1}), a_\sigma(\omega), \dots, a_\sigma(\omega+k-1)\}$$

of ∞ . Thus σ is not a winning k -tactic. □

Once the discrete space is larger than ω_1 , knowing the round number is not sufficient to construct a limited information strategy, due to a similar argument.

Theorem 1.1.8. $O \not\uparrow_{k\text{-mark}} Clus_{O,P}^*(\omega_2^*, \infty)$. ◇

Proof. Let σ be a k -mark for \mathcal{O} in $Clus_{O,P}(\omega_2^*, \infty)$. By the closing-up lemma, the set

$$C_\sigma = \{\alpha < \omega_2 : s \in \alpha^{<\omega} \Rightarrow \omega_2^* \setminus \sigma \circ \mu_k(s) \subset \alpha\}$$

is closed and unbounded. Let $a_\sigma : \omega_2 \rightarrow C_\sigma$ be an order isomorphism.

Choose $\beta_0 < \dots < \beta_{k-1} < \omega_1$ such that for each $i < k$:

$$a_\sigma(\beta_i) \in \bigcup_{n < \omega} \sigma(\langle a_\sigma(\beta_0), \dots, a_\sigma(\beta_{i-1}), a_\sigma(\omega_1 + i), \dots, a_\sigma(\omega_1 + k - 1) \rangle, n)$$

Finally, observe that the legal counterattack

$$\langle a_\sigma(\beta_0), \dots, a_\sigma(\beta_{k-1}), a_\sigma(\omega_1), \dots, a_\sigma(\omega_1+k-1), a_\sigma(\beta_0), \dots, a_\sigma(\beta_{k-1}), a_\sigma(\omega_1), \dots, a_\sigma(\omega_1+k-1), \dots \rangle$$

has a range outside the open neighborhood

$$\omega_2^* \setminus \{a_\sigma(\beta_0), \dots, a_\sigma(\beta_{k-1}), a_\sigma(\omega_1), \dots, a_\sigma(\omega_1+k-1)\}$$

of ∞ . Thus σ is not a winning k -mark. \square

1.2 Sigma-products

Knowing the status of W -games in simpler spaces yields insight to larger spaces.

Proposition 1.2.1. *Suppose $S \subseteq Y \subseteq X$, \uparrow_{limit} is any of \uparrow , $\uparrow_{k\text{-tact}}$, or $\uparrow_{k\text{-mark}}$, and $G(X, S)$ is any of $\text{Con}_{O,P}(X, S)$, $\text{Con}_{O,P}^*(X, S)$, $\text{Clus}_{O,P}(X, S)$, or $\text{Clus}_{O,P}^*(X, S)$.*

Then $\mathcal{O} \uparrow_{\text{limit}} G(X, S)$ implies $\mathcal{O} \uparrow_{\text{limit}} G(Y, S)$. \diamond

Proof. Simply intersect the output of the winning strategy in $G(X, S)$ with Y . \square

A natural superspace of a Fort space is the sigma-product of a discrete cardinal.

Definition 1.2.2. Let $\Sigma_y X^\kappa$ be a *sigma product* of X with dimension κ for each $y \in X^\kappa$, the subset of the usual Tychonoff product space X^κ such that $x \in \Sigma_y X^\kappa$ if and only if $\{\alpha < \kappa : x(\alpha) \neq y(\alpha)\}$ is countable.

For homogeneous spaces X containing 0, y is usually assumed to be the zero vector $\vec{0}$ and the sigma product is written ΣX^κ . \diamond

Proposition 1.2.3. κ^* is homeomorphic to the space

$$\{x \in \Sigma 2^\kappa : x(\alpha) = 0 \text{ for all but one } \alpha < \kappa\}$$

\diamond

Proof. Map $\alpha < \kappa$ to x_α such that

$$x_\alpha(\beta) = \begin{cases} 0 & \beta \neq \alpha \\ 1 & \beta = \alpha \end{cases}$$

and map ∞ to the zero vector $\vec{0}$. \square

Corollary 1.2.4. $\mathcal{O} \not\uparrow_{k\text{-tact}} \text{Clus}_{\mathcal{O},P}^*(\Sigma\mathbb{R}^{\omega_1}, \vec{0})$, $\mathcal{O} \not\uparrow_{k\text{-mark}} \text{Con}_{\mathcal{O},P}^*(\Sigma\mathbb{R}^{\omega_1}, \vec{0})$, and $\mathcal{O} \not\uparrow_{k\text{-mark}} \text{Clus}_{\mathcal{O},P}^*(\Sigma\mathbb{R}^{\omega_2}, \vec{0})$.

◇

While this closes the question on tactics and marks for high dimensional sigma- (and Tychonoff-) products of the real line, there is another type of limited information strategy to investigate.

Definition 1.2.5. For a game $G = \langle M, W \rangle$ and *coding strategy* or *code* $\sigma : M^2 \rightarrow M$, the σ -coding fog-of-war $\gamma_\sigma : M^{<\omega} \rightarrow M^{\leq 2}$ is the function defined such that

$$\gamma_\sigma(\emptyset) = \emptyset$$

and

$$\gamma_\sigma(s \frown \langle x \rangle) = \langle \sigma \circ \gamma_\sigma(s), x \rangle$$

For a coding strategy σ , its corresponding strategy is $\sigma \circ \gamma_\sigma$. For a game G , if $\sigma \circ \gamma_\sigma$ is a winning strategy for \mathcal{A} , then σ is a winning coding strategy and we write $\mathcal{A} \uparrow_{\text{code}} G$. ◇

Intuitively, a σ -coding fog-of-war converts perfect information of the game into the last moves of both the player and her opponent, so a player has a winning coding strategy when she only needs to know the move of her opponent and her own last move. The term “coding” comes from the fact that a player may encode information about the history of the game into her own moves, and use this encoded information in later rounds.

Coding strategies have been studied since the earliest days of the Banach-Mazur game.

Theorem 1.2.6. (TODO: Cite the precise version of the BM game where a winning strategy implies a coding strategy.) ◇

The hard and normal versions of the W games are all equivalent with regards to coding strategies since \mathcal{O} may always ensure her new move is a subset of her previous move. For Fort spaces, the question is immediately closed.

Proposition 1.2.7. $\mathcal{O} \uparrow_{code} Con_{O,P}(\kappa^*, \infty)$. \diamond

Proof. Let $\sigma(\emptyset) = \kappa^*$, $\sigma(\langle U, \alpha \rangle) = U \setminus \{\alpha\}$ for $\alpha < \kappa$, and $\sigma(\langle U, \infty \rangle) = U$. \mathcal{P} cannot legally repeat non- ∞ points of the set, so her points converge to ∞ . \square

This trick does not simply extend to the $\Sigma\mathbb{R}^\kappa$ case, however. An open set may only restrict finitely many coordinates of the product, and a point in $\Sigma\mathbb{R}^\kappa$ may have countably infinite non-zero coordinates. Thus, information about the previous non-zero coordinates cannot be directly encoded into the open set.

Circumventing this takes a bit of extra bookkeeping. We proceed by defining a simpler infinite game for each cardinal κ .

Game 1.2.8. Let $PF_{F,C}(\kappa)$ denote the *point-finite game* with players \mathcal{F}, \mathcal{C} for each cardinal κ .

In round n , \mathcal{F} chooses $F_n \in [\kappa]^{<\omega}$, followed by \mathcal{C} choosing $C_n \in [\kappa \setminus \bigcup_{m \leq n} F_m]^{\leq \omega}$.

\mathcal{F} wins the game if the collection $\{C_n : n < \omega\}$ is a point-finite cover of its union $\bigcup_{n < \omega} C_n$, that is, each point in $\bigcup_{n < \omega} C_n$ is in C_n only for finitely many $n < \omega$. \diamond

Theorem 1.2.9. $\mathcal{F} \uparrow_{code} PF_{F,C}(\kappa)$ implies $\mathcal{O} \uparrow_{code} Con_{O,P}(\Sigma\mathbb{R}^\kappa, \vec{0})$. \diamond

Proof. Let σ be a coding strategy for \mathcal{F} in $PF_{F,C}(\kappa)$ such that $\sigma(F, C) \supset F$.

For $F \in [\kappa]^{<\omega}$ and $\epsilon > 0$ let $U(F, \epsilon)$ be the open set in \mathbb{R}^κ such that

$$\pi_\alpha(U(F, \epsilon)) = \begin{cases} (-\epsilon, \epsilon) & \alpha \in F \\ \mathbb{R} & \alpha \notin F \end{cases}$$

Note that F is uniquely identifiable given $U(F, \epsilon) \cap \Sigma\mathbb{R}^\kappa$.

For each point $x \in \Sigma\mathbb{R}^\kappa$ and $\epsilon \geq 0$, let $C_\epsilon(x) \in [\kappa]^{\leq \omega}$ such that $\alpha \in C_\epsilon(x)$ if and only if $|x(\alpha)| \geq \epsilon$.

TODO: Finish

\square

We aim to show that $\mathcal{F} \uparrow_{\text{code}} PF_{F,C}(\kappa)$ for all κ .

Theorem 1.2.10. *Let $cf([\kappa]^{\leq \omega}) = \kappa$. Then $\mathcal{F} \uparrow_{\text{code}} PF_{F,C}(\kappa)$.* \diamond

Proof. Let $c : \kappa \rightarrow [\kappa]^{\leq \omega}$ be a function such that for every $W \in [\kappa]^{\leq \omega}$, there exists $\gamma_W < \kappa$ with $W \subseteq c(\gamma_W)$.

Let $W \restriction n \in [\kappa]^n$ be a subset of $W \in [\kappa]^\omega$ such that $W \restriction n \subset W \restriction (n+1)$ and $\bigcup_{n < \omega} W \restriction n = W$.

Define the coding strategy σ as follows:

$$\sigma(\emptyset) = \emptyset$$

$$\sigma(\langle N, W \rangle) = N \cup (|N| + 1) \cup \{\gamma_W\} \cup \bigcup_{\alpha \in N} (c(\alpha) \restriction |N|)$$

Let $\langle W_0, W_1, \dots \rangle$ be a legal attack against σ , $N_0 = \sigma(\emptyset)$, and $N_{i+1} = \sigma(\langle N_i, W_i \rangle)$. Then $\langle N_0, W_0, N_1, W_1, \dots \rangle$ is the playthrough of the game corresponding to the coding strategy σ and attack a .

Let $\alpha \in W_i$ for some $i < \omega$. Then $\gamma_{W_i} \in N_n$ when $n > i$. Since $W_i \subseteq c(\gamma_{W_i})$, there exists some $n > i$ such that $\alpha \in c(\gamma_{W_i}) \restriction n$. Finally, we consider

$$N_{n+1} = \sigma(\langle N_n, W_n \rangle) \supseteq \bigcup_{\alpha \in N_n} (c(\alpha) \restriction |N_n|) \supseteq c(\gamma_{W_i}) \restriction |N_n|$$

and since $|N_n| \geq n$, we may conclude $\alpha \in N_{n+1}$. Thus $\alpha \notin W_j$ for $j > n$, and \mathcal{F} wins the game. \square

Theorem 1.2.11. *Let κ be the limit of cardinals κ_n such that $cf([\kappa_n]^{\leq \omega}) = \kappa_n$. Then $F \uparrow_{\text{code}} PF_{F,C}(\kappa)$.* \diamond

Theorem 1.2.12. *$F \uparrow_{\text{code}} PF_{F,C}(\kappa)$ for all cardinals κ .* \diamond

Corollary 1.2.13. *$O \uparrow_{\text{code}} Con_{O,P}(\Sigma \mathbb{R}^\kappa, \vec{0})$ for all cardinals κ .* \diamond

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