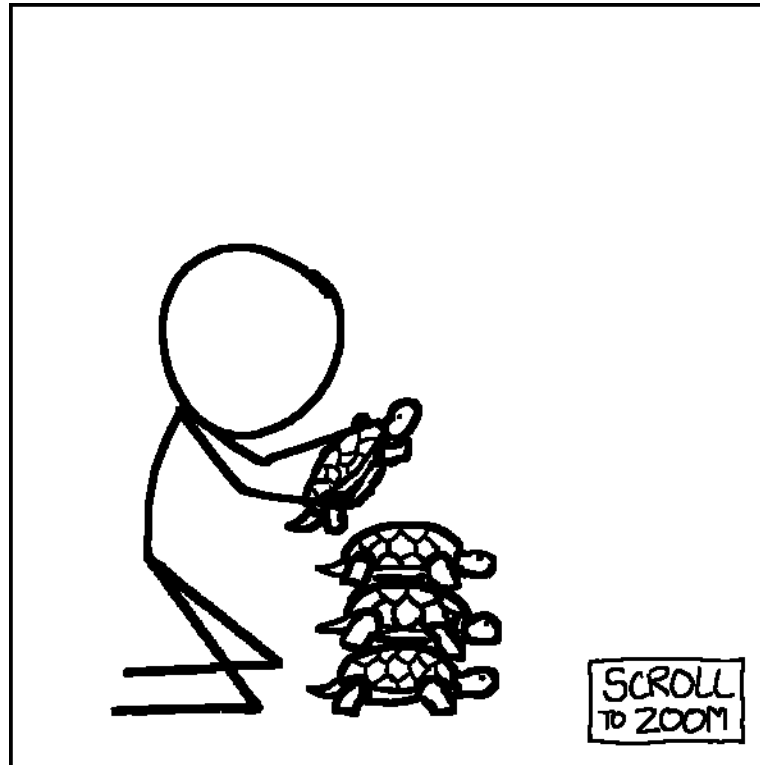


VE280 Programming and Elementary Data Structures

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Linear List; Stack



Learning Objectives

- Understand what is a linear list and what is a stack
- Know how they can be implemented
- Discover some applications of the stack data structure

Outline

- Linear List
- Stack
 - Implementation
 - Application

Linear List ADT

- Recall the IntSet ADT
 - A collection of zero or more integers, with **no duplicates**.
 - It supports insertion and removal, but by value.
- A related ADT: linear list
 - A collection of zero or more integers; **duplicates possible**.
 - $L = (e_0, e_1, \dots, e_{N-1})$
 - It supports insertion and removal **by position**.

Linear List ADT

Insertion

```
void insert(int i, int v) // if  $0 \leq i \leq N$   
// (N is the size of the list), insert v at  
// position i; otherwise, throws BoundsError  
// exception.
```

Example: $L1 = (1, 2, 3)$

$L1.insert(0, 5) = (5, 1, 2, 3);$

$L1.insert(1, 4) = (1, 4, 2, 3);$

$L1.insert(3, 6) = (1, 2, 3, 6);$

$L1.insert(4, 0)$ throws **BoundsError**

Linear List ADT

Removal

```
void remove(int i) // if  $0 \leq i < N$  (N is  
// the size of the list), remove the i-th  
// element; otherwise, throws BoundsError  
// exception.
```

Example: `L2 = (1, 2, 3)`

`L2.remove(0) = (2, 3);`

`L2.remove(1) = (1, 3);`

`L2.remove(2) = (1, 2);`

`L2.remove(3) throws BoundsError`



A linear list ADT can be implemented with:

Select all the correct answers.

- **A.** an array.
- **B.** a singly-linked list.
- **C.** a doubly-linked list.
- **D.** any container.

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Stack

- A “pile” of objects where new object is put on **top** of the pile and the top object is removed first.
 - LIFO access: last in, first out.
 - Restricted form of a **linear list**: insert and remove only at the end of the list.



Methods of Stack

- **size()** : number of elements in the stack.
- **isEmpty()** : checks if stack has no elements.
- **push(Object o)** : add object **o** to the top of stack.
- **pop()** : remove the top object if stack is not empty; otherwise, throw **stackEmpty**.
- **Object &top()** : return a reference to the top element.

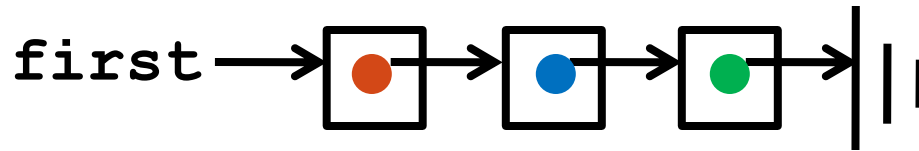
Stacks Using Arrays

Array[**MAXSIZE**] :

2	3	1	4		
---	---	---	---	--	--

- Maintain an integer **size** to record the size of the stack.
- **size():return size;**
- **isEmpty():return (size == 0);**
- **push(Object o):** add object **o** to the end of the array and increment **size**. Allocate more space if necessary.
- **pop():** If **isEmpty()**, throw **stackEmpty**; otherwise, decrement **size**.
- **Object &top():** return a reference to the top element **Array[size-1]**

Stacks Using Linked Lists

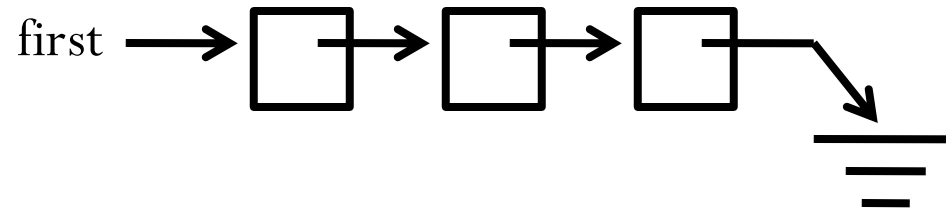


For single-ended linked list, which end is preferred to be the top? Why?

- `size() : LinkedList::size() ;`
- `isEmpty() : LinkedList::isEmpty() ;`
- `push(Object o) :` insert object at the beginning
`LinkedList::insertFirst(Object o) ;`
- `pop() :` remove the first node
`LinkedList::removeFirst() ;`
- `Object &top() :` return a reference to the object stored in the first node.

LinkedList::size()

- How to get the size of a linked list?



```
int LinkedList::size() {  
    int count = 0;  
    node *current = first;  
    while(current) {  
        count++;  
        current = current->next;  
    }  
    return count;  
}
```

Array vs. Linked List: Which is Better?

- They both have advantages and disadvantages
- Linked list
 - memory-efficient: a new item just needs extra constant amount of memory
 - not time-efficient for size operation
- Array
 - time-efficient for size operation
 - not memory-efficient: need to allocate a big enough array

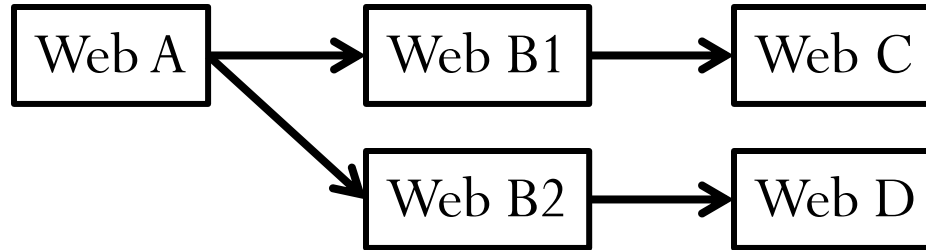
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Applications of Stacks

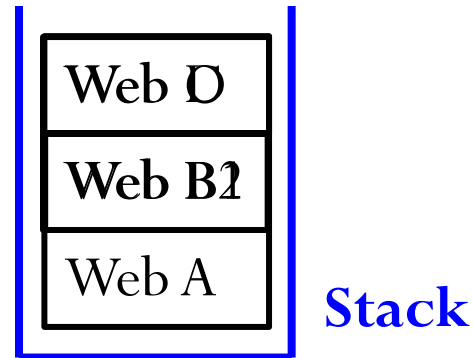
- Function calls in C++
- Web browser's “back” feature
- Parentheses Matching

Web Browser's “back” Feature



Visiting order

- Web A
- Web B1
- Web C
- Back (to Web B1)
- Back (to Web A)
- Web B2
- Web D



Parentheses Matching

- Output pairs (u, v) such that the left parenthesis at position u is matched with the right parenthesis at v .

$((a + b) * c + d - e) / (f + g)$
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18

- Output is: $(1, 5)$; $(0, 12)$; $(14, 18)$;

$(a + b)) * ((c + d)$
0 1 2 3 4 5 6 7 8 9 10 12

- Output is

$(0, 4)$;

Right parenthesis at 5 has no matching left parenthesis;

$(8, 12)$;

Left parenthesis at 7 has no matching right parenthesis

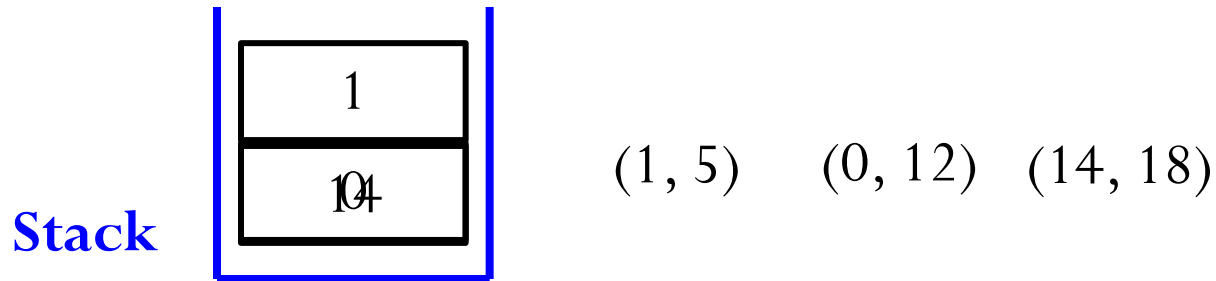
How to Realize Parentheses Matching?

((a + b) * c + d - e) / (f + g)
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18

- Scan expression from left to right.
- When a **left** parenthesis is encountered, push its position to the stack.
- When a **right** parenthesis is encountered, pop the top position from the stack, which is the position of the **matching left** parenthesis.
 - If the stack is empty, the **right** parenthesis is not matched.
- If string is scanned over but the stack is not empty, there are not-matched **left** parentheses.

Parentheses Matching

((a + b) * c + d - e) / (f + g)
0 1 2 3 4 5 6 7 8 9 10 12 14 16 18





A stack can be used:

Select all the correct answers.

- **A.** to manage any arithmetic expression.
- **B.** to undo operations (such as in a text editor).
- **C.** to reverse a list.
- **D.** to provide an efficient implementation of a queue.

Reference

- **Problem Solving with C++ (8th Edition)**, by *Walter Savitch*, Addison Wesley Publishing (2011)
 - Chapter 13.2 **Stack**