

MEDIANS AND ORDER STATISTICS

CS340

What is a median?

17	34	36	49	57
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17	34	36	49	57	102
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- The **median** of a set of numbers is the middle value, such that half the numbers are less than the median and half are greater.
- The median of an odd-length sorted array is at index $(n+1)/2$.
- An even-length sorted array has 2 medians, at $n/2$ and $n/2 + 1$
- How can we find the median of a sorted array?

What is an order statistic?

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- In a set of n distinct numbers, $(i-1)$ of the numbers are less than the i th order statistic.
- The **selection problem** is the problem of selecting an order statistic.
- For example, the 2nd order statistic of the above arrays is 34. ($i-1=1$ numbers are less than 34).

Minimum and Maximum

17	34	36	49	57	102
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MINIMUM(*A*)

```
1  min = A[1]
2  for i = 2 to A.length
3      if min > A[i]
4          min = A[i]
5  return min
```

What is the time complexity?
Is this the best we can do?

Simultaneous Min and Max

MINIMUM(*A*)

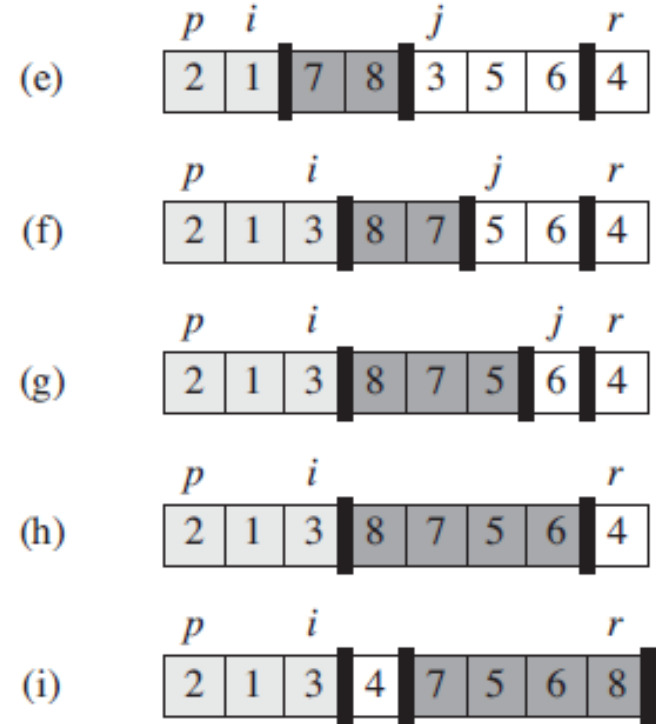
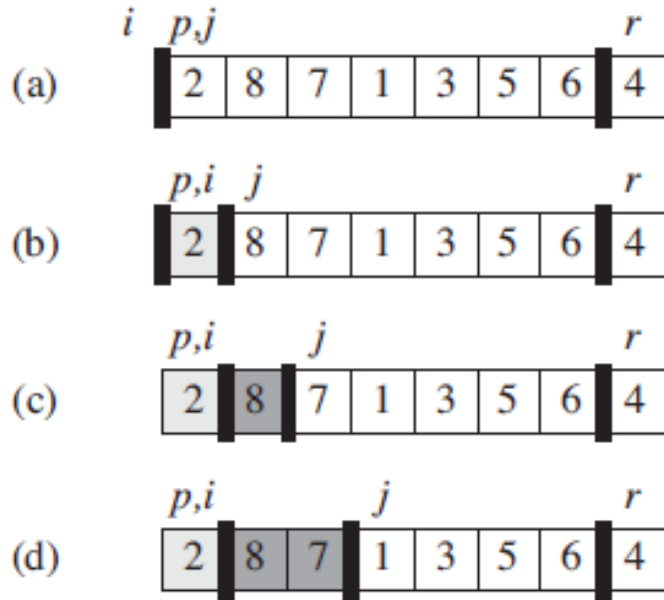
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- Can find the minimum using $n-1$ comparisons, and the max using $n-1$ comparisons = $2n-2$ comparisons.
- Can do better: pick elements in pairs and compare greater item to max and smaller item to min. 3 comparisons for every 2 items = $3n/2-2$ comparisons.

Selecting in $O(n)$ time

- Remember Quicksort Partition?



- In this example, Partition found the 4th smallest value

Selecting in $O(n)$ time

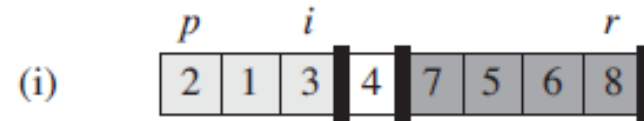
- Here we found the 4th smallest value
- What if we are looking for the 2nd smallest value?
- We can throw away half the data, and partition on the range $p..i$

(i)



What is the time complexity of partitioning?

Selecting in $O(n)$ time



RANDOMIZED-SELECT(A, p, r, i)

```
1  if  $p == r$ 
2      return  $A[p]$ 
3   $q = \text{RANDOMIZED-PARTITION}(A, p, r)$ 
4   $k = q - p + 1$ 
5  if  $i == k$            // the pivot value is the answer
6      return  $A[q]$ 
7  elseif  $i < k$ 
8      return RANDOMIZED-SELECT( $A, p, q - 1, i$ )
9  else return RANDOMIZED-SELECT( $A, q + 1, r, i - k$ )
```

What is the expected time complexity?