ARHANGELSKII'S α -PRINCIPLES AND SELECTION GAMES

STEVEN CLONTZ

Abstract. Arhangelskii's properties α_2 and α_4 defined for convergent sequences may be characterized in terms of Scheeper's selection principles. We generalize these results to hold for more general collections and consider these results in terms of selection games.

- The following characterizations were given as Definition 1 by Kocinac in [cite
- Kocinac selection principles related.

2

- **Definition 1.** Arhangelskii's α -principles $\alpha_i(\mathcal{A},\mathcal{B})$ are defined as follows for $i \in$ $\{1,2,3,4\}$. Let $A_n \in \mathcal{A}$ for all $n < \omega$; then there exists $B \in \mathcal{B}$ such that:
- α_1 : $A_n \cap B$ is cofinite in A_n for all $n < \omega$.
- α_2 : $A_n \cap B$ is infinite for all $n < \omega$.
- α_3 : $A_n \cap B$ is infinite for infinitely-many $n < \omega$.
- α_4 : $A_n \cap B$ is non-empty for infinitely-many $n < \omega$. 10
- When $(\mathcal{A}, \mathcal{B})$ is omitted, it is assumed that $\mathcal{A} = \mathcal{B}$ is the collection $\Gamma_{X,x}$ of 11
- sequences converging to some point $x \in X$, as introduced by Arhangelskii in [cite 12
- Arhangelskii frequency spectrum]. Provided A only contains infinite sets, it's easy 13
- to see that $\alpha_n(\mathcal{A}, \mathcal{B})$ implies $\alpha_{n+1}(\mathcal{A}, \mathcal{B})$. 14
- We aim to relate these to the following games. 15
- **Definition 2.** The selection game $G_1(\mathcal{A}, \mathcal{B})$ (resp. $G_{fin}(\mathcal{A}, \mathcal{B})$) is an ω -length
- game involving Players I and II. During round n, I chooses $A_n \in \mathcal{A}$, followed
- by II choosing $a_n \in A_n$ (resp. $F_n \in [A_n]^{<\aleph_0}$). Player II wins in the case that
- $\{a_n:n<\omega\}\in\mathcal{B}\ (\text{resp. }\bigcup\{F_n:n<\omega\}\in\mathcal{B}), \text{ and Player I wins otherwise.}$
- Such games are well-represented in the literature; see [cite Scheepers combi-20
- natorics ramsey] for example. We will also consider the similarly-defined games 21 $G_{<2}(\mathcal{A},\mathcal{B})$ (II chooses 0 or 1 points from each choice by I) and $G_{cf}(\mathcal{A},\mathcal{B})$ (II
- chooses cofinitely-many points). 23
- **Definition 3.** Let P be a player in a game G. P has a winning strategy for G,
- denoted $P \uparrow G$, if P has a strategy that defeats every possible counterplay by 25
- their opponent. If a strategy only relies on the round number and ignores the 26
- moves of the opponent, the strategy is said to be predetermined; the existence of a
- predetermined winning strategy is denoted $P \uparrow G$.
- We briefly note that the statement I $egthinspace{1mu}{f} G_{\star}(\mathcal{A}, \mathcal{B})$ is often denoted as the selection
- principle $S_{\star}(\mathcal{A}, \mathcal{B})$.

22

44

50

51

55

56

Definition 4. Let $\Gamma_{X,x}$ be the collection of non-trivial sequences $S \subseteq X$ converging to x, that is, infinite subsets of $X \setminus \{x\}$ such that for each neighborhood U of x, $S \cap U$ is cofinite in S.

Definition 5. Let Γ_X be the collection of open γ -covers \mathcal{U} of X, that is, infinite open covers of X such that $X \notin \mathcal{U}$ and for each $x \in X$, $\{U \in \mathcal{U} : x \in U\}$ is cofinite in \mathcal{U} .

The similarity in nomenclature follows from the observation that every nontrivial sequence in $C_p(X)$ converging to the zero function $\mathbf{0}$ naturally defines a corresponding γ -cover in X, see e.g. Theorem 4 of [Scheepers a sequential property and covering property].

The equivalence of $\alpha_2(\Gamma_{X,x}\Gamma_{X,x})$ and I $\underset{\text{pre}}{\gamma} G_1(\Gamma_{X,x},\Gamma_{X,x})$ was briefly asserted by

Sakai in the introduction of [cite Sakai sequence selection properties]; the similar equivalence of $\alpha_4(\Gamma_{X,x}\Gamma_{X,x})$ and I $\uparrow_{\text{pre}} G_{fin}(\Gamma_{X,x},\Gamma_{X,x})$ seems to be folklore. In

fact, these relationships hold in more generality.

Note that by these definitions, convergent sequences (resp. γ -covers) may be uncountable, but any infinite subset of either would remain a convergent sequence (resp. γ -cover), in particular, countably infinite subsets. We capture this idea as follows.

Definition 6. Say a collection \mathcal{A} is Γ-like if it satisfies the following for each $A \in \mathcal{A}$.

- $|A| \geq \aleph_0$.
 - If $A' \subseteq A$ and $|A'| \ge \aleph_0$, then $A' \in \mathcal{A}$.

We also require the following.

Definition 7. Say a collection \mathcal{A} is almost-Γ-like if for each $A \in \mathcal{A}$, there is $A' \subseteq A$ such that:

- $\bullet |A'| = \aleph_0.$
 - If A'' is a cofinite subset of A', then $A'' \in \mathcal{A}$.

7 So all Γ-like sets are almost-Γ-like.

We are now able to prove a few general equivalences between α -princples and selection games.

1. On
$$\alpha_2(\mathcal{A}, \mathcal{B})$$
 and $G_1(\mathcal{A}, \mathcal{B})$

Theorem 8. Let \mathcal{A} be almost- Γ -like and \mathcal{B} be Γ -like. Then $\alpha_2(\mathcal{A}, \mathcal{B})$ holds if and only if $\prod_{pre} G_1(\mathcal{A}, \mathcal{B})$.

Proof. We first assume $\alpha_2(\mathcal{A}, \mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined strategy for I. We may apply $\alpha_2(\mathcal{A}, \mathcal{B})$ to choose $B \in \mathcal{B}$ such that $|A_n \cap B| \ge \aleph_0$. We may then choose $a_n \in (A_n \cap B) \setminus \{a_i : i < n\}$ for each $n < \omega$. It follows that $B' = \{a_n : n < \omega\} \in \mathcal{B}$ since B' is an infinite subset of $B \in \mathcal{B}$; therefore A_n does not define a winning predetermined strategy for I.

Now suppose I \mathcal{V} $G_1(\mathcal{A}, \mathcal{B})$. Given $A_n \in \mathcal{A}$ for $n < \omega$, first choose $A'_n \in \mathcal{A}$ such

Now suppose I $\uparrow G_1(\mathcal{A}, \mathcal{B})$. Given $A_n \in \mathcal{A}$ for $n < \omega$, first choose $A'_n \in \mathcal{A}$ such

that $A'_n = \{a_{n,j} : j < \omega\} \subseteq A_n, \ j < k \text{ implies } a_{n,j} \neq a_{n,k}, \text{ and } A_{n,m} = \{a_{n,j} : m \leq j < \omega\} \in \mathcal{A}$. Finally choose some $\theta : \omega \to \omega$ such that $|\theta^{\leftarrow}(n)| = \aleph_0$ for each $n < \omega$.

Since playing $A_{\theta(m),m}$ during round m does not define a winning strategy for I in

 $G_1(\mathcal{A},\mathcal{B})$, II may choose $x_m \in A_{\theta(m),m}$ such that $B = \{x_m : m < \omega\} \in \mathcal{B}$. Choose

```
i_m < \omega for each m < \omega such that x_m = a_{\theta(m),i_m}, noting i_m \ge m. It follows that
      A_n \cap B \supseteq \{a_{\theta(m),i_m} : m \in \theta^{\leftarrow}(n)\}. Since for each m \in \theta^{\leftarrow}(n) there exists M \in A_n \cap B
      \theta^{\leftarrow}(n) such that m \leq i_m < M \leq i_M, and therefore a_{\theta(m),i_m} \neq a_{\theta(m),i_M} = a_{\theta(M),i_M},
      we have shown that A_n \cap B is infinite. Thus B witnesses \alpha_2(\mathcal{A}, \mathcal{B}).
          While \alpha_2(\mathcal{A}, \mathcal{B}) involves infinite intersection and G_1(\mathcal{A}, \mathcal{B}) involves single selec-
 78
      tions, the previous result is made more intuitive given the following result, shown
      for \mathcal{A} = \mathcal{B} = \Gamma_{X,x} by Nogura in [cite product of alpha spaces].
      Definition 9. \alpha'_2(\mathcal{A}, \mathcal{B}) is the following claim: if A_n \in \mathcal{A} for all n < \omega, then there
      exists B \in \mathcal{B} such that A_n \cap B is nonempty for all n < \omega.
 81
          (Note that \alpha_5 is sometimes used in the literature in place of \alpha'_2.)
 82
      Proposition 10. If A is almost-\Gamma-like, then \alpha_2(A, B) is equivalent to \alpha'_2(A, B).
 83
      Proof. The forward implication is immediate, so we assume \alpha'_2(\mathcal{A},\mathcal{B}). Given A_n \in
      \mathcal{A}, we apply the almost-\Gamma-like property to obtain A'_n = \{a_{n,m} : m < \omega\} \subseteq A_n such
      that A_{n,m} = A_n \setminus \{a_{i,j} : i, j < m\} \in \mathcal{A} \text{ for all } m < \omega.
          By applying \alpha'_2(\mathcal{A}, \mathcal{B}) to A_{n,m}, we obtain B \in \mathcal{B} such that A_{n,m} \cap B is nonempty
 87
      for all n, m < \omega. Since it follows that A_n \cap B is infinite for all n < \omega, we have
      established \alpha_2(\mathcal{A}, \mathcal{B}).
                                         2. On \alpha_4(\mathcal{A}, \mathcal{B}) and G_{fin}(\mathcal{A}, \mathcal{B})
          A similar correspondence exists between \alpha_4(\mathcal{A}, \mathcal{B}) and G_{fin}(\mathcal{A}, \mathcal{B}).
 91
      Theorem 11. Let A be almost-\Gamma-like and B be \Gamma-like. Then \alpha_4(A, B) holds if and
 92
      only if I \uparrow_{pre} G_{<2}(\mathcal{A}, \mathcal{B}) if and only if I \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B}).
 93
      Proof. We first assume \alpha_4(\mathcal{A}, \mathcal{B}) and let A_n \in \mathcal{A} for n < \omega define a predetermined
      strategy for I in G_{<2}(\mathcal{A},\mathcal{B}). We then may choose A'_n \in \mathcal{A} where A'_n = \{a_{n,j} : j < 1\}
 95
      \{\omega\} \subseteq A_n, j < k \text{ implies } a_{n,j} \neq a_{n,k}, \text{ and } A''_n = A'_n \setminus \{a_{i,j} : i, j < n\} \in \mathcal{A}.
 96
          By applying \alpha_4(\mathcal{A}, \mathcal{B}) to A_n'', we obtain B \in \mathcal{B} such that A_n'' \cap B \neq \emptyset for infintely-
 97
      many n < \omega. We then let F_n = \emptyset when A''_n \cap B = \emptyset, and F_n = \{x_n\} for some
      x_n \in A_n'' \cap B otherwise. Then we will have that B' = \bigcup \{F_n : n < \omega\} \subseteq B belongs
 99
      to \mathcal{B} once we show that B' is infinite. To see this, for m \leq n < \omega note that either
      F_m is empty (and we let j_m = 0) or F_m = \{a_{m,j_m}\} for some j_m \ge m; choose N < \omega
      such that j_m < N for all m \le n and F_N = \{x_N\}. Thus F_m \ne F_N for all m \le n
      since x_N \notin \{a_{i,j} : i, j < N\}. Thus II may defeat the predetermined strategy A_n by
103
      playing F_n each round.
          Since I \uparrow G_{<2}(\mathcal{A}, \mathcal{B}) immediately implies I \uparrow G_{fin}(\mathcal{A}, \mathcal{B}), we assume the latter.
105
      Given A_n \in \mathcal{A} for n < \omega, we note this defines a (non-winning) predetermined
106
      strategy for I, so II may choose F_n \in [A_n]^{<\aleph_0} such that B = \bigcup \{F_n : n < \omega\} \in \mathcal{B}.
      Since B is infinite, we note F_n \neq \emptyset for infinitely-many n < \omega. Thus B witnesses
108
      \alpha_4(\mathcal{A}, \mathcal{B}) since A_n \cap B \supseteq F_n \neq \emptyset for infinitely-many n < \omega.
109
          This shows that II gains no advantage from picking more than one point per
110
      round. This in fact only depends on \mathcal{B} being \Gamma-like, which we formalize in the
```

Theorem 12. Let \mathcal{B} be Γ -like. Then $I \uparrow_{pre} G_{<2}(\mathcal{A}, \mathcal{B})$ if and only if $I \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$.

following results.

112

125

126

127

128

129

130

131

132

134

135

136

139

140

142

143

144

145 146

Proof. Assume $\bigcup \mathcal{A}$ is well-ordered. Given a winning predetermined strategy A_n for I in $G_{<2}(\mathcal{A},\mathcal{B})$, consider $F_n \in [A_n]^{<\aleph_0}$. We set

$$F_n^* = \begin{cases} \emptyset & \text{if } F_n \setminus \bigcup \{F_m : m < n\} = \emptyset \\ \{\min(F_n \setminus \bigcup \{F_m : m < n\})\} & \text{otherwise} \end{cases}$$

Since $|F_n^*| < 2$, we have that $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$. In the case that $\bigcup \{F_n^* : n < \omega\}$ is finite, we immediately see that $\bigcup \{F_n : n < \omega\}$ is also finite and therefore not in \mathcal{B} . Otherwise $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$ is an infinite subset of $\bigcup \{F_n : n < \omega\}$, and thus 118 $\bigcup \{F_n : n < \omega\} \notin \mathcal{B}$ too. Therefore A_n is a winning predetermined strategy for I in $G_{fin}(\mathcal{A}, \mathcal{B})$ as well. 120

Theorem 13. Let \mathcal{B} be Γ -like. Then $I \uparrow G_{<2}(\mathcal{A}, \mathcal{B})$ if and only if $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$. 121

Proof. Assume $\bigcup A$ is well-ordered. Suppose $I \uparrow G_{<2}(A, \mathcal{B})$ is witnessed by the strategy σ . Let $\langle \rangle^* = \langle \rangle$, and for $s \cap \langle F \rangle \in ([\bigcup \mathcal{A}]^{<\aleph_0})^{<\omega} \setminus \{\langle \rangle\}$ let

$$(s^{\frown} \langle F \rangle)^{\star} = \begin{cases} s^{\star \frown} \langle \emptyset \rangle & \text{if } F \setminus \bigcup \text{range}(s) = \emptyset \\ s^{\star \frown} \langle \{\min(F \setminus \bigcup \text{range}(s))\} \rangle & \text{otherwise} \end{cases}$$

We then define the strategy τ for I in $G_{fin}(\mathcal{A}, \mathcal{B})$ by $\tau(s) = \sigma(s^*)$. Then given any counterattack $\alpha \in ([\bigcup \mathcal{A}]^{<\aleph_0})^{\omega}$ by II played against τ , we note that $\alpha^* =$ $\{(\alpha \upharpoonright n)^* : n < \omega\}$ is a counterattack to σ , and thus loses. This means $B = \{(\alpha \upharpoonright n)^* : n < \omega\}$ $||\operatorname{Jrange}(\alpha^*) \not\in \mathcal{B}.$

We consider two cases. The first is the case that $||\operatorname{Jrange}(\alpha^*)||$ is finite. Noting that $\alpha^*(m) \cap \alpha^*(n) = \emptyset$ whenever $m \neq n$, there exists $N < \omega$ such that $\alpha^*(n) = \emptyset$ for all n > N. As a result, $\bigcup \operatorname{range}(\alpha) = \bigcup \operatorname{range}(\alpha \upharpoonright n)$, and thus $\bigcup \operatorname{range}(\alpha)$ is finite, and therefore not in \mathcal{B} .

In the other case, $| \operatorname{Jrange}(\alpha^*) \notin \mathcal{B}$ is an infinite subset of $| \operatorname{Jrange}(\alpha)|$, and for I in $G_{fin}(\mathcal{A}, \mathcal{B})$.

We note that the above proof technique could be used to establish that perfectinformation and limited-information strategies for II in $G_{fin}(\mathcal{A},\mathcal{B})$ may be improved to be valid in $G_{<2}(\mathcal{A},\mathcal{B})$, provided \mathcal{B} is Γ -like. As such, $G_{<2}(\mathcal{A},\mathcal{B})$ and $G_{fin}(\mathcal{A},\mathcal{B})$ are effectively equivalent games under this hypothesis, so we will no longer consider $G_{<2}(\mathcal{A},\mathcal{B}).$

3. Perfect information and predetermined strategies

We now demonstrate the following, in the spirit of Pawlikowskii's celebrated result that a winning strategy for the first player in the Rothberger game may always be improved to a winning predetermined strategy [cite pawlikowskii].

Theorem 14. Let A be almost- Γ -like and B be Γ -like. Then

- I↑ G_{fin}(A,B) if and only if I↑ pre G_{fin}(A,B), and
 I↑ G₁(A,B) if and only if I↑ G₁(A,B).

Proof. We assume $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ and let the symbol \dagger mean $\langle \aleph_0 \rangle$ (respectively, $I \uparrow G_1(\mathcal{A}, \mathcal{B})$ and $\dagger = 1$, and for convenience we assume II plays singleton subsets of \mathcal{A} rather than elements). As \mathcal{A} is almost- Γ -like, there is a winning strategy σ

where $|\sigma(s)| = \aleph_0$ and $\sigma(s) \cap \bigcup \operatorname{range}(s) = \emptyset$ (that is, σ never replays the choices of II) for all partial plays s by II. 151

For each $s \in \omega^{<\omega}$, suppose $F_{s \mid m} \in [\bigcup A]^{\dagger}$ is defined for each $0 < m \le |s|$. Then let $s^*: |s| \to [\bigcup \mathcal{A}]^{\dagger}$ be defined by $s^*(m) = F_{s \mid m+1}$, and define $\tau': \omega^{<\omega} \to \mathcal{A}$ by $\tau'(s) = \sigma(s^*)$. Finally, set $[\sigma(s^*)]^{\dagger} = \{F_{s \cap \langle n \rangle} : n < \omega\}$, and for some bijection 154 $b:\omega^{<\omega}\to\omega$ let $\tau(n)=\tau'(b(n))$ be a predetermined strategy for I in $G_{fin}(\mathcal{A},\mathcal{B})$ (resp. $G_1(\mathcal{A}, \mathcal{B})$). 156

Suppose α is a counterattack by II against τ , so

157

$$\alpha(n) \in [\tau(n)]^{\dagger} = [\tau'(b(n))]^{\dagger} = [\sigma(b(n)^{\star})]^{\dagger}$$

It follows that $\alpha(n) = F_{b(n) \cap \langle m \rangle}$ for some $m < \omega$. In particular, there is some infinite subset $W \subseteq \omega$ and $f \in \omega^{\omega}$ such that $\{\alpha(n) : n \in W\} = \{F_{f \upharpoonright n+1} : n < \omega\}$. 158 159 Note here that $(f \upharpoonright n+1)^* = (f \upharpoonright n)^* \cap \langle F_{f \upharpoonright n+1} \rangle$. This shows that $F_{f \upharpoonright n+1} \in [\sigma((f \upharpoonright n+1)^*)]$ $[n]^*$)] is an attempt by II to defeat σ , which fails. Thus $\bigcup \{F_{f \mid n+1} : n < \omega\} = 0$ 161 $\bigcup \{\alpha(n) : n \in W\} \notin \mathcal{B}$, and since this set is infinite (as σ prevents II from repeating choices) we have $\bigcup \{\alpha(n) : n < \omega\} \notin \mathcal{B}$ too. Therefore τ is winning. 163

Note that the assumption in Theorem 14 that \mathcal{A} be almost- Γ -like cannot be 164 omitted. In [todo cite Clontz k-tactics in Gruenhage game] an example of a space 165 X^* and point $\infty \in X^*$ where $I \uparrow G_1(\mathcal{A}, \mathcal{B})$ but $I \not \uparrow G_1(\mathcal{A}, \mathcal{B})$ is given, where \mathcal{A} is the 166

set of open neighborhoods of ∞ (which are all uncountable), and \mathcal{B} is the set $\Gamma_{X^*,\infty}$ 167 of sequences converging to that point. (Note that $G_1(\mathcal{A},\mathcal{B})$ is called $Gru_{O,P}(X^*,\infty)$ 168 in that paper, and an equivalent game $Gru_{K,P}(X)$ is what is directly studied. In fact, more is shown: I has a winning perfect-information strategy, but for any 170 natural number k, any strategy that only uses the most recent k moves of II and 171 the round number can be defeated.) 172

While A is often not almost- Γ -like in general, it may satisfy that property in 173 combination with the selection principles being considered. 174

Proposition 15. Let \mathcal{B} be Γ -like, $\mathcal{B} \subseteq \mathcal{A}$, and $I \underset{pre}{\gamma} G_{fin}(\mathcal{A}, \mathcal{B})$. Then \mathcal{A} is almost-175 176

Proof. Let $A \in \mathcal{A}$, and for all $n < \omega$ let $A_n = A$. Then A_n is not a winning 177 predetermined strategy for I, so II may choose finite sets $B_n \subseteq A_n = A$ such that 178 $A' = \bigcup \{B_n : n < \omega\} \in \mathcal{B} \subseteq \mathcal{A}.$ 179

It follows that $A' \subseteq A$ and $|A'| = \aleph_0$, and for any infinite subset $A'' \subseteq A'$ (in particular, any cofinite subset), $A'' \in \mathcal{B} \subseteq \mathcal{A}$. Thus \mathcal{A} is almost- Γ -like.

Note that in the previous result, I $\gamma G_{fin}(\mathcal{A}, \mathcal{B})$ could be weakened to the choice 182

principle $\binom{\mathcal{A}}{\mathcal{B}}$: for every member of \mathcal{A} , there is some countable subset belonging to 183 184

Corollary 16. Let \mathcal{B} be Γ -like and $\mathcal{B} \subseteq \mathcal{A}$. Then 185

- I \(\gamma G_{fin}(\mathcal{A}, \mathcal{B})\) if and only if I \(\gamma G_{fin}(\mathcal{A}, \mathcal{B})\), and
 I \(\gamma G_1(\mathcal{A}, \mathcal{B})\) if and only if I \(\gamma G_1(\mathcal{A}, \mathcal{B})\).

Proof. Assuming I $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$, we have I $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ by Proposition 15 and 188

Theorem 14.

180

181

186 187

197

198

223

```
Similarly, assuming I \gamma G_1(\mathcal{A}, \mathcal{B}) \Rightarrow I \gamma G_{fin}(\mathcal{A}, \mathcal{B}), we have I \gamma G_1(\mathcal{A}, \mathcal{B}) by
190
       Proposition 15 and Theorem 14.
191
```

This corollary generalizes e.g. Theorems 26 and 30 of [cite Scheepers 1996 Ram-192 sey], Theorem 5 of [cite MR2119791], and Corollary 36 of [cite Clontz dual games]. 193 In summary, using the selection principle notation $S_{\star}(\mathcal{A},\mathcal{B})$: 194

Corollary 17. Let \mathcal{B} be Γ -like and $\mathcal{B} \subseteq \mathcal{A}$. Then 195

- I $\gamma G_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $S_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $\alpha_2(\mathcal{A}, \mathcal{B})$, and
- I $\not\uparrow G_1(\mathcal{A}, \mathcal{B})$ if and only if $S_1(\mathcal{A}, \mathcal{B})$ if and only if $\alpha_4(\mathcal{A}, \mathcal{B})$.

4. Disjoint selections

In each $\alpha_i(\mathcal{A}, \mathcal{B})$ principle, it is not required for the collection $\{A_n : n < \omega\}$ to 199 be pairwise disjoint. However, in many cases it may as well be. 200

Definition 18. For $i \in \{1, 2, 3, 4\}$ let $\alpha_{i,1}(\mathcal{A}, \mathcal{B})$ denote the claim that $\alpha_i(\mathcal{A}, \mathcal{B})$ 201 202 holds provided the collection $\{A_n : n < \omega\}$ is pairwise disjoint.

Of course, $\alpha_i(\mathcal{A}, \mathcal{B})$ implies $\alpha_{i,1}(\mathcal{A}, \mathcal{B})$. It's also immediate that $\alpha_{i,1}(\mathcal{A}, \mathcal{B})$ implies 203 $\alpha_{i,1+1}(\mathcal{A},\mathcal{B})$ for the same reason that $\alpha_i(\mathcal{A},\mathcal{B})$ implies $\alpha_{i+1}(\mathcal{A},\mathcal{B})$. 204

We take advantage of the following lemma. 205

Lemma 19 (Lemma 1.2 of (cite Nyikos 92)). Given a family $\{A_n : n < \omega\}$ of 206 infinite sets, there exist infinite subsets $A'_n \subseteq A_n$ such that $\{A'_n : n < \omega\}$ is pairwise 207 208

Proposition 20. Let \mathcal{A} be Γ -like. For $i \in \{2,3,4\}$, $\alpha_i(\mathcal{A},\mathcal{B})$ is equivalent to 209 210 $\alpha_{i,1}(\mathcal{A},\mathcal{B})$

Proof. Assume $\alpha_{i,1}(\mathcal{A},\mathcal{B})$. Let $A_n \in \mathcal{A}$. By applying the previous lemma, we have $\{A'_n:n<\omega\}$ pairwise disjoint with each A'_n being an infinite subset of A_n . Since Ais Γ -like, $A'_n \in \mathcal{A}$, so we have a witness $B \in \mathcal{B}$ such that $A'_n \cap B$ satisfies $\alpha_{i,1}(\mathcal{A},\mathcal{B})$ for all $n < \omega$. Since $A'_n \subseteq A_n$, it follows that $A_n \cap B$ satisfies $\alpha_i(\mathcal{A}, \mathcal{B})$ for all 214 215

It's also true that $\alpha_1(\Gamma_{X,x},\Gamma_{X,x})$ is equivalent to $\alpha_{1.1}(\Gamma_{X,x},\Gamma_{X,x})$, which is cap-216 tured by the following theorem. 217

Proposition 21. Let A be a Γ -like collection closed under finite unions. Then 218 219 $\alpha_1(\mathcal{A}, \mathcal{A})$ is equivalent to $\alpha_{1,1}(\mathcal{A}, \mathcal{A})$.

Proof. Let $A_n \in \mathcal{A}$ and assume $\alpha_{1.1}(\mathcal{A}, \mathcal{A})$. To apply the assumption, we will define a pairwise disjoint collection $\{A'_n : n < \omega\}$. First let 0' = 0 and $A'_0 = A_0$. Then 221 suppose $m' \ge m$ and $A'_m \subseteq A_{m'} \subseteq \bigcup_{i \le m} A'_i$ are defined for all $m \le n$. If $A_k \setminus \bigcup_{m \le n} A'_m$ is finite for k > n', let $B = \bigcup_{m \le n'} A_m \in \mathcal{A}$. This B then 222

witnesses $\alpha_1(\bar{\mathcal{A}}, \mathcal{A})$ since $A_k \setminus B$ is finite for all $k < \omega$.

224 Otherwise pick the minimal (n+1)' > n where $A'_{n+1} = A_{(n+1)'} \setminus \bigcup_{m \leq n} A'_m$ is 225 infinite. It follows that $A'_{n+1} \subseteq A_{(n+1)'} \subseteq \bigcup_{m \le n+1} A'_m$. By construction, $\{A'_n : n < a \le n \le n \le n \}$ 226 ω } is a pairwise disjoint collection of members of \mathcal{A} , and we may apply $\alpha_{1,1}(\mathcal{A},\mathcal{A})$ 227 to obtain $B \in \mathcal{A}$ where $A'_n \setminus B$ is finite for all $n < \omega$. 228

Finally let $k < \omega$. If k = n' for some $n < \omega$, then $A_k \setminus B = A_{n'} \setminus B \subseteq$ $(\bigcup_{m \le n} A'_m) \setminus B$ is finite. Otherwise, n' < k < (n+1)' for some $n < \omega$. Then

```
231 (A_k \setminus \bigcup_{m \le n} A'_m) \setminus B \subseteq A_k \setminus \bigcup_{m \le n} A'_m is finite, and (A_k \cap \bigcup_{m \le n} A'_m) \setminus B \subseteq
232 (\bigcup_{m \le n} A'_m) \setminus B is finite, showing A_k \setminus B is finite.
```

TODO add discussion of $\alpha_{1.5}$ here.

233

246

247

256

263

266

For this paragraph we adopt the conventional assumption that $\Gamma_{X,x}$ is restricted to countable sets. Nyikos showed a consistent example where $\alpha_2(\Gamma_{X,x},\Gamma_{X,x})$ fails to imply $\alpha_{1.5}(\Gamma_{X,x},\Gamma_{X,x})$, and a consistent example where $\alpha_{1.5}(\Gamma_{X,x},\Gamma_{X,x})$ fails to imply $\alpha_1(\Gamma_{X,x},\Gamma_{X,x})$ [cite Nyikos 92]. On the other hand, Dow showed that $\alpha_2(\Gamma_{X,x},\Gamma_{X,x})$ implies $\alpha_1(\Gamma_{X,x},\Gamma_{X,x})$ in the Laver model for the Borel conjecture [cite Dow 1990]; the author conjectures that this model (specifically, the fact that every ω -splitting family contains an ω -splitting family of size less than $\mathfrak b$ in this model) witnesses an affirmative answer to the following question.

Definition 22. A Γ-like collection is strongly-Γ-like if the collection is closed under finite unions and each member is countable.

Question 23. Let A, B be strongly- Γ -like and $B \subseteq A$. Is it consistent that $\alpha_2(A, B)$ implies $\alpha_1(A, B)$?

5. Conclusion

We conclude with the following easy result, and a couple questions.

Proposition 24. Let \mathcal{B} be Γ-like. Then $\alpha_1(\mathcal{A},\mathcal{B})$ holds if and only if $\prod_{pre} G_{cf}(\mathcal{A},\mathcal{B})$.

Proof. We first assume $\alpha_1(\mathcal{A},\mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined strategy for I. By $\alpha_1(\mathcal{A},\mathcal{B})$, we immediately obtain $B \in \mathcal{B}$ such that $|A_n \setminus B| < \aleph_0$.

Thus $B_n = A_n \cap B$ is a cofinite choice from A_n , and $B' = \bigcup \{B_n : n < \omega\}$ is an infinite subset of B, so $B' \in \mathcal{B}$. Thus II may defeat I by choosing $B_n \subseteq A_n$ each round, witnessing I \mathcal{Y} $G_{cf}(\mathcal{A},\mathcal{B})$.

On the other hand, let I $\gamma G_{cf}(\mathcal{A}, \mathcal{B})$. Given $A_n \in \mathcal{A}$ for $n < \omega$, we note that II may choose a cofinite subset $B_n \subseteq A_n$ such that $B = \bigcup \{B_n : n < \omega\} \in \mathcal{B}$. Then

Question 25. Is there a game-theoretic characterization of $\alpha_3(\mathcal{A},\mathcal{B})$?

B witnesses $\alpha_1(\mathcal{A}, \mathcal{B})$ since $|A_n \setminus B| \leq |A_n \setminus B_n| \leq \aleph_0$.

Noting that I $\uparrow G_1(\Gamma_X, \Gamma_X)$ if and only if I $\uparrow G_{fin}(\Gamma_X, \Gamma_X)$ [cite Kocinac], but the same is not true of $G_{\star}(\Gamma_{X,x}, \Gamma_{X,x})$ (i.e. there are α_4 spaces that are not α_2 [cite Shakhmatov convergence algebraic structure]), we also ask the following.

Question 26. Is there a natural condition on \mathcal{A}, \mathcal{B} guaranteeing $I \uparrow G_1(\mathcal{A}, \mathcal{B}) \Rightarrow$ 1 $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$?

6. Acknowledgements

The author would like to thank Dr. Alan Dow and Dr. Jared Holshouser for their discussions on the subject of the paper.

References

Department of Mathematics and Statistics, The University of South Alabama, Mo- $_{268}$ $_{\rm BILE,\ AL\ 36688}$

Email address: sclontz@southalabama.edu