## Exercise session notes - Week 10

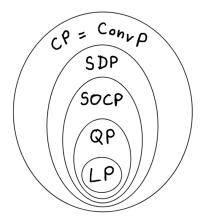
Second Order Cone Program + Robust Optimization

## 1 Convex Programs and Inclusions

Firstly we reviewed the types of convex programs that we studied so far.

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LP: \min f^{\top}x s.t Ax \le b ; Fx = g
QP: \min f^{\top}x s.t x^{\top}Ax + b^{\top}x \le c ; Fx = g
SOCP: \min f^{\top}x s.t \|Ax + b\|_2 \le c^{\top}x + d ; Fx = g
SDP: ...
CP: \min f^{\top}x s.t x \in K ; Fx = g
ConvP: \min f(x) s.t x \in C convex set ; Fx = g
```

We then pictured the diagram of all the inclusions between these sets.



The inclusions  $LP \subseteq QP$ ,  $SOCP \subseteq CP$  and  $CP \subseteq ConvP$  are trivial. Then we proved  $ConvP \subseteq CP$  and  $QP \subseteq SOCP$ , and next week we will consider also the Semi Definite Programs (SDP).

 $ConvP \subseteq CP$ . Consider a general convex program

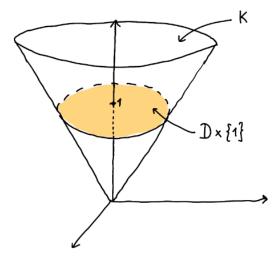
```
min f(x) s.t x \in C convex set; Fx = g
```

Then we can write an equivalent program as

```
\min \ t \quad \text{s.t} \quad t \geq f(x) \ ; \ x \in C \text{ convex set } ; \quad Fx = g \iff \min \ t \quad \text{s.t} \quad (x,t) \in \operatorname{epi}(f) \ ; \ (x,t) \in C \times \mathbb{R} \ ; \quad Fx = g
```

Let  $D := \operatorname{epi}(f) \cap (C \times \mathbb{R})$  and note that this is a convex set. Then we pick the conic hull  $K = \operatorname{cone}(D \times \{1\})$  which is

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Therefore  $(x,t) \in D \iff (x,t,1) \in E$ . Finally we can write the equivalent conic program as

min 
$$t$$
 s.t  $(x,t,s) \in K$ ;  $s=1$ ;  $Fx=g$ 

 $\mathbf{QP} \subseteq \mathbf{SOCP}$ . Consider a general quadratic program with A being a positive definite matrix

$$\min \ f^{\top}x \quad \text{s.t.} \quad x^{\top}Ax + b^{\top}x + c \le 0 \ ; \quad Fx = g$$

Recall that by the Cholesky Decomposition there exists a matrix B with  $A = B^{T}B$ . Then we can write an equivalent program as

min 
$$f^{\top}x$$
 s.t  $x^{\top}B^{\top}Bx \le -b^{\top}x - c$ ;  $Fx = g$   
min  $f^{\top}x$  s.t  $\|Bx\|_2^2 \le -b^{\top}x - c$ ;  $Fx = g$ 

Then recall the definition of the rotated Lorentz cone:

$$\mathcal{L}_{rot}^{m} = \left\{ (x, y, z) \mid \left\| x \right\|_{2}^{2} \leq 2yz \right\}$$

Therefore we can model the program as the following SOCP

min 
$$f^{\top}x$$
 s.t  $(Bx, -b^{\top}x - c, \frac{1}{2}) \in \mathcal{L}_{rot}^{m}$ ;  $Fx = g$ 

## 2 Robust Optimization

Assume we are working for an airline company and we want to assign people to fligths. We model this problem as an LP. In a perfect world we can write

max# people served s.t. assignment is feasible

Since flight disruptions may happen (flight cancellations, delays) we want to find an assignment that works for any possible scenario. Such a solution is called a Robust Solution. A non-example would be:

Rome 
$$\rightarrow$$
 Zurich (wait 15 min)  
Zurich  $\rightarrow$  London

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This solution is clearly feasible but a small delay could change entirely our assignment, therefore it is not Robust. Let's formalize this problem (Ex. 6.6 from Lecture Notes). We want to solve an LP

$$\min \quad c^{\top} x \quad \text{s.t.} \quad a^{\top} x \le b \ , \quad x \ge 0$$

but the vector  $a \in \mathbb{R}^n$  is uncertain, it varies in some Ellipse

$$a \in E = \{a_0 + Pu \mid ||u||_2 \le 1\}$$

for  $a_0$  the center and P the matrix corresponding to the stretch of the Ellipse. Our goal is to find a <u>Robust</u> solution, so a point x that is feasible for any choice of a and with minimal value. Therefore we may require

$$b \ge \sup_{a \in E} a^{\top} x$$

$$= a_0^{\top} x + \sup_{\|u\| \le 1} u^{\top} P^{\top} x$$

$$= a_0^{\top} x + \|P^{\top} x\|_2$$

Therefore we can model the robust LP as the following SOCP

min 
$$c^{\top}x$$
 s.t.  $a_0^{\top}x + \|P^{\top}x\|_2 \le b$ ,  $x \ge 0$