Kernel code

```
global__
          void histogram(unsigned int* buffer, unsigned int* bin, int size) {
 int stride = blockDim.x * gridDim.x;
 int i = blockIdx.x * blockDim.x + threadIdx.x;
   _shared__ unsigned int pr_bin[NUM_BINS];
 for (int j = i; j < NUM_BINS; j += stride)
   bin[j] = 0;
 for (int u = threadIdx.x; u < NUM_BINS; u += blockDim.x)</pre>
   pr_bin[u] = 0;
   _syncthreads();
 for (int j = i; j < size; j += stride)
   ① atomicAdd(&(pr_bin[buffer[j]]), 1);
   _syncthreads();
 for (int u = threadIdx.x; u < NUM_BINS; u += blockDim.x)
   ② atomicAdd(&bin[u], pr_bin[u]);
   _syncthreads();
}
 _global__ void histogramSaturate(unsigned int* bin) {
 int i = blockIdx.x * blockDim.x + threadIdx.x;
 bin[i] = min(bin[i], 127);
```

Host code

```
//@@ Allocate GPU memory here
 cudaMalloc((void**)&deviceInput, inputLength * sizeof(unsigned int));
 cudaMalloc((void**)&deviceBins, NUM_BINS * sizeof(unsigned int));
  CUDA_CHECK(cudaDeviceSynchronize());
//@@ Copy memory to the GPU here
 cudaMemcpy(deviceInput, hostInput, inputLength * sizeof(unsigned int), cudaMemcpyHostToDevice);
 cudaMemcpy(deviceBins, hostBins, NUM_BINS * sizeof(unsigned int), cudaMemcpyHostToDevice);
  CUDA_CHECK(cudaDeviceSynchronize());
//@@ Launch kernel
 int blockSize = 256;
 int maxGridSize = 8192;
 int gridSize = min(maxGridSize, (inputLength + blockSize - 1) / blockSize);
 histogram<<<gridSize, blockSize>>>(deviceInput, deviceBins, inputLength);
 int satGridSize = NUM_BINS / blockSize;
 histogramSaturate<<<satGridSize, blockSize>>>(deviceBins);
//@@ Copy the GPU memory back to the CPU here
 \verb|cudaMemcpy| (hostBins, deviceBins, NUM\_BINS * size of (\textit{unsigned int}), cudaMemcpyDeviceToHost); \\
  CUDA_CHECK(cudaDeviceSynchronize());
//@@ Free the GPU memory here
 cudaFree(deviceBins);
 cudaFree(deviceInput);
```

- 1) atomicAdd is called exactly inputLength(N) time throughout kernels.
- 2) atomicAdd is called NUM_BINS = 4096 times per thread block.

In this case, B = # of thread block is initialized into min $(B_{\text{max}}, \lceil N/T \rceil)$ in host code,

where T(blockSize) = 256 and Bmax(maxBlockSize) = 8192 is modifiable.

Therefore, Atomic operations are called $N + 4096 \times \min(B_{\text{max}}, [N/T])$ times.

Problem 2

The contention of ② atomicAdd is constant regardless of data.

However, contention of (1) atomicAdd will be maximized when every element has the same value.

For each thread block, addition to same shared memory is completely serialized.

If $N \leq B_{max} \times T$, the loop of ① called only once, so T of them are serialized.

Otherwise, $B=B_{max}$, then maximum $T\times \left\lceil N/_{B_{max}T} \right\rceil$ of 1 atomicAdd are serialized.

Problem 3

The more random the data is, especially in 256 chunks, the more different the input values are, the less content is. In this case, (1) atomicAdd is closer to full parallel.

Problem 4

Refer to the previous page.

Problem 5

Refer to the below table. Each time is millisecond unit.

N	16	1024	513	511	1	500000	800000	1200000
GPU alloc	0.136262	0.141	0.145751	0.147649	0.14361	0.161956	0.394259	0.332044
H->D	0.044791	0.045716	0.043647	0.043303	0.041688	0.313888	0.449736	0.57497
Kernel	0.024034	0.025683	0.023922	0.025047	0.02356	0.076281	0.11287	0.148933
D->H	0.02734	0.025967	0.029283	0.024645	0.022598	0.019647	0.025409	0.028175
GPU free	0.122255	0.119247	0.118937	0.121361	0.119485	0.14266	0.265325	0.249115