The recursion principle is the rule that states that $\operatorname{rec}_{A\times B}(f): A\times B\to C$ is well defined, and it is taken as a primitive notion. We are asked to instead take the well-definedness of pr_1 and pr_2 as primitive instead and to derive the recursion principle.

Definitions of $\mathsf{rec}_{A \times B}, \, \mathsf{pr}_1, \, \mathsf{and} \, \mathsf{pr}_2$ are given below for convenience.

$$\begin{split} \operatorname{rec}_{A\times B} &: (A\to B\to C)\to A\times B\to C\\ \operatorname{rec}_{A\times B}(f)((a,b)) &:= f(a)(b) \\ \\ \operatorname{pr}_1 &: A\times B\to A\\ \operatorname{pr}_1 &:= \operatorname{rec}_{A\times B}(\lambda\, a\, b\mapsto a) \\ \\ \operatorname{pr}_2 &: A\times B\to B\\ \operatorname{pr}_2 &:= \operatorname{rec}_{A\times B}(\lambda\, a\, b\mapsto b) \end{split}$$

(In agreement with the notation used in the text, $\lambda \, a \, b \mapsto \Phi$ is merely shorthand for $\lambda \, a \mapsto (\lambda \, b \mapsto \Phi)$. We can think of such an expression as a "two-variable" curried function [mmmm, curry].)

We need primitive definitions of pr_1 and pr_2 .

$$\begin{aligned} \operatorname{pr}_1: A \times B &\to A \\ \operatorname{pr}_1(a,b) := a \\ \\ \operatorname{pr}_2: A \times B &\to B \\ \operatorname{pr}_2(a,b) := b \end{aligned}$$

We define

$$\begin{split} \operatorname{rec}_{A\times B}': (A\to B\to C) \to A\times B \to C \\ \operatorname{rec}_{A\times B}'(f)((a,b)) &:= f(\operatorname{pr}_1(a,b))(\operatorname{pr}_2(a,b)) \end{split}$$

We have

$$\begin{split} \operatorname{rec}_{A\times B}(f)(a,b) &= f(a)(b) \\ &= f(\operatorname{pr}_1(a,b))(\operatorname{pr}_2(a,b)) \\ &= \operatorname{rec}'_{A\times B}(f)(a,b) \end{split}$$

So $\operatorname{rec}_{A\times B} = \operatorname{rec}'_{A\times B}$.