

RADIX SORT

CS340

Radix Sort

- Sorting based on the digits in a number
 - In decimal, there are 10 digits
 - Based on this, you could divide numbers into bins based on their most significant digit, sort the bins, and combine.
 - This means 9/10 bins must be put aside to sort each of the bins.
 - That's a lot of intermediate piles to keep track of.

Radix Sort

- Radix Sort counterintuitively sorts based on the least significant digit first.
- Then combines all numbers into a single bin.
- Then sorts that bin based on second-least significant digit.
- And so on.
- If the numbers being sorted have d digits, then only d passes through the numbers are required to sort.
- In order for this to work, the sort used at each pass must be stable.

Radix Sort

329	720	720	329
457	355	329	355
657	436	436	436
839	457	839	457
436	657	355	657
720	329	457	720
355	839	657	839

RADIX-SORT(A, d)

- 1 **for** $i = 1$ **to** d
- 2 use a stable sort to sort array A on digit i

Radix Sort

- Counting sort is a good choice for the sort on line 2.
- Time complexity of Radix Sort?

```
RADIX-SORT( $A, d$ )  
1  for  $i = 1$  to  $d$   
2      use a stable sort to sort array  $A$  on digit  $i$ 
```

Time complexity

- $\Theta(d(n + k))$
- With a computer, the number is binary.
- The binary number consists of b -bits
- We choose how many binary numbers make a digit, r
- Time complexity is then
 $\Theta((b/r)(n + 2^r))$

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Example of time complexity

- In the binary world, this time complexity is:
 - $\Theta((b/r)(n + 2^r))$
 - b = number of bits, r = number of bits in a digit
- Eg we are sorting a 32 bit word ($b=32$), which we consider as having 4 8bit digits ($r=8$).
- The number of passes was d , $d = (b/r) = 4$.
- Because there are r bits in a digit, the biggest number is $k = 2^r$.

Example of time complexity

- $\Theta((b/r)(n + 2^r))$
 - b = number of bits, r = number of bits in a digit
 - Eg an ordinary number has $b=32$, and we choose $r=8$.
- This formulation is handy because it removes k .
- n and b are given. Choose r to minimize time complexity
- If $r=b$, then we have
 - $\Theta((b/b)(n + 2^b)) = \Theta(n + 2^b) = \Theta(n)$

Radix Sort

- Is it preferable to other sorts like Quicksort?
 - Hard to say. $\Theta(n + 2^b)$ vs $\Theta(n \lg n)$
 - Radix sort appears better.
 - Constant factors with Radix Sort may be large.
 - Quicksort makes more passes over the n items, but counting sort passes take longer.
 - Quicksort sorts in place.

```
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1  for  $i = 1$  to  $d$   
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```

How about an example?

- COW
- DOG
- SEA
- RUG
- ROW
- MOB
- BOX

Radix Sort

- It's a loop
 - What is the loop invariant?
 - Initialization, maintenance, termination proofs?

RADIX-SORT(A, d)

1 **for** $i = 1$ **to** d

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Sorting in general

- Is Radix Sort the limit?
- Can it be faster than $O(n)$ time?