CS 3205 COMPUTER NETWORKS

JAN-MAY 2020

LECTURES 14: 9TH MAR 2020

Text book and section(s) covered in this lecture: Book Kurose and Ross – Sections 2.5

DNS: domain name system

people: many identifiers:

SSN, name, passport #

Internet hosts, routers:

- IP address (32 bit) used for addressing datagrams
- "name", e.g., www.yahoo.com used by humans
- Q: how to map between IP address and name, and vice versa?

Domain Name System:

- distributed database implemented in hierarchy of many name servers
- application-layer protocol: hosts, name servers communicate to resolve names (address/name translation)
 - note: core Internet function, implemented as applicationlayer protocol
 - complexity at network's "edge"

DNS: services, structure

DNS services

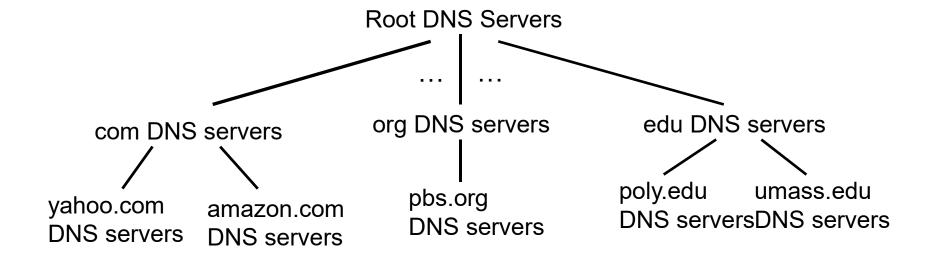
- hostname to IP address translation
- host aliasing
 - canonical, alias names
- mail server aliasing
- load distribution
 - replicated Web servers: many IP addresses correspond to one name

why not centralize DNS?

- single point of failure
- traffic volume
- distant centralized database
- maintenance

A: doesn't scale!

DNS: a distributed, hierarchical database

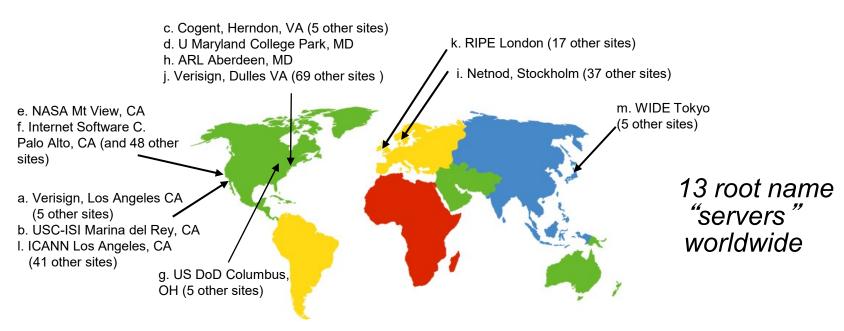


client wants IP for www.amazon.com; Ist approx:

- client queries root server to find com DNS server
- client queries .com DNS server to get amazon.com DNS server
- client queries amazon.com DNS server to get IP address for www.amazon.com

DNS: root name servers

- contacted by local name server that can not resolve name
- root name server:
 - contacts authoritative name server if name mapping not known
 - gets mapping
 - returns mapping to local name server



TLD, authoritative servers

top-level domain (TLD) servers:

- responsible for com, org, net, edu, aero, jobs, museums, and all top-level country domains, e.g.: uk, fr, ca, jp
- Network Solutions maintains servers for .com TLD
- Educause for .edu TLD

authoritative DNS servers:

- organization's own DNS server(s), providing authoritative hostname to IP mappings for organization's named hosts
- can be maintained by organization or service provider

Local DNS name server

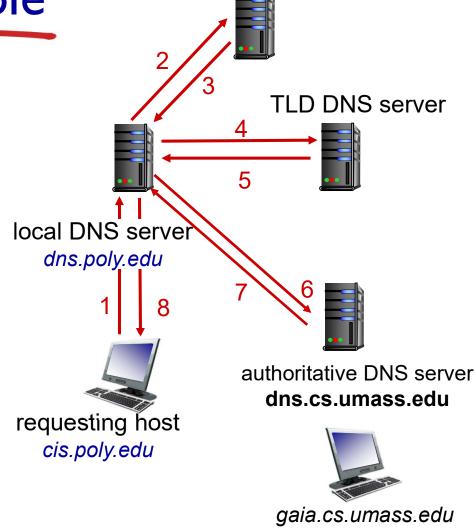
- does not strictly belong to hierarchy
- each ISP (residential ISP, company, university) has one
 - also called "default name server"
- when host makes DNS query, query is sent to its local DNS server
 - has local cache of recent name-to-address translation pairs (but may be out of date!)
 - acts as proxy, forwards query into hierarchy

DNS name resolution example

 host at cis.poly.edu wants IP address for gaia.cs.umass.edu

iterated query:

- contacted server replies with name of server to contact
- "I don't know this name, but ask this server"

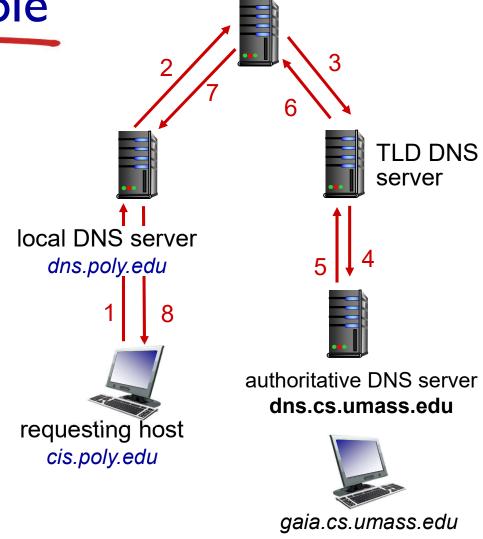


root DNS server

DNS name resolution example

recursive query:

- puts burden of name resolution on contacted name server
- heavy load at upper levels of hierarchy?



root DNS server

DNS: caching, updating records

- once (any) name server learns mapping, it caches mapping
 - cache entries timeout (disappear) after some time (TTL)
 - TLD servers typically cached in local name servers
 - thus root name servers not often visited
- cached entries may be out-of-date (best effort name-to-address translation!)
 - if name host changes IP address, may not be known Internet-wide until all TTLs expire
- update/notify mechanisms proposed IETF standard
 - RFC 2136

DNS records

DNS: distributed db storing resource records (RR)

RR format: (name, value, type, ttl)

type=A

- name is hostname
- value is IP address

type=NS

- name is domain (e.g., foo.com)
- value is hostname of authoritative name server for this domain

type=CNAME

- name is alias name for some "canonical" (the real) name
- www.ibm.com is really servereast.backup2.ibm.com
- value is canonical name

type=MX

 value is name of mailserver associated with name

DNS protocol, messages

* query and reply messages, both with same message format

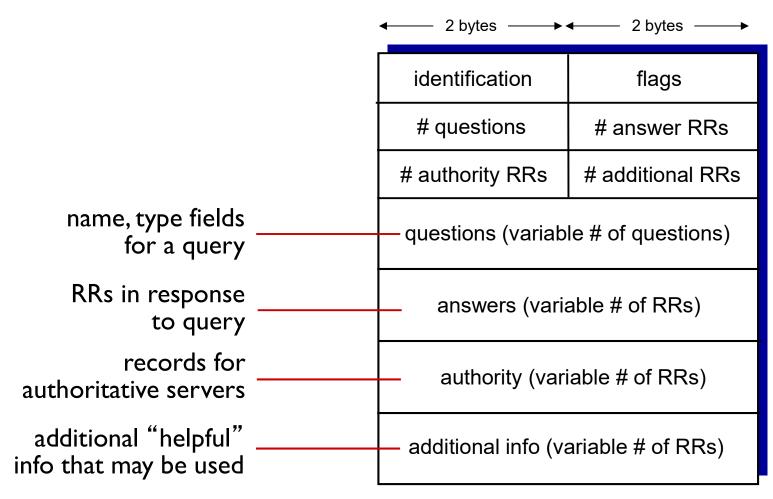
msg header

- identification: 16 bit # for query, reply to query uses same #
- flags:
 - query or reply
 - recursion desired
 - recursion available
 - reply is authoritative

| 2 Bytes | 2 5 7 6 5 |
|-------------------------------------|------------------|
| identification | flags |
| # questions | # answer RRs |
| # authority RRs | # additional RRs |
| questions (variable # of questions) | |
| answers (variable # of RRs) | |
| authority (variable # of RRs) | |
| additional info (variable # of RRs) | |

— 2 bytes → ◆ 2 bytes → →

DNS protocol, messages



Inserting records into DNS

- example: new startup "Network Utopia"
- register name networkuptopia.com at DNS registrar (e.g., Network Solutions)
 - provide names, IP addresses of authoritative name server (primary and secondary)
 - registrar inserts two RRs into .com TLD server: (networkutopia.com, dns1.networkutopia.com, NS) (dns1.networkutopia.com, 212.212.212.1, A)
- create authoritative server type A record for www.networkuptopia.com; type MX record for networkutopia.com

Attacking DNS

DDoS attacks

- Bombard root servers with traffic
 - Not successful to date
 - Traffic Filtering
 - Local DNS servers cache IPs of TLD servers, allowing root server bypass
- Bombard TLD servers
 - Potentially more dangerous

Redirect attacks

- Man-in-middle
 - Intercept queries
- DNS poisoning
 - Send bogus relies to DNS server, which caches

Exploit DNS for DDoS

- Send queries with spoofed source address: target IP
- Requires amplification