

NJIT

New Jersey's Science &
Technology University

THE EDGE IN KNOWLEDGE

CS 280
Programming Language Concepts

Languages – Why and How

How, when and why did programming languages evolve?

- All processors support a very small set of simple instructions
- Instructions are represented as bit patterns that control the operation of the machine
- VERY early on, the notion of a stored program (as compared to wired in or fed in from cards or magnetic tape) developed
- VERY early on, simple mnemonic languages were developed to make programming a little bit easier

Assembler Language

- Machine code has a series of possible operations and operands, represented by unique bit patterns
- Some mnemonic representation of the operations helps make things readable
- Symbolic names for resources like processor registers helps a lot.
- Symbolic names for memory locations helps even more

Assembler example

Count the number of 1 bits in a byte at memory location 0x2000

```

                LXI H 2000H; ;HL points at location 2000H
                MOVA A, M; ;Loads A with the content of M
                MVI C, 08H; ;Sets up counter for number of bits
                MVI B, 00H; ;Sets counter to count number of ones
jump2:          RAL; ;Rotates accumulator left through carry
                JNC jump1; ;On no carry jumps to jump1:
                INR B; ;Increases counter B by one
jump1:          DCR C; ;Decreases counter C by one
                JNZ jump2; ;When C is not zero jumps to jump2:
                HLT; ;Terminates the program

```

Assembler example

```

; -----
; Writes "Hello, World" to the console using only system calls. Runs on 64-bit Linux only.
; To assemble and run:
;
;     nasm -felf64 hello.asm && ld hello.o && ./a.out
; -----

                global _start

_start:         section .text
                mov     rax, 1 ; system call for write
                mov     rdi, 1 ; file handle 1 is stdout
                mov     rsi, message ; address of string to output
                mov     rdx, 13 ; number of bytes
                syscall ; invoke operating system to do the write
                mov     rax, 60 ; system call for exit
                xor     rdi, rdi ; exit code 0
                syscall ; invoke operating system to exit

                section .data
message:        db      "Hello, World", 10 ; note the newline at the end

```

Higher Levels of Language

- A higher level of abstraction
- Easier to read and write
- Easier to develop systems for more complex uses

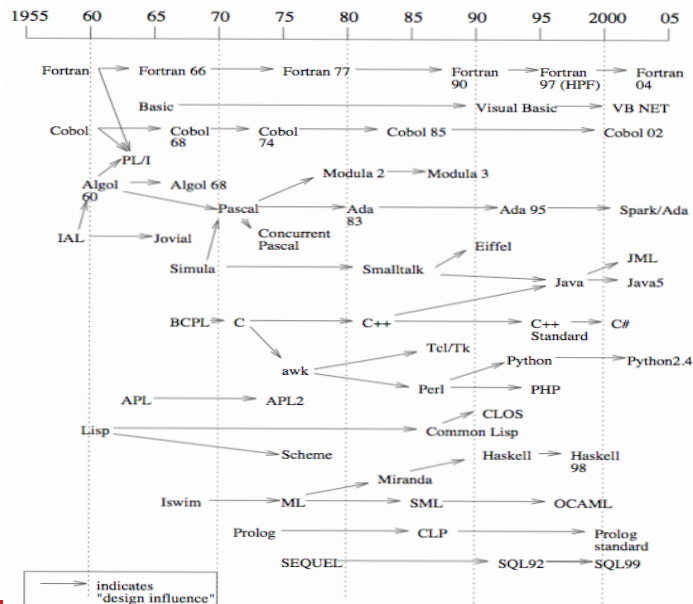
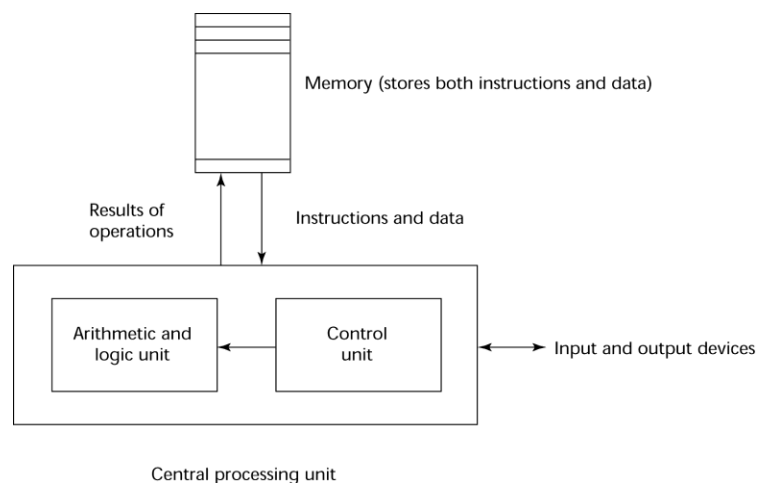


Figure 1.2: A Snapshot of Programming Language History

What Is Common

- All Languages will have
 - a syntax – what is and is not valid in the language
 - semantics – what is the meaning of elements of the language
 - names – what are the rules for names for things in the language
 - types – what values can things in the language take on
- All Languages run on some sort of machine

The von Neumann Architecture



The von Neumann Architecture

- Fetch–execute–cycle (on a von Neumann architecture computer)

```
initialize the program counter
repeat forever
    fetch the instruction pointed by the counter
    increment the counter
    decode the instruction
    execute the instruction
end repeat
```

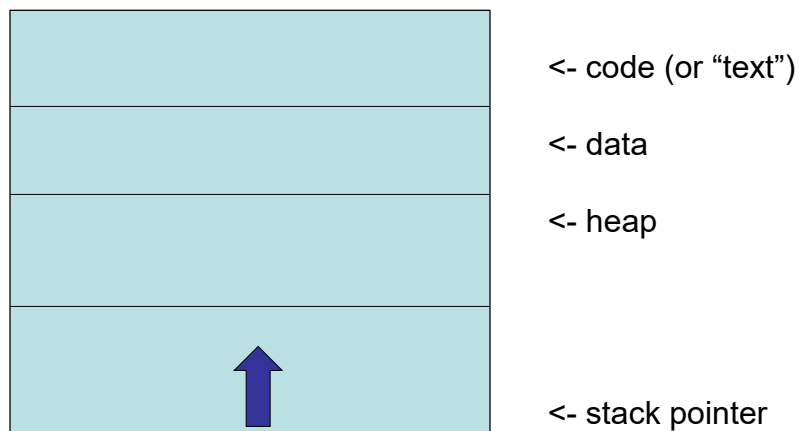
Von Neumann Bottleneck

- Connection speed between a computer's memory and its processor determines the speed of a computer
- Program instructions often can be executed much faster than the speed of the connection; the connection speed thus results in a *bottleneck*
- Known as the *von Neumann bottleneck*; it is the primary limiting factor in the speed of computers

Memory Usage

- Programs use memory in several ways:
 - the actual executable machine instructions
 - global variables
 - constants (such as strings)
 - memory that is used to support function calls
 - this is usually done in a “stack”, at runtime
 - memory that is dynamically allocated and freed as needed
 - this is usually done in a “heap”, at runtime
 - the “new” operator is one way to dynamically allocate memory

Memory Layout



Computer Architecture Influences Languages

- Imperative languages are most dominant, because of von Neumann computers
 - Data and programs stored in memory
 - Memory is separate from CPU
 - Instructions and data are piped from memory to CPU
 - This is the basis for imperative languages
 - Variables model memory cells
 - Assignment statements model piping
 - Iteration is efficient
 - A language that supports changing memory and conditional branching is said to be “Turing Complete”

Programming Methodologies over Time

- 1950s and early 1960s: Simple applications; worry about machine efficiency
- Late 1960s: People efficiency became important; readability, better control structures
 - structured programming
 - top-down design and step-wise refinement
- Late 1970s: Process-oriented to data-oriented
 - data abstraction
- Middle 1980s: Object-oriented programming
 - Data abstraction + inheritance + polymorphism

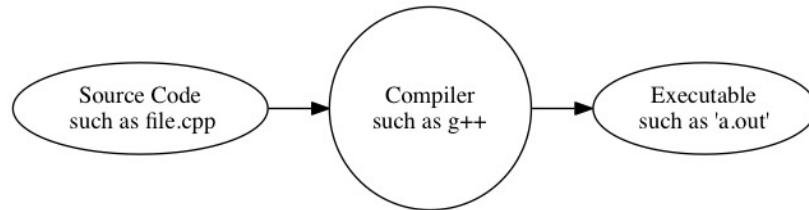
Language Design Trade-Offs

- **Reliability vs. cost of execution**
 - Example: Java demands all references to array elements be checked for proper indexing, which leads to increased execution costs
- **Readability vs. writeability**
 - Example: APL provides many powerful operators (and a large number of new symbols), allowing complex computations to be written in a compact program but at the cost of poor readability
- **Writeability (flexibility) vs. reliability**
 - Example: C++ pointers are powerful and very flexible but can be unreliable

Implementation Methods

- **Compilation**
 - Programs are translated into machine language; includes JIT systems
 - Use: Large commercial applications
- **Pure Interpretation**
 - Programs are interpreted by another program known as an interpreter
 - Use: Small programs or when efficiency is not an issue
- **Hybrid Implementation Systems**
 - A compromise between compilers and pure interpreters
 - Use: Small and medium systems when efficiency is not the first concern

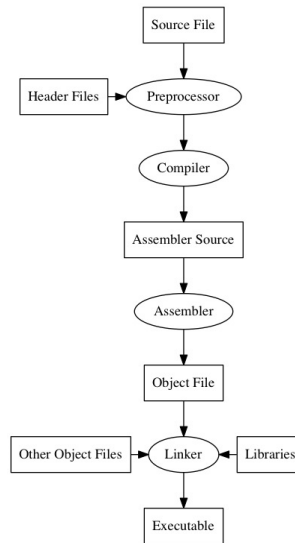
Compiling



Compilation

- Translate high-level program (source language) into machine code (machine language)
- This is specific to the machine you are compiling for
- Compiling on one machine to run on another is called “cross-compilation”
- Slow translation, fast execution

Steps In Compile



Preprocessors

- Preprocessor macros (instructions) are commonly used to specify that code from another file is to be included
- A preprocessor processes a program immediately before the program is compiled to expand embedded preprocessor macros
- A well-known example: C preprocessor
 - expands `#include`, `#define`, and similar macros

Object Files

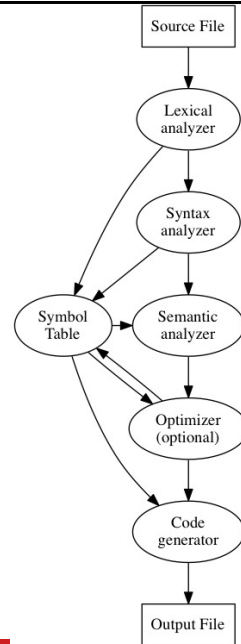
- Each Object File contains:
 - executable code
 - global variables
 - constants
 - “linkage information”
 - what code and data do I provide that others may use
 - what code and data do I require that I need others to provide for me
- A library is just a collection of object files (sometimes an “archive” or a “.a” or a “library” or a “.lib”)

Linking

- The “linker” connects all the object files together to make a single executable
- Every required connection must be resolved or the link fails
- Some environments use a “shared library” scheme
 - On windows this is a .dll, a Dynamically Linked Library
 - On linux it's sometimes a .so, a Shared Object
- With shared libraries, some of the linking happens at run time

Inside a Compiler

- lexical analyzer converts characters in the source program into lexical units
- syntax analyzer transforms lexical units into *parse trees* which represent the syntactic structure of program
- Semantics analyzer enforces the semantic rules of the language
- Optimizer improves the code
- Code generator creates the output (assembler or byte code)



Pure Interpretation

- No translation
- Easier implementation of programs (run-time errors can easily and immediately be displayed)
- Slower execution (10 to 100 times slower than compiled programs)
- Often requires more space
- Now rare for traditional high-level languages
- Significant comeback with some Web scripting languages (e.g., JavaScript, PHP)

Hybrid Implementation Systems

- A compromise between compilers and pure interpreters
- A high-level language program is translated to an intermediate language that allows easy interpretation
- Faster than pure interpretation
- Examples
 - Perl programs are partially compiled to detect errors before interpretation
 - Initial implementations of Java were hybrid; the intermediate form, *byte code*, provides portability to any machine that has a byte code interpreter and a run-time system (together, these are called *Java Virtual Machine*)

Just-in-Time Implementation Systems

- Initially translate programs to an intermediate language
- Then compile the intermediate language of the subprograms into machine code when they are called
- Machine code version is kept for subsequent calls
- JIT systems are widely used for Java programs
- .NET languages are implemented with a JIT system
- In essence, JIT systems are delayed compilers

Programming Environments

- The compiler/interpreter comes with standard libraries and a collection of other tools used in software development. This is sometimes called the “toolchain”
- Example tools:
 - Linkers
 - Archivers (to make libraries)
 - Symbol readers
 - Debuggers
- Debuggers allow for breakpointing and single stepping through your program, examining the value of variables and monitoring the flow of the program