Tools and Techniques, Lecture 3 Performance

- Algorithmic complexity
- Memory access patterns
- Benchmarking





CERN School of Computing 2025, Lund Sten Åstrand, Lund University

Why?







How does it scale?

ALICE Event Processing Nodes Cluster

24640 CPU cores 2800 GPUs







How do you scale?



Man making Python plots, color photograph



Image from imgflip.com





Algorithmic complexity

Theoretical metric for estimating resource consumption (in our case runtime - time complexity).

"The time complexity of an algorithm represents the number of steps it has to take to complete."

- Complexity Theory. Brilliant.org. Retrieved 13:58, June 18, 2025, from https://brilliant.org/wiki/complexity-theory/

Also useful: Devopedia. 2022. "Algorithmic Complexity." Version 8, February 19. Accessed 2024-06-25. https://devopedia.org/algorithmic-complexity



Big O notation



 \Rightarrow algorithm is on the order of f(n) or O(f(n))

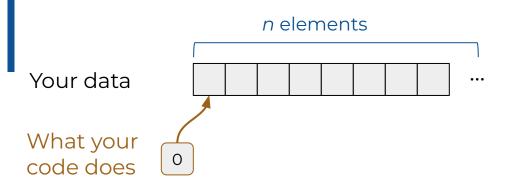
Example: if an algorithm's runtime grows linearly with the input size, the algorithm is O(n)

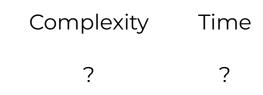
Strictly: $O(f(n)) \Leftrightarrow$ "asymptotically bounded by f(n) up to some constant" (see Bachmann-Landau notation)

Big O Notation. Brilliant.org. Retrieved 13:54, June 18, 2025, from https://brilliant.org/wiki/big-o-notation/

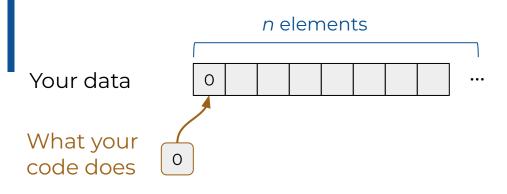


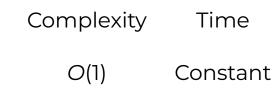




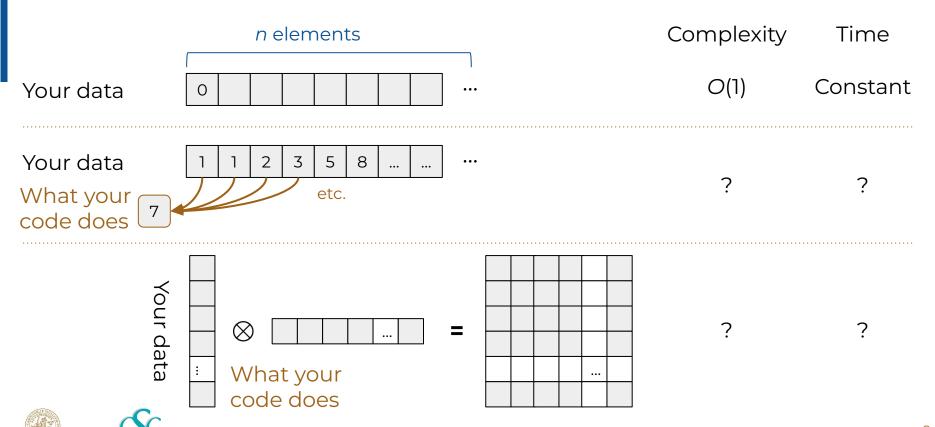


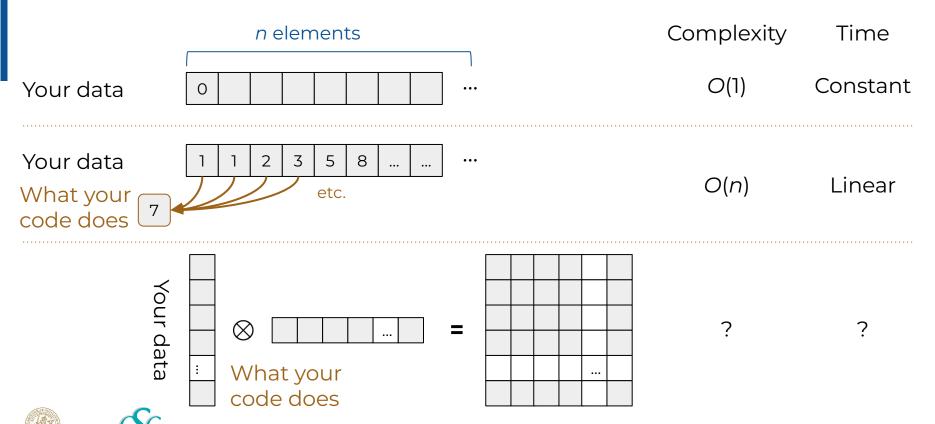


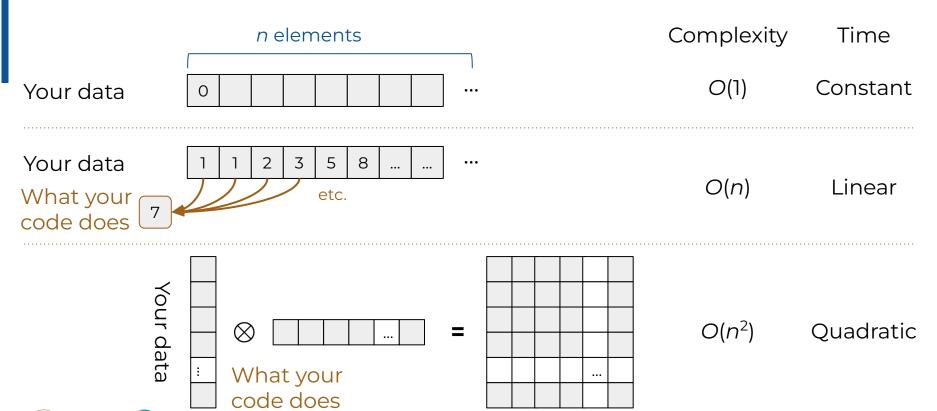
















O(n) - linear time (i.e. input size doubles = runtime doubles)

Radix sort

 $O(n^2)$ - quadratic time (input size doubles = runtimes quadruples)

- "For loop in a for loop"
- Outer product of two *n*-length vectors \Rightarrow *n*-by-*n* matrix
- Selection sort, insertion sort



O(n) - linear time (i.e. input size doubles = runtime doubles)

Radix sort aktshually O(k·n)

 $O(n^2)$ - quadratic time (input size doubles = runtimes quadruples)

- "For loop in a for loop"
- Outer product of two *n*-length vectors \Rightarrow *n*-by-*n* matrix
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O(log, n) - logarithmic time

Binary search of a sorted array

 $O(n \log_2 n)$ - log-linear or "linearithmic" time (close to O(n) for "reasonably sized" n)

Many sorting algorithms, e.g. merge sort, heap sort

$O(n^3)$ - cubic time

- "For loop in a for loop in a for loop"
- "Schoolbook" matrix multiplication

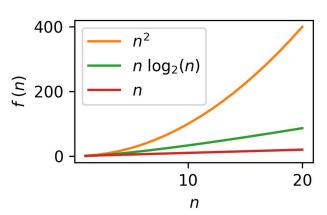




O(log, n) - logarithmic time

Binary search of a sorted array

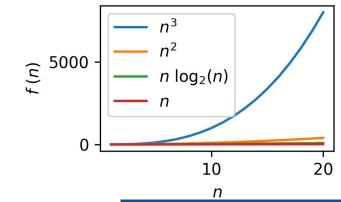
O(n log, n) - log-linear or "linearithmic" time (close to O(n) for "reasonably sized" n)



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 $O(n^3)$ - cubic time

- "For loop in a for loop in a for loop"
- "Schoolbook" matrix multiplication







Two really Big O's (to avoid if possible)

 $O(2^n)$ and O(n!) - exponential time and factorial time

- Brute-force search, combinatorics, NP-hard problems
- Finding prime numbers



Matrix multiplication

- "Schoolbook" matrix multiplication $O(n^3)$
- The Strassen algorithm (1969) $O(n^{2.805})$
- State of the art (2023) $O(n^{2.373})$

For 4 x 4 matrices:

Strassen algorithm: 49 multiplications

Google Deepmind AlphaEvolve algorithm: 48 multiplications

AlphaEvolve: A Gemini-powered coding agent for designing advanced algorithms https://deepmind.google/discover/blog/alphaevolve-a-gemini-powered-coding-agent-for-designing-advanced-algorithms/





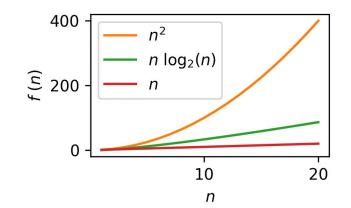
Algorithmic complexity - conclusion

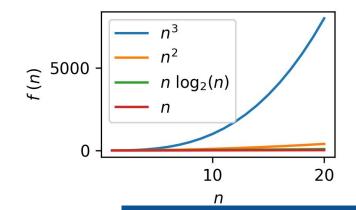
```
int binary search(const std::vector<int>& v, int target) {
    int left = 0, right = v.size() - 1;
    while (left <= right) {</pre>
        int mid = left + (right - left) / 2;
        if (v[mid] == target)
            return mid;
        else if (v[mid] < target)</pre>
            left = mid + 1;
        else
            right = mid - 1;
    return -1; // not found
```



Common code patterns

- For loop in a for loop
 - Solve differently?
 - Combine loops?
- Searching through data
 - o Consider sorting first?
 - Use look-up table?
- Recomputing values
 - Compute once, keep in a table?







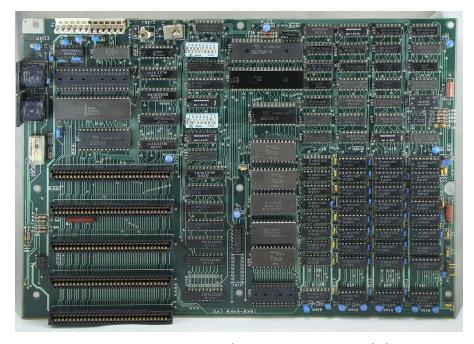


Pit stop





Memory access patterns



An IBM Personal Computer Model 5150 motherboard, 1981 - the first "PC". Image credit: user GermanX on Wikimedia Commons, under license https://doi.org/10.1007/journal.com/





Memory access patterns

Image credit: user Gravislizard on Wikimedia Commons, under license Attribution-ShareAlike 2.0 France

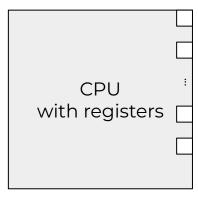


An IBM Personal Computer Model 5150 motherboard, 1981 - the first "PC". Image credit: user GermanX on Wikimedia Commons, under license Attribution-ShareAlike 2.5 Generic





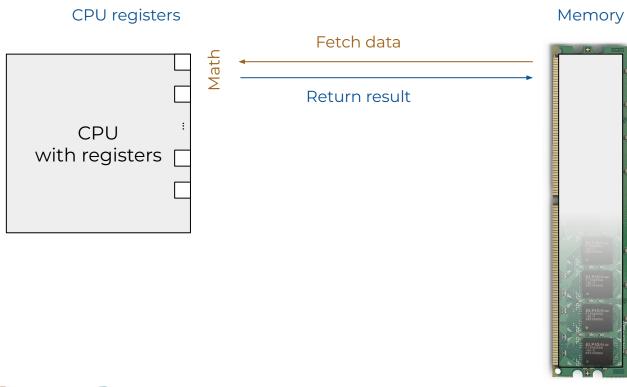
CPU registers



Memory

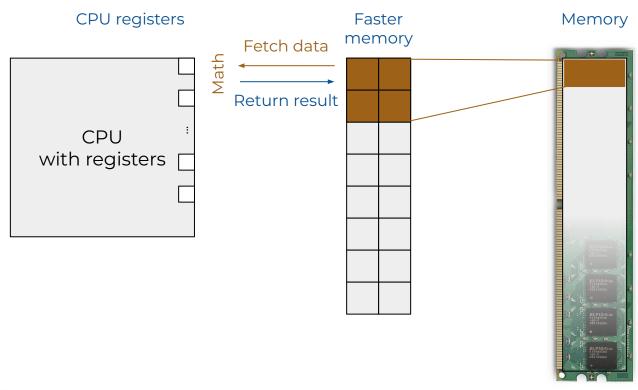






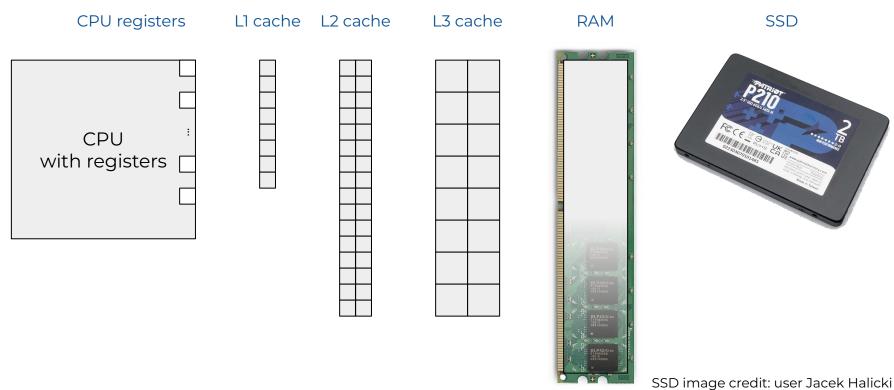
















on Wikipedia, under

Tools and Techniques, Lecture 3 - Performance,

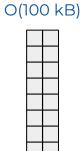
L3 cache

CPU-to-memory layout

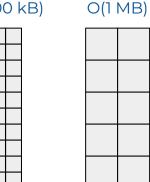
CPU registers 8 bytes

CPU with registers

L1 cache O(10 kB)



L2 cache



RAM O(10 GB)



Typical sizes of different memory hardware

> **SSD** O(1 TB)





SSD image credit: user Jacek Halicki on Wikipedia, under

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CPU-to-memory layout

Typical sizes of different memory hardware

CPU registers L1 cache L2 cache L3 cache 8 bytes O(10 kB) O(100 kB) O(1 MB) CPU with registers 1ns Typical access times 10ns 4ns



RAM





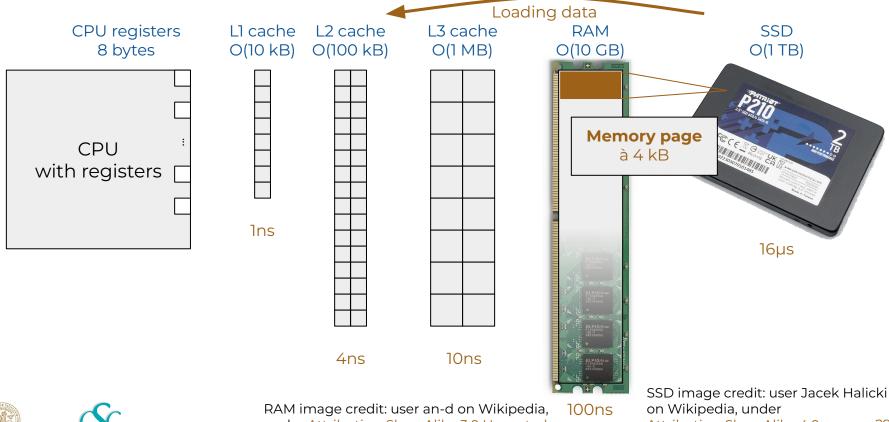


RAM image credit: user an-d on Wikipedia, under <u>Attribution-ShareAlike 3.0 Unported</u>

100ns

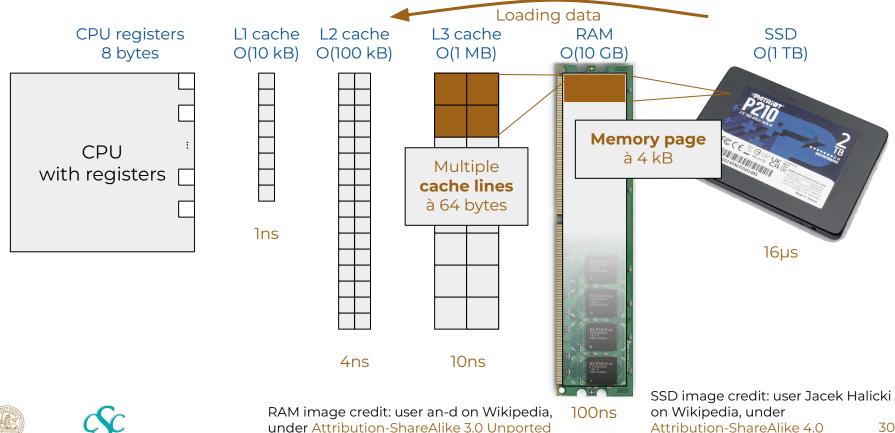
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SSD image credit: user Jacek Halicki



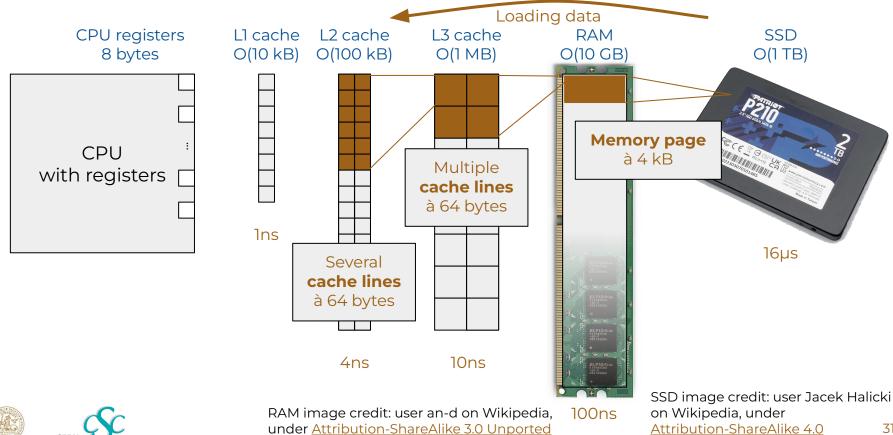






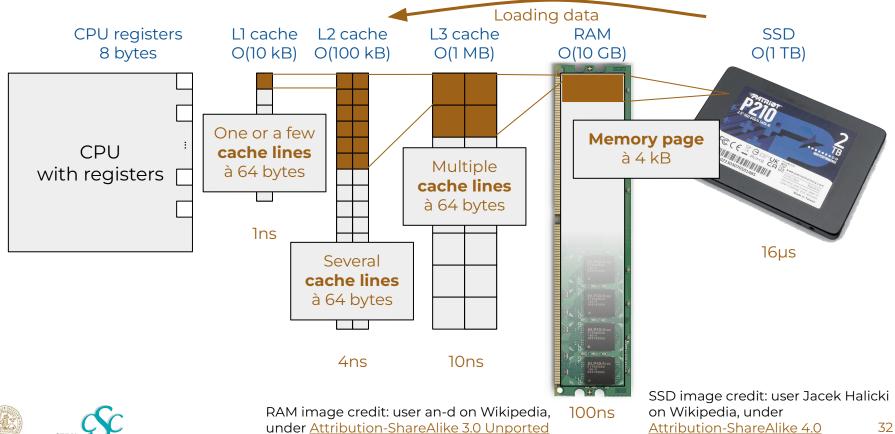


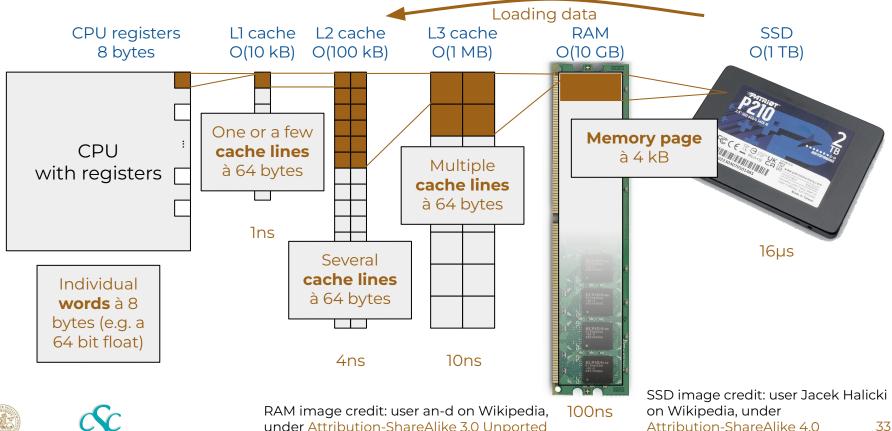




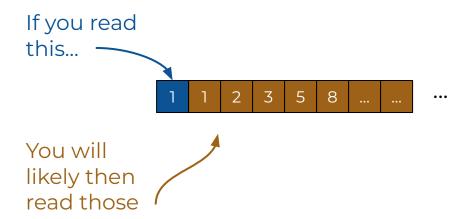








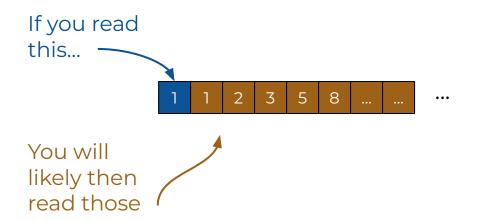
"Principle of locality" or "data locality" - data access often happens on many elements close to each other.







"Principle of locality" or "data locality" - data access often happens on many elements close to each other.



- space locality
- time locality

Pre-fetching: loading data that is likely to be needed soon





Memory access patterns

Amount that fits in L1 cache (i.e. amount that is pre-fetched)

Your data





Amount that fits in L1 cache (i.e. amount that is pre-fetched)

Your data What your code does

"cache hits"

1 ns



• • •

"cache hits"

1 ns

Amount that fits in L1 cache
(i.e. amount that is pre-fetched)

Your data

What your code does

"cache miss"

4 ns





Amount that fits in L1 cache (i.e. amount that is pre-fetched) Your data What your code does "cache hits" "cache miss" 1 ns 4 ns Your data • • • What your etc. code does

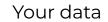




Amount that fits in L1 cache (i.e. amount that is pre-fetched)







What your code does

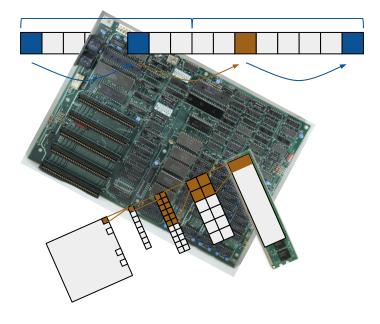


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Memory access patterns - conclusions



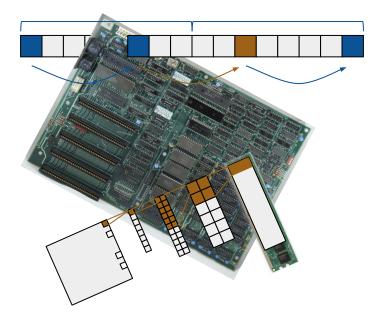


Pit stop 2





Benchmarking





Benchmarking, pitfalls of

...but measuring is hard.

- Compiler optimizations
- Programs running in parallel ⇒ variance
- CPUs boost clock frequencies
- ...

Takeaway: healthy scepticism and reasonable expectations





Linux time command

```
int main() {
        int temp;
        for (int i = 0; i < 1000000000; i++) {
            temp = i;
        }
        return temp;
}</pre>
```

```
$ g++ -03 time_example.cpp -o time_example
$ time ./time_example

real     0m0.009s
user     0m0.002s
sys     0m0.002s
```

- **real:** total time until your program finished
- user: time spent executing your program
- sys: time the system spent on behalf of your program





Note: blunt tool

Linux time command

```
int main() {
    int temp;
    for (int i = 0; i < 10000000000; i++) {
        temp = i;
    }
    return temp;
}</pre>

    godbolt.org Compiler Explorer:
    main:
        mov eax, 999999999
    ret
}
```

```
$ g++ -03 time_example.cpp -o time_example
$ time ./time_example

real    0m0.009s
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```

- real: total time until your program finished
- user: time spent executing your program
- sys: time the system spent on behalf of your program





Note: blunt tool

Compiler pitfalls

```
int main() {
    int size = 1000;
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```

```
int main() {
    int size = std::stoi(string_from_user);
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```



Compiler pitfalls

```
int main() {
    int size = 1000;
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
        index = i % size;
        // ... vector operations
    }
}</pre>
```

```
Array size: 1000
Steps taken: 1000000000
Time per step: 1.06009 ns
```

```
standard string-to-int
                   conversion
int main() {
    int size = std::stoi(string_from_user);
    int index = 0;
    for (int i = 0; i < 1000000000; i++) {
            index = i % size;
           // ... vector operations
Array size: 1000
Steps taken: 1000000000
Time per step: 2.02596 ns
```

⇒ loop 2x slower



Compiler pitfalls

```
int main() {
     int size = 1000;
     int index = 0;
     for (int i = 0; i < 1000000000; i++) {
            index = i % size;
             // ... vector operations
                   movabs rsi, 2361183241434822607
                   .L4:
                   mov rdx, rcx
                   shr rdx
                   mov rax, rdx
                   mul rsi
                   mov rax, rcx
                   shr rdx, 4
                   imul rdx, rdx, 250
                   sub rax, rdx
                   movsx rax, DWORD PTR [rbp+0+rax*4]
                   add ebx, DWORD PTR [rbp+0+rax*4]
```

```
standard string-to-int
                     conversion
int main() {
     int size = std::stoi(string_from_user);
     int index = 0;
     for (int i = 0; i < 10000000000; i++) {
            l index = i % size;
            // ... vector operations
 mov rax, rcx
 cgo
 idiv rbp
 movsx rax, DWORD PTR [rbx+rdx*4]
 add esi, DWORD PTR [rbx+rax*4]
 mov r12d, esi
```

⇒ loop 2x slower

Benchmarking tools

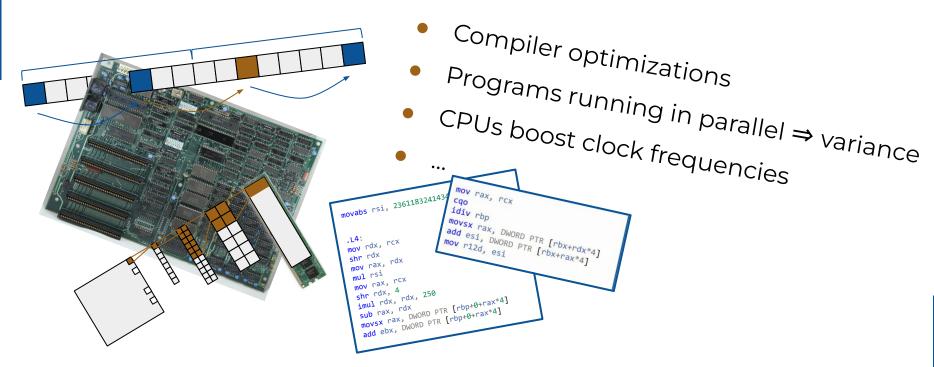
Some more sophisticated benchmarking tools:

- C/C++
 - Catch2
 - Google Benchmark
- Rust
 - o <u>bench</u>
- Python
 - <u>pytest-benchmark</u>
 - o <u>timeit</u>
- General purpose
 - <u>hyperfine</u>
 - <u>perf</u> (Linux)





Benchmarking - conclusions



Healthy scepticism and reasonable expectations





Conclusions - conclusions

Tools and Techniques, Lecture 3 Performance

- Algorithmic complexity
- Memory access patterns
- Benchmarking





Backup





Catch2

```
#include <catch2/catch test macros.hpp>
#include <catch2/benchmark/catch benchmark.hpp>
int arithmetic_sum (int upper) {
  volatile int sum = 0;
   for (int i = 1; i <= upper; i++) {
      sum += i;
   return sum;
TEST CASE ("Benchmarking Basics") {
  BENCHMARK ("Arithmetic sum 1 to 100") {
      return arithmetic sum (100);
```





Catch2

```
#include <catch2/catch test macros.hpp>
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int arithmetic sum (int upper) {
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      sum += i;
   return sum;
TEST CASE ("Benchmarking Basics") {
   BENCHMARK ("Arithmetic sum 1 to 100") {
       return arithmetic sum (100);
```

including Catch2 library

function to benchmark

defining a set of tests or benchmarks

configuring a benchmark





Catch2

```
including Catch2
#include <catch2/catch test macros.hpp>
                                                              library
#include <catch2/benchmark/catch benchmark.hpp>
int arithmetic sum (int upper) {
  volatile int sum = 0;
   for (int i = 1; i <= upper; i++) {
                                                            function to
                                                            benchmark
       sum += i;
   return sum;
                                                       defining a set of tests
                                                          or benchmarks
TEST CASE ("Benchmarking Basics") {
   BENCHMARK ("Arithmetic sum 1 to 100") {
                                                              configuring a
       return arithmetic sum (100);
                                                               benchmark
```





Catch2 output

```
Benchmarking Basics
/eos/user/k/kaastran/schools/CSC2025/prep/exercises/src/ex0.0 benchmarking basics.cpp:15
benchmark name
                                       low mean
                                                      high mean
                              std dev low std dev high std dev
Arithmetic sum 1 to 100
                               38.5609 ns 36.9933 ns 40.8075 ns
                               9.47688 ns 7.24546 ns 12.492 ns
test cases: 1 | 1 passed
assertions: - none -
```











