

The *recursion principle* is the rule that states that  $\text{rec}_{A \times B}(f) : A \times B \rightarrow C$  is well defined, and it is taken as a primitive notion. We are asked to instead take the well-definedness of  $\text{pr}_1$  and  $\text{pr}_2$  as primitive instead and to derive the recursion principle.

Definitions of  $\text{rec}_{A \times B}$ ,  $\text{pr}_1$ , and  $\text{pr}_2$  are given below for convenience.

$$\begin{aligned}\text{rec}_{A \times B} &: (A \rightarrow B \rightarrow C) \rightarrow A \times B \rightarrow C \\ \text{rec}_{A \times B}(f)((a, b)) &:= f(a)(b)\end{aligned}$$

$$\begin{aligned}\text{pr}_1 &: A \times B \rightarrow A \\ \text{pr}_1 &:= \text{rec}_{A \times B}(\lambda a b \mapsto a)\end{aligned}$$

$$\begin{aligned}\text{pr}_2 &: A \times B \rightarrow B \\ \text{pr}_2 &:= \text{rec}_{A \times B}(\lambda a b \mapsto b)\end{aligned}$$

(In agreement with the notation used in the text,  $\lambda a b \mapsto \Phi$  is merely shorthand for  $\lambda a \mapsto (\lambda b \mapsto \Phi)$ . We can think of such an expression as a “two-variable” curried function [mmmm, curry].)

We need primitive definitions of  $\text{pr}_1$  and  $\text{pr}_2$ .

$$\begin{aligned}\text{pr}_1 &: A \times B \rightarrow A \\ \text{pr}_1(a, b) &:= a\end{aligned}$$

$$\begin{aligned}\text{pr}_2 &: A \times B \rightarrow B \\ \text{pr}_2(a, b) &:= b\end{aligned}$$

We define

$$\begin{aligned}\text{rec}'_{A \times B} &: (A \rightarrow B \rightarrow C) \rightarrow A \times B \rightarrow C \\ \text{rec}'_{A \times B}(f)((a, b)) &:= f(\text{pr}_1(a, b))(\text{pr}_2(a, b))\end{aligned}$$

We have

$$\begin{aligned}\text{rec}_{A \times B}(f)(a, b) &= f(a)(b) \\ &= f(\text{pr}_1(a, b))(\text{pr}_2(a, b)) \\ &= \text{rec}'_{A \times B}(f)(a, b)\end{aligned}$$

So  $\text{rec}_{A \times B} = \text{rec}'_{A \times B}$ .