

INTRO AND INSERTION SORT

CS340

Your experience

- You have had
 - CS 234 Database and Web Development
 - CS 286 Computer Architecture
 - CS 314 Operating Systems
 - CS 325 Software Engineering
 - CS 321 Human-Computer Interaction
 - Math 224 Discrete Math
 - Calculus and other math?
 - Artificial Intelligence
 - Networking
 - Other?

Your goals

- Professions
 - Software Engineer
 - Application Developer
 - IT / Support
 - Database Administrator
 - Web Developer
 - Information Security Analyst
 - Researcher / Educator
 - Other?

Your skills and attributes

- I know Python or Java?
- I know more than 4 programming languages.
- I am a good Object-Oriented programmer.
- I know databases / database programming.
- I am a good communicator.
- I am a good technical writer.
- I like to test software.
- I am comfortable making presentations.
- I am good at math.
- I am comfortable writing a proof.
- I have a job and work more than 20 hours per week.
- I have had professional programming experience or an internship.

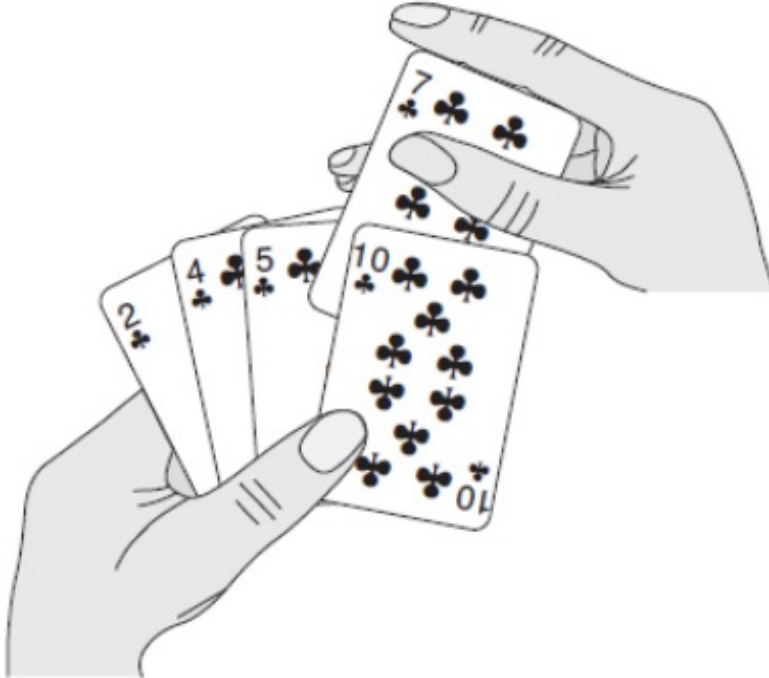
Algorithms

- An algorithm is any well-defined computational procedure that takes some value, or set of values, as **input** and produces some value, or set of values, as **output**.

Sorting: Insertion Sort

- Input
 - a sequence of n numbers $\langle a_1, a_2, a_3, \dots, a_n \rangle$
- Output
 - a reordering $\langle a_1', a_2', a_3', \dots, a_n' \rangle$ of the input sequence such that $a_1' \leq a_2' \leq a_3' \leq \dots \leq a_n'$
 - Often the numbers are stored in an array, A

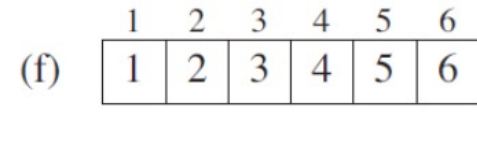
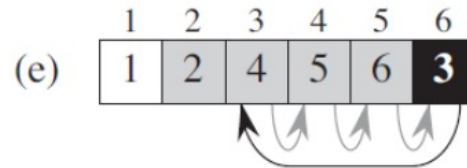
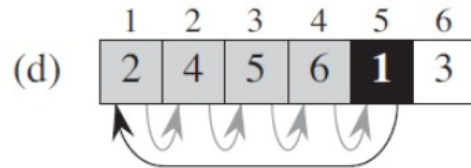
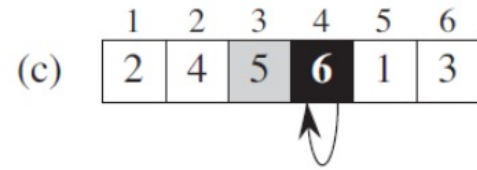
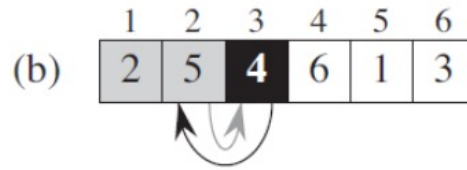
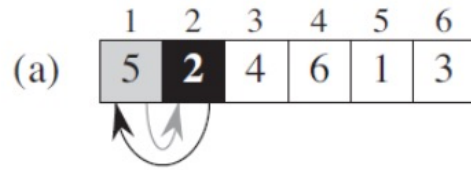
Insertion Sort



Insertion Sort

- The idea is to insert the current element into a sub-array that is assumed to be sorted.
- What we need:
 - A loop L1 going through the elements $A[i]$
 - Within L1 we need to perform insertion of
 - the current element $A[i]$ within already sorted sub-array $A[1], \dots, A[i-1]$

Insertion Sort



- (c) no elements are moved
 (d) all elements are moved

Insertion Sort

- If $A[i]$ is already in the “correct position”
 - No extra “insertion” needed
 - What it means: $A[i] > A[i-1]$
 - Namely, when $A[i]$ is the maximum of the subarray
- On the other hand, if $A[i]$ is the minimum of the subarray
 - Namely, when $A[i] < A[1]$
 - Then, $A[i]$ needs to be inserted into the first index.
 - **Entire subarray must be right-shifted.**

Insertion Sort

- The larger $A[i]$ is in comparison to subarray $A[1] \dots A[i-1]$
 - The smaller the amount of right-shifting.
- Whereas L1 goes left to right through the $A[i]$
 - There is an inner loop L2
 - used for insertion into subarray
 - which goes right to left

Insertion Sort

INSERTION-SORT(A)

```
1  for  $j = 2$  to  $A.length$ 
2       $key = A[j]$ 
3      // Insert  $A[j]$  into the sorted sequence  $A[1..j-1]$ .
4       $i = j - 1$ 
5      while  $i > 0$  and  $A[i] > key$ 
6           $A[i + 1] = A[i]$ 
7           $i = i - 1$ 
8       $A[i + 1] = key$ 
```

Proving correctness with a loop invariant

3 Things must be shown

1. **Initialization:** It is true prior to the first iteration of the loop.
2. **Maintenance:** If it is true before an iteration of the loop, it remains true before the next iteration.
3. **Termination:** When the loop terminates, the invariant—usually along with the reason that the loop terminated—gives us a useful property that helps show that the algorithm is correct.

Loop invariant for insertion sort

- Loop invariant: At the start of each iteration of the “outer” for loop—the loop indexed by j —the subarray $A[1.. j-1]$ consists of the elements originally in $A[1.. j-1]$ but in sorted order.
- In other words, $A[1..j-1]$ is a sorted subarray.
- Notice that this tells us about the algorithm and gives us a goal:
 - If the sorted subarray can consist of all items in the array, then the array will be sorted!

Correctness of insertion sort

- **Initialization**

- Just before the first iteration, $j = 2$. The subarray $A[1.. j-1]$ is the single element $A[1]$, which is **trivially** sorted.

- **Maintenance:**

- The body of the inner **while** loop puts *key* (which has the value that started out in $A[j]$) into the correct sorted position. The subarray $A[1..j]$ then consists of the elements originally in $A[1..j]$, but in sorted order.

- **Termination:**

- The outer **for** loop ends when $j > n$, which occurs when $j=n+1$. Therefore, $j-1=n$. Plugging n in for $j-1$ in the loop invariant, the subarray $A[1..n]$ consists of the elements originally in $A[1..n]$ but in sorted order. In other words, the entire array is sorted.

Interview Questions

- Write a proof of correctness for LINEAR-SEARCH, which scans through a sequence, looking for value v . Use a loop invariant to prove that your algorithm is correct. (Make sure that your loop invariant fulfills the three necessary properties – initialization, maintenance, termination.)

14	3	7	2	27
A[1]	A[2]	A[3]	A[4]	A[5]

```
LINEAR-SEARCH(A, v)
for i = 1 to A.length
    if A[i] == v
        return i
return NIL
```