**Definition 1.** A uniform space  $\langle X, \mathcal{D} \rangle$  is a set X paired with a filter  $\mathcal{D}$  (called its uniformity) of relations (called **entourages**) on X such that for each entourage  $D \in \mathcal{D}$ :

- D is reflexive, i.e., the diagonal  $\Delta \subseteq D$ .
- Its inverse  $D^{-1} = \{ \langle y, x \rangle : \langle x, y \rangle \in D \} \in \mathcal{D}$ .
- There exists  $\frac{1}{2}D \in \mathcal{D}$  such that

$$2(\frac{1}{2}D) = \frac{1}{2}D \circ \frac{1}{2}D = \{\langle x, z \rangle : \exists y(\langle x, y \rangle, \langle y, z \rangle \in \frac{1}{2}D)\} \subseteq D$$

Note that since  $\mathcal{D}$  is a filter, for each  $D \in \mathcal{D}$ , the symmetric relation  $D \cap D^{-1} \in \mathcal{D}$ .

**Proposition 2.** For each  $D \in \mathcal{D}$  and  $n < \omega$  there exists  $\frac{1}{2^{n+1}}D \in \mathcal{D}$  such that

$$2(\frac{1}{2^{n+1}}D) = \frac{1}{2^{n+1}}D \circ \frac{1}{2^{n+1}}D \subseteq \frac{1}{2^n}D$$

and if  $2E \subseteq \frac{1}{2^n}D$ , then  $E \subseteq \frac{1}{2^{n+1}}D$ .

**Definition 3.** For an entourage  $D \in \mathcal{D}$ , let  $D[x] = \{y : (x,y) \in D\}$  be the D-neighborhood of x. The uniform topology for a uniform space  $\langle X, \mathcal{D} \rangle$  is generated by the base  $\{D[x] : x \in X, D \in \mathcal{D}\}$ .

**Theorem 4.** A space X is uniformizable (its topology is the uniform topology for some uniformity) if and only if X is completely regular  $(T_{3\frac{1}{\alpha}})$ .

**Proposition 5.** If X is a uniform space, then for all  $x \in X$  and symmetric entourages D:

$$x \in \frac{1}{2}D[y] \text{ and } y \in \frac{1}{2}D[z] \Rightarrow x \in D[z]$$

and

$$\frac{1}{2}D[x]\subseteq\overline{\frac{1}{2}D[x]}\subseteq D[x]$$

*Proof.* The first is by definition of  $\frac{1}{2}D$ .

If  $z \in \overline{\frac{1}{2}D[x]}$ , it follows that there is  $y \in \overline{\frac{1}{2}D[x]} \cap \overline{\frac{1}{2}D[z]}$  since  $\overline{\frac{1}{2}D[z]}$  is an open neighborhood of z. Thus  $(x,z) \in D \Rightarrow z \in D[x] \Rightarrow \overline{\frac{1}{2}D[x]} \subseteq D[x]$ .

**Definition 6.** For a uniform space X, Bell's proximity game proceeds as follows.

In round 0,  $\mathscr{D}$  chooses an entourage  $D_0$ , followed by  $\mathscr{P}$  choosing a point  $p_0 \in X$ .

In round n+1,  $\mathscr{D}$  chooses an entourage  $D_{n+1} \subseteq D_n$ , followed by  $\mathscr{P}$  choosing a point  $p_{n+1} \in 4D_n[p_n]$ .

Player  $\mathscr{D}$  wins if either  $\bigcap_{n<\omega} 4D_n[p_n] = \emptyset$  or  $\langle p_0, p_1, \ldots \rangle$  converges.

proximity.tex - Updated on February 27, 2014

**Definition 7.** For a uniform space X, the simplified proximal game  $Prox_{D,P}(X)$  can be defined as follows:

In round 0,  $\mathscr{D}$  chooses a symmetric entourage  $D_0$ , followed by  $\mathscr{P}$  choosing a point  $p_0 \in X$ .

In round n+1,  $\mathscr{D}$  chooses a symmetric entourage  $D_{n+1}$ , followed by  $\mathscr{P}$  choosing a point  $p_{n+1} \in \left(\bigcap_{m \leq n} D_m\right)[p_n]$ .

Player 
$$\mathscr{D}$$
 wins if either  $\bigcap_{n<\omega}\left(\bigcap_{m\leq n}D_m\right)[p_n]=\emptyset$  or  $\langle p_0,p_1,\ldots\rangle$  converges.

**Theorem 8.**  $\mathscr{D}$  has a winning perfect-information strategy in Bell's game if and only if  $\mathscr{D} \uparrow Prox_{D,P}(X)$ .

*Proof.* Let  $\sigma$  be a winning perfect information strategy for  $\mathscr{D}$  in Bell's game. We define a perfect information strategy  $\tau$  in the simplified game to yield symmetric entourages  $\tau(p \upharpoonright n) = \sigma(p \upharpoonright n) \cap (\sigma(p \upharpoonright n))^{-1}$  for all partial attacks  $p \upharpoonright n$ . Note that  $\tau(p \upharpoonright n) = \bigcap_{m \le n} \tau(p \upharpoonright m)$ .

If p attacks  $\tau$  in the simplified game,  $p(n+1) \in \left(\bigcap_{m \leq n} \tau(p \upharpoonright m)\right)[p(n)] = \tau(p \upharpoonright n)[p(n)] \subseteq \sigma(p \upharpoonright n)[p(n)] \subseteq 4\sigma(p \upharpoonright n)[p(n)]$ , so p attacks  $\sigma$  in Bell's game. Thus either p converges, or

$$\emptyset = \bigcap_{n < \omega} 4\sigma(p \upharpoonright n)[p(n)] \supseteq \bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)] = \bigcap_{n < \omega} \left(\bigcap_{m \le n} \tau(p \upharpoonright n)\right)[p(n)]$$

For the other direction, let  $\sigma$  be a winning perfect information strategy for  $\mathscr{D}$  in the simplified game such that  $\sigma(p \upharpoonright n) = \bigcap_{m \leq n} \sigma(p \upharpoonright m)$ . Define the perfect information strategy  $\tau$  in Bell's Game such that  $4\tau(p \upharpoonright n) \subseteq \sigma(p \upharpoonright n)$  and  $\tau(p \upharpoonright n) = \bigcap_{m \leq n} \tau(p \upharpoonright m)$  for all partial attacks  $p \upharpoonright n$ .

If p attacks  $\tau$  in Bell's game,  $p(n) \in 4\tau(p \upharpoonright n) \subseteq \sigma(p \upharpoonright n) = \bigcap_{m \le n} \sigma(p \upharpoonright m)$ , so p attacks  $\sigma$  in the simplified game. Thus either p converges, or

$$\emptyset = \bigcap_{n < \omega} \left( \bigcap_{m \le n} \sigma(p \upharpoonright n) \right) [p(n)] = \bigcap_{n < \omega} \sigma(p \upharpoonright n) [p(n)] \supseteq \bigcap_{n < \omega} 4\tau(p \upharpoonright n) [p(n)] \supseteq \bigcap_{n < \omega} \tau(p \upharpoonright n) [p(n)]$$

**Proposition 9.**  $\mathscr{P}$  has a winning perfect-information strategy in Bell's game if and only if  $\mathscr{P} \uparrow Prox_{D,P}(X)$ .

*Proof.* Similar to the previous. 
$$\Box$$

**Definition 10.** A uniform space is **proximal** if  $\mathcal{D} \uparrow Prox_{D,P}(X)$ .

**Definition 11.** For a space X and a point  $x \in X$ , the W-convergence-game  $Con_{O,P}(X,x)$  proceeds as follows.

In round 0,  $\mathscr{O}$  chooses a neighborhood  $U_n$  of x, followed by  $\mathscr{P}$  choosing a point  $p_n \in \bigcap_{m \leq n} U_m$ .

Player  $\mathcal{O}$  wins if  $\langle p_0, p_1, \ldots \rangle$  converges.

**Definition 12.** A space is W if  $\mathcal{O} \uparrow Con_{O,P}(X,x)$  for all  $x \in X$ .

**Definition 13.** For each finite tuple  $(m_0, \ldots, m_{n-1})$ , we define the k-tactical fog-of-war

$$T_k(\langle m_0,\ldots,m_{n-1}\rangle) = \langle m_{n-k},\ldots,m_{n-1}\rangle$$

and the k-Marköv fog-of-war

$$M_k(\langle m_0, \dots, m_{n-1} \rangle) = \langle \langle m_{n-k}, \dots, m_{n-1} \rangle, n \rangle$$

So  $P \uparrow G$  if and only if there exists a winning strategy for P of the form  $\sigma \circ T_k$ , and  $P \uparrow G$  if and only if there exists a winning strategy of the form  $\sigma \circ M_k$ .

**Theorem 14.** For all  $x \in X$ :

- $\mathscr{D} \uparrow Prox_{D,P}(X) \Rightarrow \mathscr{O} \uparrow Con_{O,P}(X,x)$
- $\bullet \ \mathscr{D} \ {\underset{2k-tact}{\uparrow}} \ Prox_{D,P}(X) \Rightarrow \mathscr{O} \ {\underset{k-tact}{\uparrow}} \ Con_{O,P}(X,x)$
- $\mathscr{D} \underset{2k-mark}{\uparrow} Prox_{D,P}(X) \Rightarrow \mathscr{O} \underset{k-mark}{\uparrow} Con_{O,P}(X,x)$

Proof. Let  $\sigma$  witness  $\mathscr{D} \underset{2k\text{-tact}}{\uparrow} Prox_{D,P}(X)$  (resp.  $\mathscr{D} \underset{2k\text{-mark}}{\uparrow} Prox_{D,P}(X), \mathscr{D} \uparrow Prox_{D,P}(X)$ ). We define the k-tactical (resp. k-Marköv, perfect info) strategy  $\tau$  such that

$$\tau \circ L_k(p) = \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(|p|-1)\rangle)[x] \cap \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(|p|-1), x\rangle)[x]$$

where  $L_{2k}$  is the 2k-tactical fog-of-war (resp. 2k-Marköv fog-of-war, identity) and  $L_k$  is the k-tactical fog-of-war (resp. k-Marköv fog-of-war, identity).

Let p attack  $\tau$ . Consider the attack q against the winning strategy  $\sigma$  such that q(2n) = x and q(2n+1) = p(n), and let  $D_n = \sigma \circ L_{2k}(q)$  and  $E_n = \bigcap_{m \leq n} D_n$ .

Certainly,  $x \in E_{2n}[x] = E_{2n}[q(2n)]$  for any  $n < \omega$ . Note also for any  $n < \omega$  that

$$p(n) \in \bigcap_{m \le n} \tau \circ L_k(p \upharpoonright n)$$

$$= \bigcap_{m \le n} (\sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(m-1) \rangle)[x] \cap \sigma \circ L_{2k}(\langle x, p(0), \dots, x, p(m-1), x \rangle)[x])$$

$$= \bigcap_{m \le n} (D_{2m}[x] \cap D_{2m+1}[x]) = \bigcap_{m \le 2n+1} D_m[x] = E_{2n+1}[x]$$

so by the symmetry of  $E_{2n+1}$ ,  $x \in E_{2n+1}[p(n)] = E_{2n+1}[q(2n+1)]$ . Thus  $x \in \bigcap_{n < \omega} E_n[q(n)] \neq \emptyset$ , and since  $\sigma$  is a winning strategy, the attack q converges. Since q(2n) = x, q must converge to x. Thus its subsequence p converges to x, and  $\tau$  is a winning strategy in  $Con_{O,P}(X,x)$ .

Corollary 15. For all  $x \in X$ :

• 
$$\mathscr{D} \underset{k\text{-tact}}{\uparrow} Prox_{D,P}(X) \Rightarrow \mathscr{O} \underset{k\text{-tact}}{\uparrow} Con_{O,P}(X,x)$$

$$\bullet \ \mathscr{D} \underset{k\text{-}mark}{\uparrow} Prox_{D,P}(X) \Rightarrow \mathscr{O} \underset{k\text{-}mark}{\uparrow} Con_{O,P}(X,x)$$

Corollary 16. All proximal spaces are W-spaces.

**Theorem 17.** Let  $X \cup \{\infty\}$  be a uniformizable space such that X is discrete. Then

• 
$$\mathscr{O} \uparrow Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathscr{D} \uparrow Prox_{D,P}(X \cup \{\infty\})$$

• 
$$\mathscr{O} \underset{k\text{-tact}}{\uparrow} Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathscr{D} \underset{k\text{-tact}}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$$

• 
$$\mathscr{O} \underset{k-mark}{\uparrow} Con_{O,P}(X \cup \{\infty\}, \infty) \Rightarrow \mathscr{D} \underset{k-mark}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$$

*Proof.* Note that the topology on  $X \cup \{\infty\}$  is induced by the uniformity with equivalence relation entourages  $D(U) = \Delta \cup U^2$  for each open neighborhood U of  $\infty$ .

Let  $\sigma$  witness  $\mathscr{D} \uparrow Con_{O,P}(X \cap \{\infty\}, \infty)$  (resp.  $\mathscr{D} \uparrow Con_{O,P}(X \cap \{\infty\}, \infty)$ ),  $\mathscr{D} \uparrow Con_{O,P}(X \cap \{\infty\}, \infty)$ ). We define the k-tactical (resp. k-Marköv, perfect info) strategy  $\tau$  such that

$$\tau \circ L(p) = D(\sigma \circ L(p))$$

where L is the k-tactical fog-of-war (resp. k-Marköv fog-of-war, identity).

Let 
$$p \in (X \cup \{\infty\})^{\omega}$$
 attack  $\tau$  such that  $\bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)] \neq \emptyset$ .

If  $\infty \in \bigcap_{n < \omega} \tau(p \upharpoonright n)[p(n)]$ , it follows that p is an attack on  $\sigma$ . Since  $\sigma$  is a winning strategy, it follows that q and its subsequence p must coverge to  $\infty$ .

Otherwise,  $\infty \notin \tau(p \upharpoonright N)[p(N)]$  for some  $N < \omega$ , and then  $\tau(p \upharpoonright N)[p(N)] = \{p(N)\}$  implies  $p \to p(N)$ .

Thus 
$$\tau \circ L$$
 is a winning strategy.

**Corollary 18.** Let  $X \cup \{\infty\}$  be a uniformizable space such that X is discrete. Then

- $\mathscr{O} \uparrow Con_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathscr{D} \uparrow Prox_{D,P}(X \cup \{\infty\})$
- $\mathscr{O} \underset{k\text{-tact}}{\uparrow} Con_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathscr{D} \underset{k\text{-tact}}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$
- $\mathscr{O} \underset{k-mark}{\uparrow} Con_{O,P}(X \cup \{\infty\}, \infty) \Leftrightarrow \mathscr{D} \underset{k-mark}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$

**Proposition 19.** For any  $x \in X$  and  $k \ge 1$ ,

- $\mathscr{O} \underset{k\text{-}tact}{\uparrow} Con_{O,P}(X,x) \Leftrightarrow \mathscr{O} \underset{tact}{\uparrow} Con_{O,P}(X,x)$
- $\bullet \ \, \mathscr{O} \underset{k-mark}{\uparrow} Con_{O,P}(X,x) \Leftrightarrow \mathscr{O} \underset{mark}{\uparrow} Con_{O,P}(X,x)$

*Proof.* If  $\sigma$  witnesses  $\mathcal{O} \underset{k\text{-tact}}{\uparrow} Con_{O,P}(X,x)$ , let  $\tau(\emptyset) = \sigma(\emptyset)$  and

$$\tau(\langle q \rangle) = \bigcap_{i < k} \sigma(\langle \underbrace{x, \dots, x}_{k-i-1}, q, \underbrace{x, \dots, x}_{i} \rangle)$$

This is easily verified to be a winning strategy. The proof for  $\mathscr{O} \underset{k\text{-mark}}{\uparrow} Con_{O,P}(X,x)$  is analogous.

**Corollary 20.** Let  $X \cup \{\infty\}$  be a uniformizable space such that X is discrete, and  $k \geq 1$ . Then

- $\mathscr{D} \underset{k\text{-tact}}{\uparrow} Prox_{D,P}(X \cup \{\infty\}) \Leftrightarrow O \underset{tact}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$
- $\mathscr{D} \underset{k-mark}{\uparrow} Prox_{D,P}(X \cup \{\infty\}) \Leftrightarrow O \underset{mark}{\uparrow} Prox_{D,P}(X \cup \{\infty\})$

**Proposition 21.** For any uniform space X,

- $\mathscr{O} \underset{k\text{-tact}}{\uparrow} Prox_{D,P}(X) \Leftrightarrow \mathscr{O} \underset{2\text{-tact}}{\uparrow} Prox_{D,P}(X)$
- $\mathscr{O} \underset{k-mark}{\uparrow} Prox_{D,P}(X) \Leftrightarrow \mathscr{O} \underset{2-mark}{\uparrow} Prox_{D,P}(X)$

*Proof.* If  $\sigma$  witnesses  $\mathscr{O} \uparrow_{k\text{-tact}} Con_{O,P}(X,x)$ , let  $\tau(\emptyset) = \sigma(\emptyset)$  and

$$\tau(\langle q \rangle) = \bigcap_{i < k} \sigma(\langle \underbrace{q, \dots, q}_{i} \rangle)$$

proximity.tex - Updated on February 27, 2014

$$\tau(\langle q, q' \rangle) = \bigcap_{i < k} \sigma(\langle \underbrace{q, \dots, q}_{k-i}, \underbrace{q', \dots, q'}_{i} \rangle)$$

This is easily verified to be a winning strategy. The proof for  $\mathscr{O} \uparrow_{k\text{-mark}} Con_{O,P}(X,x)$  is analogous.  $\Box$ 

**Definition 22.** The absolute proximal game  $aProx_{D,P}(X)$  is analogous to  $Prox_{D,P}(X)$ , except  $\mathscr{D}$  may only win if p converges.

**Definition 23.** A uniformly locally compact space is a uniformizable space with a uniformly compact entourage M where  $\overline{M[x]}$  is compact for all x.

**Theorem 24.** For any uniformly locally compact space X,  $\mathscr{D} \uparrow Prox_{D,P}(X) \Leftrightarrow \mathscr{D} \uparrow aProx_{D,P}(X)$ 

*Proof.* Let M be a uniformly locally compact entourage. Let  $\sigma$  witness  $\mathscr{D} \uparrow Prox_{D,P}(X)$  such that  $\sigma(a) \subseteq M$  always (so  $\overline{\sigma(a)[x]} \subseteq \overline{M[x]}$  is compact), and  $a \supseteq b$  implies  $\sigma(a) \subseteq \frac{1}{4}\sigma(b)$ .

Let  $\tau(p \upharpoonright n) = \frac{1}{2}\sigma(p \upharpoonright n)$ . If p attacks  $\tau$  in  $aProx_{D,P}(X)$ , then

$$p(n+1) \in \tau(p \upharpoonright n)[p(n)] = \frac{1}{2}\sigma(p \upharpoonright n)[p(n)]$$

and for

$$x \in \overline{\sigma(p \upharpoonright (n+1))[p(n+1)]} \subseteq \overline{\frac{1}{4}\sigma(p \upharpoonright n)[p(n+1)]} \subseteq \frac{1}{2}\sigma(p \upharpoonright n)[p(n+1)]$$

we can conclude  $x \in \sigma(p \upharpoonright n)[p(n)]$ . Thus

$$\sigma(p \upharpoonright (n+1))[p(n+1)] \subseteq \overline{\sigma(p \upharpoonright (n+1))[p(n+1)]} \subseteq \sigma(p \upharpoonright n)[p(n)]$$

Finally, note that p attacks the winning strategy  $\sigma$  in  $Prox_{D,P}(X)$ , but since the intersection of a chain of nonempty compact sets is nonempty:

$$\bigcap_{n<\omega}\sigma(p\upharpoonright n)[p(n)]=\bigcap_{n<\omega}\overline{\sigma(p\upharpoonright n)[p(n)]}\neq\emptyset$$

We conclude that p converges.

**Corollary 25.** A uniformaly locally compact space X is proximal if and only if  $\mathscr{D} \uparrow aProx_{D,P}(X)$ .

**Theorem 26.** For any uniformly locally compact proximal space X,  $\mathscr{O} \uparrow Clus_{O,P}(X,H)$  for all compact  $H \subseteq X$ .

*Proof.* Let  $\sigma$  witness  $\mathscr{D} \uparrow aProx_{D,P}(X)$  such that  $p \supseteq q$  implies  $\sigma(p) \subseteq \frac{1}{4}\sigma(q)$ .

Let o(t) be the subsequence of t consisting of its odd-indexed terms.

We define  $T(\emptyset)$ , etc. as follows:

- Let  $\emptyset \in T(\emptyset)$ .
- Choose  $m_{\emptyset} < \omega$ ,  $h_{\emptyset,i} \in H$  for  $i < m_{\emptyset}$ , and  $h_{\emptyset,i,j} \in H \cap \frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}]$  for  $i, j < m_{\emptyset}$  such that

$$\{\frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}] : i < m_{\emptyset}\}$$

is a cover for H and such that for each  $i < m_{\emptyset}$ 

$$\left\{ \frac{1}{4} \sigma(\langle h_{\emptyset,i} \rangle) [h_{\emptyset,i,j}] : j < m_{\emptyset} \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(\emptyset)[h_{\emptyset,i}]}$ .

• Let  $\langle i \rangle \in T(\emptyset)$ ,  $\langle i, h_{\emptyset,i} \rangle \in T(\emptyset)$ , and  $\langle i, h_{\emptyset,i}, j \rangle \in T(\emptyset)$  for  $i, j < m_{\emptyset}$ .

Suppose T(a), etc. are defined. We then define  $T(a \land \langle x \rangle)$ , etc. for

$$x \in \bigcup_{s \cap \langle i, h_{s,i}, j \rangle \in \max(T(a))} \frac{1}{4} \sigma(o(s) \cap \langle h_{s,i} \rangle) [h_{s,i,j}]$$

as follows:

- Let  $T(a) \subseteq T(a^{\widehat{}}\langle x \rangle)$ .
- Choose  $t = s^{\widehat{}}\langle i, h_{s,i}, j, x \rangle$  such that  $s^{\widehat{}}\langle i, h_{s,i}, j \rangle \in \max(T(a))$  and  $x \in \frac{1}{4}\sigma(o(s)^{\widehat{}}\langle h_{s,i}\rangle)[h_{s,i,j}].$
- Note that, assuming  $o(s) \cap \langle h_{s,i} \rangle$  is a legal partial attack against  $\sigma$ , then

$$x \in \frac{1}{4}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}] \subseteq \frac{1}{4}\sigma(o(s))[h_{s,i,j}]$$

and

$$h_{s,i,j} \in \overline{\frac{1}{4}\sigma(o(s))[h_{s,i}]} \subseteq \frac{1}{2}\sigma(o(s))[h_{s,i}]$$

implies

$$x \in \sigma(o(s))[h_{s,i}]$$

and thus  $o(s)^{\hat{}}\langle h_{s,i}, x \rangle = o(t)$  is a legal partial attack against  $\sigma$ .

• Choose  $m_t < \omega$ ,  $h_{t,k} \in H \cap \frac{1}{4}\sigma(o(s) \cap \langle h_{s,i} \rangle)[h_{s,i,j}]$  for  $k < m_t$ , and  $h_{t,k,l} \in H \cap \frac{1}{4}\sigma(t)[h_{t,k}]$  for  $k, l < m_t$  such that

$$\{\frac{1}{4}\sigma(o(t))[h_{t,k}]: k < m_t\}$$

is a cover for  $H \cap \frac{1}{4}\sigma(o(s)^{\hat{}}(h_{s,i}))[h_{s,i,j}]$  and such that for each  $k < m_t$ 

$$\{\frac{1}{4}\sigma(o(t)^{\frown}\langle h_{t,k}\rangle)[h_{t,i,j}]: l < m_t\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(o(t))[h_{t,k}]}$ .

• Note that, assuming o(t) is a legal partial attack against  $\sigma$ , then

$$h_{t,k} \in \overline{\frac{1}{4}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}]} \subseteq \frac{1}{2}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}]$$

and

$$x \in \frac{1}{4}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}]$$

implies

$$h_{t,k} \in \sigma(o(s) \widehat{\ } \langle h_{s,i} \rangle)[x]$$

and thus  $o(t)^{\hat{}}\langle h_{t,k}\rangle$  is a legal partial attack against  $\sigma$ .

- Let  $t \in T(a^{\ }\langle x \rangle)$ ,  $t^{\ }\langle k \rangle \in T(a^{\ }\langle x \rangle)$ ,  $t^{\ }\langle k, h_{t,k} \rangle \in T(a^{\ }\langle x \rangle)$ , and  $t^{\ }\langle k, h_{t,k}, l \rangle \in T(a^{\ }\langle x \rangle)$  for  $k, l < m_t$ .
- Note that assuming

$$\{\frac{1}{4}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}]: s^{\frown}\langle i, h_{s,i}, j\rangle \in \max(T(a))\}$$

covers H, then since

$$\{\frac{1}{4}\sigma(o(t)^{\frown}\langle h_{t,k}\rangle)[h_{t,k,l}]: s^{\frown}\langle i, h_{s,i}, j, x, k, h_{t,k}, l\rangle \in \max(T(a^{\frown}\langle x\rangle)) \setminus \max(T(a))\}$$

covers  $H \cap \frac{1}{4}\sigma(o(s)^{\frown}\langle h_{s,i}\rangle)[h_{s,i,j}]$ , we have that

$$\{\frac{1}{4}\sigma(o(t)^{\frown}\langle h_{t,k}\rangle)[h_{t,k,l}]:t^{\frown}\langle k,h_{t,k},l\rangle\in\max(T(a^{\frown}\langle x\rangle))\}$$

covers H.

With this we may define the perfect information strategy  $\tau$  for  $\mathscr{O}$  in  $Con_{O,P}(X,H)$  such that:

$$\tau(p \upharpoonright n) = \bigcup_{s \cap \langle i, h_{s,i}, j \rangle \in \max(T(p \upharpoonright n))} \frac{1}{4} \sigma(o(s) \cap \langle h_{s,i} \rangle) [h_{s,i,j}]$$

If p attacks  $\tau$ , then it follows that  $T(p \upharpoonright n)$  is defined for all  $n < \omega$ , so let  $T(p) = \bigcup_{n < \omega} T(p \upharpoonright n)$ . We note T(p) is an infinite tree with finite levels:

- $\emptyset$  has exactly  $m_{\emptyset}$  successors  $\langle i \rangle$ .
- $s^{\hat{}}\langle i\rangle$  has exactly one successor  $t^{\hat{}}\langle i, h_{s,i}\rangle$
- $s^{\frown}\langle i, h_{s,i}\rangle$  has exactly  $m_s$  successors  $t^{\frown}\langle i, h_{s,i}, j\rangle$
- $s \cap \langle i, h_{s,i}, j \rangle$  has either no successors or exactly one successor  $t \cap \langle i, h_{s,i}, j, x \rangle$

•  $t = s^{\hat{}}\langle i, h_{s,i}, j, x \rangle$  has exactly  $m_t$  successors  $t^{\hat{}}\langle k \rangle$ 

Let  $q' = \langle i_0, h_0, j_0, x_0, i_1, h_1, j_1, x_1, \ldots \rangle$  correspond to this infinite branch in T(p), and let  $q = o(q') = \langle h_0, x_0, h_1, x_1, \ldots \rangle$ . Note that by the construction of T(p), q is an attack on the winning strategy  $\sigma$  in  $aProx_{D,P}(X)$ , so it must converge. Since every other term of q is in H, it must converge to H. Then since q is a subsequence of p, p must cluster at H.  $\square$ 

**Corollary 27.** For any uniformly locally compact proximal space,  $\mathcal{O} \uparrow Con_{O,P}(X,H)$  for all compact  $H \subseteq X$ .

*Proof.*  $\mathscr{O} \uparrow Con_{O,P}(X,H)$  if and only if  $\mathscr{O} \uparrow Clus_{O,P}(X,H)$ .

Corollary 28. A compact uniform space X is Corson compact if and only if it is proximal.

*Proof.* A characterization of Corson compact is having a W-set diagonal. If X is proximal compact, then  $X^2$  is proximal compact, and its compact diagonal is a W-set.

**Theorem 29.**  $\mathscr{O} \uparrow_{pre} Con_{O,P}(X,H)$  if and only if there exists a countable base around H.

*Proof.* Let  $\{U_n : n < \omega\}$  be a countable base around H. We define the predetermined strategy  $\sigma(n) = \bigcap_{m \leq n} U_m$ . Let p attack  $\sigma(n)$  - then if U is any neighborhood of H, we may choose  $H \subseteq U_m \subseteq U$ , and note that  $\sigma(n) \subseteq U_m$  for  $n \geq m$ , and thus  $p(n) \in U_m \subseteq U$  for all  $n \geq m$ . Thus  $\sigma$  is a winning strategy.

For the other direction, suppose there does not exist a countable base around H, and let  $\sigma(n)$  be an arbitrary predetermined strategy. Since  $\{\bigcap_{m\leq n}\sigma(m):n<\omega\}$  is not a countable base around H, we may choose an open set U around H such that  $\bigcap_{m\leq n}\sigma(m)\not\subseteq U$  for all  $n<\omega$ . We may easily verify that if  $p(n)\in\bigcap_{m\leq n}\sigma(m)\setminus U$  for all  $n<\omega$ , then p is a successful counterattack to  $\sigma$ .

Corollary 30. X is first countable if and only if  $\mathcal{O} \underset{pre}{\uparrow} Con_{O,P}(X,x)$  for all  $x \in X$ 

Corollary 31.  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$  implies X is first countable.

**Definition 32.** Scattered Eberlein compact spaces are known as **strong Eberlein compact** spaces.

**Theorem 33** (folklore). Scattered compact first-countable spaces are metrizable.

**Corollary 34.** If X is scattered compact and  $\mathcal{O} \uparrow_{pre} Con_{O,P}(X,x)$  for all  $x \in X$  (or  $\mathcal{D} \uparrow_{pre} Prox_{D,P}(X)$ ), then X is metrizable.

Example 35.  $\mathscr{D} \bigwedge_{\text{pre}} Prox_{D,P}(\omega_1^*)$ 

*Proof.* There does not exist a countable base around  $\infty$ , so  $\mathscr{O} \bigwedge_{\mathrm{pre}} Con_{O,P}(X,\omega_1)$ .

**Example 36.**  $\mathscr{O} \underset{\text{tact}}{\uparrow} Con_{O,P}(\kappa^*, \infty)$  and  $\mathscr{D} \underset{\text{tact}}{\uparrow} Prox_{D,P}(\kappa^*)$  for all cardinals  $\kappa$ 

*Proof.* For  $Con_{O,P}(\kappa^*,\infty)$ , let  $\sigma()=\sigma(\infty)=\kappa^*$  and  $\sigma(x)=\kappa^*\setminus\{x\}$  otherwise.

**Theorem 37.** If H is a closed subset of X, then  $\mathscr{Q} \uparrow_{limit} Prox_{D,P}(X) \Rightarrow \mathscr{Q} \uparrow_{limit} Prox_{D,P}(H)$  where  $\uparrow_{limit}$  is any of  $\uparrow$ ,  $\uparrow_{k-tact}$ , or  $\uparrow_{k-mark}$ .

*Proof.* Let  $\sigma \circ L$  witness  $\mathscr{D} \uparrow_{\text{limit}} Prox_{D,P}(X)$ . We define  $\tau \circ L$  for  $\mathscr{D}$  in  $Prox_{D,P}(H)$  as follows:

$$\tau \circ L(p \upharpoonright n) = \sigma \circ L(p \upharpoonright n) \cap H^2$$

Let p attack  $\tau \circ L$ . p also attacks the winning strategy  $\sigma \circ L$ , so either

$$\bigcap_{n<\omega}\left(\bigcap_{m\leq n}\tau\circ L(p\upharpoonright n)\right)[p(n)]\subseteq\bigcap_{n<\omega}\left(\bigcap_{m\leq n}\sigma\circ L(p\upharpoonright n)\right)[p(n)]=\emptyset$$

or p converges in X, and thus converges in H.

**Theorem 38.** If  $\mathscr{D} \uparrow_{\underset{limit}{limit}} Prox_{D,P}(X_i)$  for  $i < \omega$ , then  $\mathscr{D} \uparrow_{\underset{limit}{limit}} Prox_{D,P}(\prod_{i < \omega} X_i)$ , where  $\uparrow_{\underset{limit}{limit}} is \ either \uparrow \ or \ \uparrow_{\underset{k-mark}{l}}$ .

*Proof.* A subbase for  $\prod_{i<\omega} X_i$  is

$$\{\pi_i^{-1}(D): i < \omega, D \in \mathcal{D}_i\}$$

where  $\pi_i$  is the natural projection from  $\left(\prod_{i<\omega}X_i\right)^2$  onto  $X_i^2$ . (See Bell.)

For 
$$p \in (\prod_{i < \omega} X_i)^{\omega}$$
, let  $p_i \in X_i^{\omega}$  such that  $p_i(n) = p(n)(i)$ .

Let  $\sigma_i \circ L$  witness  $\mathscr{D} \uparrow \underset{\text{limit}}{\uparrow} Prox_{D,P}(X_i)$  for  $i < \omega$ , and assume without loss of generality that  $\sigma_i \circ L$  always yields  $X_i^2$  before round i.

Then we define the strategy  $\tau \circ L$  for  $\mathscr{D}$  in  $Prox_{D,P}(\prod_{i<\omega} X_i)$  as follows:

$$\tau \circ L(p \upharpoonright n) = \bigcap_{i \le n} \pi_i^{-1}(\sigma_i \circ L(p_i \upharpoonright n))$$

Let p attack  $\tau \circ L$ . If  $\bigcap_{n < \omega} \left( \bigcap_{m \le n} \sigma_i(p_i \upharpoonright n) \right) [p_i(n)] = \emptyset$  for any  $i < \omega$ , it easily follows that  $\bigcap_{n < \omega} \left( \bigcap_{m \le n} \tau(p \upharpoonright n) \right) [p(n)] = \emptyset$ .

Otherwise, we assume that for each  $i < \omega$ ,  $p_i$  converges to some  $x_i \in X_i$ . Thus p converges to  $x = \langle x_0, x_1, \ldots \rangle$ .

Note: I expect I should be able to do some clever things assuming  $S(\kappa, \omega, \omega)$  to get a similar result for sigma products of dimension  $\kappa$ .

proximity.tex - Updated on February 27, 2014

Example 39. 
$$\mathscr{D} \underset{\text{mark}}{\uparrow} Prox_{D,P}((\kappa^*)^{\omega})$$

Proof. 
$$\mathscr{D} \uparrow_{\text{tact}} Prox_{D,P}(\kappa^*) + \text{previous result}$$

**Theorem 40.** For any predetermined absolutely proximal space X,  $\mathscr{O} \uparrow_{pre} Con_{O,P}(X,H)$  for all compact  $H \subseteq X$ .

*Proof.* Let  $\sigma(n)$  be a winning predetermined strategy for  $\mathscr{D}$  in the absolutely proximal game such that  $\sigma(n+1) \subseteq \frac{1}{4}\sigma(n)$ . For a given tree T, let  $\max(T)$  denote its maximal nodes.

First we define  $T(0) \subseteq \omega^{\leq 2}$ .

- Let  $\emptyset \in T(0)$ .
- Choose  $m_{\emptyset} < \omega$ ,  $h_{\langle i \rangle} \in H$  for  $i < m_{\emptyset}$ , and  $h_{\langle i,j \rangle} \in H \cap \overline{\frac{1}{4}\sigma(0)[h_{\langle i \rangle}]}$  for  $i,j < m_{\emptyset}$  such that

$$\left\{ \frac{1}{4}\sigma(0)[h_{\langle i\rangle}] : i < m_{\emptyset} \right\}$$

is a cover for H and such that for each  $i < m_{\emptyset}$ 

$$\left\{ \frac{1}{4}\sigma(1)[h_{\langle i,j\rangle}] : j < m_{\emptyset} \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(0)[h_{\langle i\rangle}]}$ .

• Let  $\langle i \rangle$  and  $\langle i, j \rangle$  be in T(0) for  $i, j < m_{\emptyset}$ .

Now suppose  $T(n) \subseteq \omega^{\leq 2n+2}$  is defined. We then define  $T(n+1) \subseteq \omega^{\leq 2n+4}$  as follows:

- Let  $T(n) \subseteq T(n+1)$ .
- For each  $t \in \max(T(n))$ , choose  $m_t < \omega$ ,  $h_{t \cap \langle i \rangle} \in H \cap \frac{1}{4}\sigma(2n+2)[h_t]$  for  $i < m_t$ , and  $h_{t \cap \langle i,j \rangle} \in H \cap \frac{1}{4}\sigma(2n+3)[h_{t \cap \langle i \rangle}]$  for  $i,j < m_t$  such that

$$\left\{ \frac{1}{4}\sigma(2n+2)[h_{t} | \langle i \rangle] : i < m_t \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(2n+1)[h_t]}$  and such that for each  $i < m_t$ 

$$\left\{ \frac{1}{4}\sigma(2n+3)[h_{t^{\frown}\langle i,j\rangle}] : i < m_t \right\}$$

is a cover for  $H \cap \overline{\frac{1}{4}\sigma(2n+2)[h_{t^{\frown}\langle i\rangle}]}$ .

• For each  $t \in \max(T(n))$  and each  $i, j < m_t$ , put  $t \cap \langle i \rangle$  and  $t \cap \langle i, j \rangle$  in T(n+1).

We now define the predetermined strategy  $\tau$  for  $\mathscr{O}$  in  $Clus_{O,P}(X,H)$  such that:

$$\tau(n) = \bigcup_{t \in \max(T(n))} \frac{1}{4} \sigma(2n+1)[h_t]$$

In order for  $\tau$  to be a legal strategy, we must show that  $\tau(n)$  contains H. For n=0,

$$\tau(0) = \bigcup_{i,j < m_{\emptyset}} \frac{1}{4} \sigma(1) [h_{\langle i,j \rangle}]$$

Since  $\left\{\frac{1}{4}\sigma(1)[h_{\langle i,j\rangle}]: j < m_{\emptyset}\right\}$  is a cover for  $H \cap \overline{\frac{1}{4}\sigma(0)[h_{\langle i\rangle}]}$ , and since  $\left\{\frac{1}{4}\sigma(0)[h_{\langle i\rangle}]: i < m_{\emptyset}\right\}$  is a cover for H,  $\tau(0)$  must contain H. A similar argument follows for  $\tau(n)$ .

Let p be an attack against  $\tau$  such that  $p(n) \in \bigcap_{m \le n} \tau(m)$ . If we can construct an attack q against  $\sigma$  which shares a subsequence with p, then p must cluster. To find such a q, we construct a new tree T'.

We begin by setting

$$T'(0) = \{s : s \le \langle i, h_{\langle i \rangle}, j \rangle \text{ for } i, j < m_{\emptyset} \}$$

Since

$$p(0) \in \frac{1}{4}\sigma(1)[h_{\langle i,j\rangle}]$$

for some  $i, j < m_{\emptyset}$ , we may find  $t \in \max(T'(0))$  such that  $p(0) \in \frac{1}{4}\sigma(|o(t)|)[h_{e(t)}]$ .

Assume that T'(n) is defined such that there is some  $t \in \max(T'(n))$  where  $p(n) \in \frac{1}{4}\sigma(|o(t)|)[h_{e(t)}]$ . Then let

$$T'(n+1) = T'(n) \cup \{s : s \le t \widehat{\ } \langle p(n), i, h_{e(t) \widehat{\ } \langle i \rangle}, j \rangle \text{ for } i, j < m_{e(t)} \}$$

Note that

$$p(n+1) \in \bigcap_{m \le n+1} \tau(m) = \bigcap_{m \le n+1} \left( \bigcup_{t \in \max(T(m))} \frac{1}{4} \sigma(2m+1)[h_t] \right)$$

(Need details on why there is some t in  $\max(T'(n+1))$  such that  $p(n+1) \in \frac{1}{4}\sigma(|o(t)|)[h_{e(t)}]$ .)

(Follow this up with choosing an infinite branch  $q' \in T'$ , and showing that q = o(q') is an attack on  $\sigma$  similar to perfect info result.)

**Example 41.** Let  $X = I \times 2$  be the Alexandrov double interval. Then  $\mathscr{D} \uparrow_{\text{pre}} Prox_{D,P}(X)$ , but  $\mathscr{D} \uparrow_{\text{mark}} Prox_{D,P}(X)$ .

*Proof.* We assume that the uniformity on X is given by entourages

$$D(\epsilon, F) = \{ \langle x, 0 \rangle, \langle y, 0 \rangle : |x - y| < \epsilon \} \cup \{ \langle x, 1 \rangle, \langle y, 0 \rangle : |x - y| < \epsilon \lor x \not\in F \}$$
$$\cup \{ \langle x, 0 \rangle, \langle y, 1 \rangle : |x - y| < \epsilon \lor y \not\in F \} \cup \{ \langle x, 1 \rangle, \langle y, 1 \rangle : x = y \}$$

That is, points are  $D(\epsilon, F)$ -close if they are the same point, or the first coordinates are within  $\epsilon$  of each other while neither second coordinate is in F.

Suppose  $\mathscr{D}$  had a predetermined winning strategy  $\sigma(n) = D(\epsilon_n, F_n)$ . Then  $\mathscr{P}$  can choose  $x \notin \bigcup_{n < \omega} F_n$ , and play  $\langle x, 1 \rangle$  during even rounds, and  $\langle x_{2n+1}, 0 \rangle$  where  $|x - x_{2n+1}| < \epsilon_{2n}$  during odd rounds, preventing convergence.

However, assume  $\mathscr{D}$  uses the Marköv strategy  $\sigma(x,n) = D(2^{-n},\{x\})$ . If  $\mathscr{P}$  repeats a point of the form  $\langle x,1\rangle$ , then since  $D(2^{-n},\{x\})[\langle x,1\rangle] = \{\langle x,1\rangle\}$ ,  $\mathscr{P}$  must repeat  $\langle x,1\rangle$  for the rest of the game, and  $\mathscr{D}$  wins. Otherwise,  $\mathscr{P}$  cannot repeat points played in  $I \times \{1\}$ , and as the first coordinates form a Cauchy sequence and converge to some z, any open set about  $\langle z,0\rangle$  contains all but finitely many points of  $\mathscr{P}$ 's sequence, and  $\mathscr{D}$  wins.

**Theorem 42.** For any uniformly locally compact space X,  $\mathscr{Q} \uparrow_{pre} Prox_{D,P}(X) \Leftrightarrow \mathscr{Q} \uparrow_{pre} aProx_{D,P}(X)$ 

*Proof.* Let M be a uniformly locally compact entourage. Let  $\sigma$  witness  $\mathscr{D} \uparrow Prox_{D,P}(X)$  such that  $\sigma(n) \subseteq M$  always (so  $\overline{\sigma(a)[x]} \subseteq \overline{M[x]}$  is compact),  $\sigma(n+1) \subseteq \frac{1}{4}\sigma(n)$ .

Let  $\tau(n) = \frac{1}{2}\sigma(n)$ . If p attacks  $\tau$  in  $aProx_{D,P}(X)$ , then

$$p(n+1) \in \tau(n)[p(n)] = \frac{1}{2}\sigma(n)[p(n)]$$

and for

$$x \in \overline{\sigma(n+1)[p(n+1)]} \subseteq \overline{\frac{1}{4}\sigma(n)[p(n+1)]} \subseteq \frac{1}{2}\sigma(n)[p(n+1)]$$

we can conclude  $x \in \sigma(n)[p(n)]$ . Thus

$$\sigma(n+1)[p(n+1)] \subseteq \overline{\sigma(n+1)[p(n+1)]} \subseteq \sigma(n)[p(n)]$$

Finally, note that p attacks the winning strategy  $\sigma$  in  $Prox_{D,P}(X)$ , but since the intersection of a chain of nonempty compact sets is nonempty:

$$\bigcap_{n<\omega}\sigma(n)[p(n)]=\bigcap_{n<\omega}\overline{\sigma(n)[p(n)]}\neq\emptyset$$

We conclude that p converges.

**Proposition 43.** If  $\mathscr{D} \uparrow_{pre} Prox_{D,P}(X)$ , then X has a  $G_{\delta}$  diagonal.

*Proof.* If  $\mathscr{D} \uparrow_{\text{pre}} Prox_{D,P}(X)$  with strategy  $\sigma$ , then consider  $\langle x,y \rangle \in \bigcap_{n<\omega} \sigma(n)$ . It follows that  $\langle x,y,x,y,\ldots \rangle$  attacks  $\sigma$ , and  $\{x,y\} \subseteq \bigcap_{n<\omega} \sigma(n)[x] \cap \bigcap_{n<\omega} \sigma(n)[y] \neq 0$  so it must converge, and x=y. Thus  $\bigcap_{n<\omega} \sigma(n) = \Delta$  is  $G_{\delta}$ .

**Example 44.** The Sorgenfrey line S has a  $G_{\delta}$  diagonal but  $\mathscr{P} \uparrow Prox_{D,P}(S)$ .

Corollary 45. For X with uniformity  $\mathbb{D}$  inducing the compact Hausdorff topology  $\tau$ , the following are equivalent:

- (a)  $\mathscr{D} \uparrow_{pre} Prox_{D,P}(X)$
- (b)  $\mathscr{D} \uparrow_{pre} aProx_{D,P}(X)$
- (c) X has a  $G_{\delta}$  diagonal
- (d)  $\mathbb{D}$  is metrizable
- (e)  $\tau$  is metrizable

*Proof.* For compact Hausdorff spaces, it is well known that there is exactly one uniformity inducing the topology. Thus  $(d) \Leftrightarrow (e)$ . Since X is uniformly locally compact,  $(a) \Leftrightarrow (b)$ . Also, compact spaces with a  $G_{\delta}$  diagonal are metrizable, so  $(c) \Rightarrow (e)$ . Bell noted  $(d) \Rightarrow (a)$  for arbitrary uniform spaces, and the previous proposition shows  $(a) \Rightarrow (c)$ .

**Theorem 46.** A uniformly locally compact space with a  $G_{\delta}$  diagonal is metrizable.

*Proof.* Based on several folklore results.

Uniformly locally compact implies the topological sum of  $\sigma$ -compact spaces implies paracompact. Locally compact plus  $G_{\delta}$  diagonal implies locally metrizable. Locally metrizable plus paracompact characterizes metrizable.

**Corollary 47.** If X is uniformly locally compact, then  $\mathscr{D} \uparrow_{pre} Prox_{D,P}(X)$  implies X's topology is metrizable.

**Example 48.** Let R be the Michael Line. Then  $\mathscr{P} \uparrow Prox_{D,P}(X)$ .

*Proof.* Let  $\pi_n$  be the *n*th digit of  $\pi$ .

During round 0,  $\mathscr{P}$  may choose m(0)=0 and  $p(0)=\pi_0$ , and during round n+1,  $\mathscr{P}$  may choose m(n+1)>m(n) and  $p(n+1)=p(n)+\frac{\pi_{n+1}}{10^{m(n+1)}}$  such that p is a legal attack.

It follows that p "converges" to  $x = \sum_{n < \omega} \frac{\pi_n}{10^{m(n)}}$ , except x is an irrational number.  $\square$