# Game-theoretic strengthenings of Menger's property

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## **Abstract**

A certain topological game introduced by Hurewicz characterizes Menger's covering property whenever the first player lacks a winning strategy. It was later shown by Telgarsky (and later by Scheepers using a different argument) that for metric spaces, the second player having a winning strategy characterizes the stronger property of  $\sigma$ -compactness.

We factor out Scheepers' proof to show that for regular spaces, the second player having a winning 1-Marköv strategy characterizes  $\sigma$ -compactness; also, for second-countable spaces, the presence of a winning perfect information strategy for the second player implies the existence of a winning 1-Marköv strategy for that player. We also show the that for all k, the existence of a winning k-Marköv strategy for the second player implies the existence of a winning 2-Marköv strategy, and give a space which permits a winning 2-Marköv strategy but not a winning 1-Marköv strategy.

# The Menger property

## Definition

A space X is Menger if for every sequence  $\langle \mathcal{U}_0, \mathcal{U}_1, \ldots \rangle$  of open covers of X there exists a sequence  $\langle \mathcal{F}_0, \mathcal{F}_1, \ldots \rangle$  such that  $\mathcal{F}_n \subseteq \mathcal{U}_n, \, |\mathcal{F}_n| < \omega$ , and  $\bigcup_{n < \omega} \mathcal{F}_n$  is a cover of X.

## Proposition

*X* is  $\sigma$ -compact  $\Rightarrow$  *X* is Menger  $\Rightarrow$  *X* is Lindelöf.

# The Menger game

## Game

Let  $Cov_{C,F}(X)$  denote the *Menger game* with players  $\mathscr{C}$ ,  $\mathscr{F}$ . In round n,  $\mathscr{C}$  chooses an open cover  $\mathcal{C}_n$ , followed by  $\mathscr{F}$  choosing a finite subcollection  $\mathcal{F}_n \subseteq \mathcal{C}_n$ .

 $\mathscr{F}$  wins the game, that is,  $\mathscr{F} \uparrow Cov_{C,F}(X)$  if  $\bigcup_{n<\omega} \mathcal{F}_n$  is a cover for the space X, and  $\mathscr{C}$  wins otherwise.

#### **Theorem**

*X* is Menger if and only if  $\mathscr{C} \uparrow Cov_{C,F}(X)$ . (Hurewicz 1926, effectively)

Menger suspected that the subsets of the real line with his property were exactly the  $\sigma$ -compact spaces; however:

#### Theorem

There are ZFC examples of non- $\sigma$ -compact subsets of the real line which are Menger. (Fremlin, Miller 1988)

But metrizable non- $\sigma$ -compact Menger spaces will be *undetermined* for the Menger game.

#### Theorem

Let X be metrizable.  $\mathscr{F} \uparrow Cov_{C,F}(X)$  if and only if X is  $\sigma$ -compact. (Telgarsky 1984, Scheepers 1995)

Note that for Lindelöf spaces, metrizability is characterized by regularity and secound countability.

## Sketch of Scheeper's proof:

- Using second-countability and the winning strategy for  $\mathscr{F}$ , construct certain subsets  $R_s$  for  $s \in \omega^{<\omega}$  such that  $X = \bigcup_{s \in \omega < \omega}$ .
- Using regularity, show that each R<sub>s</sub> is compact.
- The result follows since  $|\omega^{<\omega}| = \omega$ .

By considering winning *limited-information strategies*, we'll be able to factor out this proof a bit.



# Limited information strategies

#### **Definition**

A (perfect information) strategy has knowledge of all the past moves of the opponent.

## Definition

A *k-tactical strategy* has knowledge of only the past *k* moves of the opponent.

## Definition

A *k-Marköv strategy* has knowledge of only the past *k* moves of the opponent and the round number.



## Obviously,

$$\mathscr{A} \underset{k\text{-tact}}{\uparrow} G \Rightarrow \mathscr{A} \underset{k\text{-mark}}{\uparrow} G \Rightarrow \mathscr{A} \underset{\text{(perfect)}}{\uparrow} G$$

But tactical strategies aren't interesting for the Menger game.

## Proposition

For any 
$$k < \omega$$
,  $\mathscr{F} \underset{k\text{-tact}}{\uparrow} Cov_{C,F}(X)$  if and only if  $X$  is compact.

Effectively,  $\mathscr{F}$  needs some sort of seed to prevent from being stuck in a loop: there's nothing stopping  $\mathscr{C}$  from playing the same open cover during every round of the game.

Comparitively, Marköv strategies are very powerful.

## **Proposition**

If X is 
$$\sigma$$
-compact, then  $\mathscr{F} \underset{1-mark}{\uparrow} Cov_{C,F}(X)$ .

## Proof.

Let  $X = \bigcup_{n < \omega} K_n$ . During round n,  $\mathscr{F}$  picks a finite subcollection of the last open cover played by  $\mathscr{C}$  (the only one  $\mathscr{F}$  remembers) which covers  $K_n$ .

Without assuming regularity, we can't quite reverse the implication, but we can get close.

## Definition

A subset Y of X is *relatively compact* if for every open cover for X, there exists a finite subcollection which covers Y.

## Proposition

If X is  $\sigma$ -relatively-compact, then  $\mathscr{F} \underset{1-mark}{\uparrow} Cov_{C,F}(X)$ .

## Proposition

For regular spaces,  $Y \subseteq X$  is relatively compact if and only if  $\overline{Y}$  is compact. So  $\sigma$ -relatively-compact regular spaces are exactly the  $\sigma$ -compact regular spaces.

### Theorem

#### Proof.

Let  $\sigma(\mathcal{U}, n)$  represent a 1-Marköv strategy. For every open cover  $\mathcal{U} \in \mathfrak{C}$ ,  $\sigma(\mathcal{U}, n)$  witnesses relative compactness for the set

$$R_n = \bigcap_{\mathcal{U} \in \mathfrak{C}} \bigcup \sigma(\mathcal{U}, n)$$

If X is not  $\sigma$ -relatively compact, fix  $x \notin R_n$  for any  $n < \omega$ . Then  $\mathscr C$  can beat  $\sigma$  by choosing  $\mathcal U_n \in \mathfrak C$  during each round such that  $x \notin \bigcup \sigma(\mathcal U_n, n)$ .

So for regular spaces, a winning strategy for  $\mathscr{F}$  in the Menger game isn't sufficient to characterize  $\sigma$ -compactness, but a winning 1-Marköv strategy does the trick.

We can complete Telgarsky's/Scheeper's result by showing the following:

#### Theorem

For second countable spaces X,  $\mathscr{F} \uparrow Cov_{C,F}(X)$  if and only if  $\mathscr{F} \uparrow Cov_{C,F}(X)$ .

## **Proof**

Let  $\sigma$  be a perfect information strategy. Since X is a second-countable space, we may pretend that there are only countably many finite collections of open sets. Thus for  $s \in \omega^{<\omega}$ , we may define open covers  $\mathcal{U}_{s \frown \langle n \rangle}$  such that for each open cover  $\mathcal{U}$ , there is some  $n < \omega$  where

$$\sigma(\mathcal{U}_{s \upharpoonright 1}, \dots, \mathcal{U}_{s}, \mathcal{U}) = \sigma(\mathcal{U}_{s \upharpoonright 1}, \dots, \mathcal{U}_{s}, \mathcal{U}_{s \frown \langle n \rangle})$$

Let  $t: \omega \to \omega^{<\omega}$  be a bijection. During round n and seeing only the latest open cover  $\mathcal{U}$ ,  $\mathscr{F}$  may play the finite subcollection

$$\tau(\mathcal{U}, n) = \sigma(\mathcal{U}_{t(n) \mid 1}, \dots, \mathcal{U}_{t(n)}, \mathcal{U})$$



# Proof (cont.)

Suppose there exists a counter-attack  $\langle \mathcal{V}_0, \mathcal{V}_1, \ldots \rangle$  which defeats the 1-Marköv strategy  $\tau$ . Then there exists  $f: \omega \to \omega$  such that, if  $\mathcal{V}^n = \mathcal{V}_{t^{-1}(f \upharpoonright n)}$ 

$$\begin{array}{ll}
x & \notin & \bigcup \tau(\mathcal{V}^n, t^{-1}(f \upharpoonright n)) \\
& = & \bigcup \sigma(\mathcal{U}_{f \upharpoonright 1}, \dots, \mathcal{U}_{f \upharpoonright n}, \mathcal{V}^n) \\
& = & \bigcup \sigma(\mathcal{U}_{f \upharpoonright 1}, \dots, \mathcal{U}_{f \upharpoonright n}, \mathcal{U}_{f \upharpoonright (n+1)})
\end{array}$$

Thus  $\langle \mathcal{U}_{f|1}, \mathcal{U}_{f|2}, \ldots \rangle$  is a successful counter-attack by  $\mathscr{C}$  against the perfect information strategy  $\sigma$ .

Unlike the Banach-Mazur game, we can immediately see that knowledge of more than two previous moves of  $\mathscr{F}$ 's opponent must be infinite to be of any use.

### Theorem

If 
$$\mathscr{F} \underset{k\text{-mark}}{\uparrow} Cov_{C,F}(X)$$
, then  $\mathscr{F} \underset{2\text{-mark}}{\uparrow} Cov_{C,F}(X)$ .

## Proof.

$$\tau(\langle \mathcal{U}, \mathcal{V} \rangle, n+1) = \bigcup_{m < k+2} \sigma(\langle \underbrace{\mathcal{U}, \dots, \mathcal{U}}_{k+1-m}, \underbrace{\mathcal{V}, \dots, \mathcal{V}}_{m+1} \rangle, (n+1)(k+2) + m)$$

foo

Introduction
Menger Spaces and the Menger Game
1-Marköv Strategies
k-Marköv strategies for k > 2

Questions? Thanks for having me!

