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A Study of the Synchronisation and Concurrency Issues in the Dining Philosophers’ Problem completed using the ThreadMentor Visualisation Tool

**Submitted by: Habiba Nour, B00151078**

**Steven Kelly, B00150588**

**Rochelle Mullen, B00156311**

**Piotr Momat, B00156112**

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Author: Habiba Nour Dated: 18/MARCH/2024

Author: Steven Kelly Dated: 18/MARCH/2024

Author: Rochelle Mullen Dated: 18/MARCH/2024

Author: Piotr Momat Dated: 18/MARCH/2024

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Note: Please make sure that you use proper citation within the body of the text of your report + include the reference cited in the References section at the end. It is **not** sufficient simply to list the references at the end. See:

<https://libguides.ucd.ie/harvardstyle/introduction>

# Introduction (Steven)

## Background

* Concurrency and Synchronisation Issues in OSes (Steven)
* Mutex Locks (Rochelle)
* Semaphores (Habiba)

Semaphores can be viewed as an extension to mutex locks. “Semaphores used to solve the critical section problems and to active process synchronization in multiprocessing environment”. It has two methods “Wait” and “Signal”(atomic), a private integer counter, and a private queue (of threads).

Two possibilities when “Wait” is executed by a thread

- The counter of S (Semaphore) is positive: in this case, the counter is decreased by one and the thread resumes its execution.

- The counter of S is zero: in this case, the thread is suspended and put into the private queue of S.

Two possibilities when “Signal” is executed by a thread

- The queue has no waiting thread: the counter of S increased by one and the thread resumes its execution.

- The queue of S has a waiting thread: in this case, the counter of S must be zero. one of the waiting threads will be allowed to leave the queue and resume its execution.

When the element is one that means that the element could be used but if it is zero that means that the element has to wait. All the elements are supposed to be initialized to 1.

As shown above that “Wait” and “Signal” are atomic which means once the activities of “Wait” start, they will continue with no interruption. There are many steps for “Wait” and “Signal”, these steps are considered as a single non-interruption instruction. The same thing applies to “Signal”. If more than one threads try to execute “Wait” or “Signal”, only one of them will succeed.

**How “wait” and “Signal” are working**

“Wait”: out of many threads, only one thread can successfully execute "Wait." This will cause the counter to decrease by 1 and enter the critical section. Once the thread enters the critical section, the counter becomes 0, and as a result, all subsequent attempts at executing the wait will be blocked.

“Signal”: when the thread exits the critical section, it executes "Signal." If there are threads waiting, only one of them will release and enter the critical section, but the counter will not increase by 1, the “Signal” will not increase, and the “Wait” will not decrease. If there are no waiting threads, the execution of “Signal” causes the value of the counter to increase by 1. Then the next thread that executes “Wait” can enter the critical section.

This type of semaphore is called a “binary semaphore,” and the counter is either 1 or 0. Binary semaphore can be used to control access to a single resource which can be used by one process.

* ThreadMentor (Piotr)

## The Dining Philosophers Problem (Habiba)

A diagram of a circular object with text and symbols

Description automatically generatedDining philosophers problem happens when we have food and a chopstick in front of each philosopher. When one of the philosophers gets hungry, he will pick up his chopsticks and start eating. All the philosophers have to have two chopsticks to eat, and in our case, only two philosophers will eat.

e.g., if philosophers P5 and P2 want to eat, P5 will pick up chopsticks one and five and P2 will pick up chopsticks two and three, leaving P1, P3, and P4 with either one or no chopstick, and they have to wait until one of them finishes their food and puts down his chopstick.

This scenario causes problems because we want all of them to eat in order. To solve the philosopher's problem, we are using either mutex locks or semaphore.

## Outline/Layout of your Report

* In section 2, we describe “this”. In section 3, we describe “that”…
* Tie the sections together: *briefly* describe how they are related

# <Your Solution>

* Theory/How it works (Rochelle)
* ThreadMentor: (Piotr)
  + Makefile(Piotr)
  + Explanation of “Tags” and other ThreadMentor issues related to your solution (Piotr)

# Results and Analysis (Piotr & Steven)

Screenshots of a **single** program run, to include:

* Several History Graph screenshots + proper captions
* Screenshots of highlighted code corresponding to various ThreadMentor tags for Philosopher threads in the History graph above + proper captions
* Screenshots of main ThreadMentor window showing relevant and corresponding information relating to the above.
* Screenshots of Thread Status window showing relevant and corresponding information relating to the above.
* Detailed descriptions in the text of what is happening to each of the Philosopher threads making reference to the screenshots.

# Conclusions (Habiba & Rochelle)

Technical conclusions about your solution; for example:

* “our solution avoids a specified problem”, or
* “our does had a specified problem”
* Advantages/Disadvantages of our solution

Advantages

Disadvantages:

* Semaphore can be complex, as any mistake or error will cause synchronisation problems.
* Might need more resources, such as additional memory or CPU.
* All the threads might be racing over which read we should access the critical section first, and that can lead to unpredictable behaviour.
* Can cause deadlock as all the threads could be initialised to zero, causing them to be waiting forever.

The conclusions should **not**include things like: “I loved this project!”, “I hated this project!”, “I learned x, y and z on this project”. These things should go in the Personal Reflections section.

# References

The references listed here **must** be cited within the text of your report.

It is not good enough simply to list the references here and have no citations in the text – marks **will** be lost if you leave out the citations or if you use Wikipedia as a reference source.

# Appendix: Personal Reflections

* A brief summary of what you learned
* What you liked about the project
* What you didn’t like about the project
* What would you have done differently if you could do it again
* Any other recommendations/feedback for the Lecturer and/or future Students

# Appendix: Project Planning and Management

(Ask if we can put Diary here)

* Gantt chart
* Description of your Gantt chart
* How you managed the project on a week-to-week basis