

Optimization of snow plowing operations in Skinnarila area, Lappeenranta

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- 1 Problem description
- 2 Graph Theory and the Chinese Postman Problem (CPP)
- 3 Dealing with the real problem
- 4 Heuristic Alternative
- 5 Conclusions
- 6 Results

Motivation



What is the shortest snow plowing route in Lappeenranta?



An instance of an old problem...

Chinese Postman Problem

To find the shortest closed path that visits every edge of a connected undirected graph.



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Chinese Postman Problem

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- In our problem, **nodes** are intersections and **edges** are **road-segments**.
- We want to find the route with the shortest total distance.



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The Lappeenranta road network is NOT Eulerian!



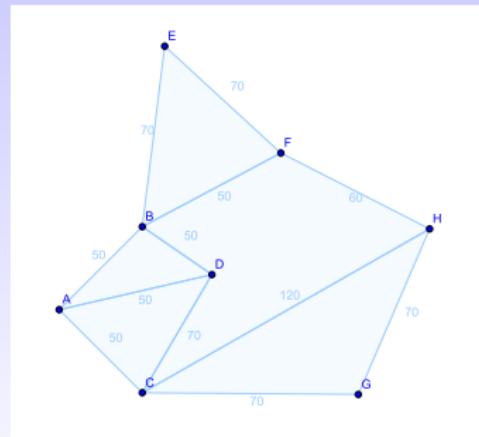
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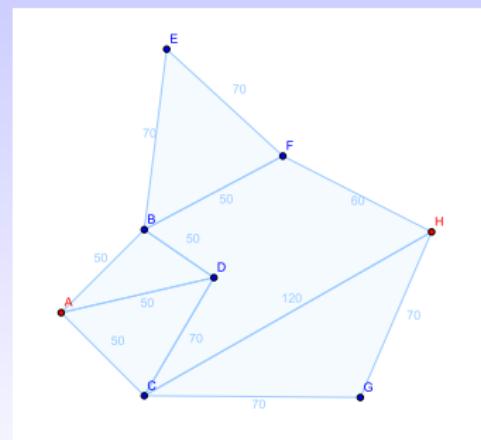
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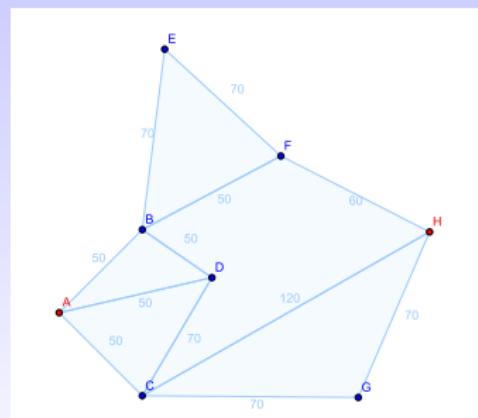
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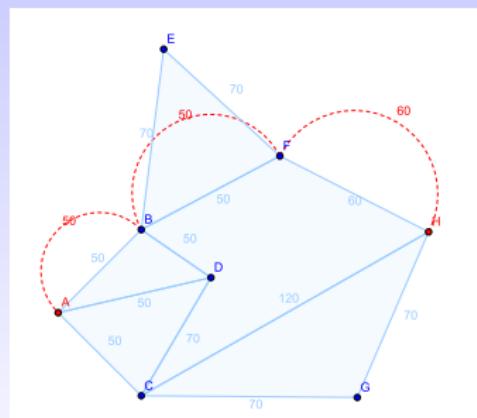
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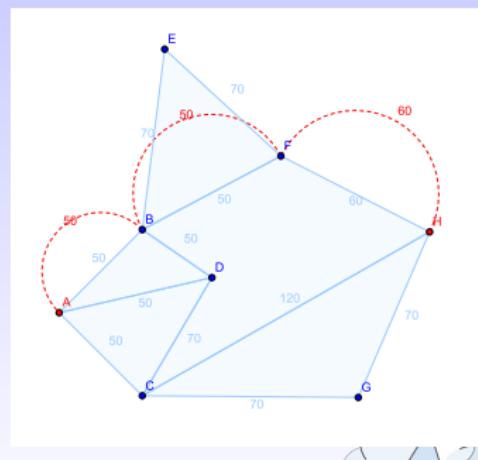
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- 5 Add those edges to the original graph G' .



Python Implementation

How to find the edges that must be added to form G' while minimizing the extra distance?

Example: for two pairs of odd nodes A, B, C and D there are 3 possible combinations:

- 1 AB + CD
- 2 AC + BD
- 3 BC + AD

Python Implementation

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Example: for two pairs of odd nodes A, B, C and D there are 3 possible combinations:

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But for big networks the amount of combinations increases following a factorial rule. For N nodes:

$$\text{Combinations} = \binom{N}{2} \binom{N-2}{2} \cdots \binom{4}{2} = N!!$$



Python Implementation: Blossom Algorithm

In 1973 Edmund and Johnson proposed the Blossom algorithm which solves this problem in polynomial time. For a map with 20 odd nodes:

$$\text{Combinations} = 20! \approx 3,715,891,200$$



Python Implementation: Blossom Algorithm

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$$\text{Combinations} = 20! \cdot 1! = 3,715,891,200$$

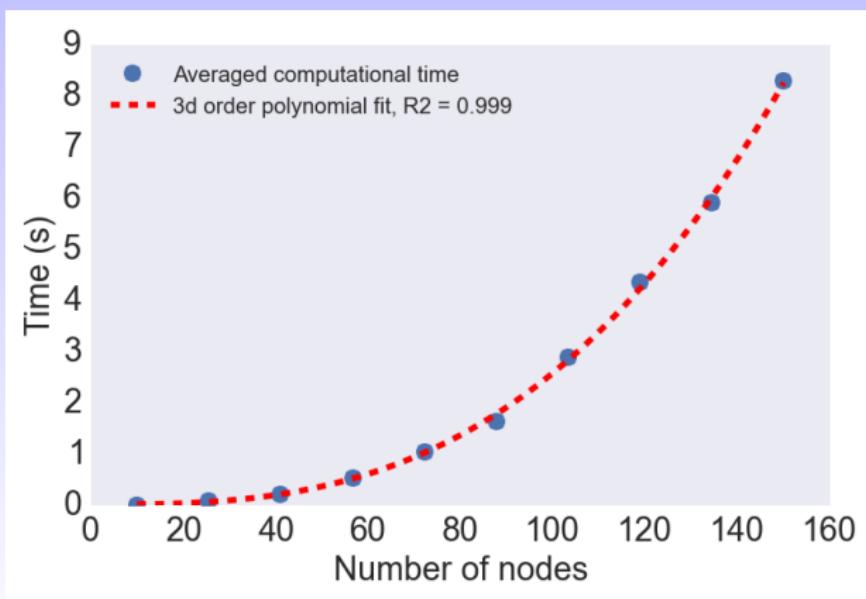
On a normal computer we get the following results:

- Searching through all possibilities: possibly days of computational time.
- Blossom algorithm: 0.3 seconds.



Python Implementation: Scaling

How does the computational time scale with graph size?



Python Implementation

- Pros
 - Scales well with size.
 - Optimality guarantee in finite time.
 - Can be used for graphs with one-way and two-way streets.
- Cons
 - Not flexible. Hard to add priorities, multiple machines, etc.



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LIUKKAUDEN TORJUNTA

Yleensä käsitellään kerrallaan koko pituudeltaan liukkauksen esiintyessä. Väylät hiekoitetaan pääsääntöisesti traktorivetoisilla hiekoittimilla. Hiekoitus suontetaan hiekoittimen levyisenä kaistana ja pyritään tekemään välyn keskelle.

TYÖN SUORITUSTAPA

Raitit aurataan joko heti täyteen leveyteen tai liikennöitävään leveyteen ja levitetään myöhempinä.

Työt arkitseen aloitetaan tarvittaessa jo klo 3.00 - klo 7.00. Viikonloppuisin panostetaan tarpeen mukaan pääliikenneräiltien kunnossapitoon.

A lot of Finnish words!

AURAUSURAKKA
LÄNSIALUE
Katulista

Lappeenrannan kaupunki
Tekninen toimi, Kadut ja ympäröivät alueet
25.5.2016

KADUT	pituus m
I LK	58 574

I LK	2 168
LAVOLANTIE	1 277
MUNTERONKATU	198
RUOHALAMMENKATU	560
SKINNARILANKATU	133

KADUT	pituus m
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II LK	10 059
HIETAKALLIONKATU	164
HOPEAMÄENRAITTI	166
JUPITERINKATU	201
JÄRVINIITUNKATU	296
KAARTINKATU	110
KALLIOPELLONKATU	264
KARJAKATU	209
KIERTOKATU	116

Dealing with the real problem

We were provided with the following data:

- 1 A Description of the task (in Finnish;)
- 2 A map of Lappeenranta;
- 3 A list of street names and priorities (which streets should be plowed first, second etc.)
- 4 We also used Google maps;)

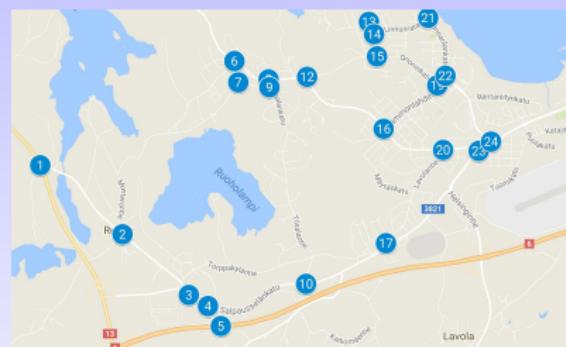
- We focused on streets with the highest priority.



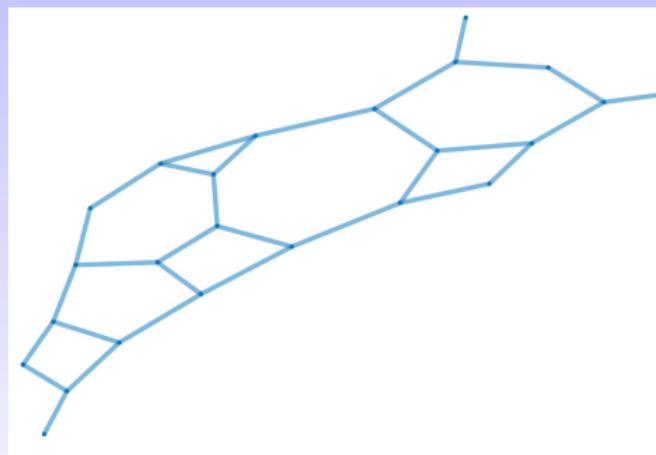
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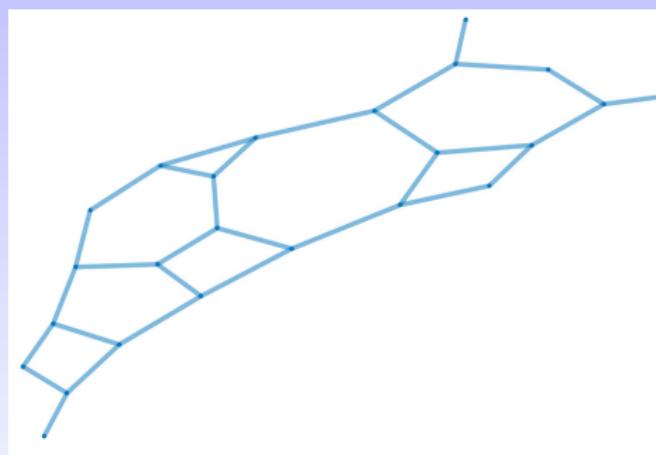
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- The edges are road segments.
- The distances are the lengths of the segments.

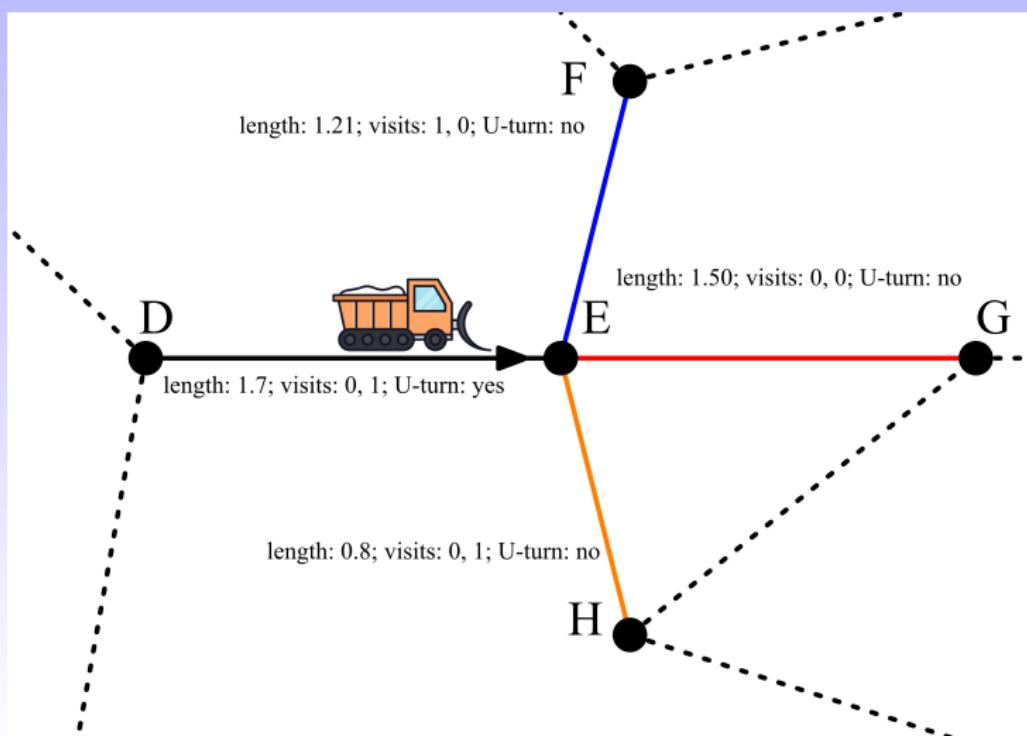


Penalty Scout Algorithm

- Heuristic methods for finding a path in an undirected connected graph.
- The idea is to create a walk along the graph and calculate the fitness.
- It uses a penalty term whenever a possible path choice has to be made.
- It allows us to compare different generated paths.

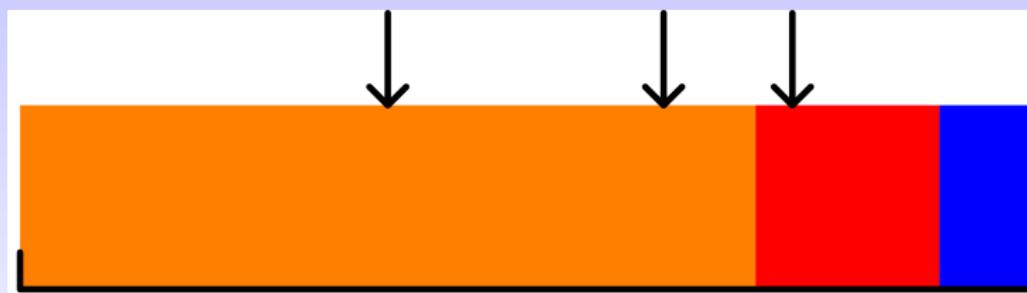


A pictorial example



A pictorial example (cont'd)

- Sort these points by penalty (the lower the penalty, the better)



Probability for each edge to be chosen

Results: Penalty Scout Algorithm

- One-way graph (24.2355 km total length)
 - One vehicle
 - 30.5275 km found within 1e6 attempts (30 minutes)
 - 35.0465 km found within 1e3 attempts (2 seconds)
 - Two vehicles
 - 41.1874 km found within 1e5 attempts (3 minutes)
 - Three vehicles
 - 58.8465 km found within 1e5 attempts (3 minutes)



Results: Penalty Scout Algorithm

- Two-way graph (48.471 km total length)
 - One vehicle
 - 53.005 km found within 1e6 attempts (60 minutes)
 - 53.627 km found within 1e4 attempts (40 seconds)
 - Two vehicles
 - 60.245 km found within 1e5 attempts (6 minutes)
 - Three vehicles
 - 89.137 km found within 1e5 attempts (9 minutes)



Solution using Penalty Scout Algorithm

■ Pros

- Fast implementation
- Terminating
- Easy to set priorities of different parameters
- Simple fitness calculation (length/time)
- Fast expandable
- Multiple vehicles considerable

■ Cons

- Decreasing compilation speed for more complicated graphs
(exponential, but still acceptable)
- No optimality guarantee (but usually performs well)



Results: Exact Solution to the CPP

The animation says it best...



