

# CMSC303 Introduction to Theory of Computation, VCU

## Spring 2017, Assignment 1 Solutions

Total marks: 26 marks + 3 bonus marks for LaTeX

NOTE: As this is a warmup assignment intended to refresh your memory on background material, it will be marked only for completeness, *not* correctness. It is your responsibility to compare your answers with the solutions (to be posted after the due date) to gauge how well you understand the concepts on this assignment.

### 1 Exercises

1. (6 marks) Sipser, Ex. 0.3: Let  $A$  be the set  $\{x, y, z\}$  and  $B$  be the set  $\{x, y\}$ .

- (a) (1 mark) Is  $A$  a subset of  $B$ ?
- (b) (1 mark) Is  $B$  a subset of  $A$ ?
- (c) (1 mark) What is  $A \cup B$ ?
- (d) (1 mark) What is  $A \cap B$ ?
- (e) (1 mark) What is  $A \times B$ ?
- (f) (1 mark) What is the power set of  $B$ ?

Solution: (a) No. (b) Yes. (c)  $A \cup B = \{x, y, z\}$ . (d)  $A \cap B = \{x, y\}$ . (e)  $A \times B = \{(x, x), (x, y), (y, x), (y, y), (z, x), (z, y)\}$ . (f)  $P(B) = \{\emptyset, \{x\}, \{y\}, \{x, y\}\}$ .

2. (2 marks) Sipser, Ex. 0.4: If  $A$  has  $a$  elements and  $B$  has  $b$  elements, how many elements are in  $A \times B$ ? Explain your answer.

Solution:  $|A \times B| = ab$ . Proof: The elements of  $A \times B$  are obtained by considering all possible pairings between some  $x \in A$  and some  $y \in B$ . Since  $A$  and  $B$  have sizes  $a$  and  $b$ , respectively, the total number of such pairings is  $ab$ .

3. (2 marks) Sipser, Ex. 0.5: If  $C$  is a set with  $c$  elements, how many elements are in the power set of  $C$ ? Explain your answer.

Solution:  $P(C) = 2^c$ . Proof: The elements of  $P(C)$  are obtained by considering all possible subsets of  $C$ . To generate all possible subsets, for each element  $x \in C$ , we must consider two options: Either we include  $x$  in the subset, or we do not. Thus, the total number of ways of generating a subset is  $2^c$ , since there are  $c$  elements, and each offers us two distinct ways to create a subset. (Alternatively, one can view this as assigning a bit to each  $x \in C$ . If and only if the bit  $i$  is set to 1, we think of  $x_i$  as being in our subset. The number of distinct settings for  $c$  bits is then  $2^c$ .)

4. (7 marks) Sipser, Ex. 0.6: Let  $X$  be the set  $\{1, 2, 3, 4, 5\}$  and  $Y$  be the set  $\{6, 7, 8, 9, 10\}$ . The unary function  $f : X \mapsto Y$  and the binary function  $g : X \times Y \mapsto Y$  are described in the following tables.

$n$	$f(n)$	$g$	6	7	8	9	10
1	6	1	10	10	10	10	10
2	7	2	7	8	9	10	6
3	6	3	7	7	8	8	9
4	7	4	9	8	7	6	10
5	6	5	6	6	6	6	6

- (a) (1 mark) What is the value of  $f(2)$ ?

- (b) (2 marks) What are the domain and co-domain of  $f$ ?
- (c) (1 mark) What is the value of  $g(2, 10)$ ?
- (d) (2 marks) What are the domain and co-domain of  $g$ ?
- (e) (1 mark) What is the value of  $g(4, f(4))$ ?

Solution: (a)  $f(2) = 7$ . (b)  $\text{Co-Domain}(f) = \{6, 7, 8, 9, 10\} = Y$ .  $\text{Domain}(f) = X$ . (c)  $g(2, 10) = 6$ . (d)  $\text{Co-Domain}(g) = Y$ .  $\text{Domain}(g) = X \times Y$ . (e)  $g(4, f(4)) = 8$ .

5. (3 marks) Sipser, Ex. 0.7: For each part, give a relation that satisfies the condition.

- (a) (1 mark) Reflexive and symmetric but not transitive. Solution: E.g.,  $\text{differenceIsZeroOrOdd}(x, y)$ , which given  $x, y \in \mathbb{Z}$ , asks whether  $x - y$  is either 0 or odd. Reflexivity and symmetry are trivial, transitivity breaks because  $(1, 2)$  and  $(2, 3)$  are in the relation, but  $(1, 3)$  is not. Another example is  $\text{haveHadLunchTogether}(x, y)$ , which asks if person  $x$  and person  $y$  have ever had lunch together. (many possible answers)
- (b) (1 mark) Reflexive and transitive but not symmetric. Solution: E.g.,  $\text{lessThanOrEqualTo}(x, y)$ , which asks whether (say) for  $x, y \in \mathbb{Z}$ , is  $x \leq y$ ? (many possible answers)
- (c) (1 mark) Symmetric and transitive but not reflexive. Solution: E.g., let  $D = \{a, b, c\}$ . Then, we can define a relation  $R : D \times D \mapsto \{\text{TRUE}, \text{FALSE}\}$  as  $\{(a, b), (b, a), (a, c), (c, a), (b, c), (c, b)\}$ .

## 2 Problems

1. (2 marks) Sipser, Prob. 0.12 (0.11 in 2nd edition): Find the error in the following proof that all horses are the same color.

CLAIM: In any set of  $h$  horses, all horses are the same color.

PROOF: By induction on  $h$ .

**Base Case:** For  $h = 1$ . In any set containing just one horse, all horses clearly are the same color.

**Induction Step:** For  $k \geq 1$  assume that the claim is true for  $h = k$  and prove that it is true for  $h = k + 1$ . Take any set  $H$  of  $k + 1$  horses. We show that all the horses in this set are the same color. Remove one horse from this set to obtain the set  $H_1$  with just  $k$  horses. By the induction hypothesis, all the horses in  $H_1$  are the same color. Now replace the removed horse and remove a different one to obtain the set  $H_2$ . By the same argument, all the horses in  $H_2$  are the same color. Therefore all horses in  $H$  must be the same color, and the proof is complete.

Solution: The problem is that the inductive step does *not* work for all  $h \geq 2$ ; in particular, it fails for  $h = 2$ . This is because when  $h = 2$ , we have  $H_1 \cap H_2 = \emptyset$ , and the argument here works only if we can guarantee  $H_1 \cap H_2 \neq \emptyset$  (can you see why?).

2. (4 marks) Prove using induction that

$$\sum_{m=0}^n m = \frac{n(n+1)}{2}.$$

Solution: Observation: Since when  $m = 0$ , no value is contributed to the sum, we can instead evaluate the sum  $\sum_{m=1}^n m = \frac{n(n+1)}{2}$ .

**Base Case** ( $n = 1$ ): Here,  $\sum_{m=1}^1 m = 1$ , which also equals  $1 \cdot (1 + 1)/2$ .

**Inductive Hypothesis:** Assume the claim holds for  $n = k$  for any  $k \geq 1$ .

**Inductive Step:** We prove the claim holds for  $n = k + 1$ . Specifically,

$$\sum_{m=1}^{k+1} m = (1 + 2 + \cdots + k) + k + 1 = \frac{k(k+1)}{2} + k + 1 = \frac{k^2 + k + 2k + 2}{2} = \frac{(k+1)(k+2)}{2},$$

as required, where the second equality holds by the Induction Hypothesis. Hence, by the principle of Mathematical Induction, the claim holds for all  $n \geq 1$ .