

# PrimeLister v2.9: Hybrid-Optimized Prime Generation via Segmented Marking

Tobias Jung  
`info@scyrus.de`

January 30, 2026

## Abstract

This paper documents *PrimeLister v2.9*, a C++ implementation that generates all primes up to a given limit  $N$ . The focus is on a hybrid, segmented marking approach (cache-friendly segmentation, OpenMP parallelization) and on providing the full source code for reproducibility.

## 1 Introduction

Prime tables are used in number theory, cryptography, and performance benchmarking. This document describes the approach implemented in `Version2.9.txt` and provides reproducibility notes.

## 2 Algorithmic Approach (Overview)

The code operates on segments of an index space and marks composite candidates using multiple formulas/steps. The segment size is chosen adaptively to improve cache locality and load balancing (with OpenMP).

### 2.1 Index Mapping (Index to Candidate)

Internally, the algorithm does not represent all integers. Instead it enumerates the  $6k \pm 1$  candidates only. In the C++ implementation this is done via a 0-based index  $m$  and the mapping

$$p(m) = \begin{cases} 3m + 5, & m \text{ even} \\ 3m + 4, & m \text{ odd} \end{cases}$$

which yields the ordered sequence  $5, 7, 11, 13, 17, 19, \dots$  (and finally 2 and 3 are prepended explicitly). Marking is performed in this index space: a marked index means the corresponding candidate  $p(m)$  is composite.

### 2.2 Why the Offsets Appear (and a Correctness Sketch)

The three start indices used in the ImageJ prototype (and mirrored in the C++ version) were found empirically by the author (manual derivation with printed candidate lists and pen-and-paper). They are closed-form descriptions of where products of already-generated candidates land in the index space. Concretely, the offsets  $7i + 10$ ,  $3i^2 + 8i + 4$ , and  $3i^2 + 10i + 7$  are the first

index positions (for alternating parity classes of  $i$ ) at which the elimination sequences should start so that stepping by  $2 \cdot A[i]$  hits exactly those indices whose mapped values are multiples of the current candidate. Because every prime  $> 3$  is of the form  $6k \pm 1$ , every composite number that survives the trivial pre-marking (multiples of 5 and 7 in the C++ code) still has a smallest prime factor of the same form. Thus, for each composite candidate there exists an index  $i$  corresponding to its smallest factor, and one of the three marking schemes reaches the index of that composite (with the square-handling step covering the remaining square cases), ensuring that all composites in the candidate set are eventually marked.

### 2.3 Core Idea

The original prototype was developed as an ImageJ macro; this version is often easier to read for understanding the core idea. The approach was then ported to C++ with the help of Gemini and optimized for performance (including OpenMP parallelization and segmented bit marking).

Both variants (C++ and the ImageJ macro) share the same basic construction: only candidates of the form  $6k \pm 1$  are considered (2 and 3 are added separately). Composite candidates are then removed/marketed in three systematic steps:

1. Remove certain multiples from the subsequence  $E_0, E_2, E_4, \dots$  (macro: start index  $7i + 10$  with step size  $2 \cdot A[i]$ ).
2. Remove certain multiples from the subsequence  $E_1, E_3, E_5, \dots$  (macro: start index  $3i^2 + 8i + 4$  with step size  $2 \cdot A[i]$ ).
3. Remove the remaining squares (macro: index  $3i^2 + 10i + 7$ ) and their subsequent elements.

The C++ implementation translates these steps into cache-friendly segmented bit marking and parallelizes segment processing using OpenMP.

## 3 Implementation Details

### 3.1 Parallelization

If OpenMP is available, segments are processed in parallel and results are collected per segment.

### 3.2 File Output

Results are written to `PrimeList.txt` as a table with `cols` columns.

## 4 Experimental Setup and Results

Benchmark values in Table 1 are rounded means from 5 program runs per configuration. The program was restarted between runs, and no other programs were running in the background. Across repeated runs on the same machine, the program suggested the same number of threads.

The reported throughput in Table 1 ( $\text{primes/s}$ ) refers to the core prime computation only (as measured and printed by the program), excluding the file output step (writing `PrimeList.txt`). Existing benchmarks (multiple systems) are included as slides in `Benchmark 002.pptx`.

## 4.1 Benchmark Systems

- **Slow System:** Intel(R) Core(TM) i7-9750H CPU, 2.60GHz, 6 cores, 12 threads; RAM: 16 GB; OS: Windows 11 (64-bit).
- **Medium System:** Intel(R) Core(TM) Ultra 7 155U, 1.70 GHz, 12 cores, 14 threads; RAM: 64 GB; OS: Windows 11 (64-bit).
- **Fast System:** Intel(R) Xeon(R) w7-3455, 2.50 GHz, 24 cores, 48 threads; RAM: 512 GB; OS: Windows 11 (64-bit).

$N$	Slow System (primes/s)	Medium System (primes/s)	Fast System (primes/s)
1,000,000	35,279,999	22,620,598	16,016,731
100,000,000	75,872,411	113,673,142	68,046,243
1,000,000,000	61,019,767	121,854,039	145,904,389
2,000,000,000	46,757,637	103,569,420	120,028,926
5,000,000,000	21,305,698	62,130,239	79,255,786
10,000,000,000	11,990,755	35,989,927	72,354,973

Table 1: Benchmarks from `Benchmark fÃ¶r ArchiveX.csv` (PrimeLister v2.9, “Calculate primes below”). Values are in *primes/s*.

## 5 Discussion and Related Work

At a high level, PrimeLister follows the same core ideas as established high-performance prime generators: a segmented sieve for cache efficiency, pre-sieving of small prime factors, and parallelization across segments. These ingredients are well-known performance techniques for prime sieving.[\[1, 2, 3\]](#)

The distinctive aspect of the present work is the closed-form index arithmetic used in the ImageJ prototype (three start indices and step sizes) and its translation into an indexed marking scheme in the C++ implementation. A targeted literature search for the exact offset formulas used here did not reveal an earlier publication that matches these specific expressions; however, functionally related approaches (segmented sieves and wheel factorization / pre-sieving) are widely used.[\[1, 2\]](#)

### 5.1 Prior Dissemination

The underlying idea was presented publicly in a three-part YouTube video series by the author (Part 1/2: video IDs 58EMZtSLT9Q and N053-sZHmj4; PrimeLister 2.9 update: video ID dpT2nV7t1p0).

## 6 Limitations and Future Work

The current benchmark table reports throughput in primes/s but does not yet record the exact compiler flags (e.g., optimization level, OpenMP settings) or the run protocol (number of runs, warm-up, and aggregation statistic). Future work includes adding a standardized benchmarking harness (fixed number of repetitions and reporting median/variance), documenting build flags precisely, and extending the evaluation to additional platforms/compiler.

## 7 Reproducibility

- Source code: see Appendix A (C++) and Appendix B (ImageJ macro).
- Repository: GitHub [StrangestThings/PrimeLISTER-2.9-Multicore](#).
- Build environment: Microsoft Visual Studio 2026 (MSVC toolset 14.44.35207, Release x64).
- Key flags (excerpt): /O2 /Ot /GL /arch:AVX2 /std:c++20 /permissive- /MD /openmp:llvm /fp:precise; linker: /LTCG /OPT:REF /OPT:ICF.
- Full compiler and linker command lines are provided below.
- Benchmark slides: [Benchmark 002.pptx](#) (measurements on multiple systems).
- Inputs: limit  $N$  and the number of columns `cols`.
- Output: `PrimeList.txt`.
- License: CC0 1.0 Universal (public domain dedication).

### 7.1 MSVC Build Log (Command Lines)

Compiler (CL.exe):

```
C:\Program Files\Microsoft Visual Studio\18\Community\VC\Tools\MSVC\14.44.35207\bin\HostX64\x64\CL.exe /c /I"C:\Users\user\source\repos\PrimeGeneratorMulticore\vcpkg\installed\x64-windows\include" /Zi /nologo /W3 /WX- /diagnostics:column /sdl /MP /O2 /Oi /Ot /GL /D NDEBUG /D _CONSOLE /D _UNICODE /D UNICODE /Gm- /EHsc /MD /GS /Gy /arch:AVX2 /fp:precise /Zc:wchar_t /Zc:forScope /Zc:inline /openmp /std:c++20 /permissive- /Fo"x64\Release\" /Fd"x64\Release\vc143.pdb" /external:W3 /Gd /TP /FC /errorReport:prompt /openmp:llvm
```

Linker (link.exe):

```
C:\Program Files\Microsoft Visual Studio\18\Community\VC\Tools\MSVC\14.44.35207\bin\HostX64\x64\link.exe /ERRORREPORT:PROMPT /OUT:"C:\Users\user\source\repos\PrimeGeneratorMulticore\x64\Release\PrimeGeneratorMulticore.exe" /NOLOGO /LIBPATH:"C:\Users\user\source\repos\PrimeGeneratorMulticore\vcpkg\installed\x64-windows\lib" /LIBPATH:"C:\Users\user\source\repos\PrimeGeneratorMulticore\vcpkg\installed\x64-windows\lib\manual-link" kernel32.lib user32.lib gdi32.lib winspool.lib comdlg32.lib advapi32.lib shell32.lib ole32.lib oleaut32.lib uuid.lib odbc32.lib odbc32.lib "C:\Users\user\source\repos\PrimeGeneratorMulticore\vcpkg\installed\x64-windows\lib\*.lib" /MANIFEST /MANIFESTUAC:"level='asInvoker' uiAccess='false'" /manifest:embed /DEBUG /PDB:"C:\Users\user\source\repos\PrimeGeneratorMulticore\x64\Release\PrimeGeneratorMulticore.pdb" /SUBSYSTEM:CONSOLE /OPT:REF /OPT:ICF /LTCG /LTCGOUT:"x64\Release\PrimeGeneratorMulticore.iobj" /TLBID:1 /DYNAMICBASE /NXCOMPAT /IMPLIB:"C:\Users\user\source\repos\PrimeGeneratorMulticore\x64\Release\PrimeGeneratorMulticore.lib" /MACHINE:X64 x64\Release\PrimeGeneratorMulticore.obj
```

## References

- [1] GeeksforGeeks. *Segmented Sieve*.
- [2] Kim Walisch. *primesieve* (documentation / manual).
- [3] Stephan Brumme. *Parallel Prime Sieve*.

## 8 License and Attribution

The author dedicates the source code and benchmark data in this project to the public domain under the CC0 1.0 Universal dedication.

The underlying formulas and the original implementation were developed by Tobias Jung in ImageJ; the C++ port and parallelization were produced with the assistance of Gemini-3.

## A Source Code

```
1 #include <iostream>
2 #include <vector>
3 #include <string>
4 #include <algorithm>
5 #include <cmath>
6 #include <cstdint>
7 #include <cstdlib>
8 #include <fstream>
9 #include <iomanip>
10 #include <charconv>
11 #include <cstring>
12 #include <thread>
13
14 // ---- OpenMP Setup ----
15 #ifdef _OPENMP
16 #include <omp.h>
17 #else
18 #include <chrono>
19 inline double omp_get_wtime() {
20     static const auto t0 = std::chrono::steady_clock::now();
21     return std::chrono::duration<double>(std::chrono::steady_clock::now()
22         - t0).count();
23 }
24 inline int omp_get_num_procs() { return (int)std::thread::
25     hardware_concurrency(); }
26 #endif
27
28 #if defined(_WIN32)
29 #ifndef NOMINMAX
30 #define NOMINMAX
31 #endif
32 #include <windows.h>
33 #include <intrin.h>
34 #endif
35
36 using u64 = unsigned long long;
37
38 // ----- Optimierte Helfer -----
39 static inline u64 ceil_div_u64(u64 a, u64 b) {
40     return (a + b - 1) / b;
41 }
42
43 static inline unsigned ctz64(uint64_t x) {
```

```

42 #if defined(_MSC_VER)
43     unsigned long idx;
44     _BitScanForward64(&idx, x);
45     return (unsigned)idx;
46 #else
47     return (unsigned)_builtin_ctzll(x);
48 #endif
49 }

50
51 // ----- Datei-Ausgabe -----
52 static inline void buf_append_u64(std::string& buf, u64 x) {
53     char tmp[32];
54     auto res = std::to_chars(tmp, tmp + 32, x);
55     buf.append(tmp, (size_t)(res.ptr - tmp));
56 }
57
58 static void write_prime_table(const std::vector<u64>& P, u64 cols) {
59     std::ofstream out("PrimeList.txt", std::ios::out | std::ios::binary);
60     if (!out) return;
61
62     std::string buf;
63     buf.reserve(16 * 1024 * 1024);
64
65     u64 count = 0;
66     for (size_t i = 0; i < P.size(); ++i) {
67         if (count == 0) {
68             buf_append_u64(buf, (u64)(i + 1));
69             buf.push_back('-');
70             buf_append_u64(buf, (u64)(std::min)((size_t)(i + cols), P.size()
71                                         ()));
72             buf.push_back('\t');
73         }
74         buf_append_u64(buf, P[i]);
75         buf.push_back('\t');
76         if (++count == cols) {
77             buf.push_back('\n');
78             count = 0;
79         }
80         if (buf.size() >= 15 * 1024 * 1024) {
81             out.write(buf.data(), (std::streamsize)buf.size());
82             buf.clear();
83         }
84     }
85     if (count != 0) buf.push_back('\n');
86     out.write(buf.data(), (std::streamsize)buf.size());
87 }

88 // ----- Kern-Algorithmus v2.9 (HYBRID) -----
89 static std::vector<u64> primes_formeln(u64 N, u64 cols, bool silent) {
90     u64 bereich = (N + 2ull) / 3ull + (15ull * (u64)(std::max)((size_t)1,
91                                         (size_t)cols));
92     if (bereich == 0) return {};
93
94     int T = 1;

```

```

94 #ifdef _OPENMP
95     T = omp_get_max_threads();
96 #endif
97
98 // ADAPTIVE: mehr Segmente fuer besseres Load-Balancing bei grossen N
99 u64 target_segments = (N > 1000000000ull) ? ((u64)T * 6ull) : ((u64)T
100    * 4ull);
100 u64 SEG_IDX = (bereich + target_segments - 1) / target_segments;
101
102 // Cache-optimierte Grenzen
103 const u64 MIN_SEG = 1ull << 23; // 1 MB (kleiner fuer kleine N)
104 const u64 MAX_SEG = 1ull << 27; // 16 MB
105 if (SEG_IDX < MIN_SEG) SEG_IDX = MIN_SEG;
106 if (SEG_IDX > MAX_SEG) SEG_IDX = MAX_SEG;
107 if (SEG_IDX > bereich) SEG_IDX = bereich;
108
109 u64 nsegs = (bereich + SEG_IDX - 1) / SEG_IDX;
110 const long long M = (long long)berreich;
111
112 long long end1 = (M - 11) / 7;
113 if (end1 < 0) end1 = -1;
114
115 long long end2 = (M <= 4) ? -1 : (long long)std::floor((-8.0L + std::
116    sqrt(16.0L + 12.0L * (long double)M)) / 6.0L);
116 long long end3 = (M <= 7) ? -1 : (long long)std::floor((-10.0L + std::
117    sqrt(16.0L + 12.0L * (long double)M)) / 6.0L);
118
119 std::vector<std::vector<u64>> buckets((size_t)nsegs);
120
121 double density = 1.0 / std::log((double)N + 1.0);
122 size_t estimated_per_seg = (size_t)(SEG_IDX * 3.0 * density * 1.15);
123 for (auto& b : buckets) {
124     b.reserve(estimated_per_seg);
125 }
126
127 #pragma omp parallel
128 {
129     std::vector<uint64_t> mark_words;
130     mark_words.reserve((size_t)((SEG_IDX + 63ull) >> 6));
131
132 #pragma omp for schedule(guided, 1)
133     for (long long s = 0; s < (long long)nsegs; ++s) {
134         u64 base = (u64)s * SEG_IDX;
135         u64 end = (u64)(std::min)((size_t)berreich, (size_t)(base +
136             SEG_IDX));
137         size_t seg_len = (size_t)(end - base);
138         size_t nwords = (seg_len + 63u) >> 6;
139
140         mark_words.resize(nwords);
141         std::memset(mark_words.data(), 0, nwords * sizeof(uint64_t));
142
143         long long Mend = (long long)end;
144         long long e1seg = (long long)(std::min)((size_t)end1, (size_t)
145             ((Mend - 11) / 7));

```

```

143     long long e2seg = (Mend <= 4) ? -1 : (long long)(std::min)((
144         size_t)end2, (size_t)std::floor((-8.0L + std::sqrt(16.0L +
145             12.0L * (long double)Mend)) / 6.0L));
146     long long e3seg = (Mend <= 7) ? -1 : (long long)(std::min)((
147         size_t)end3, (size_t)std::floor((-10.0L + std::sqrt(16.0L +
148             12.0L * (long double)Mend)) / 6.0L));
149
150 #define SETBIT(idx) mark_words[(idx) >> 6] |= (1ull << ((idx) & 63))
151
152     u64 r10 = base % 10ull;
153     u64 r14 = base % 14ull;
154
155     u64 i1 = (r10 == 0 ? base : base + (10 - r10));
156     if (i1 == 0 && base == 0) i1 = 10;
157     for (u64 i = i1; i < end; i += 10) SETBIT((size_t)(i - base));
158
159     for (u64 i = base + (r10 <= 7 ? 7 - r10 : 17 - r10); i < end;
160         i += 10)
161         SETBIT((size_t)(i - base));
162
163     for (u64 i = base + (r14 <= 10 ? 10 - r14 : 24 - r14); i < end
164         ; i += 14)
165         SETBIT((size_t)(i - base));
166
167     u64 i14 = base + (r14 <= 1 ? 1 - r14 : 15 - r14);
168     if (i14 == 1 && base == 0) i14 = 15;
169     for (u64 i = i14; i < end; i += 14) SETBIT((size_t)(i - base))
170         ;
171
172 // Formel 1: Simpler Loop (kein Unrolling bei grossem Overhead
173 // )
174     for (long long i = 0; i <= e1seg; i += 2) {
175         u64 iu = (u64)i;
176         u64 step = 6ull * iu + 10ull, j0 = 7ull * iu + 10ull;
177         if (j0 < base) j0 += ceil_div_u64(base - j0, step) * step;
178         for (u64 j = j0; j < end; j += step) SETBIT((size_t)(j -
179             base));
180     }
181
182 // Formel 2
183     for (long long i = 1; i <= e2seg; i += 2) {
184         u64 iu = (u64)i;
185         u64 step = 6ull * iu + 8ull;
186         u64 j0 = 3ull * iu * iu + 8ull * iu + 4ull;
187         if (j0 < base) j0 += ceil_div_u64(base - j0, step) * step;
188         for (u64 j = j0; j < end; j += step) SETBIT((size_t)(j -
189             base));
190     }
191
192 // Formel 3: Branch-Optimierung
193     for (long long i = 0; i <= e3seg; i += 2) {
194         u64 iu = (u64)i;
195         u64 step = 6ull * iu + 10ull;
196         u64 idx = 3ull * iu * iu + 10ull * iu + 7ull;

```

```

187
188     if (idx < base) {
189         u64 j0 = idx + ceil_div_u64(base - idx, step) * step;
190         for (u64 j = j0; j < end; j += step) SETBIT((size_t)(j
191             - base));
192     }
193     else if (idx < end) {
194         SETBIT((size_t)(idx - base));
195         for (u64 j = idx + step; j < end; j += step) SETBIT((
196             size_t)(j - base));
197     }
198 #undef SETBIT
199
200     std::vector<u64>& local = buckets[(size_t)s];
201
202     for (size_t wi = 0; wi < nwords; ++wi) {
203         uint64_t inv = ~mark_words[wi];
204         while (inv) {
205             unsigned b = ctz64(inv);
206             size_t off = (wi << 6) + b;
207             if (off >= seg_len) break;
208             u64 idx_v = base + off;
209             u64 p = ((idx_v & 1ull) == 0ull) ? (3ull * idx_v + 5
210                 ull) : (3ull * idx_v + 4ull);
211             if (p <= N) local.push_back(p);
212             inv &= (inv - 1ull);
213         }
214     }
215 }
216
217     std::vector<u64> primes;
218     size_t total_p = 1 + (N >= 3 ? 1 : 0);
219     for (const auto& b : buckets) total_p += b.size();
220     primes.reserve(total_p);
221
222     primes.push_back(2);
223     if (N >= 3) primes.push_back(3);
224
225     for (const auto& b : buckets) {
226         primes.insert(primes.end(), b.begin(), b.end());
227     }
228
229     return primes;
230 }
231
232 // ----- MAIN -----
233 int main() {
234     std::ios::sync_with_stdio(false); std::cin.tie(NULL);
235
236 #if defined(_WIN32)
237     SetPriorityClass(GetCurrentProcess(), HIGH_PRIORITY_CLASS);

```

```

238 #endif
239
240     std::cout << "*****\n";
241     std::cout << "*  PrimeLister v2.9 Hybrid-Optimiert (2026) *\n";
242     std::cout << "*****\n\n";
243     ;
244
245     int max_t = omp_get_num_procs();
246     int best_t = 1; double min_t = 1e30;
247     u64 calib_N = 1000000000;
248
249     std::cout << "[Step 1] Kalibriere Hardware mit N = 1.000.000.000...\n";
250     ;
251     std::vector<int> test_configs = { 1, max_t / 4, max_t / 2, (3 * max_t)
252                                     / 4, max_t };
253
254     for (int t : test_configs) {
255         if (t < 1 || t > max_t) continue;
256 #ifdef _OPENMP
257         omp_set_num_threads(t);
258 #endif
259         double start = omp_get_wtime();
260         auto d = primes_formeln(calib_N, 1, true);
261         double end = omp_get_wtime() - start;
262         std::cout << " - Threads " << std::setw(2) << t << ":" "
263                 << std::fixed << std::setprecision(4) << end << "s\n";
264         if (end < min_t) { min_t = end; best_t = t; }
265     }
266 #ifdef _OPENMP
267         omp_set_num_threads(best_t);
268 #endif
269         std::cout << ">> Optimal: " << best_t << " Threads ("
270             << std::fixed << std::setprecision(2) << (min_t * 1000) << " ms).\n";
271
272         u64 N; u64 cols;
273         std::cout << "[Step 2] Parameter eingeben\nLimit N: "; std::cin >> N;
274         std::cout << "Spalten: "; std::cin >> cols;
275
276         std::cout << "\n[Step 3] Hauptberechnung laeuft...\n";
277         double t0 = omp_get_wtime();
278         auto P = primes_formeln(N, cols, false);
279         double t1 = omp_get_wtime();
280
281         std::cout << "\nFirst 100 primes:\n";
282         size_t limit_first = (P.size() < 100) ? P.size() : 100;
283         for (size_t i = 0; i < limit_first; ++i) std::cout << P[i] << ',';
284
285         std::cout << "\n\nLast 10 primes under limit:\n";
286         size_t start_last = (P.size() > 10) ? (P.size() - 10) : 0;
287         for (size_t i = start_last; i < P.size(); ++i) std::cout << P[i] << ',';
288
289         std::cout << "\n\nOverall primes found: " << P.size() << "\n";

```

```

287     std::cout << "Overall calculation time [s]: " << std::fixed << std::
288         setprecision(6) << (t1 - t0) << "\n";
289
290     if (t1 - t0 > 0) {
291         double rate = (double)P.size() / (t1 - t0);
292         std::cout << "~ " << (u64)rate << " primes/s\n";
293     }
294
295     std::cout << "\nWriting PrimeList.txt ... \n";
296     double two = omp_get_wtime();
297     write_prime_table(P, cols);
298     double tw1 = omp_get_wtime();
299     std::cout << "Write time: " << std::fixed << std::setprecision(3) << (
300         tw1 - two) << "s\n";
301
302     std::cout << "\nDone. Press Enter to exit... ";
303     std::cin.ignore(1000, '\n'); std::cin.get();
304     return 0;
305 }
```

## B Source Code (ImageJ)

```

1 time=getTime();
2 print("//\Clear");
3
4 // 16GB-Ram: Not more than 100 Mio.
5
6 // Calculates about 0.7 sec per Million
7
8 bereich=100000;
9
10 A=newArray(bereich);
11
12
13 print("Array was created: " + IJ.freeMemory());
14
15 // Fill Array A with prime candidates
16 k=1;
17 for (i = 1; i < A.length; i=i+2) {A[i-1]=6*k-1; A[i]=6*k+1; k++;}
18 // ****
19
20 print("Array filled: " + IJ.freeMemory());
21
22 print(" Delete all the multiples of the single elements (E0, E2, E4...) ");
23
24 for(i = 0; (7*i+10) < A.length; i=i+2){
25     while (A[i]<1){i=i+2;}
26     step=2*A[i];
27     for (j = 7*i+10; j < A.length; j=j+step) {A[j]=0;}
28 }
29
30 print(" Delete all the multiples of the single elements (E1, E3, E5...) ");
```

```

31
32 for(i = 1; (3*i*i+8*i+4) < A.length; i=i+2){
33 while (A[i]<1){i=i+2;}
34 step=2*A[i];
35 for (j = 3*i*i+8*i+4; j < A.length; j=j+step) {A[j]=0;}
36 }
37
38 print(" Now delete all the remaining squares of the single elements (E0,
      E2, E4...)");
39
40 for (i = 0; (3*i*i+10*i+7) < A.length; i=i+2) {
41 if(A[3*i*i+10*i+7]>0){step=2*sqrt(A[3*i*i+10*i+7]); A[3*i*i+10*i+7]=0; for
      (j =(3*i*i+10*i+7); j < A.length; j=j+step) {A[j]=0;}}
42 }
43
44 // Delete Zeros
45
46 A=Array.deleteValue(A, 0);
47
48 // Attach 2 and 3
49 A=Array.concat(A,2);
50 A=Array.concat(A,3);
51
52 A=Array.sort(A);
53
54 time2=getTime();
55 zeit=(time2-time)/1000;
56
57 print("Range: " + bereich);
58 print("Primes identified: " + A.length);
59 print("Calculation time [sec]: " + zeit);
60 print("Primes per second: " + A.length/zeit);
61
62
63 print("*****");
64
65 results = getInfo("log");
66 results=replace(results,".",",");
67 print("\Clear");
68 print(results);
69
70 // Output of 100 primes per line with the according counter
71
72
73 spalten=100;
74
75 for (i = 0; i < A.length-spalten; i=i+spalten) {
76 line=toString((1+i) + "-" + (i+spalten) + ":" + "\t");
77 for (s = i; s < i+spalten; s++) {line=line+A[s]+\t;};
78 print(line);
79 }

```