

HEAPSORT

CS340

Heapsort Properties

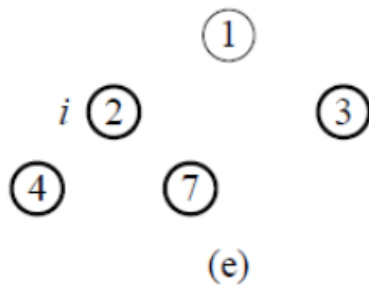
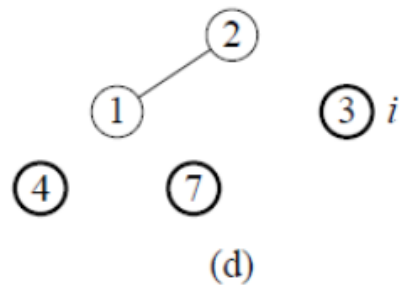
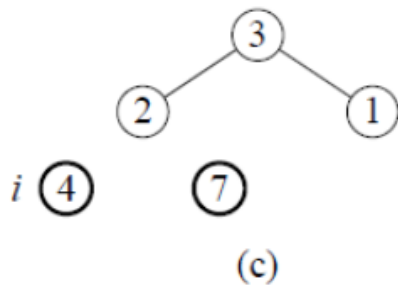
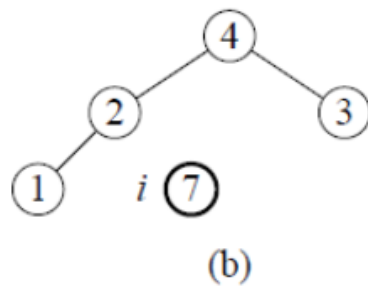
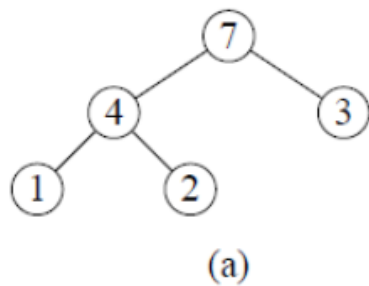
- Combines the best properties of insertion sort and merge sort
 - $O(n \lg n)$ worst case
 - Sorts in place (no extra storage needed)

Heapsort

1. Build a max-heap from the array ($O(n)$)
2. Place max element in its correct position by swapping it with the last item ($O(1)$)
3. Decrease heap size by 1
4. Call max-heapify on root $O(\lg n)$ each time, n times

```
HEAPSORT(A)  
1  BUILD-MAX-HEAP(A)  
2  for i = A.length downto 2  
3      exchange A[1] with A[i]  
4      A.heap-size = A.heap-size - 1  
5      MAX-HEAPIFY(A, 1)
```

Heapsort example



A

1	2	3	4	7
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Priority Queues

- A **priority queue** is a data structure for maintaining a set S of elements, each with an associated value called a **key**.

Max-Priority Queue

- A max-priority queue supports the following operations:
 - **INSERT(A, x)** inserts the element x into the array A
 - **MAXIMUM(A)** returns the element of A with the largest key.
 - **EXTRACT-MAX(A)** removes and returns the element of A with the largest key.
 - **INCREASE-KEY(A, i, k)** increases the value of element i 's key to the new value k , which is assumed to be at least as large as i 's current key value.

Priority Queue methods

```
<T> T maximum(T[ ] A) {
```

```
}
```

```
<T> T extractMax(T[ ] A) {
```

```
}
```

Priority Queue methods

HEAP-INCREASE-KEY(A, i, key)

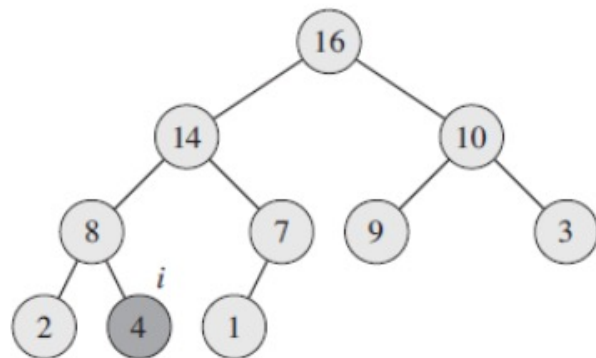
```
1  if  $key < A[i]$ 
2      error "new key is smaller than current key"
3   $A[i] = key$ 
4  while  $i > 1$  and  $A[PARENT(i)] < A[i]$ 
5      exchange  $A[i]$  with  $A[PARENT(i)]$ 
6       $i = PARENT(i)$ 
```

MAX-HEAP-INSERT(A, key)

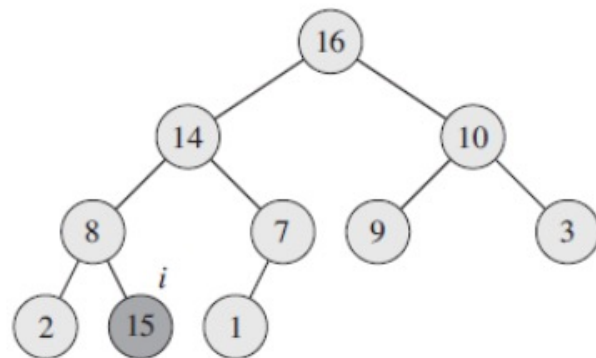
```
1   $A.heap-size = A.heap-size + 1$ 
2   $A[A.heap-size] = -\infty$ 
3  HEAP-INCREASE-KEY( $A, A.heap-size, key$ )
```

Time complexities?

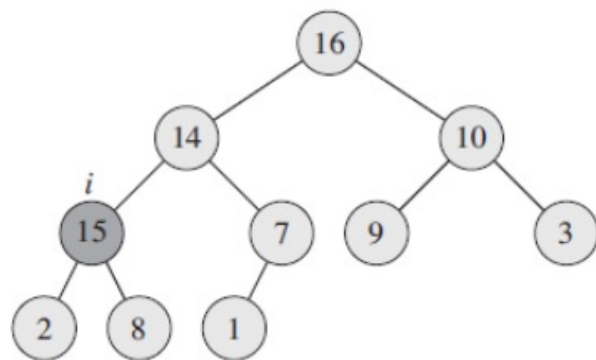
Increase-Key example



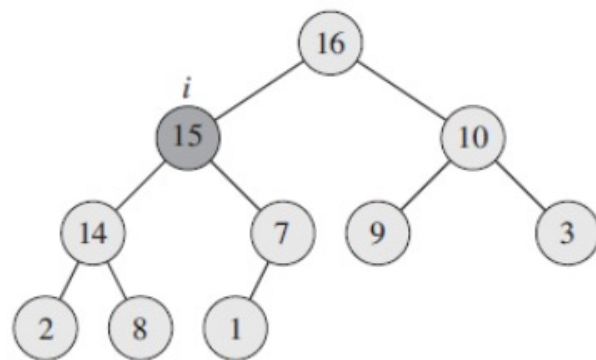
(a)



(b)



(c)



(d)