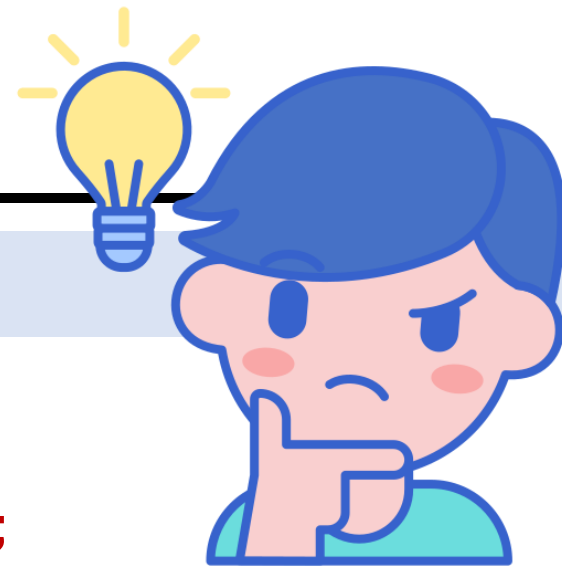


Data structure and Programming (C++)

Hash Table



About array



- Assume that we have an array

```
int a[ ] = {4, 7, 10, 0, 9, 100, 50, 80, 25, 56, 90, 17};
```

We can use for loop to access data in this array.

Suppose that we want to search for a number and tell the user whether that number is exist or not.

- A. Let's search for the number 90 in this array. How many times did you use the comparison operation?
- B. Let's search for the number 10 in this array. How many times did you use the comparison operation?
- C. Let's search for the number 99 in this array. How many times did you use the comparison operation?



By using **hash table**, we compare only 1 time in order to search.

Outline

- Time complexity
- What is Hashing? Hash table? Hash function?
- Collision
- Collision resolution techniques
 - Open hashing
 - Close hashing
- Examples

Time complexity

Time complexity is a time required for an algorithm to run.
It is denoted by a big-O notation: $O(n)$

- Linked list : $O(n)$
- Array (unsorted) : $O(n)$
- Array (sorted) : $O(\log_2 n)$

- Best case
- Average case
- **Worst case**

Searching time
complexity

Big-O Notation

□ Introduction

- Introduction about Big-O Notation: bit.ly/312x9ou

Big-O Notation:
Introduction in 5

Other Data structure

❑ Searching operation

- Linked list
- Array
- Stack
- Queue

What is the time required to search ?

$O(n)$

Hash Table

❑ Searching operation: Time Complexity

Hash Table : $O(1)$

General rules for time complexity

1. ignore constants

$$5n \rightarrow O(n)$$

2. certain terms "dominate" others

$$O(1) < O(\log n) < O(n) < O(n \log n) < O(n^2) < O(2^n) < O(n!)$$

i.e., ignore low-order terms

constant time

$O(1)$ "big oh of one"

```
x = 5 + (15 * 20);
```

independent of input size, N

```
x = 5 + (15 * 20);  
y = 15 - 2;  
print x + y;
```

total time = $O(1) + O(1) + O(1) = O(1)$
 $3 * O(1)$

linear time

$$N * O(1) = O(N)$$

```
for x in range (0, n):  
    print x; // O(1)
```

```
y = 5 + (15 * 20);           O(1)
for x in range (0, n):      }
    print x;                } O(N)
```

total time = $O(1) + O(N) = O(N)$

quadratic time

$O(N^2)$

```
for x in range (0, n):  
    for y in range (0, n):  
        print x * y; //  $O(1)$ 
```

$O(N^2)$

```
x = 5 + (15 * 20);  
for x in range (0, n):  
    print x;  
for x in range (0, n):  
    for y in range (0, n):  
        print x * y;
```

$O(1)$

$O(N)$

$O(N^2)$

Terms

❑ Hash Table Vs. Hash Function Vs. Hashing

▪ Hash Table

- Is a data structure which is used to store key-value pairs. It is implemented as array

▪ Hash Function

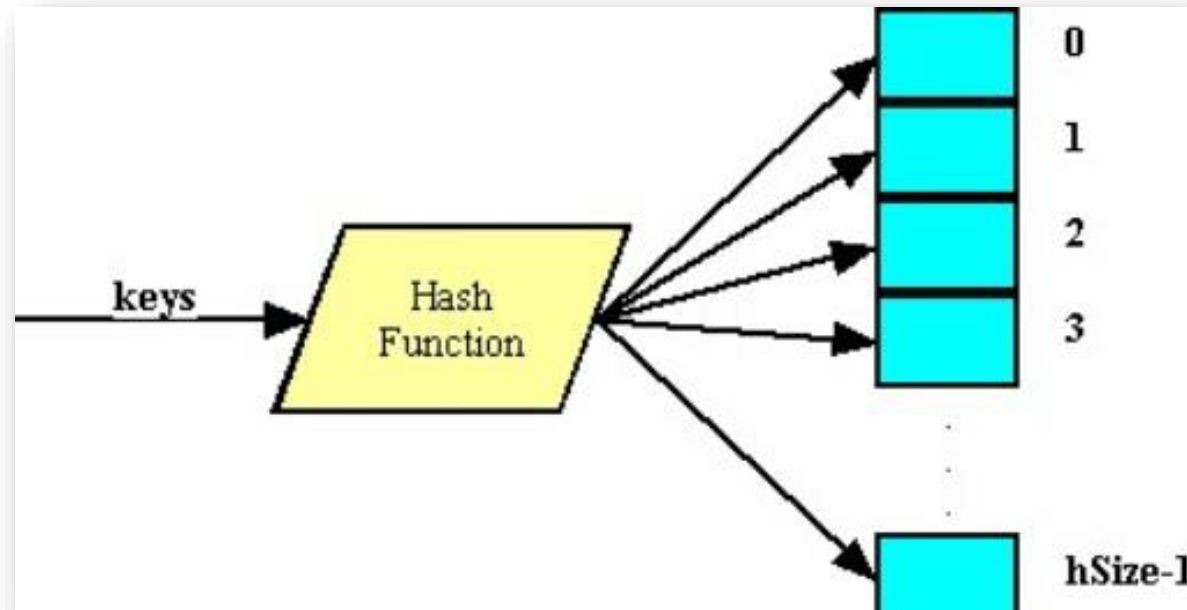
- Is a function used by hash table **to compute an index of an array** in which an element will be inserted to or be searched at

▪ Hashing

- In hashing, large keys are converted into smaller ones using hash function AND then the values are stored in hash table.

Hash Functions

- A hash function usually means a function that compresses.
 - The output is shorter than the input



An example of a hash function

❑ Definition

```
int hash(int key){  
    return key%7;  
}
```



Size of array

Example: Implementation of hash table

```
#include<iostream>
using namespace std;
```

```
const int SIZE=7;
int ht[SIZE];
```

```
void initializeArray(){
    for(int i=0; i<SIZE; i++){
        ht[i] = -999
    }
}
```

```
int hashFunction(int n){
    return n%SIZE;
}
```

```
void insertData(int value){
    int index;
    index = hashFunction(value);
    ht[index] = value;
}
```

```
void displayHT(){
    for(int i=0; i<SIZE; i++){
        cout<<i<<"\t--> ";
        cout<<ht[i]<<"\n";
    }
}
```

1

```
main(){
    insertData(7);
    insertData(8);
    insertData(25);
    displayHT();
}
```

2

```
main(){
    insertData(7);
    insertData(8);
    insertData(15);
    insertData(22);
    insertData(25);
    displayHT();
}
```

What is the output?

```
0 --> 7
1 --> 8
2 --> 0
3 --> 0
4 --> 0
5 --> 25
6 --> 0
```

```
0 --> 7
1 --> 22
2 --> 0
3 --> 0
4 --> 25
5 --> 0
6 --> 0
```

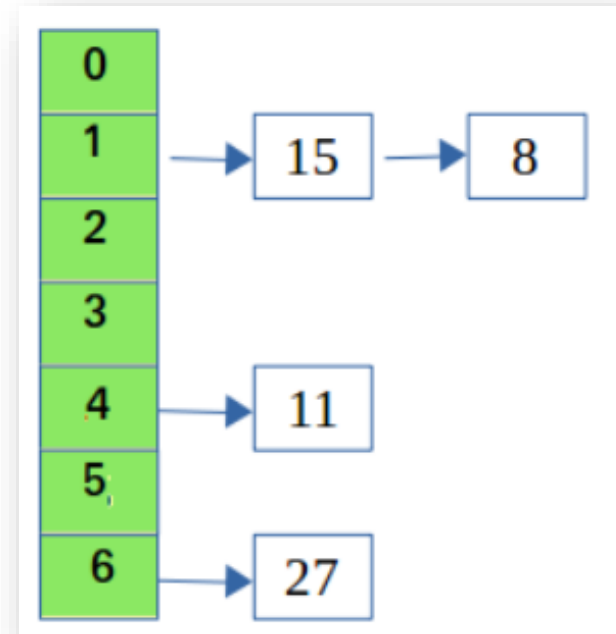
Ooop! Data is overridden?
This is known as collision!

Collision

❑ Definition

- A situation when two or more data hash to be stored in the same location in the table, is called a *hash collision*.

- Suppose we have a hash table with 7 elements
 - Hash function is to modulo with 7
 - How to insert these numbers?
 - 15, 11, 27, 8



Factors of a good hash function

- ✓ Easy to compute
- ✓ Minimize collision



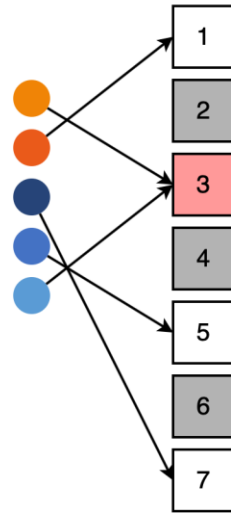
Perfect Hashing

□ Definition

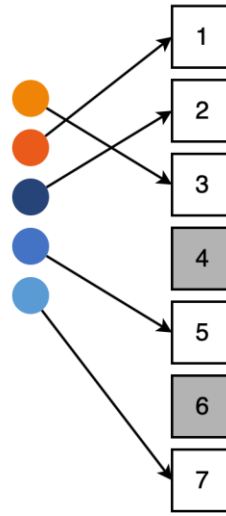
- Perfect hashing maps each valid input to a different hash value

(no collision)

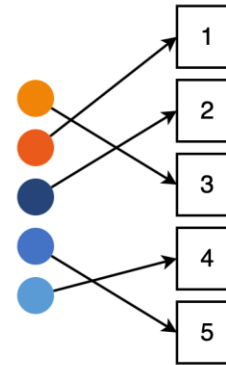
Universal



Perfect



Minimal Perfect

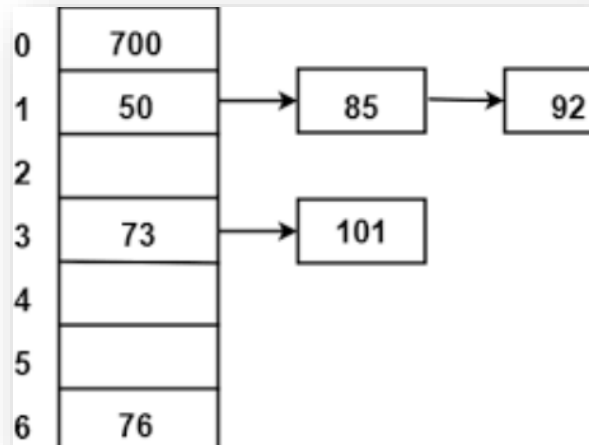


Collision Resolution Techniques

Collision Resolution Techniques

Separate Chaining (Open Hashing)

(An array of linked list implementation)



An example of using chaining

Open Addressing (Closed Hashing)

(Array-based implementation)

- Linear Probing
- Quadratic Probing
- Double Hashing

Solution for Hash Collision

❑ Open Hashing

- Open hashing defines **each slot in the hash table to be the head of a linked list**
- All records that hash to a particular slot (collision cases) are placed on that slot's **linked list**

Hash Table with Chaining resolution technique (An array of linked list)

1

```
struct Element{
    int value;
    Element *next;
};
struct List{
    int n;
    Element *head, *tail;
};
const int SIZE=7;
List *ht[SIZE];
```

2

```
List* createAnEmptylist(){
    List *L1;
    L1 = new List();
    L1->n = 0;
    L1->head=NULL;
    L1->tail=NULL;
    return L1;
}
void createEmptyAllLists(){
    for(int i=0; i<SIZE; i++){
        ht[t] = createEmptyList()
    }
}
```

3

```
void addEnd(List *ls, int value){
    Element *e;
    e= new Element;
    e->next=NULL;
    e->value=value;

    if(ls->n==0){
        ls->head = e;
        ls->tail = e;
        ls->n = ls->n + 1;
    }else{
        ls->tail->next = e;
        ls->tail = e;
        ls->n = ls->n + 1;
    }
}
```

Display data
in hash table

4

```
int hashFunction(int n){
    return n%SIZE;
}
```

5

```
void insertData(int value){
    int index;
    index = hashFunction(value);
    addEnd(ht[index], value);
}
```

6

```
void displayHTO(){
    Element *e;
    for(int i=0; i<SIZE; i++){
        cout<<i<<"\t-->";
        if(ht[i]!=NULL){
            e = ht[i]->head;
            while(e!=0){
                cout<<e->value<<" ";
                e=e->next;
            }
        }
        cout<<endl;
    }
}
```

7

```
int main(){
    createEmptyAllLists();
    insertData(7);
    insertData(8);
    insertData(15);
    insertData(22);
    insertData(25);
    displayHT();
}
```

Output?

```
0      -->7
1      -->8  15  22
2      -->
3      -->
4      -->25
5      -->
6      -->
```

Solution for Hash Collision

❑ Closed Hashing

1. Linear probing
2. Quadratic probing
3. Double hashing

Case Study

Try these topics and explore how each method works!

Q & A