

1 ARHANGELSKII'S α -PRINCIPLES AND SELECTION GAMES

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ABSTRACT. Arhangel'skii's convergence properties α_2 and α_4 may be characterized in terms of Scheeper's selection games. We generalize these folklore results to hold for more general collections.

3 The following characterizations were given as Definition 1 by Kocinac in [cite
4 Kocinac selection principles related].

5 **Definition 1.** *Arhangel'skii's α -principles $\alpha_i(\mathcal{A}, \mathcal{B})$ are defined as follows for $i \in$
6 $\{1, 2, 3, 4\}$. Let $A_n \in \mathcal{A}$ for all $n < \omega$; then there exists $B \in \mathcal{B}$ such that:*

- 7 α_1 : $A_n \cap B$ is cofinite in A_n for all $n < \omega$.
- 8 α_2 : $A_n \cap B$ is infinite for all $n < \omega$.
- 9 α_3 : $A_n \cap B$ is infinite for infinitely-many $n < \omega$.
- 10 α_4 : $A_n \cap B$ is non-empty for infinitely-many $n < \omega$.

11 When $(\mathcal{A}, \mathcal{B})$ is omitted, it is assumed that $\mathcal{A} = \mathcal{B}$ is the collection $\Gamma_{X,x}$ of
12 sequences converging to some point $x \in X$, as introduced by Arhangel'skii in [cite
13 Arhangel'skii frequency spectrum]. Provided \mathcal{A} only contains infinite sets, it's easy
14 to see that $\alpha_n(\mathcal{A}, \mathcal{B})$ implies $\alpha_{n+1}(\mathcal{A}, \mathcal{B})$.

15 We aim to relate these to the following games.

16 **Definition 2.** The *selection game* $G_1(\mathcal{A}, \mathcal{B})$ (resp. $G_{fin}(\mathcal{A}, \mathcal{B})$) is an ω -length
17 game involving Players I and II. During round n , I chooses $A_n \in \mathcal{A}$, followed
18 by II choosing $a_n \in A_n$ (resp. $F_n \in [A_n]^{<\aleph_0}$). Player II wins in the case that
19 $\{a_n : n < \omega\} \in \mathcal{B}$ (resp. $\bigcup \{F_n : n < \omega\} \in \mathcal{B}$), and Player I wins otherwise.

20 Such games are well-represented in the literature; see [cite Scheepers combi-
21 natorics ramsey] for example. We will also consider the similarly-defined games
22 $G_{<2}(\mathcal{A}, \mathcal{B})$ (II chooses 0 or 1 points from each choice by I) and $G_{cf}(\mathcal{A}, \mathcal{B})$ (II
23 chooses cofinitely-many points).

24 **Definition 3.** Let P be a player in a game G . P has a *winning strategy* for G ,
25 denoted $P \uparrow G$, if P has a strategy that defeats every possible counterplay by
26 their opponent. If a strategy only relies on the round number and ignores the
27 moves of the opponent, the strategy is said to be *predetermined*; the existence of a
28 predetermined winning strategy is denoted $P \uparrow_{\text{pre}} G$.

29 We briefly note that the statement $\text{I} \not\uparrow_{\text{pre}} G_\star(\mathcal{A}, \mathcal{B})$ is often denoted as the *selection*
30 *principle* $S_\star(\mathcal{A}, \mathcal{B})$.

31 The equivalence of $\alpha_2(\Gamma_{X,x}, \Gamma_{X,x})$ and $\text{I} \not\uparrow_{\text{pre}} G_1(\Gamma_{X,x}, \Gamma_{X,x})$ was briefly asserted by
32 Sakai in the introduction of [cite Sakai sequence selection properties]; the similar

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equivalence of $\alpha_4(\Gamma_{X,x}\Gamma_{X,x})$ and $\text{I} \not\preceq_{\text{pre}} G_{fin}(\Gamma_{X,x}, \Gamma_{X,x})$ seems to be folklore. In fact, these relationships hold in more generality.

1. MAIN RESULTS

Definition 4. Let $\Gamma_{X,x}$ be the collection of non-trivial sequences $S \subseteq X$ converging to x , that is, infinite subsets of $X \setminus \{x\}$ such that for each neighborhood U of x , $S \cap U$ is cofinite in S .

Definition 5. Let Γ_X be the collection of open γ -covers \mathcal{U} of X , that is, infinite open covers of X such that $X \notin \mathcal{U}$ and for each $x \in X$, $\{U \in \mathcal{U} : x \in U\}$ is cofinite in \mathcal{U} .

The similarity in nomenclature follows from the observation that every non-trivial sequence in $C_p(X)$ converging to the zero function $\mathbf{0}$ naturally defines a corresponding γ -cover in X , see e.g. Theorem 4 of [Scheepers a sequential property and covering property].

Note that by these definitions, convergent sequences (resp. γ -covers) may be uncountable, but any infinite subset of either would remain a convergent sequence (resp. γ -cover), in particular, countably infinite subsets. We capture this idea as follows.

Definition 6. Say a collection \mathcal{A} is Γ -like if it satisfies the following for each $A \in \mathcal{A}$.

- $|A| \geq \aleph_0$.
- If $A' \subseteq A$ and $|A'| \geq \aleph_0$, then $A' \in \mathcal{A}$.

We also require the following.

Definition 7. Say a collection \mathcal{A} is *almost- Γ -like* if for each $A \in \mathcal{A}$, there is $A' \subseteq A$ such that:

- $|A'| = \aleph_0$.
- If A'' is a cofinite subset of A' , then $A'' \in \mathcal{A}$.

So all Γ -like sets are almost- Γ -like.

We are now able to prove a few general equivalences between α -principles and selection games.

Theorem 8. *Let \mathcal{A} be almost- Γ -like and \mathcal{B} be Γ -like. Then $\alpha_2(\mathcal{A}, \mathcal{B})$ holds if and only if $\text{I} \not\preceq_{\text{pre}} G_1(\mathcal{A}, \mathcal{B})$.*

Proof. We first assume $\alpha_2(\mathcal{A}, \mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined strategy for I. We may apply $\alpha_2(\mathcal{A}, \mathcal{B})$ to choose $B \in \mathcal{B}$ such that $|A_n \cap B| \geq \aleph_0$. We may then choose $a_n \in (A_n \cap B) \setminus \{a_i : i < n\}$ for each $n < \omega$. It follows that $B' = \{a_n : n < \omega\} \in \mathcal{B}$ since B' is an infinite subset of $B \in \mathcal{B}$; therefore A_n does not define a winning predetermined strategy for I.

Now suppose $\text{I} \not\preceq_{\text{pre}} G_1(\mathcal{A}, \mathcal{B})$. Given $A_n \in \mathcal{A}$ for $n < \omega$, first choose $A'_n \in \mathcal{A}$ such that $A'_n = \{a_{n,j} : j < \omega\} \subseteq A_n$, $j < k$ implies $a_{n,j} \neq a_{n,k}$, and $A_{n,m} = \{a_{n,j} : m \leq j < \omega\} \in \mathcal{A}$. Finally choose some $\theta : \omega \rightarrow \omega$ such that $|\theta^{\leftarrow}(n)| = \aleph_0$ for each $n < \omega$.

Since playing $A_{\theta(m),m}$ during round m does not define a winning strategy for I in $G_1(\mathcal{A}, \mathcal{B})$, II may choose $x_m \in A_{\theta(m),m}$ such that $B = \{x_m : m < \omega\} \in \mathcal{B}$. Choose $i_m < \omega$ for each $m < \omega$ such that $x_m = a_{\theta(m),i_m}$, noting $i_m \geq m$. It follows that $A_n \cap B \supseteq \{a_{\theta(m),i_m} : m \in \theta^{\leftarrow}(n)\}$. Since for each $m \in \theta^{\leftarrow}(n)$ there exists $M \in$

75 $\theta^{\leftarrow}(n)$ such that $m \leq i_m < M \leq i_M$, and therefore $a_{\theta(m), i_m} \neq a_{\theta(m), i_M} = a_{\theta(M), i_M}$,
 76 we have shown that $A_n \cap B$ is infinite. Thus B witnesses $\alpha_2(\mathcal{A}, \mathcal{B})$. \square

77 **Theorem 9.** *Let \mathcal{A} be almost- Γ -like and \mathcal{B} be Γ -like. Then $\alpha_4(\mathcal{A}, \mathcal{B})$ holds if and*
 78 *only if $\text{I} \nVdash_{\text{pre}} G_{<2}(\mathcal{A}, \mathcal{B})$ if and only if $\text{I} \nVdash_{\text{pre}} G_{fin}(\mathcal{A}, \mathcal{B})$.*

79 *Proof.* We first assume $\alpha_4(\mathcal{A}, \mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined
 80 strategy for I in $G_{<2}(\mathcal{A}, \mathcal{B})$. We then may choose $A'_n \in \mathcal{A}$ where $A'_n = \{a_{n,j} : j <$
 81 $\omega\} \subseteq A_n$, $j < k$ implies $a_{n,j} \neq a_{n,k}$, and $A''_n = A'_n \setminus \{a_{i,j} : i, j < n\} \in \mathcal{A}$.

82 By applying $\alpha_4(\mathcal{A}, \mathcal{B})$ to A''_n , we obtain $B \in \mathcal{B}$ such that $A''_n \cap B \neq \emptyset$ for infinitely-
 83 many $n < \omega$. We then let $F_n = \emptyset$ when $A''_n \cap B = \emptyset$, and $F_n = \{x_n\}$ for some
 84 $x_n \in A''_n \cap B$ otherwise. Then we will have that $B' = \bigcup \{F_n : n < \omega\} \subseteq B$ belongs
 85 to \mathcal{B} once we show that B' is infinite. To see this, for $m \leq n < \omega$ note that either
 86 F_m is empty (and we let $j_m = 0$) or $F_m = \{a_{m,j_m}\}$ for some $j_m \geq m$; choose $N < \omega$
 87 such that $j_m < N$ for all $m \leq n$ and $F_N = \{x_N\}$. Thus $F_m \neq F_N$ for all $m \leq n$
 88 since $x_N \notin \{a_{i,j} : i, j < N\}$. Thus II may defeat the predetermined strategy A_n by
 89 playing F_n each round.

90 Since $\text{I} \nVdash_{\text{pre}} G_{<2}(\mathcal{A}, \mathcal{B})$ immediately implies $\text{I} \nVdash_{\text{pre}} G_{fin}(\mathcal{A}, \mathcal{B})$, we assume the latter.

91 Given $A_n \in \mathcal{A}$ for $n < \omega$, we note this defines a (non-winning) predetermined
 92 strategy for I, so II may choose $F_n \in [A_n]^{<\aleph_0}$ such that $B = \bigcup \{F_n : n < \omega\} \in \mathcal{B}$.
 93 Since B is infinite, we note $F_n \neq \emptyset$ for infinitely-many $n < \omega$. Thus B witnesses
 94 $\alpha_4(\mathcal{A}, \mathcal{B})$ since $A_n \cap B \supseteq F_n \neq \emptyset$ for infinitely-many $n < \omega$. \square

95 This shows that II gains no advantage from picking more than one point per
 96 round. This in fact only depends on \mathcal{B} being Γ -like, which we formalize in the
 97 following results.

98 **Theorem 10.** *Let \mathcal{B} be Γ -like. Then $\text{I} \uparrow_{\text{pre}} G_{<2}(\mathcal{A}, \mathcal{B})$ if and only if $\text{I} \uparrow_{\text{pre}} G_{fin}(\mathcal{A}, \mathcal{B})$.*

99 *Proof.* Assume $\bigcup \mathcal{A}$ is well-ordered. Given a winning predetermined strategy A_n
 100 for I in $G_{<2}(\mathcal{A}, \mathcal{B})$, consider $F_n \in [A_n]^{<\aleph_0}$. We set

$$F_n^* = \begin{cases} \emptyset & \text{if } F_n \setminus \bigcup \{F_m : m < n\} = \emptyset \\ \{\min(F_n \setminus \bigcup \{F_m : m < n\})\} & \text{otherwise} \end{cases}$$

101 Since $|F_n^*| < 2$, we have that $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$. In the case that $\bigcup \{F_n^* : n < \omega\}$
 102 is finite, we immediately see that $\bigcup \{F_n : n < \omega\}$ is also finite and therefore not in
 103 \mathcal{B} . Otherwise $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$ is an infinite subset of $\bigcup \{F_n : n < \omega\}$, and thus
 104 $\bigcup \{F_n : n < \omega\} \notin \mathcal{B}$ too. Therefore A_n is a winning predetermined strategy for I in
 105 $G_{fin}(\mathcal{A}, \mathcal{B})$ as well. \square

106 **Theorem 11.** *Let \mathcal{B} be Γ -like. Then $\text{I} \uparrow G_{<2}(\mathcal{A}, \mathcal{B})$ if and only if $\text{I} \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$.*

107 *Proof.* Assume $\bigcup \mathcal{A}$ is well-ordered. Suppose $\text{I} \uparrow G_{<2}(\mathcal{A}, \mathcal{B})$ is witnessed by the
 108 strategy σ . Let $\langle \rangle^* = \langle \rangle$, and for $s \frown \langle F \rangle \in ([\bigcup \mathcal{A}]^{<\aleph_0})^{<\omega} \setminus \{\langle \rangle\}$ let

$$(s \frown \langle F \rangle)^* = \begin{cases} s^* \frown \langle \emptyset \rangle & \text{if } F \setminus \bigcup \text{range}(s) = \emptyset \\ s^* \frown \langle \{\min(F \setminus \bigcup \text{range}(s))\} \rangle & \text{otherwise} \end{cases}$$

109 We then define the strategy τ for I in $G_{fin}(\mathcal{A}, \mathcal{B})$ by $\tau(s) = \sigma(s^*)$. Then given
 110 any counterattack $\alpha \in ([\bigcup \mathcal{A}]^{<\aleph_0})^\omega$ by II played against τ , we note that $\alpha^* =$

111 $\bigcup\{(\alpha \upharpoonright n)^* : n < \omega\}$ is a counterattack to σ , and thus loses. This means $B =$
 112 $\bigcup \text{range}(\alpha^*) \notin \mathcal{B}$.

113 We consider two cases. The first is the case that $\bigcup \text{range}(\alpha^*)$ is finite. Noting
 114 that $\alpha^*(m) \cap \alpha^*(n) = \emptyset$ whenever $m \neq n$, there exists $N < \omega$ such that $\alpha^*(n) = \emptyset$
 115 for all $n > N$. As a result, $\bigcup \text{range}(\alpha) = \bigcup \text{range}(\alpha \upharpoonright n)$, and thus $\bigcup \text{range}(\alpha)$ is
 116 finite, and therefore not in \mathcal{B} .

117 In the other case, $\bigcup \text{range}(\alpha^*) \notin \mathcal{B}$ is an infinite subset of $\bigcup \text{range}(\alpha)$, and
 118 therefore $\bigcup \text{range}(\alpha) \notin \mathcal{B}$ as well. Thus we have shown that τ is a winning strategy
 119 for I in $G_{fin}(\mathcal{A}, \mathcal{B})$. \square

120 We note that the above proof technique could be used to establish that perfect-
 121 information and limited-information strategies for II in $G_{fin}(\mathcal{A}, \mathcal{B})$ may be improved
 122 to be valid in $G_{<2}(\mathcal{A}, \mathcal{B})$, provided \mathcal{B} is Γ -like. As such, $G_{<2}(\mathcal{A}, \mathcal{B})$ and $G_{fin}(\mathcal{A}, \mathcal{B})$
 123 are effectively equivalent games under this hypothesis.

124 **Theorem 12.** *Let \mathcal{A} be almost- Γ -like and \mathcal{B} be Γ -like. Then*

- 125 • $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $I \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$, and
- 126 • $I \uparrow G_1(\mathcal{A}, \mathcal{B})$ if and only if $I \uparrow_{pre} G_1(\mathcal{A}, \mathcal{B})$.

127 *Proof.* We assume $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ and let the symbol \dagger mean $< \aleph_0$ (respectively,
 128 $I \uparrow G_1(\mathcal{A}, \mathcal{B})$ and $\dagger = 1$, and for convenience we assume II plays singleton subsets
 129 of \mathcal{A} rather than elements). As \mathcal{A} is almost- Γ -like, there is a winning strategy σ
 130 where $|\sigma(s)| = \aleph_0$ and $\sigma(s) \cap \bigcup \text{range}(s) = \emptyset$ (that is, σ never replays the choices
 131 of II) for all partial plays s by II.

132 For each $s \in \omega^{<\omega}$, suppose $F_{s \upharpoonright m} \in [\bigcup \mathcal{A}]^\dagger$ is defined for each $0 < m \leq |s|$. Then
 133 let $s^* : |s| \rightarrow [\bigcup \mathcal{A}]^\dagger$ be defined by $s^*(m) = F_{s \upharpoonright m+1}$, and define $\tau' : \omega^{<\omega} \rightarrow \mathcal{A}$ by
 134 $\tau'(s) = \sigma(s^*)$. Finally, set $[\sigma(s^*)]^\dagger = \{F_{s \upharpoonright n} : n < \omega\}$, and for some bijection
 135 $b : \omega^{<\omega} \rightarrow \omega$ let $\tau(n) = \tau'(b(n))$ be a predetermined strategy for I in $G_{fin}(\mathcal{A}, \mathcal{B})$
 136 (resp. $G_1(\mathcal{A}, \mathcal{B})$).

137 Suppose α is a counterattack by II against τ , so

$$\alpha(n) \in [\tau(n)]^\dagger = [\tau'(b(n))]^\dagger = [\sigma(b(n)^*)]^\dagger$$

138 It follows that $\alpha(n) = F_{b(n) \upharpoonright \langle m \rangle}$ for some $m < \omega$. In particular, there is some
 139 infinite subset $W \subseteq \omega$ and $f \in \omega^\omega$ such that $\{\alpha(n) : n \in W\} = \{F_{f \upharpoonright n+1} : n < \omega\}$.
 140 Note here that $(f \upharpoonright n+1)^* = (f \upharpoonright n)^* \frown \langle F_{f \upharpoonright n+1} \rangle$. This shows that $F_{f \upharpoonright n+1} \in [\sigma((f \upharpoonright$
 141 $n)^*)]^\dagger$ is an attempt by II to defeat σ , which fails. Thus $\bigcup\{F_{f \upharpoonright n+1} : n < \omega\} =$
 142 $\bigcup\{\alpha(n) : n \in W\} \notin \mathcal{B}$, and since this set is infinite (as σ prevents II from repeating
 143 choices) we have $\bigcup\{\alpha(n) : n < \omega\} \notin \mathcal{B}$ too. Therefore τ is winning. \square

144 Note that the assumption in Theorem 12 that \mathcal{A} be almost- Γ -like cannot be
 145 omitted. In [todo cite Clontz k-tactics in Gruenhage game] an example of a space
 146 X^* and point $\infty \in X^*$ where $I \uparrow_{pre} G_1(\mathcal{A}, \mathcal{B})$ but $I \not\uparrow_{pre} G_1(\mathcal{A}, \mathcal{B})$ is given, where \mathcal{A} is the
 147 set of open neighborhoods of ∞ (which are all uncountable), and \mathcal{B} is the set $\Gamma_{X^*, \infty}$
 148 of sequences converging to that point. (Note that $G_1(\mathcal{A}, \mathcal{B})$ is called $Gru_{O,P}(X^*, \infty)$
 149 in that paper, and an equivalent game $Gru_{K,P}(X)$ is what is directly studied. In
 150 fact, more is shown: I has a winning perfect-information strategy, but for any
 151 natural number k , any strategy that only uses the most recent k moves of II and
 152 the round number can be defeated.)

While \mathcal{A} is often not almost- Γ -like in general, it may satisfy that property in combination with the selection principles being considered.

Proposition 13. *Let \mathcal{B} be Γ -like, $\mathcal{B} \subseteq \mathcal{A}$, and $\mathbf{I} \nVdash_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$. Then \mathcal{A} is almost- Γ -like.*

Proof. Let $A \in \mathcal{A}$, and for all $n < \omega$ let $A_n = A$. Then A_n is not a winning predetermined strategy for \mathbf{I} , so \mathbf{II} may choose finite sets $B_n \subseteq A_n = A$ such that $A' = \bigcup \{B_n : n < \omega\} \in \mathcal{B} \subseteq \mathcal{A}$.

It follows that $A' \subseteq A$ and $|A'| = \aleph_0$, and for any infinite subset $A'' \subseteq A'$ (in particular, any cofinite subset), $A'' \in \mathcal{B} \subseteq \mathcal{A}$. Thus \mathcal{A} is almost- Γ -like. \square

Note that in the previous result, $\mathbf{I} \nVdash_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$ could be weakened to the choice principle (\mathcal{A}_B^A) : for every member of \mathcal{A} , there is some countable subset belonging to \mathcal{B} .

Corollary 14. *Let \mathcal{B} be Γ -like and $\mathcal{B} \subseteq \mathcal{A}$. Then*

- $\mathbf{I} \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $\mathbf{I} \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$, and
- $\mathbf{I} \uparrow G_1(\mathcal{A}, \mathcal{B})$ if and only if $\mathbf{I} \uparrow_{pre} G_1(\mathcal{A}, \mathcal{B})$.

Proof. Assuming $\mathbf{I} \nVdash_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$, we have $\mathbf{I} \nVdash G_{fin}(\mathcal{A}, \mathcal{B})$ by Proposition 13 and Theorem 12.

Similarly, assuming $\mathbf{I} \nVdash_{pre} G_1(\mathcal{A}, \mathcal{B}) \Rightarrow \mathbf{I} \nVdash_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$, we have $\mathbf{I} \nVdash G_1(\mathcal{A}, \mathcal{B})$ by Proposition 13 and Theorem 12. \square

This corollary generalizes e.g. Theorems 26 and 30 of [cite Scheepers 1996 Ramsey], Theorem 5 of [cite MR2119791], and Corollary 36 of [cite Clontz dual games].

In summary, using the selection principle notation $S_*(\mathcal{A}, \mathcal{B})$:

Corollary 15. *Let \mathcal{B} be Γ -like and $\mathcal{B} \subseteq \mathcal{A}$. Then*

- $\mathbf{I} \nVdash G_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $S_{fin}(\mathcal{A}, \mathcal{B})$ if and only if $\alpha_2(\mathcal{A}, \mathcal{B})$, and
- $\mathbf{I} \nVdash G_1(\mathcal{A}, \mathcal{B})$ if and only if $S_1(\mathcal{A}, \mathcal{B})$ if and only if $\alpha_4(\mathcal{A}, \mathcal{B})$.

2. CONCLUSION

We conclude with the following easy result, and a couple questions.

Proposition 16. *Let \mathcal{B} be Γ -like. Then $\alpha_1(\mathcal{A}, \mathcal{B})$ holds if and only if $\mathbf{I} \nVdash_{pre} G_{cf}(\mathcal{A}, \mathcal{B})$.*

Proof. We first assume $\alpha_1(\mathcal{A}, \mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined strategy for \mathbf{I} . By $\alpha_1(\mathcal{A}, \mathcal{B})$, we immediately obtain $B \in \mathcal{B}$ such that $|A_n \setminus B| < \aleph_0$. Thus $B_n = A_n \cap B$ is a cofinite choice from A_n , and $B' = \bigcup \{B_n : n < \omega\}$ is an infinite subset of B , so $B' \in \mathcal{B}$. Thus \mathbf{II} may defeat \mathbf{I} by choosing $B_n \subseteq A_n$ each round, witnessing $\mathbf{I} \nVdash_{pre} G_{cf}(\mathcal{A}, \mathcal{B})$.

On the other hand, let $\mathbf{I} \nVdash_{pre} G_{cf}(\mathcal{A}, \mathcal{B})$. Given $A_n \in \mathcal{A}$ for $n < \omega$, we note that \mathbf{II} may choose a cofinite subset $B_n \subseteq A_n$ such that $B = \bigcup \{B_n : n < \omega\} \in \mathcal{B}$. Then B witnesses $\alpha_1(\mathcal{A}, \mathcal{B})$ since $|A_n \setminus B| \leq |A_n \setminus B_n| \leq \aleph_0$. \square

Question 17. *Is there a game-theoretic characterization of $\alpha_3(\mathcal{A}, \mathcal{B})$?*

190 Noting that $I \uparrow G_1(\Gamma_X, \Gamma_X)$ if and only if $I \uparrow G_{fin}(\Gamma_X, \Gamma_X)$ [cite Kocinac], but
 191 the same is not true of $G_*(\Gamma_{X,x}, \Gamma_{X,x})$ (i.e. there are α_4 spaces that are not α_2
 192 [cite Arhangel'skii]), we also ask the following.

193 **Question 18.** *Is there an elegant condition on \mathcal{A}, \mathcal{B} guaranteeing $I \uparrow G_1(\mathcal{A}, \mathcal{B}) \Rightarrow$*
 194 *$I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$?*

195 REFERENCES

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