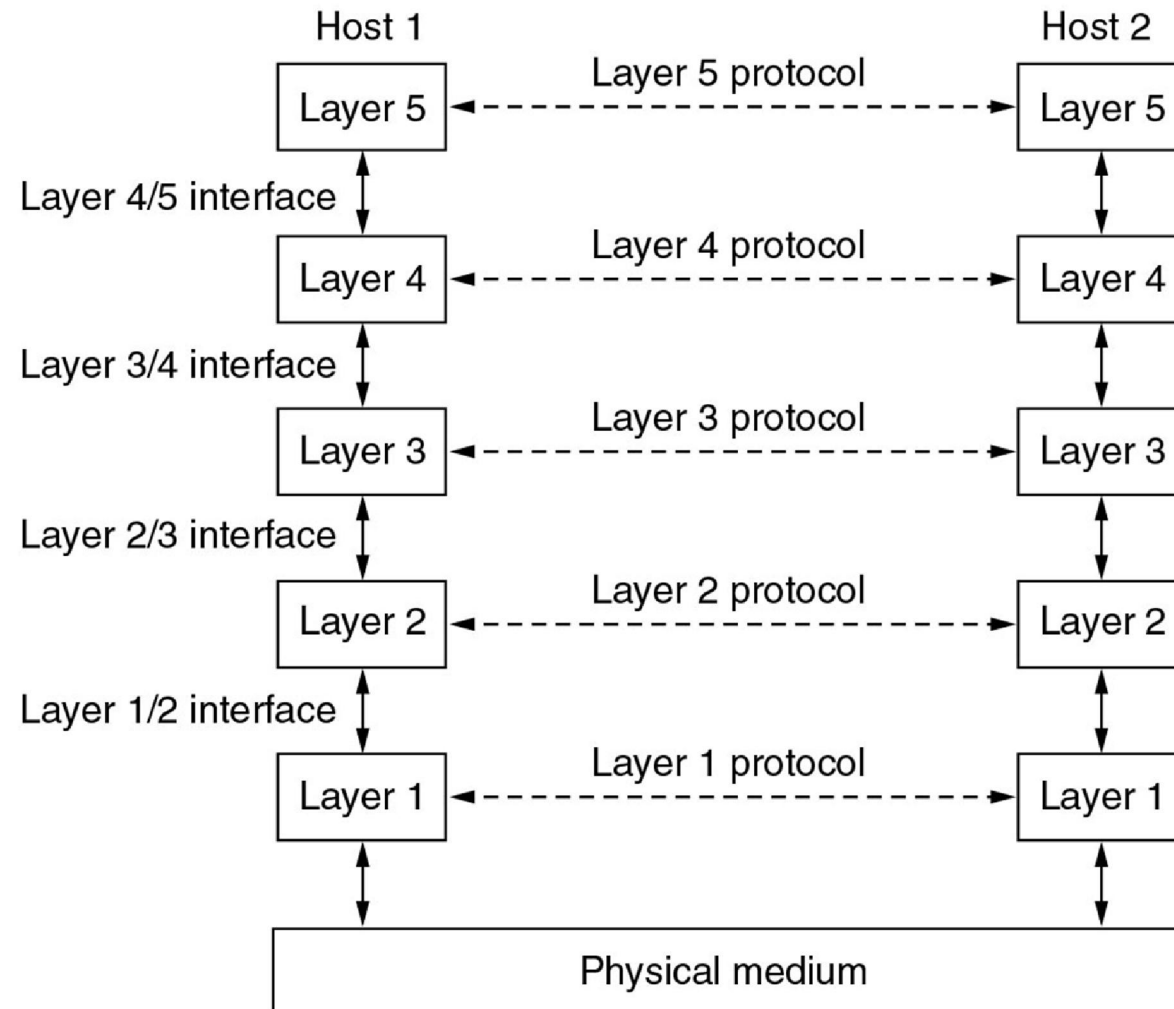


Protocol Hierarchies/Network Software

- Computer networks are generally comprised of numerous pieces of hardware and software
- To simplify network design most networks are organized as a stack of layers of hardware **and/or** software
- The purpose of each layer is to offer *services* to higher layers
- This concept is common in programming viz. objects or libraries which perform specific operations
 - Importantly the object/library function keeps details of its internal state and algorithms hidden from the main program
- The example five-layer network illustrates the layers

Protocol Hierarchies/Network Software



Protocol Hierarchies/Network Software

- *Peer* entities exist at corresponding layers on the source and destination stations
 - These may be processes, hardware devices, or even human beings
- These *peer entities* communicate with each other across a layer
 - The communications rules and conventions used between *peer entities* are collectively known as the *Layer N protocol*
 - A protocol is an *agreement* between communicating parties on how communication should take place

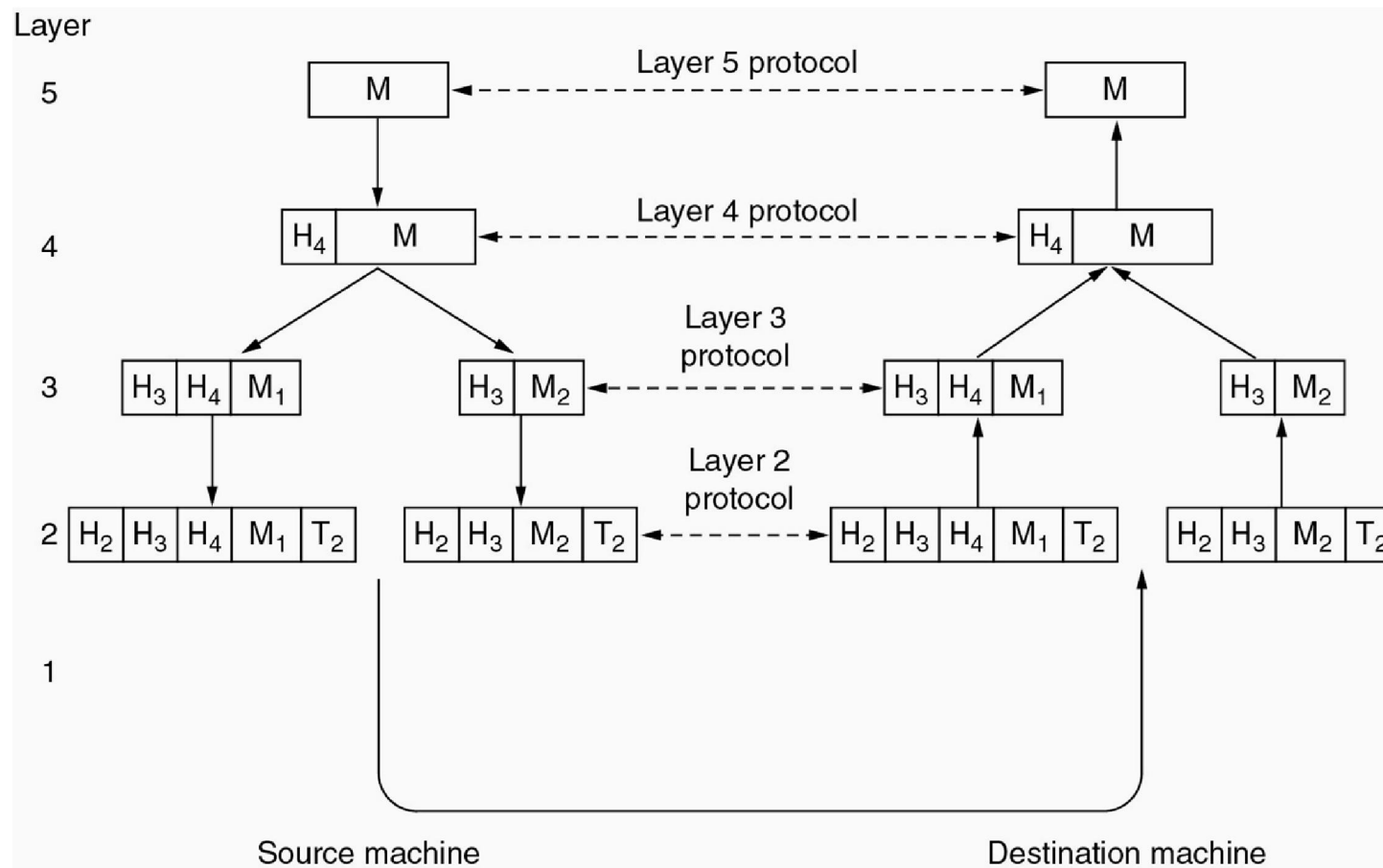
Protocol Hierarchies/Network Software

- In reality, no data is passed directly between peer entities
 - Instead data passes upwards towards the *network applications* or downwards towards the *physical medium*
- *Virtual* communications between peer entities are shown as dotted lines
- *Physical* communications are shown with solid lines

Protocol Hierarchies/Network Software

- Between each pair of adjacent layers is an *interface*
 - A layer's interface reveals the specific functions performed within the layer
- Well defined *interfaces* are an essential design feature of protocol software because:
 - They minimize the amount of information passed between layers
 - They make it easier to replace *layer entities* without affecting the ability of the hosts to communicate with each other

Information flow supporting virtual communication in layer 5



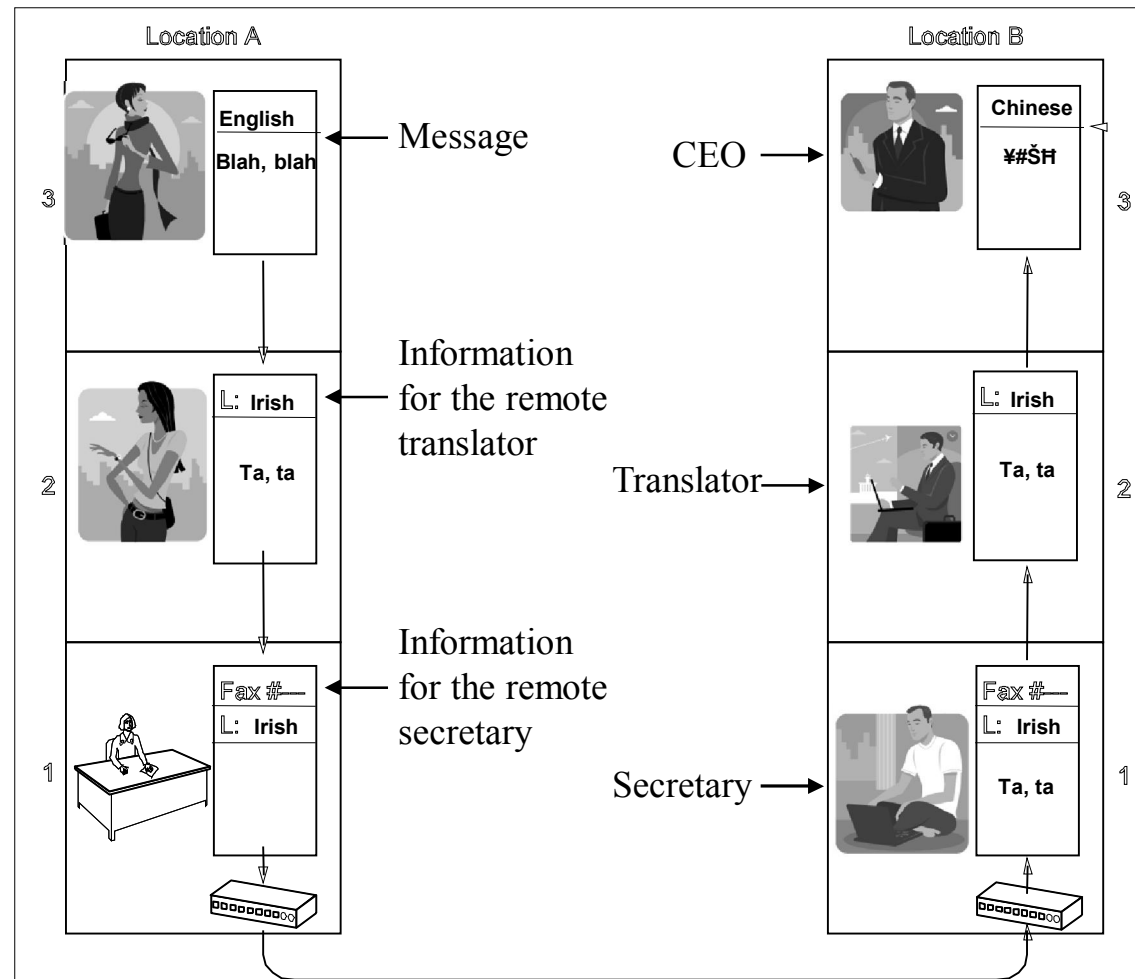
The *Layering Principle*

- When designing layered protocol software it is important to adhere to the *layering principle*:
 - “*Layer N software on the Destination machine must receive the exact message sent by Layer N software on the Source machine*”
- Any transformations or additions made by layers below Layer N on the sending side must be reversed/removed before a message is passed to Layer N software on the receiving side.

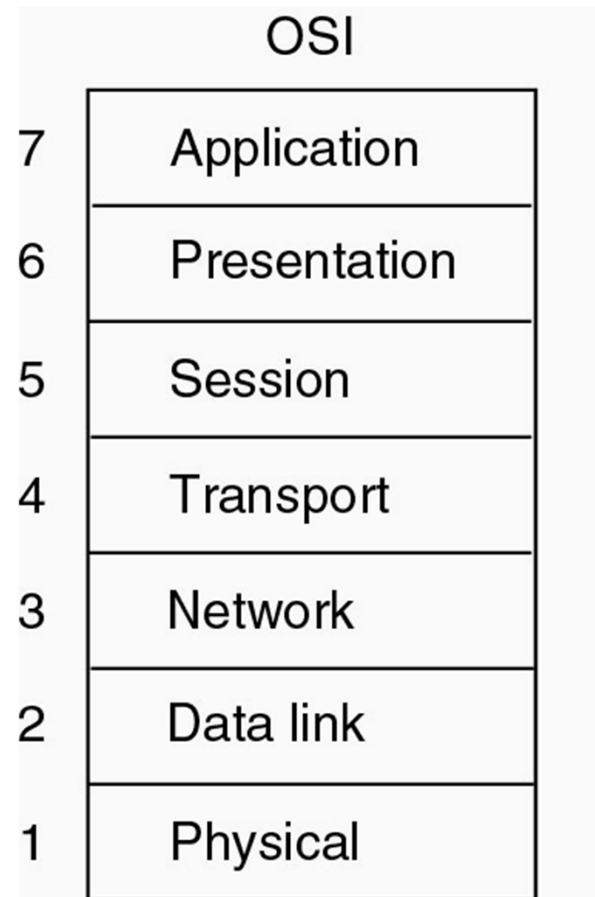
Network Architecture

- Some terminology:
 - A set or stack of layers and protocols is called a *protocol architecture*
 - The architecture only specifies the functions associated with each layer. It does not specify the *implementation* or the *interfaces* of the layer
 - The set of protocols (one per layer) used within an actual network system is called a *protocol stack*
- An analogy may help explain the idea of multilayer communication

The CEO-Translator-Secretary Scenario



Reference Models - The *ISO OSI Model*



Reference Models - The OSI ISO Model

- The “International Standards Organisation Open Systems Interconnect” model:
 - AKA: “*The OSI model*”, “ *The 7-layer model*”,
 - It deals with connecting open systems, Systems that are open for communication with other systems
- The OSI Reference model is a Protocol Architecture Reference Model :
- It is not a Protocol Stack:
 - It does not specify how functionality is to be provided only what the functionality of each layer should be,
 - It is not something one can purchase or install on a networked machine.

Reference Models - The OSI ISO Model

- The principles that were applied to arrive at the seven layers are as follows:
 - Each layer was created when a different level of abstraction was required
 - Each layer performs a well-defined function with each function chosen carefully to facilitate internationally standardized protocols
 - The layer interfaces (boundaries) were carefully defined to *minimize* information flow across the interfaces
 - The number of layers chosen was sufficient enough to ensure that distinct functions were not lumped together without becoming unwieldy

The *ISO Reference Model* - Layer-by-layer

- The *Application* Layer:
 - Contains a variety of protocols commonly used on the Internet
 - E.g. HTTP (HyperText Transfer Protocol) is the underlying protocol for the World Wide Web
 - Other protocols include FTP, E-mail etc.

The *ISO Reference Model* - Layer-by-layer

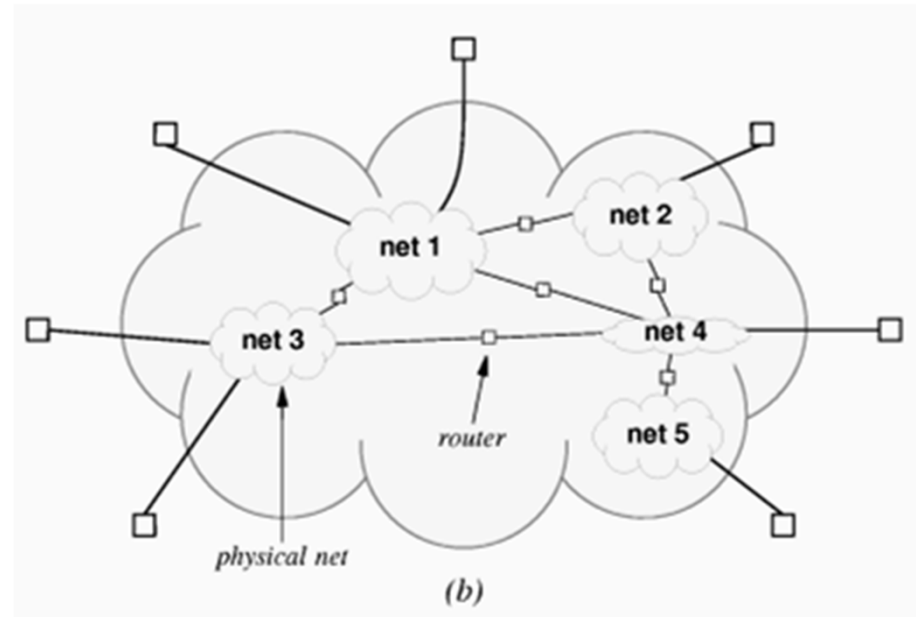
- The *Presentation* Layer: Concerned with the *syntax* and *semantics* of the information transmitted
 - Facilitates communication between *big-endian* computers e.g. Sun Sparcs and *little-endian* computers e.g. Windows machines
- The *Session* Layer: Facilitates the use of *sessions* between end stations. During a session the user and the computer system engage in a *dialogue*
 - The session layer establishes and maintains dialogues
 - It also determines:
 - The type of control to be used i.e. two-way simultaneous communication, two-way alternate comm. or one-way comm.
 - *Re-synchronization* of the dialogue after a crash

The *ISO Reference Model* - Layer-by-layer

- The *Transport* Layer: THIS IS KEY LAYER from our perspective.
- It is a true end-to-end layer in that it operates between end-hosts.
- Its function is to isolate the applications from the underlying network hardware technology:
 - It uses the network as a reliable ‘deliverer’ of data
- It splits the **source** data into manageable chunks and passes them to the network layer for onward delivery.

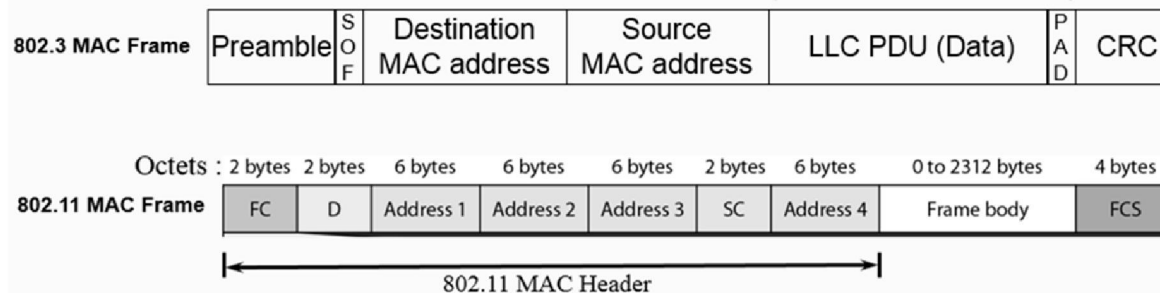
The *ISO Reference Model* - Layer-by-layer

- The *Network Layer*:
- Concerned with controlling the operation of the *sub-network* (subnet)
 - Deals with the routing of *packets* from the *source* station towards the *destination* station across ***sub-networks***
 - It handles the different sub-net addressing formats
 - Essentially this layer is responsible for interconnecting *heterogeneous* networks

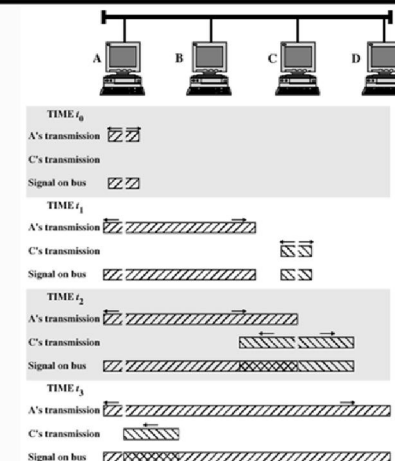


The *ISO Reference Model* - Layer-by-layer

- The *Data Link* Layer: Concerned with getting data across an individual link
 - Essentially it transforms a raw transmission facility into a *data communications channel* that *appears* free of transmission errors
 - Breaks up the data into *data frames*. Also deals with flow control, controlling access to a shared channel etc.

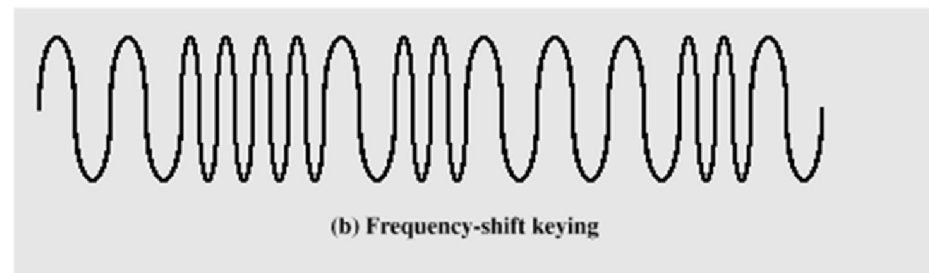
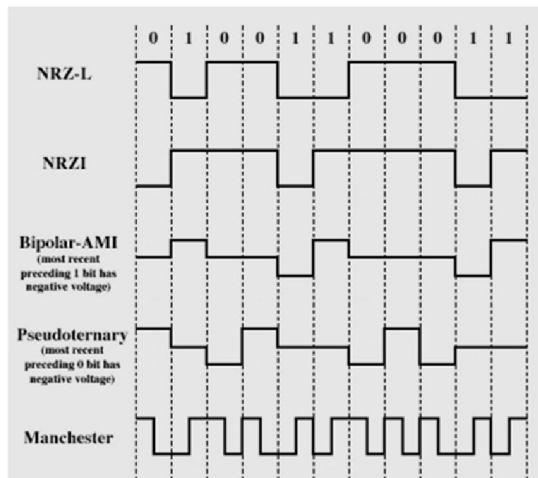


CSMA Operation

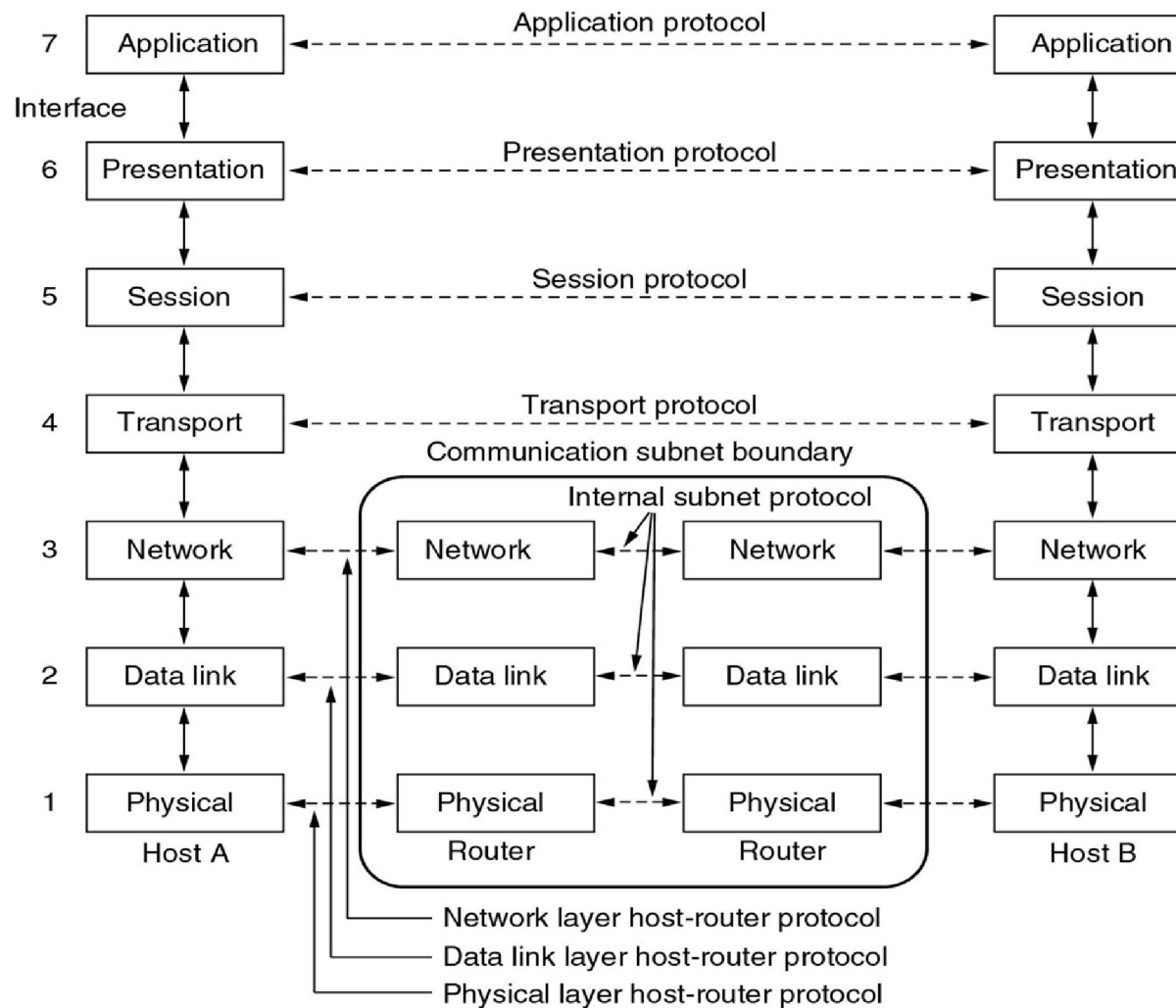


The *ISO Reference Model* - Layer-by-layer

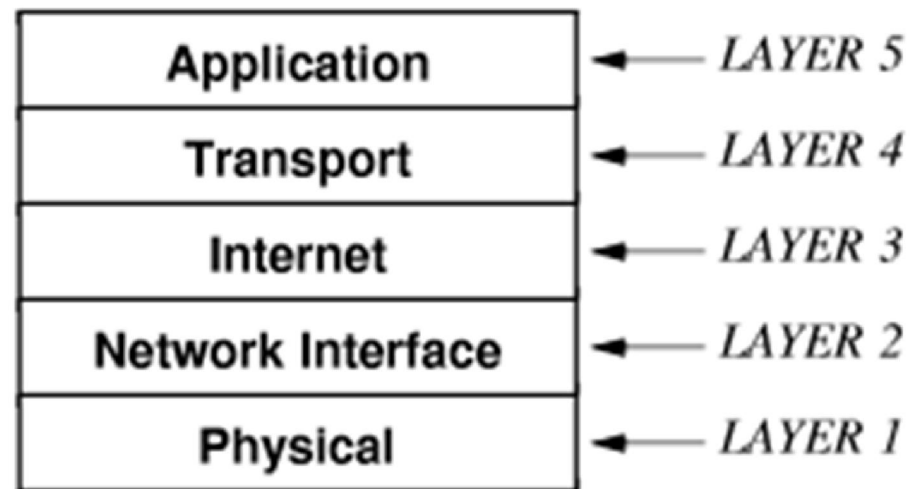
- The *Physical* Layer: Concerned with transmitting raw bits over a *communication channel*. Must ensure that when a binary 1 is sent it is received as such by the *receiver*
 - Deals with voltage levels used, bit duration etc.
 - Design issues deal with *mechanical*, *electrical*, and *timing interfaces*, and the physical *transmission medium*



The *ISO OSI Model* – Layers 1-3 Versus Layers 4-7



Reference Models – The *TCP/IP Reference Model*



The *TCP/IP Reference Model*

- Transport Control Protocol / Internetwork Protocol Reference Model:
 - AKA: “The TCP/IP Model”
- The protocols upon which this model is based (*TCP and IP*) fuelled the early growth of the *Internet*:
 - TCP and IP were adapted because they were available,
 - The ISO protocols on the other hand were still being developed
- The *TCP/IP Reference Model* was developed after the protocols
 - This is similar to retrospectively drawing plans for a house after it is built.

The *TCP/IP Reference Model*

- Specific design goals of this *Reference Model* included:
 - Ability to survive the loss of subnet hardware. Specifically in a “Theatre of War” (a Battlefield),
 - Ability to handle multiple types of data including files and real-time speech.
- These requirements led to the adoption of a connectionless packet-switching network within the *internet* layer

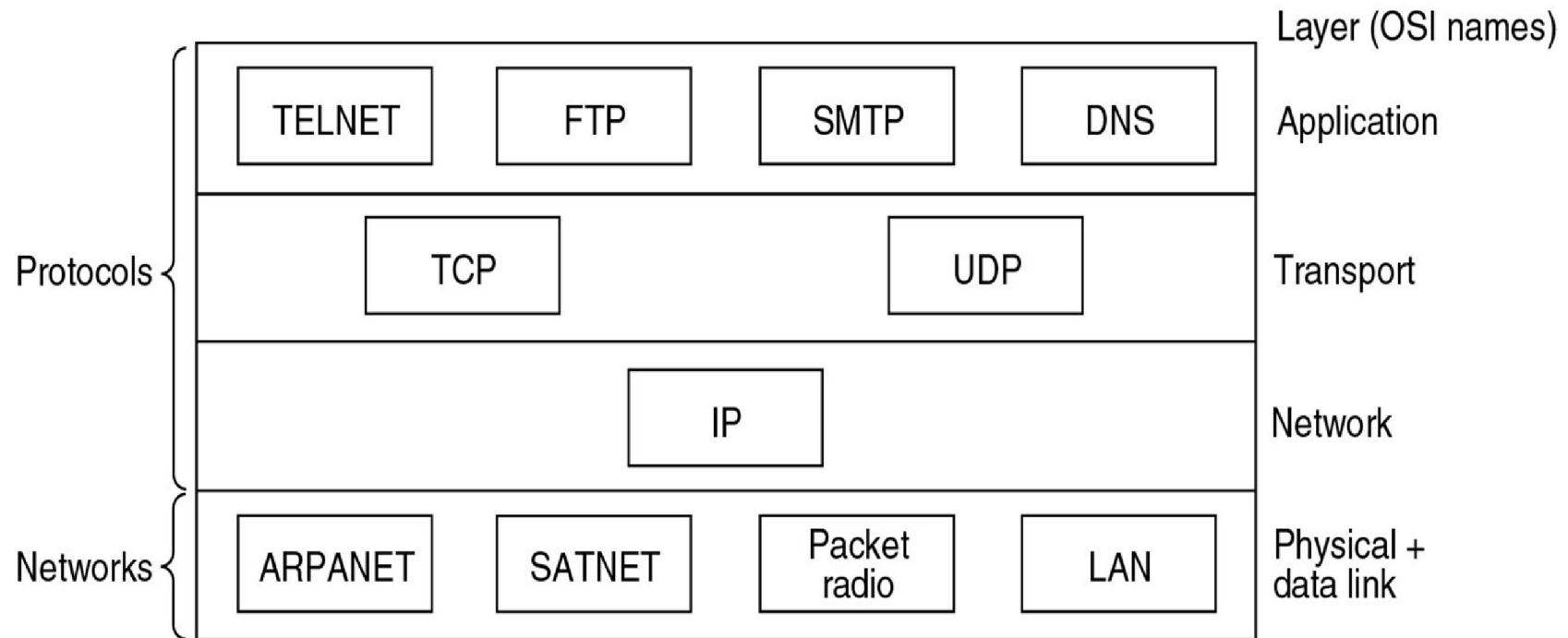
The *TCP/IP Reference Model* - Layer-by-layer

- The *Application* Layer: This layer contains all of the higher-level protocols including *FTP*, *E-mail*, the *Domain Name System (DNS)* and *HTTP*.
- The *Transport* Layer: Facilitates *end-to-end* communication between the *Source* and *Destination* hosts
- Two end-to-end transport protocols have been defined:
 - *TCP (Transmission Control Protocol)*: This is a reliable, connection-oriented protocol that allows a byte stream originating on one machine to be delivered without error to any other machine in the internet.
 - *UDP (User Datagram Protocol)*: This is an unreliable, connectionless protocol for applications that provide their own *sequencing* and *flow control* functionality

The *TCP/IP Reference Model* - Layer-by-layer

- The *Internet* Layer: This layer is key to the whole architecture
 - It facilitates hosts injecting *packets* into any network,
 - It ensures correct *routing* of packets to the *Destination* station
- The *Host-to-Network* Layer: This layer is meant to deal with hosts connecting to the network in order to transmit packets
 - It is not well defined within the TCP/IP reference model

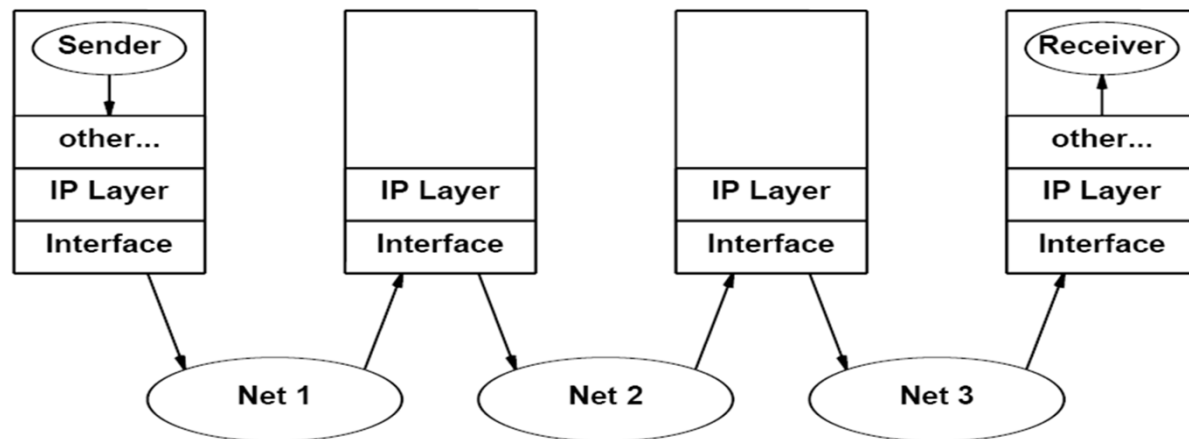
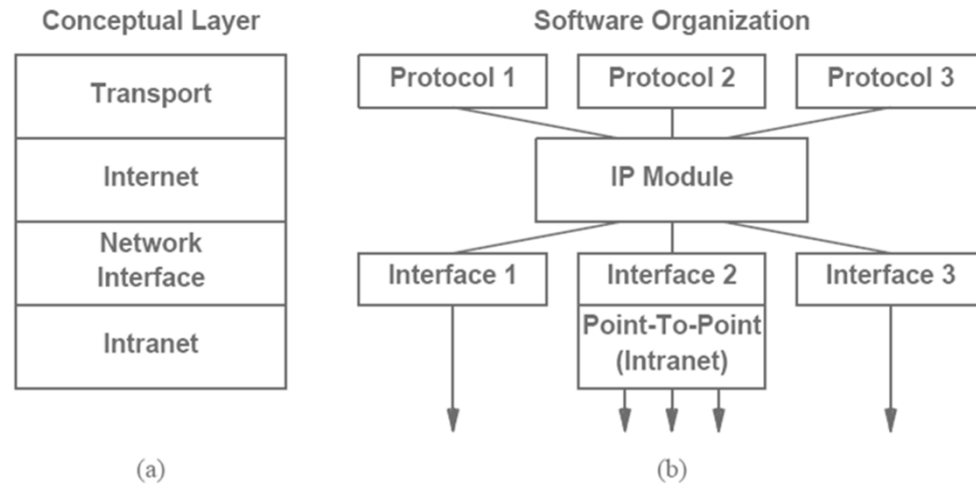
The *TCP/IP Reference Model* – Relationship between *TCP*, *UDP* and *IP*



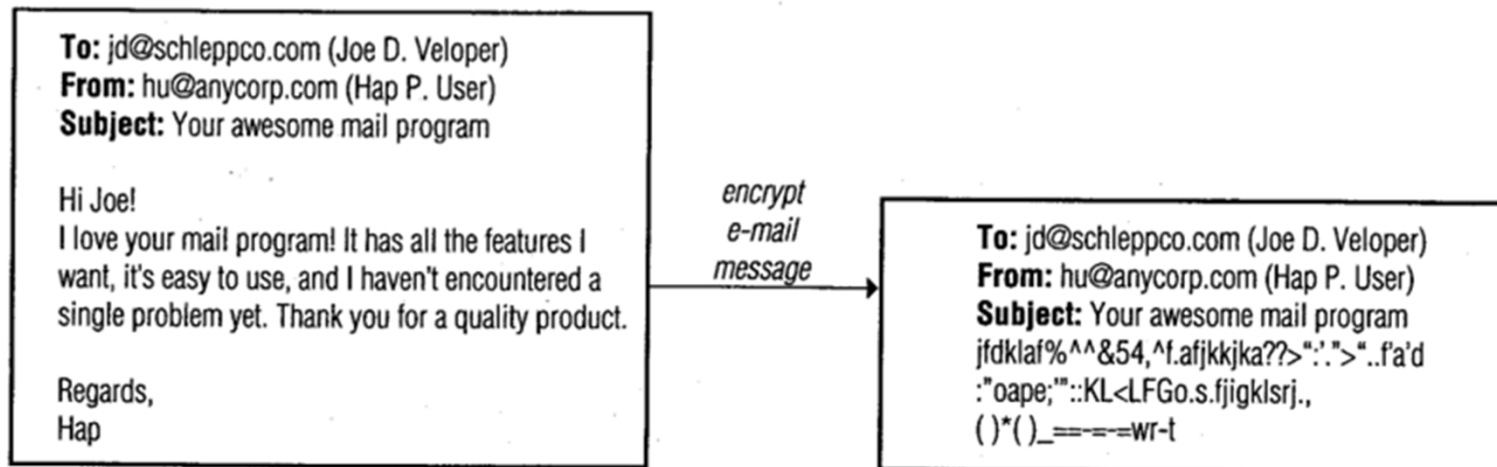
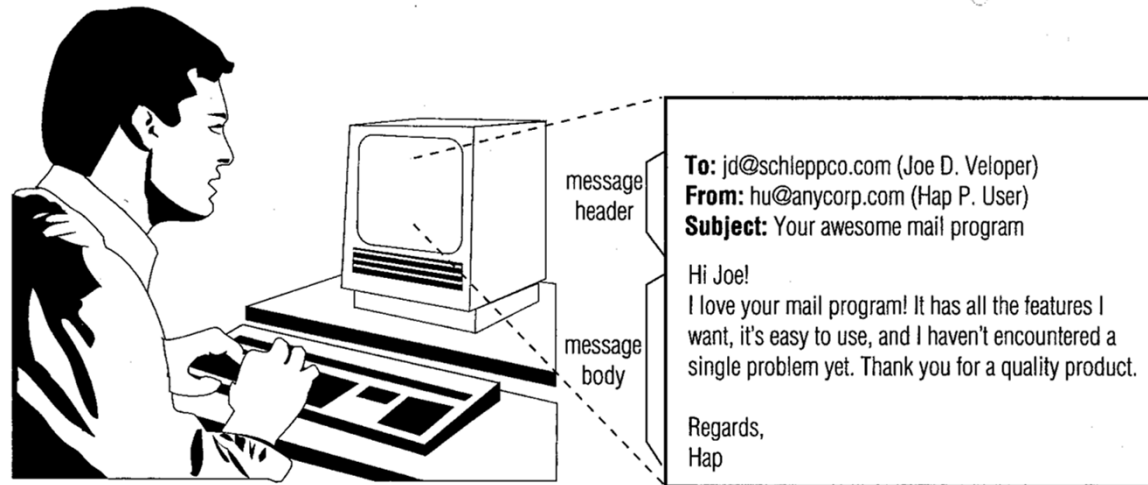
A Comparison of OSI and TCP/IP Reference Models

- Both models are similar in many ways viz.:
 - Both use the concept of a stack of independent protocols
 - Both transport layers provide an end-to-end, network-independent transport service to applications
- However, there are some notable differences as follows:
 - *Number of layers*: OSI model has **7** layers, TCP/IP has **5** layers
 - *Services versus Interfaces/Protocols*:
 - OSI clearly defines what each layer does using *service definitions*
 - TCP/IP did not originally clearly distinguish between *service*, *interface* and *protocol*. This hindered switching-out protocols to facilitate technological change
 - *Timing*:
 - OSI model was developed before the protocols were invented
 - With TCP/IP the protocols came first and then the model

Conceptual versus Realistic view of protocol layering



The *Layers* in operation – Application and Presentation Layers



The *Layers* in operation – The Session Layer

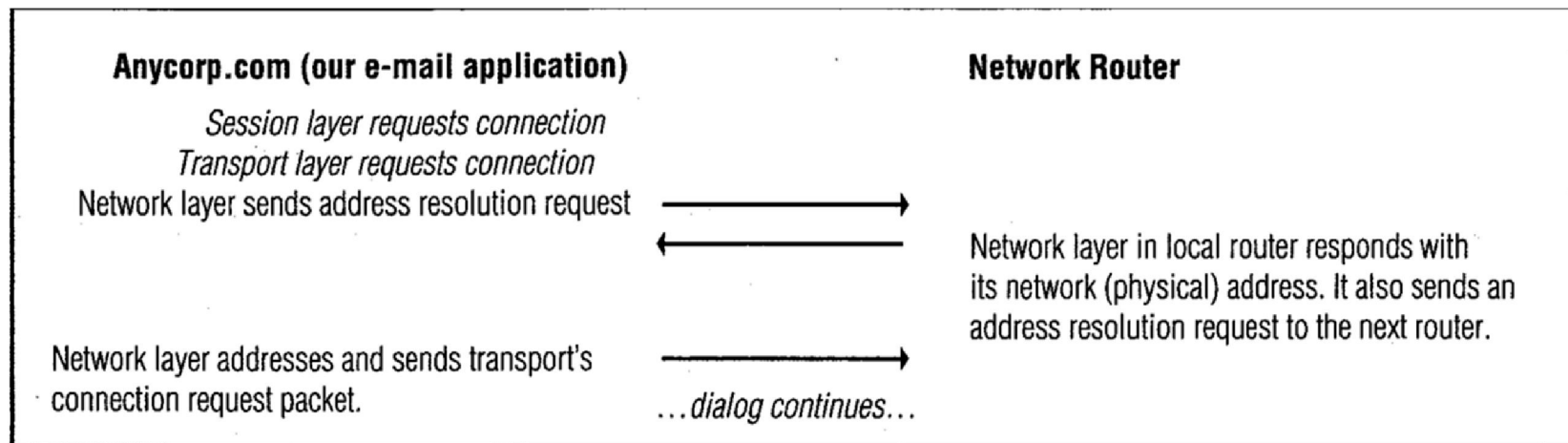
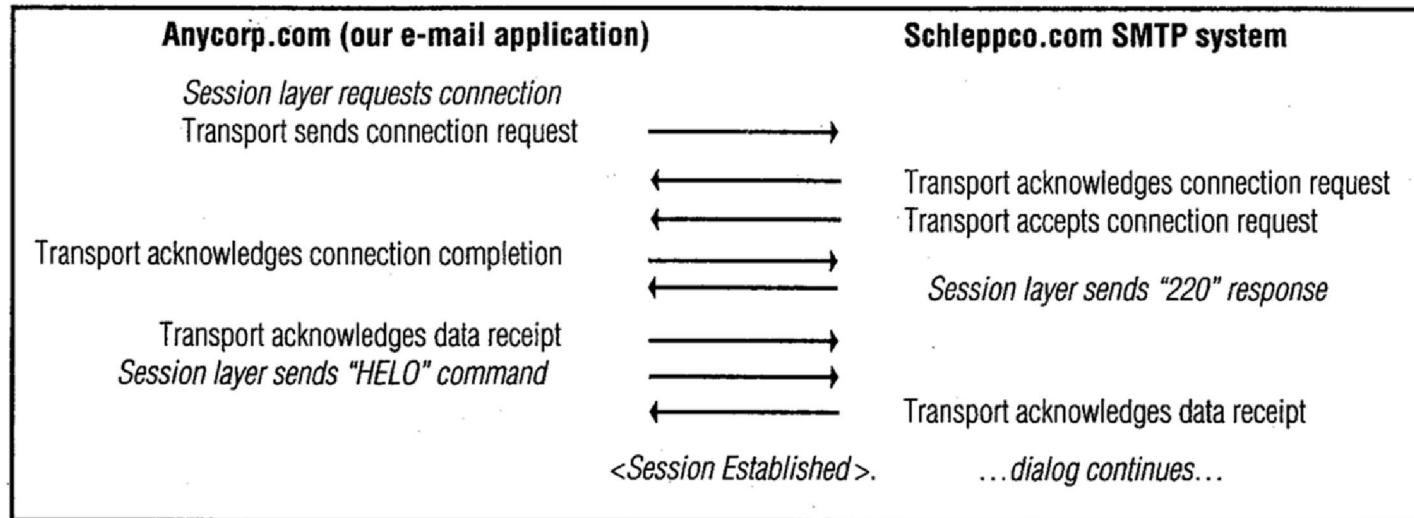
```
Server: 220 Schleppco.COM Simple Mail Transfer Protocol (SMTP) Server ready to serve you!
App:   HELO AnyCorp.COM
Server: 250 Ok to proceed!
       <Session Established>
App:   MAIL FROM:<hu@anycorp.com>
Server: 250 Ok to proceed!

App:   RCPT TO:<jdv@schleppco.com>
Server: 250 Ok to proceed!

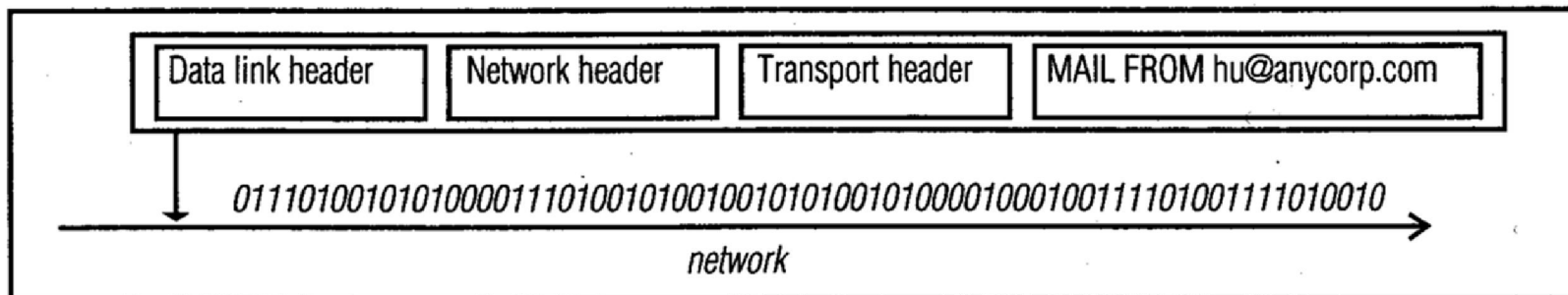
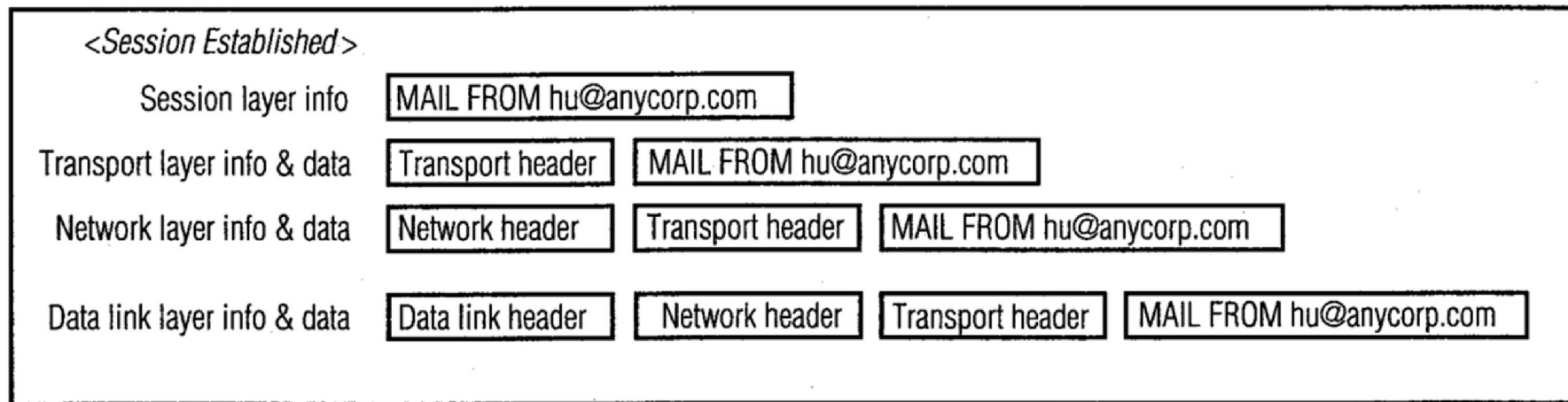
App:   DATA
Server: 354 Send mail message and end it with '.' alone on a line
App:   Date: 24 June 94 14:20:00 EST
App:   jfdklaf%^&54,^f.afjkkjka??>:"..f'a' d
       (our encrypted & compressed mail message)
App:   : "oape;"::KL<LFG0.s.fjigklsrj.,
App:   ()*( )_ = ---wr-t
App:   .
Server: 250 Ok to proceed!

App:   QUIT
Server: 221 Schleppco.COM over and out  Thanks for the visit!
       <Session Ended>
```

The *Layers* in operation – Transport and Network Layers



The *Layers* in operation – Data Link and Physical Layers



Design Issues for Layered Software

- There are a number of key design issues common to a several layers:
 - *Addressing* – Each layer needs to be able to identify *senders* and *receivers*. Some form of *addressing* is required
 - *Error control* - The *receiver* must be able to tell the *sender* which messages have been correctly received and which have not
 - *Sequencing* - The protocol software on the the receiver must be able to resequence incoming messages
 - *Flow Control* – The receiver must be able to control the flow of information from the sender
- The following are two examples of network *models* namely the **OSI** and **TCP/IP** reference models
 - These form the basis for many of today's *network architectures*

Operation of TCP/IP - Sender

1. Preparing the data. The application protocol prepares a block of data for transmission. For example, an email message (SMTP), a file (FTP), or a block of user input (TELNET).

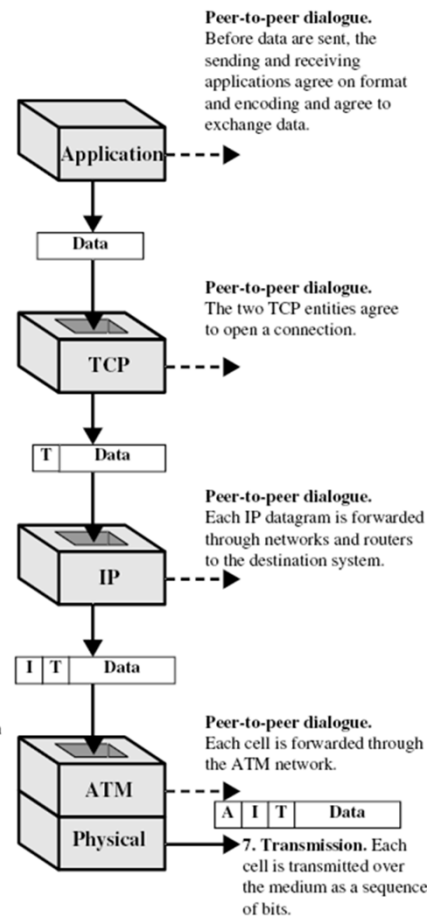
2. Using a common syntax. If necessary, the data are converted to a form expected by the destination. This may include a different character code, the use of encryption, and/or compression.

3. Segmenting the data. TCP may break the data block into a number of segments, keeping track of their sequence. Each TCP segment includes a header containing a sequence number and a frame check sequence to detect errors.

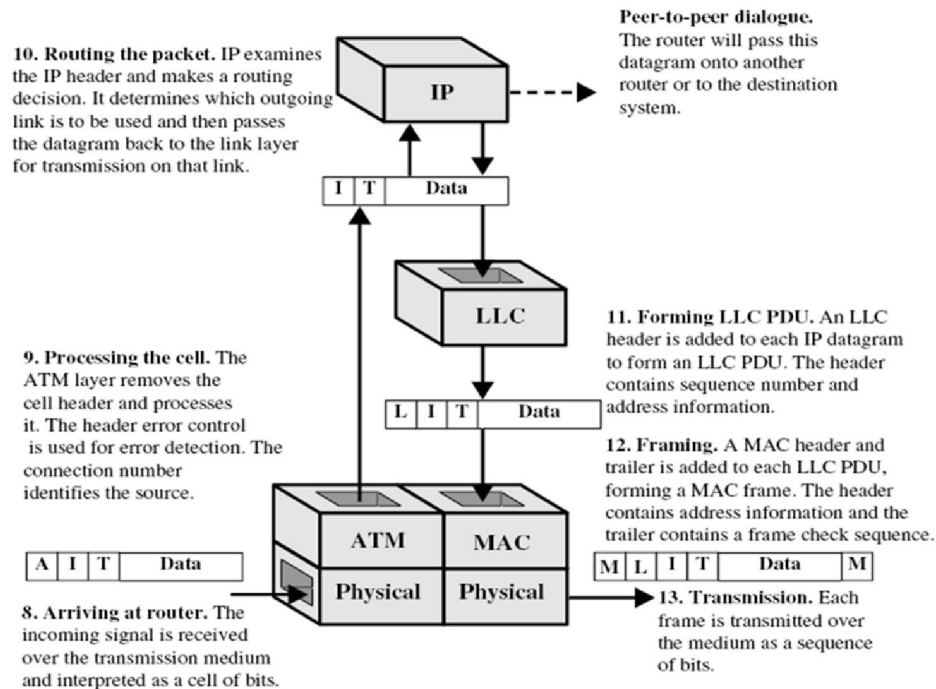
4. Duplicating segments. A copy is made of each TCP segment, in case the loss or damage of a segment necessitates retransmission. When an acknowledgment is received from the other TCP entity, a segment is erased.

5. Fragmenting the segments. IP may break a TCP segment into a number of datagrams to meet size requirements of the intervening networks. Each datagram includes a header containing a destination address, a frame check sequence, and other control information.

6. Framing. An ATM header is added to each IP datagram to form an ATM cell. The header contains a connection identifier and a header error control field



Operation of TCP/IP – Router and Receiver



20. Delivering the data. The application performs any needed transformations, including decompression and decryption, and directs the data to the appropriate file or other destination.

19. Reassembling user data. If TCP has broken the user data into multiple segments, these are reassembled and the block is passed up to the application.

18. Processing the TCP segment. TCP removes the header. It checks the frame check sequence and acknowledges if there is a match and discards for mismatch. Flow control is also performed.

17. Processing the IP datagram. IP removes the header. The frame check sequence and other control information are processed.

16. Processing the LLC PDU. The LLC layer removes the header and processes it. The sequence number is used for flow and error control.

15. Processing the frame. The MAC layer removes the header and trailer and processes them. The frame check sequence is used for error detection.

14. Arriving at destination. The incoming signal is received over the transmission medium and interpreted as a frame of bits.

