# Tools and Techniques, Lecture 3 Performance

- Algorithmic complexity
- Memory access patterns
- Benchmarking





CERN School of Computing 2025, Lund Sten Åstrand, Lund University

# Why?







#### How does it scale?

ALICE Event Processing Nodes Cluster

24640 CPU cores 2800 GPUs







## How do you scale?



Man making Python plots, color photograph



Image from imgflip.com





### Algorithmic complexity

Theoretical metric for estimating resource consumption (in our case runtime - time complexity).

"The time complexity of an algorithm represents the number of steps it has to take to complete."

- Complexity Theory. Brilliant.org. Retrieved 13:58, June 18, 2025, from <a href="https://brilliant.org/wiki/complexity-theory/">https://brilliant.org/wiki/complexity-theory/</a>

Also useful: Devopedia. 2022. "Algorithmic Complexity." Version 8, February 19. Accessed 2024-06-25. <a href="https://devopedia.org/algorithmic-complexity">https://devopedia.org/algorithmic-complexity</a>



#### Big O notation



 $\Rightarrow$  algorithm is on the order of f(n) or O(f(n))

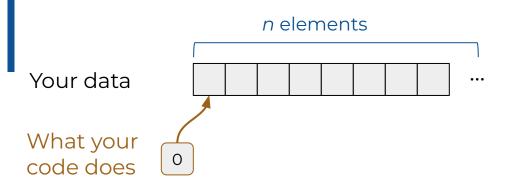
Example: if an algorithm's runtime grows linearly with the input size, the algorithm is O(n)

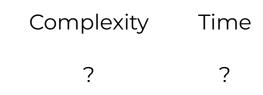
Strictly:  $O(f(n)) \Leftrightarrow$  "asymptotically bounded by f(n) up to some constant" (see Bachmann-Landau notation)

Big O Notation. Brilliant.org. Retrieved 13:54, June 18, 2025, from https://brilliant.org/wiki/big-o-notation/

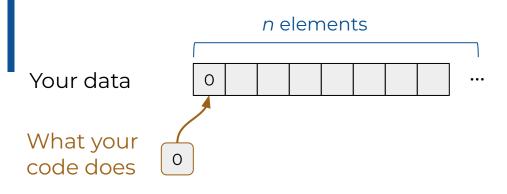


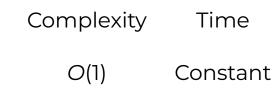




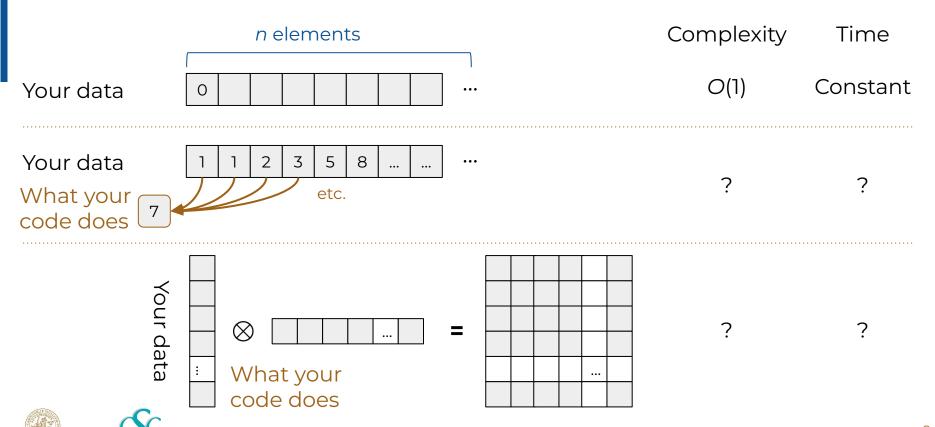


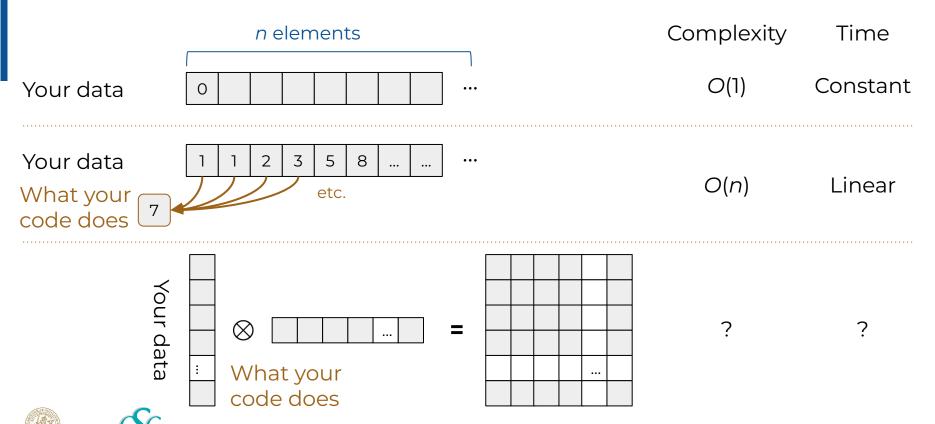


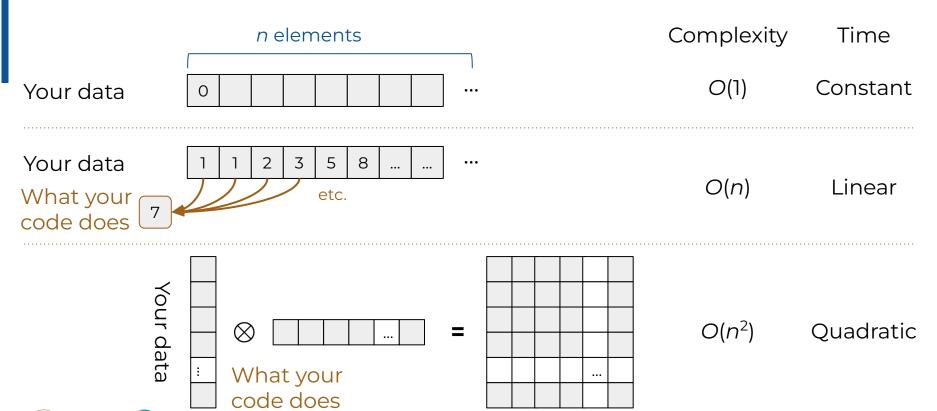
















**O(n)** - linear time (i.e. input size doubles = runtime doubles)

Radix sort

 $O(n^2)$  - quadratic time (input size doubles = runtimes quadruples)

- "For loop in a for loop"
- Outer product of two *n*-length vectors  $\Rightarrow$  *n*-by-*n* matrix
- Selection sort, insertion sort



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Radix sort aktshually O(k·n)

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- "For loop in a for loop"
- Outer product of two *n*-length vectors  $\Rightarrow$  *n*-by-*n* matrix
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#### O(log, n) - logarithmic time

Binary search of a sorted array

 $O(n \log_2 n)$  - log-linear or "linearithmic" time (close to O(n) for "reasonably sized" n)

Many sorting algorithms, e.g. merge sort, heap sort

#### $O(n^3)$ - cubic time

- "For loop in a for loop in a for loop"
- "Schoolbook" matrix multiplication

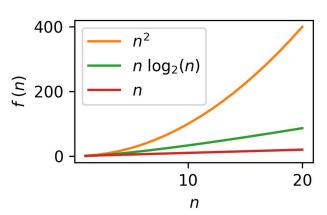




O(log, n) - logarithmic time

Binary search of a sorted array

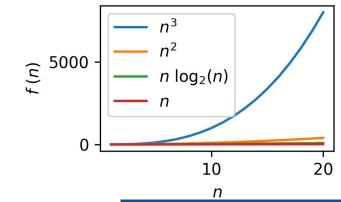
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 $O(n^3)$  - cubic time

- "For loop in a for loop in a for loop"
- "Schoolbook" matrix multiplication







## Two really Big O's (to avoid if possible)

 $O(2^n)$  and O(n!) - exponential time and factorial time

- Brute-force search, combinatorics, NP-hard problems
- Finding prime numbers



## Matrix multiplication

- "Schoolbook" matrix multiplication  $O(n^3)$
- The Strassen algorithm (1969)  $O(n^{2.805})$
- State of the art (2023)  $O(n^{2.373})$

For 4 x 4 matrices:

Strassen algorithm: 49 multiplications

Google Deepmind AlphaEvolve algorithm: 48 multiplications

AlphaEvolve: A Gemini-powered coding agent for designing advanced algorithms <a href="https://deepmind.google/discover/blog/alphaevolve-a-gemini-powered-coding-agent-for-designing-advanced-algorithms/">https://deepmind.google/discover/blog/alphaevolve-a-gemini-powered-coding-agent-for-designing-advanced-algorithms/</a>





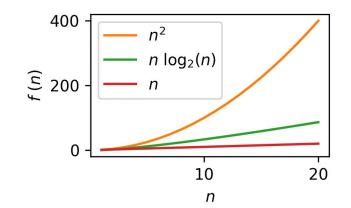
### Algorithmic complexity - conclusion

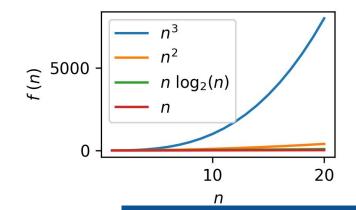
```
int binary search(const std::vector<int>& v, int target) {
    int left = 0, right = v.size() - 1;
    while (left <= right) {</pre>
        int mid = left + (right - left) / 2;
        if (v[mid] == target)
            return mid;
        else if (v[mid] < target)</pre>
            left = mid + 1;
        else
            right = mid - 1;
    return -1; // not found
```



#### Common code patterns

- For loop in a for loop
  - Solve differently?
  - Combine loops?
- Searching through data
  - o Consider sorting first?
  - Use look-up table?
- Recomputing values
  - Compute once, keep in a table?







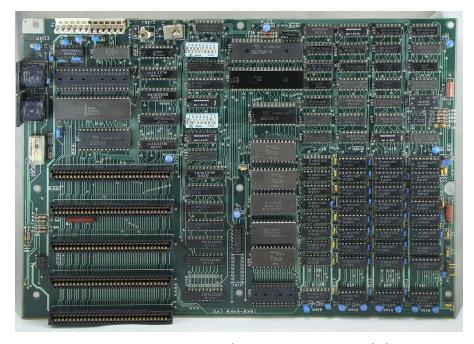


# Pit stop





#### Memory access patterns

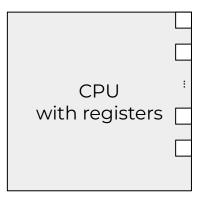


An IBM Personal Computer Model 5150 motherboard, 1981 - the first "PC". Image credit: user GermanX on Wikimedia Commons, under license <a href="https://doi.org/10.1007/journal.com/">https://doi.org/10.1007/journal.com/</a>





**CPU** registers

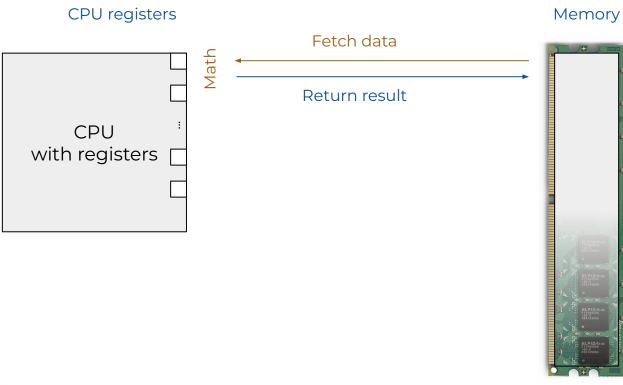


Memory



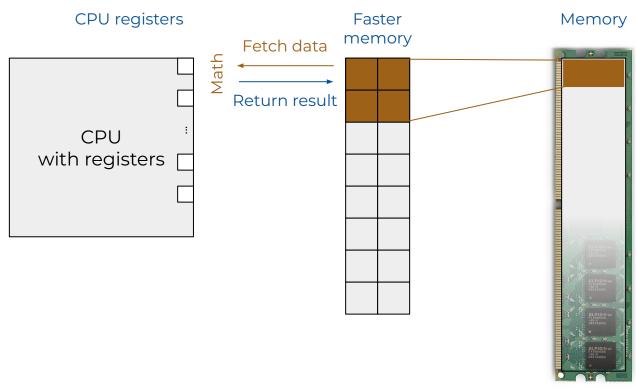






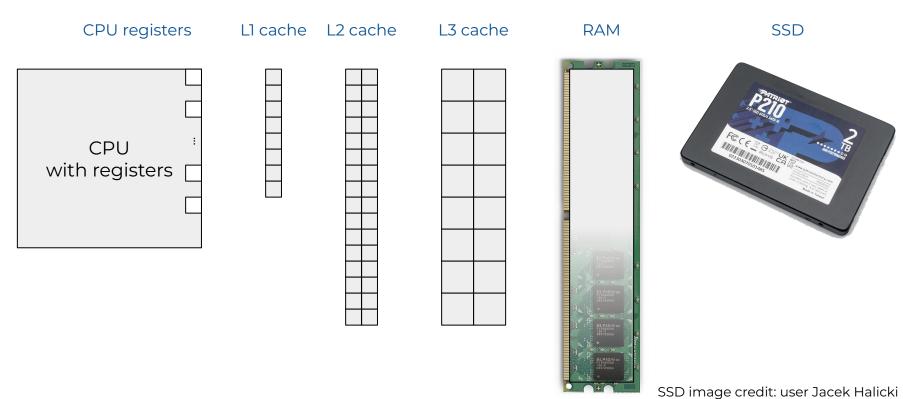
















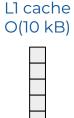
on Wikipedia, under

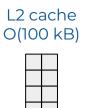
Tools and Techniques, Lecture 3 - Performance,

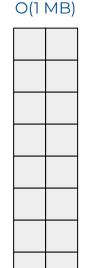
## CPU-to-memory layout

CPU registers 8 bytes

CPU with registers







L3 cache



RAM SSD O(10 GB) O(1 TB)







RAM image credit: user an-d on Wikipedia, under <u>Attribution-ShareAlike 3.0 Unported</u>

on Wikipedia, under

SSD image credit: user Jacek Halicki

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## CPU-to-memory layout

Typical sizes of different memory hardware

**CPU** registers L1 cache L2 cache L3 cache 8 bytes O(10 kB) O(100 kB) O(1 MB) CPU with registers 1ns Typical access times 10ns 4ns



**RAM** 







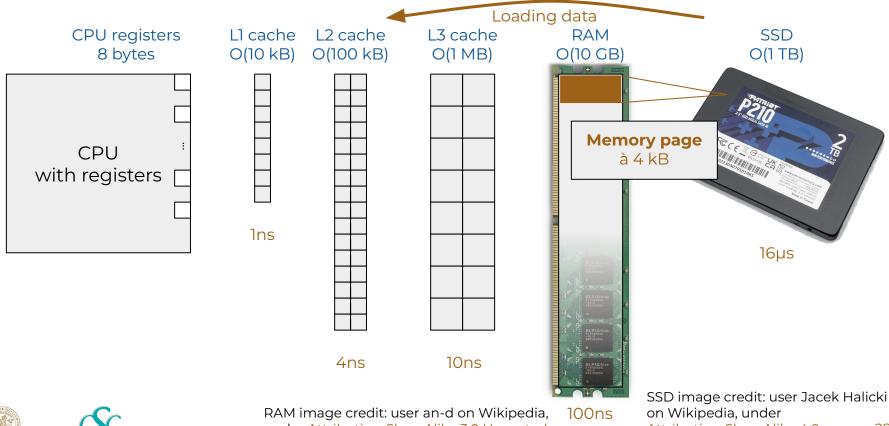
RAM image credit: user an-d on Wikipedia, under <u>Attribution-ShareAlike 3.0 Unported</u>

100ns

on Wikipedia, under Attribution-ShareAlike 4.0

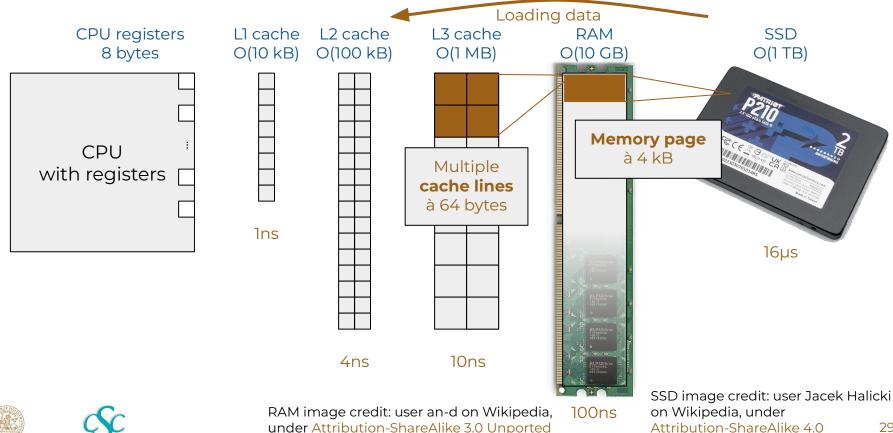
SSD image credit: user Jacek Halicki

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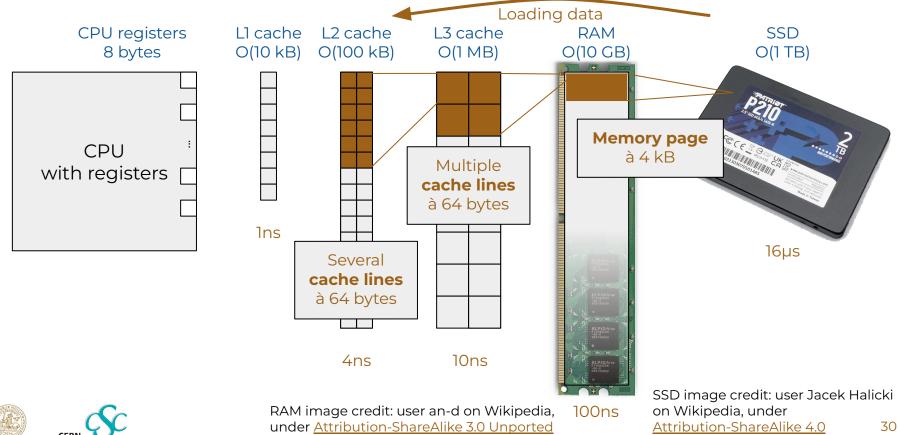






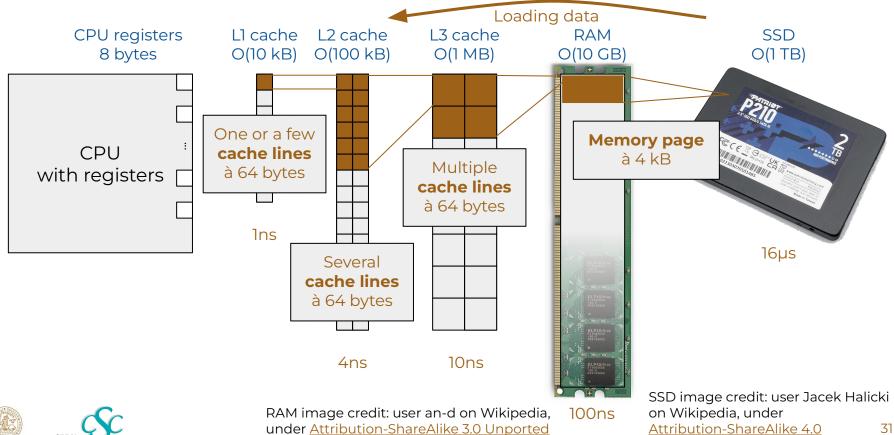


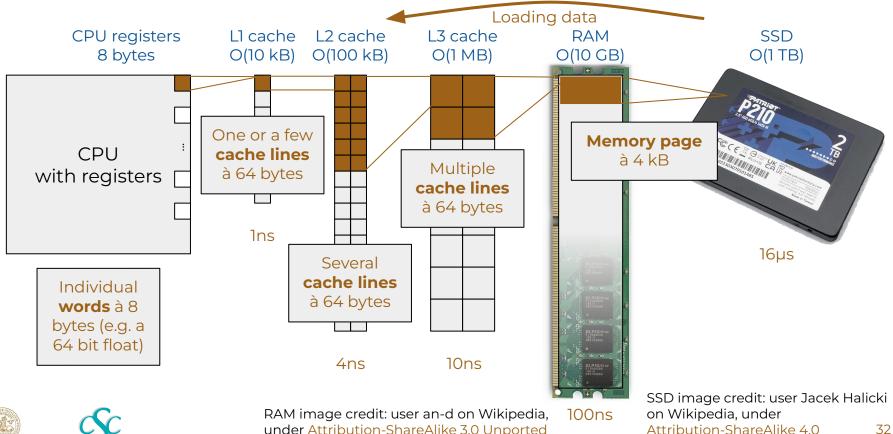




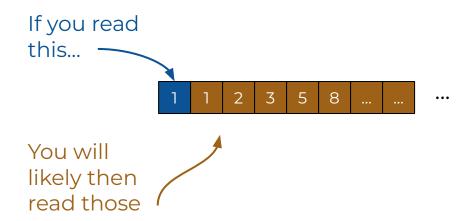








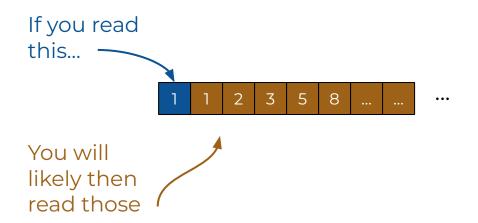
"Principle of locality" or "data locality" - data access often happens on many elements close to each other.







"Principle of locality" or "data locality" - data access often happens on many elements close to each other.



- space locality
- time locality

**Pre-fetching**: loading data that is likely to be needed soon





## Memory access patterns

Amount that fits in L1 cache (i.e. amount that is pre-fetched)

Your data





## Memory access patterns

1 ns

Amount that fits in L1 cache
(i.e. amount that is pre-fetched)

Your data

What your code does

"cache hits"



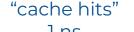


## Memory access patterns

Amount that fits in L1 cache (i.e. amount that is pre-fetched)



What your code does



1 ns

"cache miss"

4 ns





• • •

# Memory access patterns

1 ns

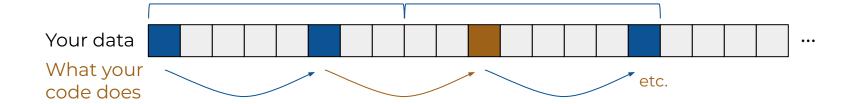
Amount that fits in L1 cache

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Your data

What your code does

"cache hits" "cache miss"

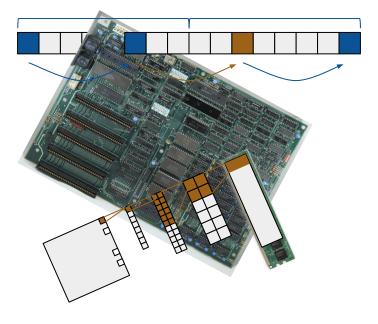


4 ns





## Memory access patterns - conclusions



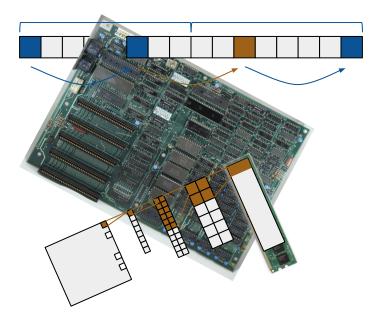


# Pit stop 2





# Benchmarking





## Benchmarking, pitfalls of

...but measuring is hard.

- Compiler optimizations
- Programs running in parallel ⇒ variance
- CPUs boost clock frequencies
- •

Takeaway: healthy scepticism and reasonable expectations





#### Linux time command

```
int main() {
        int temp;
        for (int i = 0; i < 1000000000; i++) {
            temp = i;
        }
        return temp;
}</pre>
```

```
$ g++ -03 time_example.cpp -o time_example
$ time ./time_example

real    0m0.009s
user    0m0.002s
sys    0m0.002s
```

- **real:** total time until your program finished
- user: time spent executing your program
- sys: time the system spent on behalf of your program





Note: blunt tool

#### Linux time command

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- real: total time until your program finished
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Note: blunt tool

# Compiler pitfalls

```
int main() {
    int size = 1000;
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```

```
int main() {
    int size = std::stoi(string_from_user);
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```



# Compiler pitfalls

```
int main() {
    int size = 1000;
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```

```
Array size: 1000
Steps taken: 1000000000
Time per step: 1.06009 ns
```

```
standard string-to-int
                   conversion
int main() {
    int size = std::stoi(string_from_user);
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
            index = i % size;
           // ... vector operations
Array size: 1000
Steps taken: 1000000000
```

⇒ loop 2x slower

Time per step: 2.02596 ns



# Compiler pitfalls

```
int main() {
     int size = 1000;
     int index = 0;
     for (int i = 0; i < 1000000000; i++) {
            index = i % size;
             // ... vector operations
                   movabs rsi, 2361183241434822607
                   .L4:
                   mov rdx, rcx
                   shr rdx
                   mov rax, rdx
                   mul rsi
                   mov rax, rcx
                   shr rdx, 4
                   imul rdx, rdx, 250
                   sub rax, rdx
                   movsx rax, DWORD PTR [rbp+0+rax*4]
                   add ebx, DWORD PTR [rbp+0+rax*4]
```

```
standard string-to-int
                     conversion
int main() {
     int size = std::stoi(string_from_user);
     int index = 0;
     for (int i = 0; i < 10000000000; i++) {
            l index = i % size;
            // ... vector operations
 mov rax, rcx
 cgo
 idiv rbp
 movsx rax, DWORD PTR [rbx+rdx*4]
 add esi, DWORD PTR [rbx+rax*4]
 mov r12d, esi
```

loop 2x slower





## Benchmarking tools

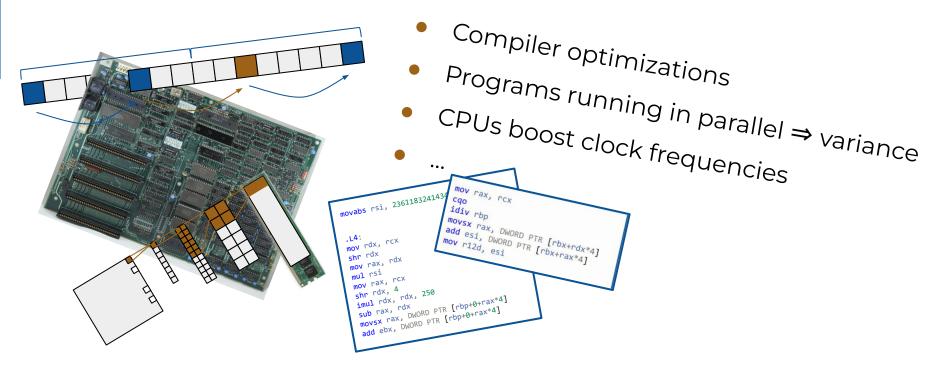
#### Some more sophisticated benchmarking tools:

- C/C++
  - Catch2
  - Google Benchmark
- Rust
  - o <u>bench</u>
- Python
  - <u>pytest-benchmark</u>
  - o <u>timeit</u>
- General purpose
  - <u>hyperfine</u>
  - <u>perf</u> (Linux)





## Benchmarking - conclusions



#### **Healthy scepticism and reasonable expectations**





### Conclusions - conclusions

# Tools and Techniques, Lecture 3 Performance

- Algorithmic complexity
- Memory access patterns
- Benchmarking





# Backup





#### Catch2

```
#include <catch2/catch test macros.hpp>
#include <catch2/benchmark/catch benchmark.hpp>
int arithmetic_sum (int upper) {
  volatile int sum = 0;
   for (int i = 1; i <= upper; i++) {
      sum += i;
   return sum;
TEST CASE ("Benchmarking Basics") {
  BENCHMARK ("Arithmetic sum 1 to 100") {
      return arithmetic sum (100);
```

#### Catch2

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#include <catch2/benchmark/catch benchmark.hpp>
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       return arithmetic sum (100);
```

including Catch2 library

function to benchmark

defining a set of tests or benchmarks

configuring a benchmark





#### Catch2

```
including Catch2
#include <catch2/catch test macros.hpp>
                                                              library
#include <catch2/benchmark/catch benchmark.hpp>
int arithmetic sum (int upper) {
  volatile int sum = 0;
   for (int i = 1; i <= upper; i++) {
                                                            function to
                                                            benchmark
       sum += i;
   return sum;
                                                       defining a set of tests
                                                          or benchmarks
TEST CASE ("Benchmarking Basics") {
   BENCHMARK ("Arithmetic sum 1 to 100") {
                                                              configuring a
       return arithmetic sum (100);
                                                               benchmark
```





## Catch2 output

```
Benchmarking Basics
/eos/user/k/kaastran/schools/CSC2025/prep/exercises/src/ex0.0 benchmarking basics.cpp:15
benchmark name
                                       low mean
                                                      high mean
                              std dev low std dev high std dev
Arithmetic sum 1 to 100
                               38.5609 ns 36.9933 ns 40.8075 ns
                               9.47688 ns 7.24546 ns 12.492 ns
test cases: 1 | 1 passed
assertions: - none -
```



