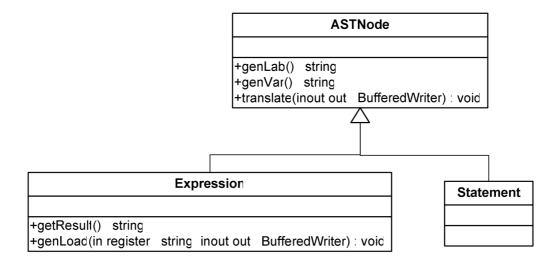
Uputstvo za sve grupe

Kreirati CUP specifikaciju koja za prosleđeni ulazni kod kreira apstraktno sintaksno stablo (AST) na osnovu koga se generiše međukod niskog nivoa. Kao osnovu je moguće koristiti CUP specifikaciju kreiranu u 3. laboratorijskoj vežbi (verziju bez korišćenja *error* simbola)

Za implementaciju AST-a koristiti paket AST koji je dat na sajtu predmeta. Paket AST dopuniti dodatnom klasom (ili klasama) potrebnom za konkretnu gramatiku. Za klasu koju sami dodajete u paket AST obavezno najpre definisati u dodatnom tekstualnom fajlu izgled generisanog međukoda niskog nivoa koji treba da generiše funkcija translate. Taj fajl predati zajedno sa projektom.

Predstavljanje čvorova u apstraktnom sintaksnom stablu



Instrukcije međukoda niskog nivoa

Instrukcija			Značenje
Load_Const	Rn,	С	(Rn) = c
Load_Mem	Rn,	x	(Rn) = x
Load_Arr	Rn,	x[Rm]	(Rn) = x[(Rm)]
Store	Rn,	x	(x) = (Rn)
Store_Arr	Rn,	x[Rm]	x[(Rm)] = (Rn)
Add	Rn,	Rm	(Rn) = (Rn) + (Rm)
Sub	Rn,	Rm	(Rn) = (Rn) - (Rm)
Mul	Rn,	Rm	(Rn) = (Rn) * (Rm)
Div	Rn,	Rm	(Rn) = (Rn) / (Rm)
Compare_Equal	Rn,	Rm	(Rn) = (Rn) == (Rm)
Compare_Greater	Rn,	Rm	(Rn) = (Rn) > (Rm)
Compare_Less	Rn,	Rm	(Rn) = (Rn) < (Rm)
And	Rn,	Rm	(Rn) = (Rn) ∧ (Rm)
Or	Rn,	Rm	$(Rn) = (Rn) \vee (Rm)$
Jump	lab		skok na naredbu sa oznakom lab
JumpIfZero	Rn,	lab	skok na naredbu sa oznakom lab ukoliko je (Rn) == 0
JumpIfNotZero	Rn,	lab	skok na naredbu sa oznakom lab ukoliko je (Rn) ≠ 0

```
Program → main Block exit

Block → Declarations Expressions

Declarations → Declarations Declaration | Declaration

Declaration → Type ID;

Type → int | float | bool

Expressions → Expressions; Expression | Expression

Expression → Assignment | ApplyExpression

Assignment → ID := ArithmeticExpression

ArithmeticExpression → ArithmeticExpression + TermExpression

| ArithmeticExpression - TermExpression

| TermExpression → CONST | ID

ApplyExpression → for ID in [ NameList ] apply Expression

NameList → NameList , ID | ID
```

Napomena: ApplyExpression se izvršava za sve promenljive u NameList-i. ID u for delu predstavlja privremenu promenljivu u koju se smešta vrednost tekuće promenljive iz NameList-e koja se obrađuje.

```
Program → main () Block
Block \rightarrow \{ Variables Statements \}
Variables → Variables Variable | Variable
Variable \rightarrow Type ID;
Type \rightarrow int \mid real \mid boolean
Statements → Statements; Statement | Statement
Statement → Assignment | IfStatement
If Statement \rightarrow if (RelExpression): Block ElsePart
\textit{ElsePart} 
ightarrow \textit{ElifList ElseStatement} \mid \textit{ElifList} \mid \textit{ElseStatement} \mid \epsilon
ElifList → ElifList Elif | Elif
Elif \rightarrow elif (RelExpression): Block
ElseStatement \rightarrow else: Block
RelExpression → Term RelOp Term | Term
Term \rightarrow ID \mid CONST
RelOp \rightarrow < | <= | == | <> | > | >=
Assignment \rightarrow ID := Term
```

Napomena: Prvi Block se izvršava ukoliko je uslov u **if** delu ispunjen. Ukoliko uslov iz **if**-a nije ispunjen a postoji jedna ili više Elif naredba izvršava se Block iz Elif naredbe čiji je uslov ispunjen. Ukoliko nijedan od prethodnih uslova nije ispunjen izvršava se Block iz Else naredbe ukoliko postoji.

```
Program → program Block return

Block → begin VarList Statments end

VarList → VarList Var | Var

Var → ID: Type;

Type → integer | char | string | file

Statements → Statements; Statement | Statement

Statement → Assignment | ReadExpression

Assignment → ID = Expression

Expression → ArithmeticExpression | open (PrimaryExpression)

ArithmeticExpression → ArithmeticExpression + PrimaryExpression

| ArithmeticExpression - PrimaryExpression

| PrimaryExpression

PrimaryExpression → ID | CONST

ReadExpression → read (ID in ID) do Block
```

Napomena: U *ReadExpression*-u drugi **ID** predstavlja identifikator fajla. Pretpostaviti da se izvršenjem **open** naredbe sadržaj fajla učita u memoriju (nije potrebno generisati međukod za naredbu **open**) tako da se podaci iz fajla prenose u niz čiji je identifikator taj **ID**. Na poziciji 0 u tom nizu je broj preostalih podataka u nizu, a stvarni podaci iz fajla kreću od pozicije 1. U *ReadExpression*-u treba korišćenjem naredbe **Load_Arr** učitavati elemente niza počevši od pozicije 1.

Program → **program** *Block* **end**

 $Block \rightarrow \{ Declarations Statements \}$

Declarations → *Declaration* | *Declaration*

 $Declaration \rightarrow Variable Declaration \mid Function Declaration$

 $Variable Declaration \rightarrow ID : Type;$

 $Type \rightarrow int \mid float \mid char$

FunctionDeclaration \rightarrow **ID** (Parameters) => Expression;

Parameters → Parameters , Parameter | Parameter

 $Parameter \rightarrow ID : Type \mid ID : Type = CONST$

Statements → Statements; Assignment | Assignment

Assignment \rightarrow **ID** = Expression

Expression → **CONST** | **ID** | *FunctionCall*

FunctionCall → ID (ArgumentsList)

ArgumentsList → ArgumentsList , Expression | Expression

Napomena: FunctionDeclaration se prevodi u Expression zadat tom funkcijom, a početak tog koda je označen labelom koja odgovara nazivu funkcije. FunctionCall treba da izvrši skok na tu labelu.

 $Y \rightarrow Term$ $Z \rightarrow Term$

Experiment → **experiment** *Body* ~**experiment** Body → Declarations Statements Requirements Execution *Declarations* → VariableDeclaration $Variable Declaration \rightarrow Variable Declaration$; $Variable \mid Variable$ $Variable \rightarrow Type ID$; $Type \rightarrow int \mid double \mid string$ Statements → Statements; Statement | Statement Statement → Assignment | IfStatement If Statement → if (RelExpression): {Statements} else {Statements} RelExpression → Term RelOp Term | Term $Term \rightarrow ID \mid CONST$ $RelOp \rightarrow < | == | >$ Assignment \rightarrow **ID** = Expression Requirements → requirements NodeNumber ~requirements *NodeNumber* → **nodes CONST**; *Execution* → **execution** *NodeList* ~**execution** NodeList → NodeList; NodeDef | NodeDef *NodeDef* → **node** *NodeName* , *Route* ~**node** $NodeName \rightarrow name ID$ Route \rightarrow [Waypoints] Waypoints → Waypoints; Waypoint | Waypoint Waypoint \rightarrow **WP** < X, Y, Z > $X \rightarrow Term$

Expression → Expression + Term | Expression * Term | Term

Napomena: U If naredbi, lista naredbi iz then dela se izvršava ukoliko je uslov različit od 0, a iz else dela u suprotnom. Potrebno je modifikovati ovo pravilo u odnosu na prethodne vežbe, što podrazumeva izmene u leksičkom i sintaksnom analizatoru.

Model → **model** *Body* ~**model**

Body → Declarations Statements Deployment

Declarations → VariableDeclaration

 $Variable Declaration \rightarrow Variable Declaration; Variable | Variable$

 $Variable \rightarrow ID : Type ;$

 $Type \rightarrow int \mid double \mid string$

Statements → Statements; Statement | Statement

Statement → Assignment | WhileStatement

WhileStatement → while (RelExpression): {Statements} default { Statement }

RelExpression → Term RelOp Term | Term

 $Term \rightarrow ID \mid CONST$

 $RelOp \rightarrow less \mid equal \mid greater$

Assignment \rightarrow **ID** := Expression

Deployment → deployment TaskList ServerList ~deployment

ServerList → ServerList; ServerDef | ServerDef

ServerDef
ightarrow server ServerName , Capacity $ilde{\ }$ server

ServerName → serverId ID

Capacity \rightarrow Term

TaskList → TaskList; TaskDef | TaskDef

TaskDef → task TaskName , Demand , Mapping ~task

TaskName → taskId ID

 $Demand \rightarrow Term$

 $Mapping \rightarrow executedOn ID$

Expression → Expression + Term | Expression * Term | Term

Napomena: Lista naredbi u While petlji se ponavlja sve dok je uslov različit od 0. Kada postane 0, izvršava se samo naredba u default delu. Potrebno je modifikovati ovo pravilo u odnosu na prethodne vežbe, što podrazumeva izmene u leksičkom i sintaksnom analizatoru.

Diagram → **diagram** *Body* ~**diagram**

Body → Declarations Statements Deployment

Declarations → VariableDeclaration

 $Variable Declaration \rightarrow Variable Declaration; Variable | Variable$

 $Variable \rightarrow Type ID$;

Type → int | double | string | bool

Statements → Statements; Statement | Statement

Statement → Assignment | DoStatement

DoStatement \rightarrow do (Statements) while (RelExpression) times (CONST)

RelExpression → Term RelOp Term | Term

 $Term \rightarrow ID \mid CONST$

 $ReIOp \rightarrow < | == | >$

Assignment \rightarrow **ID** := Expression

Deployment → **deployment** *ServiceList ServerList* ~**deployment**

ServerList → ServerList; ServerDef | ServerDef

ServerDef
ightarrow server Name , Instances , Environment ~server

ServerName → serverName ID

Instances → **numInstances** *Term*

ServiceList → ServiceList; ServiceDef | ServiceDef

ServiceDef → service ServiceName Environment Allocation ~service

ServiceName → serviceName ID

Environment → **cloud** | **edge**

Allocation → executedBy ID

 $Expression \rightarrow Expression + Term \mid Expression * Term \mid Term$

Napomena: Lista naredbi u Do-While petlji se izvršava ukoliko je uslov true, a izvršenje ponavlja maksimalno puta koliko je vrednost celobrojne konstante u times delu. Potrebno je modifikovati ovo pravilo u odnosu na prethodne vežbe, što podrazumeva izmene u leksičkom i sintaksnom analizatoru.

Strategy → strategy Body ~strategy

Body → Declarations Statements ServiceList

Declarations → declaration VariableDeclaration ~declaration

 $Variable Declaration \rightarrow Variable Declaration; Variable | Variable$

 $Variable \rightarrow Type ID$;

 $Type \rightarrow int \mid double \mid string \mid bool \mid char$

Statements → Statements; Statement | Statement

Statement → Assignment | WhileStatement

WhileStatement → repeat (Term) { Statements }

RelExpression → Term RelOp Term | Term

 $Term \rightarrow ID \mid CONST$

 $RelOp \rightarrow less \mid equal \mid greater$

Assignment \rightarrow **ID** = Expression

Instances → **numInstances** *Term*

ServiceList → ServiceList; ServiceDef | ServiceDef

ServiceDef
ightarrow service ServiceName Instances Allocation AdaptationRule "service"

ServiceName → serviceName ID

Allocation → executedBy ID

AdaptationRule → if Condition then Response

Condition → RelExpression

Response → scale Term | redeployOn ID | optimize

Expression → Expression + Term | Expression * Term | Term

Napomena: Izraz u uslovu Repeat petlje može biti integer (**modifikacija u odnosu na prethodnu vežbu**). Kolika je vrednost uslova, toliko puta se lista naredbi za Repeat ponavlja. Potrebno je modifikovati ovo pravilo u odnosu na prethodne vežbe, što podrazumeva izmene u leksičkom, sintaksnom i semantičkom analizatoru.

```
Program → main () Block

Block → { Declarations StatementList }

Declarations → Declarations VarDecl | \varepsilon

VarDecl → Type NameList;

NameList → ID | NameList, ID

Type → int | char | float

StatementList → StatementList Statement | Statement

Statement → CaseStatement | ID = Expression; | Block

CaseStatement → case (Expression) { WhenStatementList }

WhenStatementList → WhenStatementList WhenStatement | WhenStatement

WhenStatement → when CONST: Statement

Expression → Expression AddOperator Term | Term

AddOperator → + | -

Term → ID | CONST | (Expression)
```

Napomena: U okviru *CaseStatement* upravljačke strukture izvršava se naredba iz prvog *WhenStatement* dela gde je konstanta **CONST** jednaka rezultatu izraza *Expression* u uslovu *CaseStatement* upravljačke strukture. Posle izvršenja te naredbe izlazi se iz *CaseStatement*-a (tj. izvršava se najviše jedan *WhenStatement* deo).

```
Program \rightarrow program Block .

Block \rightarrow begin Variables StatementList end

Variables \rightarrow Variables Declaration | \epsilon

Declaration \rightarrow NameList : Type ;

NameList \rightarrow NameList , ID | ID

Type \rightarrow integer | char | real | boolean

StatementList \rightarrow Statement | StatementList Statement

Statement \rightarrow WhileLoop | ID := Expression ; | Block

WhileLoop \rightarrow while Expression : Statement else Statement

Expression \rightarrow Expression or AndExpression | AndExpression

AndExpression \rightarrow AndExpression and RelExpression | RelExpression

RelExpression \rightarrow Term RelOp Term | Term

RelOp \rightarrow <= | == | >=

Term \rightarrow ID | CONST | (Expression)
```

Napomena: Prva naredba u *While* petlji se izvršava sve dok je ispunjen uslov petlje (*Expression*). Kada taj uslov više nije ispunjen izvršava se tačno jednom naredba u *else* delu petlje i izlazi se iz petlje.

```
Program → main () Block

Block → { Declarations StatementList }

Declarations → Declarations VarDecl | \varepsilon

VarDecl → Type NameList;

NameList → ID | NameList, ID

Type → int | char | float | bool

StatementList → StatementList Statement | Statement

Statement → RedoLoop | ID = Expression; | Block

RedoLoop → loop (Expression) { Statement redo (Expression); Statement }

Expression → Expression | AndExpression | AndExpression

AndExpression → AndExpression & RelExpression | RelExpression

RelExpression → Term RelOp Term | Term

RelOp → <= | == | >=

Term → ID | CONST | (Expression)
```

Napomena: Spoljašnja petlja (*loop*) u okviru *RedoLoop* upravljačke strukture se izvršava sve dok je ispunjen prvi uslov *Expression*. Naredba (*Statement*) pre ključne reči **redo** se izvršava sve dok je ispunjen uslov *Expression* posle ključne reči **redo**. Kada taj uslov više nije ispunjen izvršava se druga naredba (*Statement*).

```
Program \rightarrow program \ Block.
Block → begin Variables StatementList end
Variables \rightarrow Variables Declaration \mid \epsilon
Declaration → NameList: Type;
NameList \rightarrow NameList, ID | ID
Type \rightarrow integer | char | real | boolean
StatementList → Statement | StatementList Statement
Statement → SelectStatement | ID := Expression ; | Block
SelectStatement \rightarrow \mathbf{select} \ \mathbf{begin} \ CaseList \ \mathbf{end}
CaseList → CaseList Case | Case
Case → case Expression => Statement
Expression → Expression or AndExpression | AndExpression
AndExpression → AndExpression and RelExpression | RelExpression
RelExpression → Term RelOp Term | Term
RelOp \rightarrow < | == | >
Term \rightarrow ID \mid CONST \mid (Expression)
```

Napomena: U okviru *SelectStatement* upravljačke strukture izvršava se naredba iz svakog *Case* dela čiji je uslov *Expression* ispunjen.