## SELECTION GAMES AND ARHANGELSKII'S CONVERGENCE PRINCIPLES

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ABSTRACT. We prove the things.

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## 1. Clontz results **Definition 1.** Say a collection $\mathcal{A}$ is sequence-like if it satisfies the following for each $A \in \mathcal{A}$ . • $|A| \geq \aleph_0$ . • If $A' \subseteq A$ and $|A'| \ge \aleph_0$ , then $A' \in \mathcal{A}$ . **Definition 2.** Let $\Gamma_X$ be the collection of open $\gamma$ -covers $\mathcal{U}$ of X, that is, infinite open covers of X such that for each $x \in X$ , $\{U \in \mathcal{U} : x \in U\}$ is cofinite in $\mathcal{U}$ . **Definition 3.** Let $\Gamma_{X,x}$ be the collection of non-trivial sequences $S \subseteq X$ converging 11 to x, that is, infinite subsets of X such that for each neighborhood U of $x, S \cap U$ is cofinite in S. 13 It follows that $\Gamma_X, \Gamma_{X,x}$ are both sequence-like. We also require the following. **Definition 4.** Say a collection $\mathcal{A}$ is almost-sequence-like if for each $A \in \mathcal{A}$ , there is $A' \subseteq A$ such that: $\bullet |A'| = \aleph_0.$ 17 • If A'' is a cofinite subset of A', then $A'' \in \mathcal{A}$ . 18 So all sequence-like sets are almost-sequence-like. **Theorem 5.** Let $\mathcal{B}$ be sequence-like. Then $\alpha_1(\mathcal{A},\mathcal{B})$ holds if and only if I $\checkmark$ $G_{cf}(\mathcal{A},\mathcal{B}).$ *Proof.* We first assume $\alpha_1(\mathcal{A}, \mathcal{B})$ and let $A_n \in \mathcal{A}$ for $n < \omega$ define a predetermined strategy for I. By $\alpha_1(\mathcal{A}, \mathcal{B})$ , we immediately obtain $B \in \mathcal{B}$ such that $|A_n \setminus B| < \aleph_0$ . Thus $B_n = A_n \cap B$ is a cofinite choice from $A_n$ , and $B' = \bigcup \{B_n : n < \omega\}$ is an infinite subset of B, so $B' \in \mathcal{B}$ . Thus II may defeat I by choosing $B_n \subseteq A_n$ each round, witnessing I $\uparrow \atop \text{pre} G_{cf}(\mathcal{A}, \mathcal{B})$ . 26 On the other hand, let I $\uparrow G_{cf}(\mathcal{A}, \mathcal{B})$ . Given $A_n \in \mathcal{A}$ for $n < \omega$ , we note that 27

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B witnesses  $\alpha_1(\mathcal{A}, \mathcal{B})$  since  $|A_n \setminus B| \leq |A_n \setminus B_n| \leq \aleph_0$ .

holds if and only if I  $\gamma$   $G_1(\mathcal{A}, \mathcal{B})$ .

II may choose a cofinite subset  $B_n \subseteq A_n$  such that  $B = \bigcup \{B_n : n < \omega\} \in \mathcal{B}$ . Then

**Theorem 6.** Let  $\mathcal{A}$  be almost-sequence-like and  $\mathcal{B}$  be sequence-like. Then  $\alpha_2(\mathcal{A}, \mathcal{B})$ 

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*Proof.* We first assume  $\alpha_2(\mathcal{A}, \mathcal{B})$  and let  $A_n \in \mathcal{A}$  for  $n < \omega$  define a predetermined strategy for  $\mathscr{I}$ . We may apply  $\alpha_2(\mathcal{A}, \mathcal{B})$  to choose  $B \in \mathcal{B}$  such that  $|A_n \cap B| \geq \aleph_0$ . We may then choose  $a_n \in (A_n \cap B) \setminus \{a_i : i < n\}$  for each  $n < \omega$ . It follows that  $B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ is an infinite subset of } B \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B} \text{ since } B' \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ does } B' = \{a_n : n < \omega\} \in \mathcal{B}; \text{ therefore } A_n \text{ doe$ not define a winning predetermined strategy for I. 37 that  $A'_n = \{a_{n,j} : j < \omega\} \subseteq A_n, j < k \text{ implies } a_{n,j} \neq a_{n,k}, \text{ and } A_{n,m} = \{a_{n,j} : m \leq a_{n,j} \neq a_{n,k}, a_{n,m} = a_{n,j} = a_{n,j$ 38  $j < \omega \} \in \mathcal{A}$ . Finally choose some  $\theta : \omega \to \omega$  such that  $|\theta^{\leftarrow}(n)| = \aleph_0$  for each  $n < \omega$ . Since playing  $A_{\theta(m),m}$  during round m does not define a winning strategy for I in  $G_1(\mathcal{A},\mathcal{B})$ , II may choose  $x_m \in A_{\theta(m),m}$  such that  $B = \{x_m : m < \omega\} \in \mathcal{B}$ . Choose 41  $i_m < \omega$  for each  $m < \omega$  such that  $x_m = a_{\theta(m),i_m}$ , noting  $i_m \ge m$ . It follows that  $A_n \cap B \supseteq \{a_{\theta(m),i_m} : m \in \theta^{\leftarrow}(n)\}$ . Since for each  $m \in \theta^{\leftarrow}(n)$  there exists  $M \in A_n \cap B$  $\theta^{\leftarrow}(n)$  such that  $m \leq i_m < M \leq i_M$ , and therefore  $a_{\theta(m),i_m} \neq a_{\theta(m),i_M} = a_{\theta(M),i_M}$ , we have shown that  $A_n \cap B$  is infinite. Thus B witnesses  $\alpha_2(\mathcal{A}, \mathcal{B})$ .

Theorem 7. Let  $\mathcal{A}$  be almost-sequence-like and  $\mathcal{B}$  be sequence-like. Then  $\alpha_4(\mathcal{A},\mathcal{B})$  holds if and only if I  $\gamma \atop pre$   $G_{<2}(\mathcal{A},\mathcal{B})$  if and only if I  $\gamma \atop pre$   $G_{fin}(\mathcal{A},\mathcal{B})$ .

48 Proof. We first assume  $\alpha_4(\mathcal{A},\mathcal{B})$  and let  $A_n \in \mathcal{A}$  for  $n < \omega$  define a predetermined 49 strategy for I in  $G_{<2}(\mathcal{A},\mathcal{B})$ . We then may choose  $A'_n \in \mathcal{A}$  where  $A'_n = \{a_{n,j} : j < \omega\} \subseteq A_n, j < k$  implies  $a_{n,j} \neq a_{n,k}$ , and  $A''_n = A'_n \setminus \{a_{i,j} : i,j < n\} \in \mathcal{A}$ .

By applying  $\alpha_4(\mathcal{A},\mathcal{B})$  to  $A_n''$ , we obtain  $B \in \mathcal{B}$  such that  $A_n'' \cap B \neq \emptyset$  for infintelymany  $n < \omega$ . We then let  $F_n = \emptyset$  when  $A_n'' \cap B = \emptyset$ , and  $F_n = \{x_n\}$  for some  $x_n \in A_n'' \cap B$  otherwise. Then we will have that  $B' = \bigcup \{F_n : n < \omega\} \subseteq B$  belongs to  $\mathcal{B}$  once we show that B' is infinite. To see this, for  $m \leq n < \omega$  note that either  $F_m$  is empty (and we let  $j_m = 0$ ) or  $F_m = \{a_{m,j_m}\}$  for some  $j_m \geq m$ ; choose  $N < \omega$  such that  $j_m < N$  for all  $m \leq n$  and  $j_m = \{a_{m,j_m}\}$ . Thus  $j_m \neq j_m = \{a_{m,j_m}\}$  for all  $j_m \leq n$  since  $j_m \neq j_m = \{a_{m,j_m}\}$ . Thus  $j_m \neq j_m = \{a_{m,j_m}\}$  for all  $j_m \leq n$  since  $j_m \neq j_m = \{a_{m,j_m}\}$ . Thus II may defeat the predetermined strategy  $j_m = \{a_{m,j_m}\}$  by playing  $j_m = \{a_{m,j_m}\}$  for some  $j_m \geq m$ ; choose  $j_m \geq m$  since  $j_m \neq j_m = \{a_{m,j_m}\}$ . Thus II may defeat the predetermined strategy  $j_m = \{a_{m,j_m}\}$  for some  $j_m \geq m$ ; choose  $j_m \geq m$  since  $j_m \neq j_m = \{a_{m,j_m}\}$ . Thus II may defeat the predetermined strategy  $j_m = \{a_{m,j_m}\}$  for some  $j_m \geq m$ ; choose  $j_m \geq m$ 

Since I  $\uparrow G_{<2}(\mathcal{A}, \mathcal{B})$  immediately implies I  $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ , we assume the latter.

Given  $A_n \in \mathcal{A}$  for  $n < \omega$ , we note this defines a (non-winning) predetermined strategy for I, so II may choose  $F_n \in [A_n]^{<\aleph_0}$  such that  $B = \bigcup \{F_n : n < \omega\} \in \mathcal{B}$ . Since B is infinite, we note  $F_n \neq \emptyset$  for infinitely-many  $n < \omega$ . Thus B witnesses  $\alpha_4(\mathcal{A},\mathcal{B})$  since  $A_n \cap B \supseteq F_n \neq \emptyset$  for infinitely-many  $n < \omega$ .

Theorem 8. Let  $\mathcal{B}$  be sequence-like. Then  $I \uparrow_{pre} G_{<2}(\mathcal{A}, \mathcal{B})$  if and only if  $I \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$ .

Proof. Assume  $\bigcup \mathcal{A}$  is well-ordered. Given a winning predetermined strategy  $A_n$  for I in  $G_{<2}(\mathcal{A},\mathcal{B})$ , consider  $F_n \in [A_n]^{<\aleph_0}$ . We set

$$F_n^* = \begin{cases} \emptyset & \text{if } F_n \setminus \bigcup \{F_m : m < n\} = \emptyset \\ \{\min(F_n \setminus \bigcup \{F_m : m < n\})\} & \text{otherwise} \end{cases}$$

Since  $|F_n| < 2$ , we have that  $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$ . In the case that  $\bigcup \{F_n^* : n < \omega\}$  is finite, we immediately see that  $\bigcup \{F_n : n < \omega\}$  is also finite and therefore not in  $\mathcal{B}$ . Otherwise  $\bigcup \{F_n^* : n < \omega\} \notin \mathcal{B}$  is an infinite subset of  $\bigcup \{F_n : n < \omega\}$ , and thus  $\bigcup \{F_n : n < \omega\} \notin \mathcal{B}$  too. Therefore  $A_n$  is a winning predetermined strategy for I in  $G_{fin}(\mathcal{A}, \mathcal{B})$  as well.

Theorem 9. Let  $\mathcal{B}$  be sequence-like. Then  $I \uparrow G_{<2}(\mathcal{A}, \mathcal{B})$  if and only if  $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ .

Proof. Assume  $\bigcup \mathcal{A}$  is well-ordered. Suppose  $I \uparrow G_{<2}(\mathcal{A}, \mathcal{B})$  is witnessed by the strategy  $\sigma$ . Let  $\langle \rangle^* = \langle \rangle$ , and for  $s \cap \langle F \rangle \in ([\bigcup \mathcal{A}]^{<\aleph_0})^{<\omega} \setminus \{\langle \rangle\}$  let

$$(s^{\frown} \langle F \rangle)^{\star} = \begin{cases} s^{\star \frown} \langle \emptyset \rangle & \text{if } F \setminus \bigcup \text{range}(s) = \emptyset \\ s^{\star \frown} \langle \{\min(F \setminus \bigcup \text{range}(s))\} \rangle & \text{otherwise} \end{cases}$$

We then define the strategy  $\tau$  for I in  $G_{fin}(\mathcal{A}, \mathcal{B})$  by  $\tau(s) = \sigma(s^*)$ . Then given any counterattack  $\alpha \in ([\bigcup \mathcal{A}]^{<\aleph_0})^{\omega}$  by II played against  $\tau$ , we note that  $\alpha^* = \bigcup \{(\alpha \upharpoonright n)^* : n < \omega\}$  is a counterattack to  $\sigma$ , and thus loses. This means  $B = \bigcup \operatorname{range}(\alpha^*) \notin \mathcal{B}$ .

We consider two cases. The first is the case that  $\bigcup \operatorname{range}(\alpha^*)$  is finite. Noting that  $\alpha^*(m) \cap \alpha^*(n) = \emptyset$  whenever  $m \neq n$ , there exists  $N < \omega$  such that  $\alpha^*(n) = \emptyset$  for all n > N. As a result,  $\bigcup \operatorname{range}(\alpha) = \bigcup \operatorname{range}(\alpha \upharpoonright n)$ , and thus  $\bigcup \operatorname{range}(\alpha)$  is finite, and therefore not in  $\mathcal{B}$ .

In the other case,  $\bigcup \operatorname{range}(\alpha^*) \notin \mathcal{B}$  is an infinite subset of  $\bigcup \operatorname{range}(\alpha)$ , and therefore  $\bigcup \operatorname{range}(\alpha) \notin \mathcal{B}$  as well. Thus we have shown that  $\tau$  is a winning strategy for I in  $G_{fin}(\mathcal{A}, \mathcal{B})$ .

Theorem 10. Let  $\mathcal B$  be sequence-like. Then II  $\uparrow G_{<2}(\mathcal A,\mathcal B)$  if and only if II  $\uparrow$   $G_{fin}(\mathcal A,\mathcal B)$ .

Theorem 11. Let  $\mathcal B$  be sequence-like. Then II  $\uparrow_{mark} G_{<2}(\mathcal A,\mathcal B)$  if and only if II  $\uparrow_{mark} G_{fin}(\mathcal A,\mathcal B)$ .

Theorem 12. Let  $\mathcal{A}$  be almost-sequence-like and  $\mathcal{B}$  be sequence-like. Then I  $\uparrow$   $G_{fin}(\mathcal{A},\mathcal{B})$  if and only if I  $\uparrow$   $G_{fin}(\mathcal{A},\mathcal{B})$ .

Proof. We assume  $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ . As  $\mathcal{A}$  is almost-sequence-like, there is a strategy  $\sigma$  witnessing  $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$  where  $|\sigma(s)| = \aleph_0$  and  $\sigma(s) \cap \bigcup \operatorname{range}(s) = \emptyset$  (that is,  $\sigma$  never replays the choices of II) for all partial plays s by II.

For each  $s \in \omega^{<\omega}$ , suppose  $F_{s \upharpoonright m} \in [\bigcup A]^{<\aleph_0}$  is defined for each  $0 < m \le |s|$ .

Then let  $s^* : |s| \to [\bigcup A]^{<\aleph_0}$  be defined by  $s^*(m) = F_{s \upharpoonright m+1}$ , and define  $\tau' : \omega^{<\omega} \to A$  by  $\tau'(s) = \sigma(s^*)$ . Finally, set  $[\sigma(s^*)]^{<\aleph_0} = \{F_{s \smallfrown \langle n \rangle} : n < \omega\}$ , and for some bijection  $b : \omega^{<\omega} \to \omega$  let  $\tau(n) = \tau'(b(n))$  be a predetermined strategy for I in  $G_{fin}(A, \mathcal{B})$ .

Suppose  $\alpha$  is a counterattack by II against  $\tau$ , so

$$\alpha(n) \in [\tau(n)]^{<\aleph_0} = [\tau'(b(n))]^{<\aleph_0} = [\sigma(b(n)^\star)]^{<\aleph_0}$$

It follows that  $\alpha(n) = F_{b(n) \cap \langle m \rangle}$  for some  $m < \omega$ . In particular, there is some infinite subset  $W \subseteq \omega$  and  $f \in \omega^{\omega}$  such that  $\{\alpha(n) : n \in W\} = \{F_{f \mid n+1} : n < \omega\}$ . Note here that  $(f \mid n+1)^* = (f \mid n)^* \cap \langle F_{f \mid n+1} \rangle$ . This shows that  $F_{f \mid n+1} \in [\sigma((f \mid n)^*)]^{<\aleph_0}$  is an attempt by II to defeat  $\sigma$ , which fails. Thus  $\bigcup \{F_{f \mid n+1} : n < \omega\} = \bigcup \{\alpha(n) : n \in W\} \notin \mathcal{B}$ , and since this set is infinite (as  $\sigma$  prevents II from repeating choices) we have  $\bigcup \{\alpha(n) : n < \omega\} \notin \mathcal{B}$  too. Therefore  $\tau$  is winning.

**Proposition 13.** Let  $\mathcal{B}$  be sequence-like,  $\mathcal{A} \subseteq \mathcal{B}$ , and I  $\underset{pre}{\not}$   $G_{fin}(\mathcal{A}, \mathcal{B})$ . Then  $\mathcal{A}$  is

11 almost-sequence-like.

Theorem 12.

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- 112 Proof. Let  $A \in \mathcal{A}$ , and for all  $n < \omega$  let  $A_n = A$ . Then  $A_n$  is not a winning predetermined strategy for I, so II may choose finite subsets  $B_n \subseteq A_n$  such that 114  $A'_n \bigcup \{B_n : n < \omega\} \in \mathcal{B} \subseteq \mathcal{A}$ .
  115 It follows that  $|A'_n| = \aleph_0$ , and for any infinite subset  $A''_n \subseteq A'_n$  (in particular, 116 any cofinite subset),  $A''_n \in \mathcal{B} \subseteq \mathcal{A}$ . Thus  $\mathcal{A}$  is almost-sequence-like.
- any cofinite subset),  $A''_n \in \mathcal{B} \subseteq \mathcal{A}$ . Thus  $\mathcal{A}$  is almost-sequence-like.  $\square$ Corollary 14. Let  $\mathcal{B}$  be sequence-like and  $\mathcal{A} \subseteq \mathcal{B}$ . Then  $I \uparrow G_{fin}(\mathcal{A}, \mathcal{B})$  if and only if  $I \uparrow_{pre} G_{fin}(\mathcal{A}, \mathcal{B})$ .
- 119 *Proof.* Assuming I  $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$ , we have I  $\uparrow G_{fin}(\mathcal{A}, \mathcal{B})$  by Proposition 13 and

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