

# CS 3205 COMPUTER NETWORKS

**JAN-MAY 2020**

**LECTURES 22: 16<sup>TH</sup> APR 2020**

Text book and section(s) covered in this lecture:  
Book Kurose and Ross – Sections 4.4

# Chapter 4: outline

## 4.1 introduction

## 4.2 virtual circuit and datagram networks

## 4.3 what's inside a router

## 4.4 IP: Internet Protocol

- datagram format
- IPv4 addressing
- ICMP
- IPv6

## 4.5 routing algorithms

- link state
- distance vector
- hierarchical routing

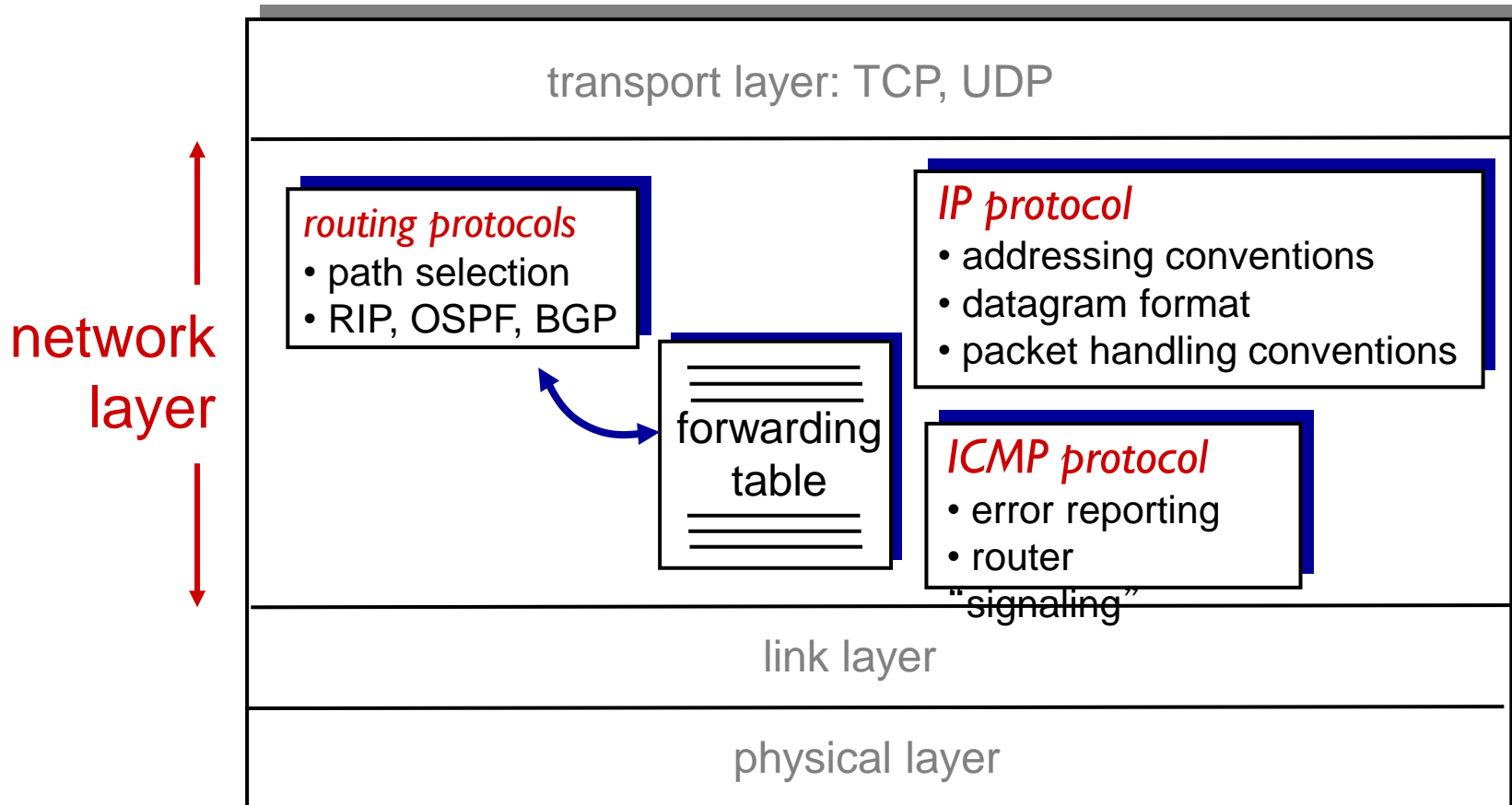
## 4.6 routing in the Internet

- RIP
- OSPF
- BGP

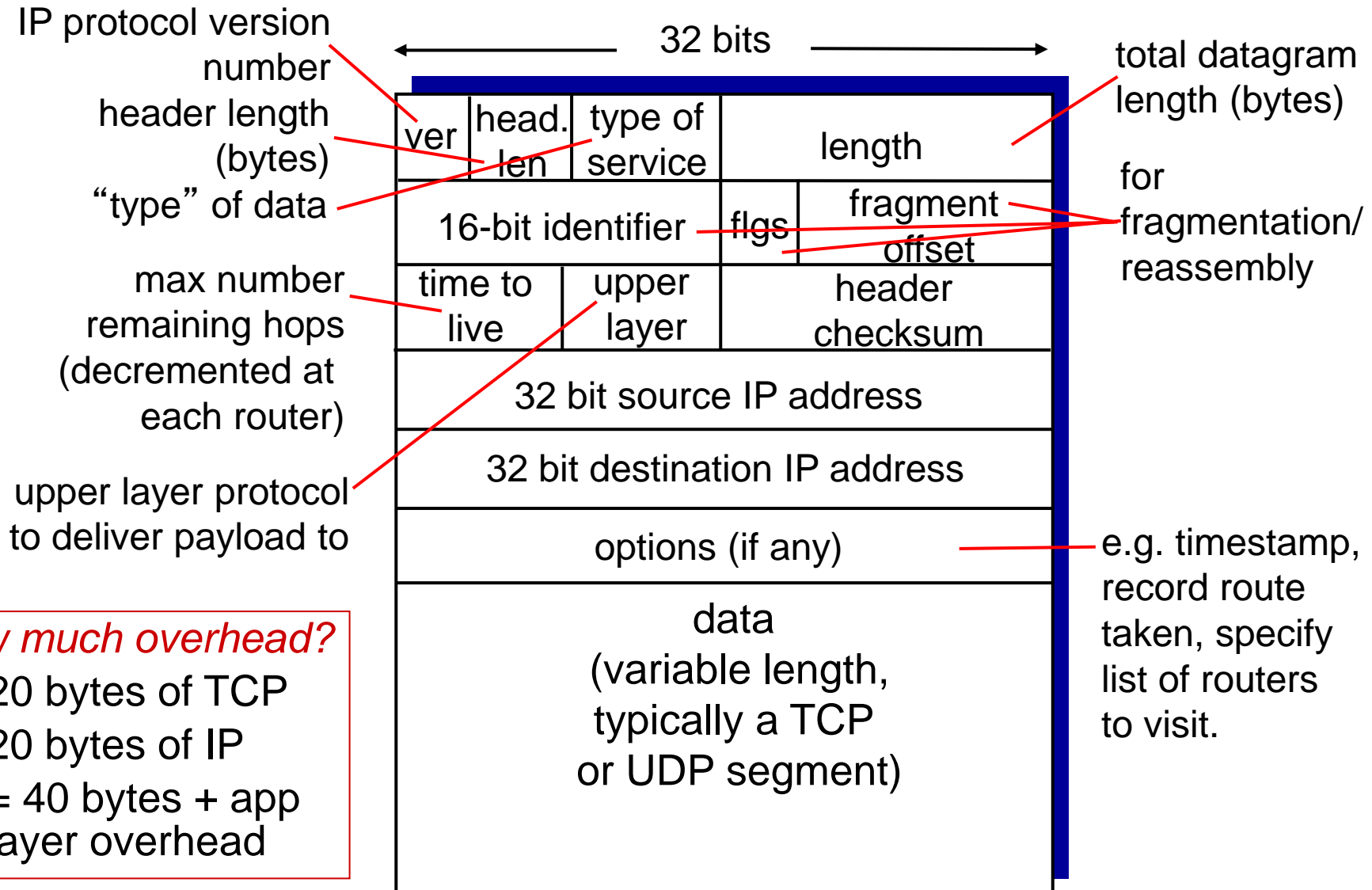
## 4.7 broadcast and multicast routing

# The Internet network layer

host, router network layer functions:

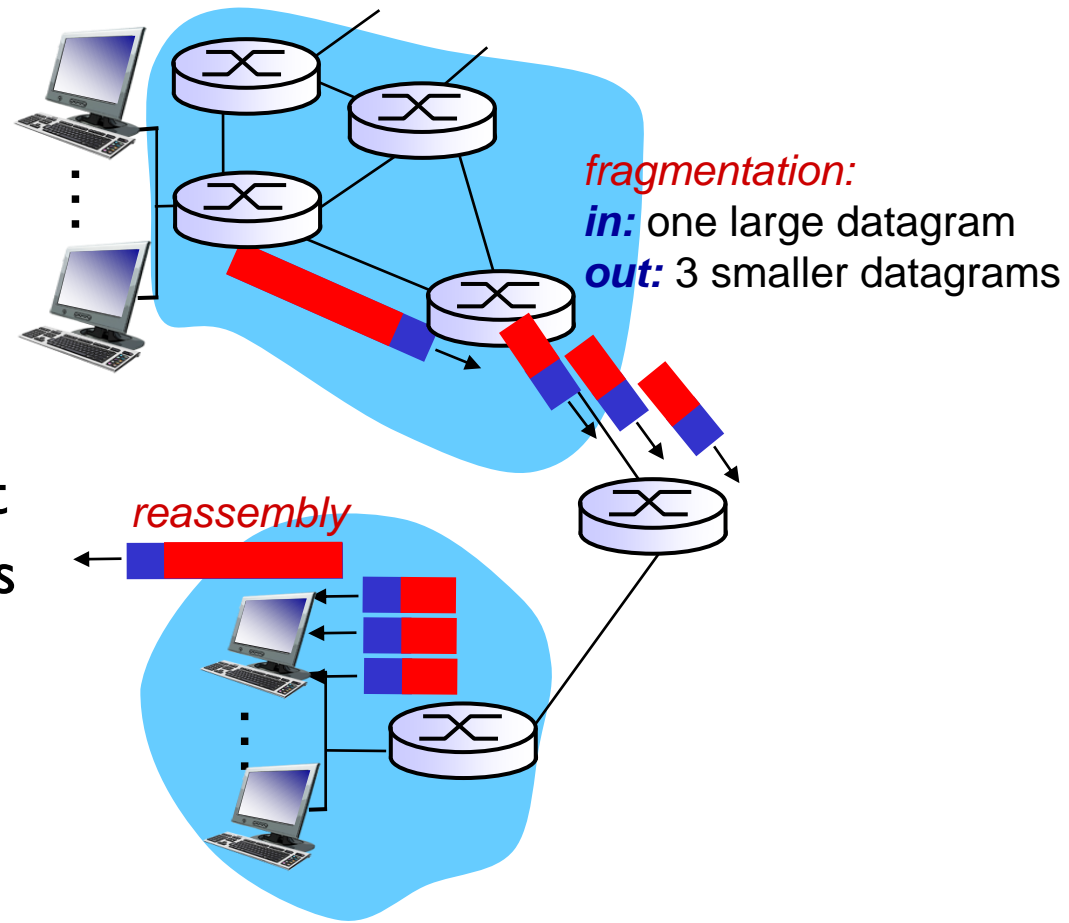


# IP datagram format



# IP fragmentation, reassembly

- ❖ network links have MTU (max.transfer size) - largest possible link-level frame
  - different link types, different MTUs
- ❖ large IP datagram divided (“fragmented”) within net
  - one datagram becomes several datagrams
  - “reassembled” only at final destination
  - IP header bits used to identify, order related fragments



# IP fragmentation, reassembly

## *example:*

- ❖ 4000 byte datagram
- ❖ MTU = 1500 bytes

	length	ID	fragflag	offset	
	=4000	=x	=0	=0	

*one large datagram becomes  
several smaller datagrams*

1480 bytes in  
data field

offset =  
 $1480/8$

	length	ID	fragflag	offset	
	=1500	=x	=1	=0	

	length	ID	fragflag	offset	
	=1500	=x	=1	=185	

	length	ID	fragflag	offset	
	=1040	=x	=0	=370	

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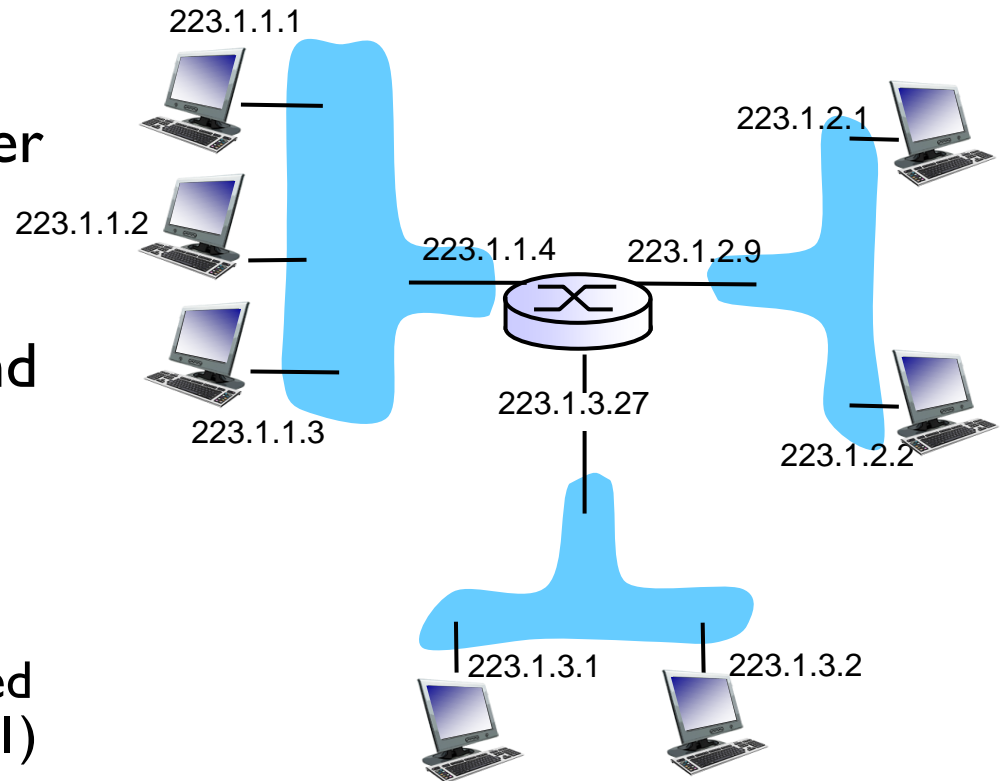
## 4.6 routing in the Internet

- RIP
- OSPF
- BGP

## 4.7 broadcast and multicast routing

# IP addressing: introduction

- ❖ **IP address:** 32-bit identifier for host, router interface
- ❖ **interface:** connection between host/router and physical link
  - router's typically have multiple interfaces
  - host typically has one or two interfaces (e.g., wired Ethernet, wireless 802.11)
- ❖ **IP addresses associated with each interface**



$$223.1.1.1 = \underbrace{11011111}_{223} \underbrace{00000001}_1 \underbrace{00000001}_1 \underbrace{00000001}_1$$



# IP addressing: introduction

CLASS	LEADING BITS	NET ID BITS	HOST ID BITS	NO. OF NETWORKS	ADDRESSES PER NETWORK	START ADDRESS	END ADDRESS
CLASS A	0	8	24	$2^7$ (128)	$2^{24}$ (16,777,216)	0.0.0.0	127.255.255.255
CLASS B	10	16	16	$2^{14}$ (16,384)	$2^{16}$ (65,536)	128.0.0.0	191.255.255.255
CLASS C	110	24	8	$2^{21}$ (2,097,152)	$2^8$ (256)	192.0.0.0	223.255.255.255
CLASS D	1110	NOT DEFINED	NOT DEFINED	NOT DEFINED	NOT DEFINED	224.0.0.0	239.255.255.255
CLASS E	1111	NOT DEFINED	NOT DEFINED	NOT DEFINED	NOT DEFINED	240.0.0.0	255.255.255.255

Class	Address Range	Supports
<b>Class A</b>	1.0.0.1 to 126.255.255.254	Supports 16 million hosts on each of 127 networks.
<b>Class B</b>	128.1.0.1 to 191.255.255.254	Supports 65,000 hosts on each of 16,000 networks.
<b>Class C</b>	192.0.1.1 to 223.255.254.254	Supports 254 hosts on each of 2 million networks.
<b>Class D</b>	224.0.0.0 to 239.255.255.255	Reserved for multicast groups.
<b>Class E</b>	240.0.0.0 to 254.255.255.254	Reserved for future use, or Research and Development Purposes.

32 positions, denoted as bunch of 8 positions  
 = abcdefgh abcdefgh abcdefgh abcdefgh  
       byte 1      byte 2      byte 3      byte 4

If "a" of byte 1 is "0" then Class A  
 If "ab" of byte 1 is "10" then Class B  
 If "abc" of byte 1 is "110" then Class C  
 If "abcd" of byte 1 is "1110" then Class D  
 If "abcd" of byte 1 is "1111" then Class E

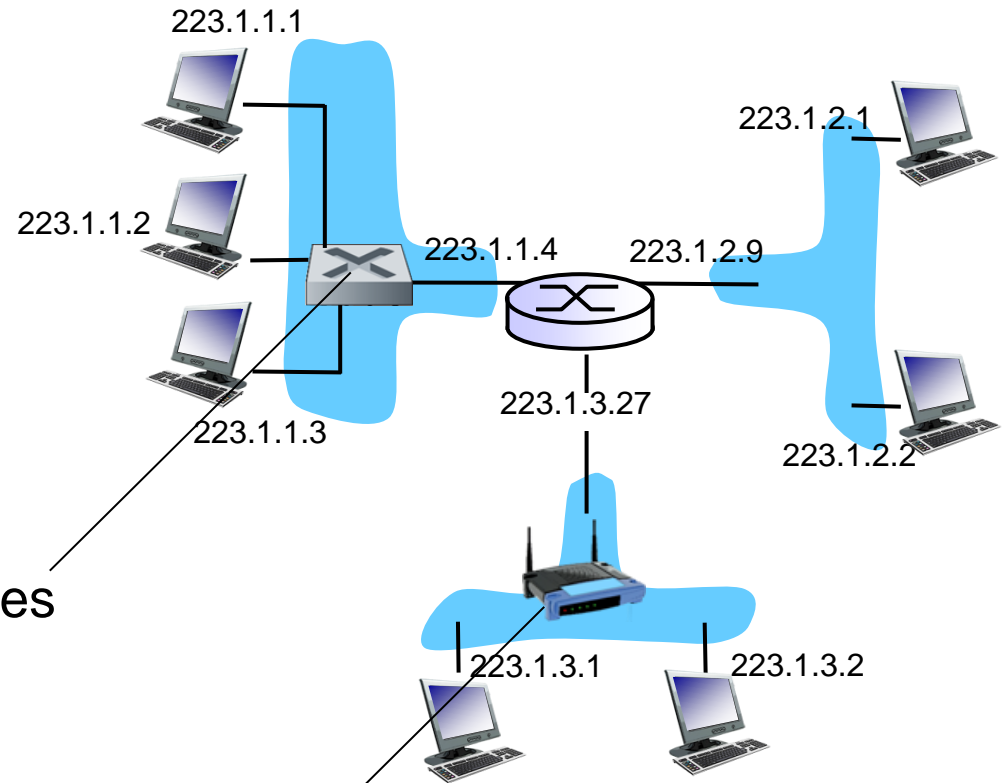
# IP addressing: introduction

*Q: how are interfaces actually connected?*

*A: we'll learn about that in chapter 5, 6.*

*A:* wired Ethernet interfaces connected by Ethernet switches

*For now:* don't need to worry about how one interface is connected to another (with no intervening router)



*A:* wireless WiFi interfaces connected by WiFi base station

To the PDF – with Details

# Subnets

**Address Range Net-0 (192.168.10.0/25)**

**Network Address**

192.168.10.00000000 = 192.168.10.0

**First Usable IP Address for Host**

192.168.10.00000001 = 192.168.10.1

**Last Usable IP Address for Host**

192.168.10.01111110 = 192.168.10.126

**Broadcast IP Address**

192.168.10.01111111 = 192.168.10.127

# Subnets

## Classful versus Classless IP Networks

IP Class	Range	Default Subnet Mask	Number of Networks	Number of Hosts per Network
A	1 – 126	255.0.0.0	126	16,777,214
B	128 – 191	255.255.0.0	16,384	65,534
C	192 – 223	255.255.255.0	2,097,152	254

**Classful**

- 255's identify the network portion
- 0's identify the host portion

**Classless**

IP Address	Subnet Mask	CIDR Notation
10.252.0.101	255.255.0.0	10.252.0.101/16

# Subnets

## Classful vs Classless Addressing

- Classful:
  - ◆ Size defined by the class (A, B, C, D, E).
  - ◆ Fixed network portion.
  - ◆ RIP & IGRP are classful routing protocols.
- Classless:
  - ◆ Network portion can be any size.
  - ◆ Protocol sends subnetting (prefix) information with routes.
    - ☞ 192.168.64.0/18
  - ◆ RIP2, EIGRP, OSPF, BGP & IS-IS.

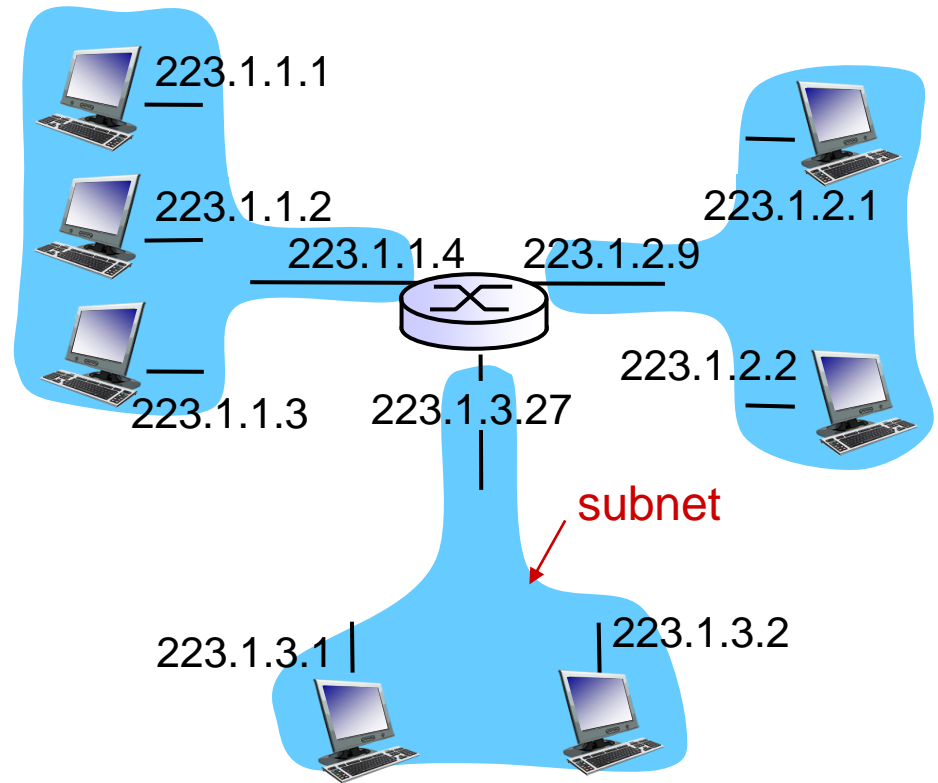
# Subnets

## ❖ IP address:

- subnet part - high order bits
- host part - low order bits

## ❖ *what 's a subnet ?*

- device interfaces with same subnet part of IP address
- can physically reach each other *without intervening router*

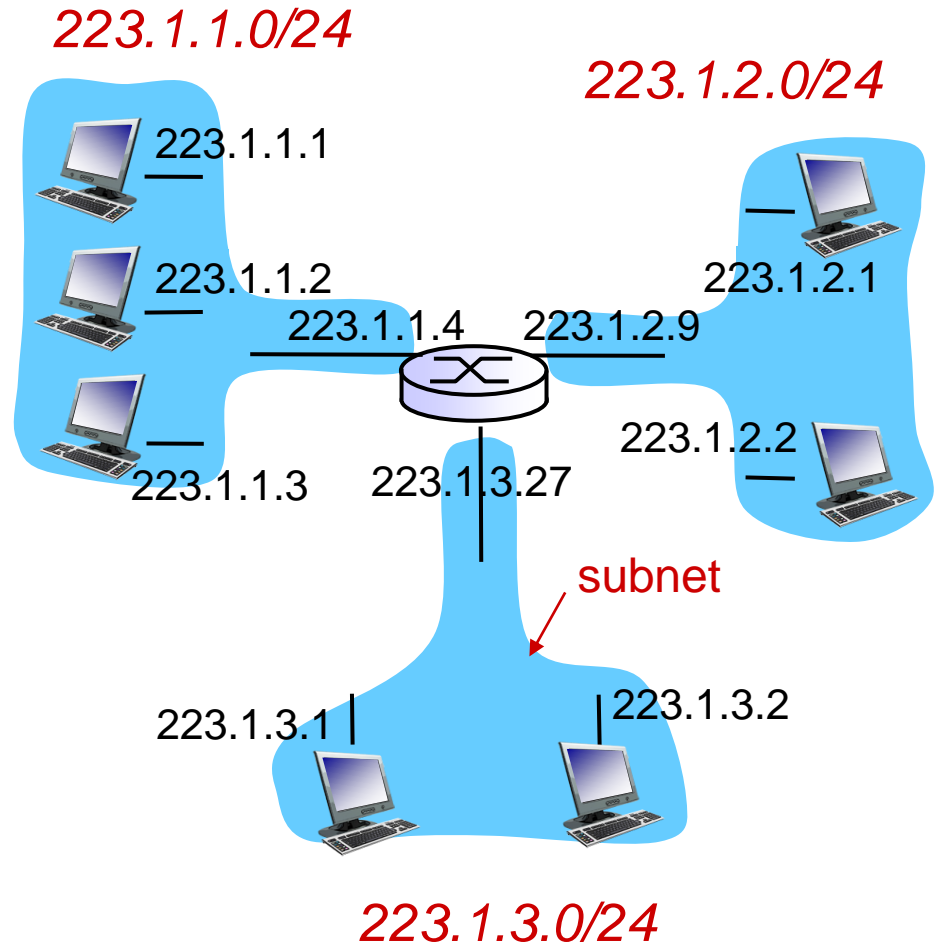


network consisting of 3 subnets

# Subnets

## *recipe*

- ❖ to determine the subnets, detach each interface from its host or router, creating islands of isolated networks
- ❖ each isolated network is called a *subnet*

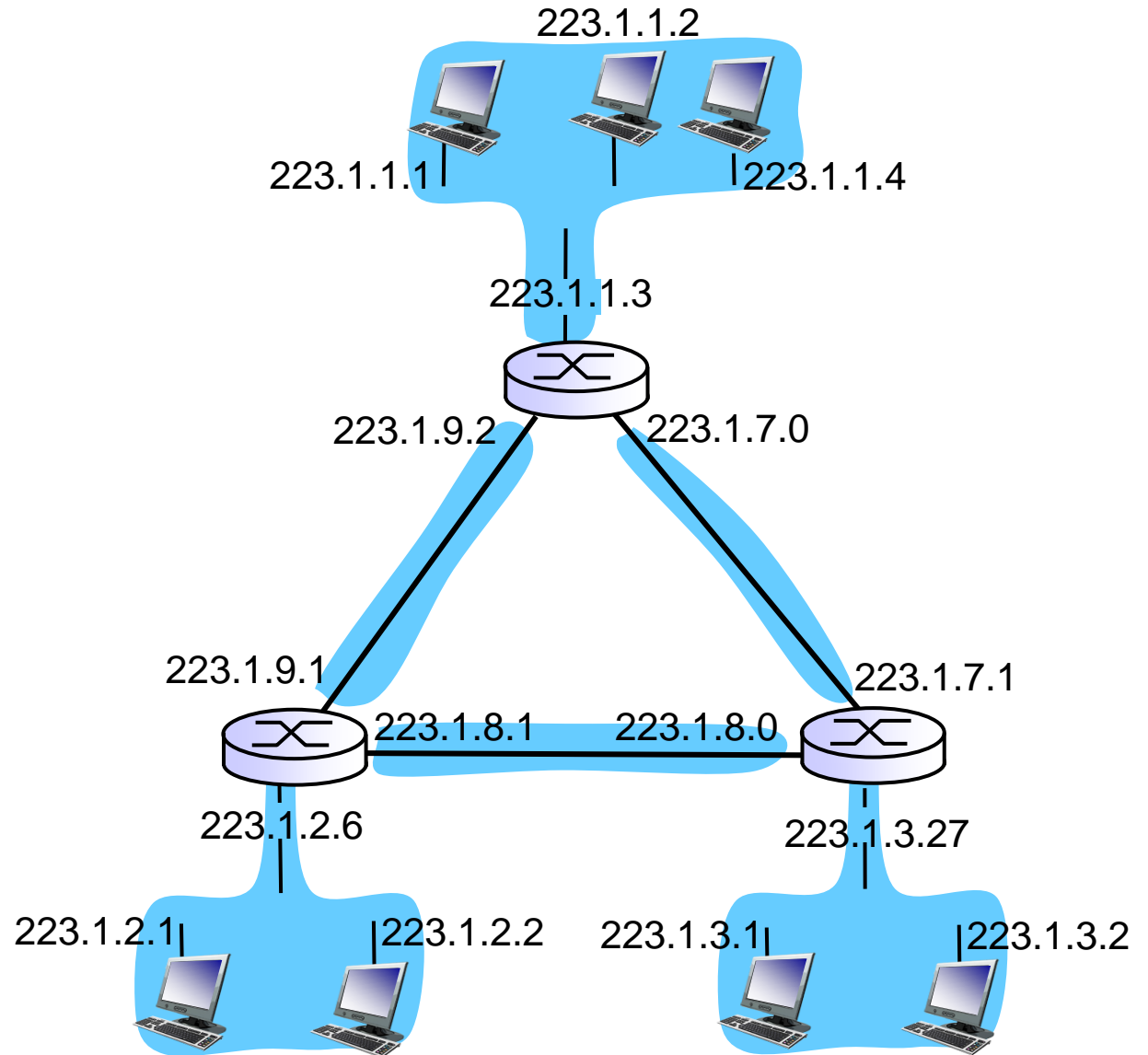


subnet mask: /24



# Subnets

how many?



# PRIVATE IP ADDRESS

(are not used anywhere on public internet, reserved for private LANs)

Network Class	Network Numbers	Network Mask	No. of Networks	No. of Hosts per Network
CLASS A	10.0.0.0	255.0.0.0	126	16,646,144
CLASS B	172.16.0.0 to 172.31.0.0	255.255.0.0	16,383	65,024
CLASS C	192.168.0.0 to 192.168.255.255	255.255.255.0	2,097,151	254
LOOPBACK (localhost)	127.0.0.0 to 127.0.0.7	255.255.255.0	-	-

# All in one place for quick ref

## Public IP Address Classes range

Class	1st Octet DEC range	1st Octet BIN	Start address	Finish address	1st Octet High order Bits	Network/ Host	Default Subnet Mask
<b>A</b>	1-126	<b>0</b> 0000001-01111110	0.0.0.0	126.255.255.255	<b>0</b>	<b>N.H.H.H</b>	255.0.0.0
<b>B</b>	128-191	<b>10</b> 000000-10111111	128.0.0.0	191.255.255.255	<b>10</b>	<b>N.N.H.H</b>	255.255.0.0
<b>C</b>	192-223	<b>110</b> 00000-11011111	192.0.0.0	223.255.255.255	<b>110</b>	<b>N.N.N.H</b>	255.255.255.0
<b>D</b>	224-239	<b>1110</b> 0000-11101111	224.0.0.0	239.255.255.255	<b>1110</b>		
<b>E</b>	240-255	<b>11110</b> 000-11111111	240.0.0.0	254.255.255.255	<b>11110</b>		

Note: Class A address **127.0.0.0 - 127.255.255.255** cannot be used and is for **LOOPBACK** and diagnostic

## Private IP Address Classes range

Class	1st Octet DEC range	1st Octet BIN	Start address	Finish address	1st Octet High order Bits	Network/ Host	Default Subnet Mask
<b>A</b>	10	<b>0</b> 0001010	10.0.0.0	10.255.255.255	<b>0</b>	<b>N.H.H.H</b>	255.0.0.0
<b>B</b>	172	<b>10</b> 101100	172.16.0.0	172.31.255.255	<b>10</b>	<b>N.N.H.H</b>	255.255.0.0
<b>C</b>	192	<b>110</b> 00000	192.168.0.0	192.168.255.255	<b>110</b>	<b>N.N.N.H</b>	255.255.255.0

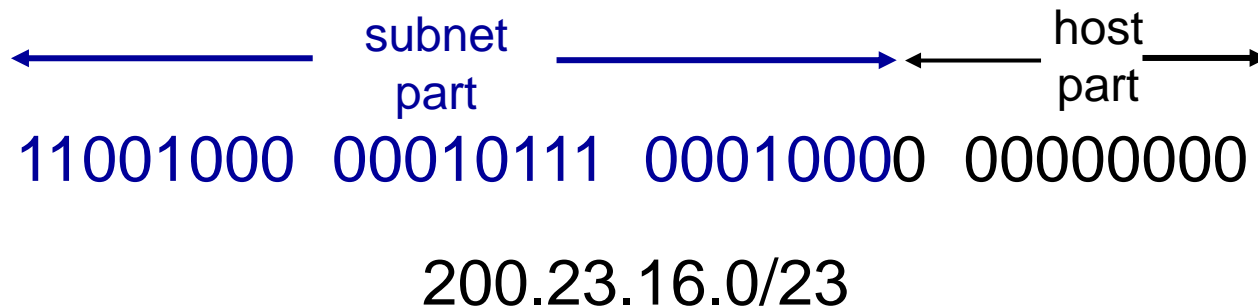
[www.ic.ims.hr](http://www.ic.ims.hr)

Number of unique IP Addresses in a network? =  $2^{(\text{Number of host bits})} - 2$

# IP addressing: CIDR

## CIDR: Classless InterDomain Routing

- subnet portion of address of arbitrary length
- address format: **a.b.c.d/x**, where x is # bits in subnet portion of address



# IP addresses: how to get one?

**Q:** How does a *host* get IP address?

- ❖ hard-coded by system admin in a file
  - Windows: control-panel->network->configuration->tcp/ip->properties
  - UNIX: /etc/rc.config
- ❖ **DHCP: Dynamic Host Configuration Protocol:** dynamically get address from as server
  - “plug-and-play”

# DHCP: Dynamic Host Configuration Protocol

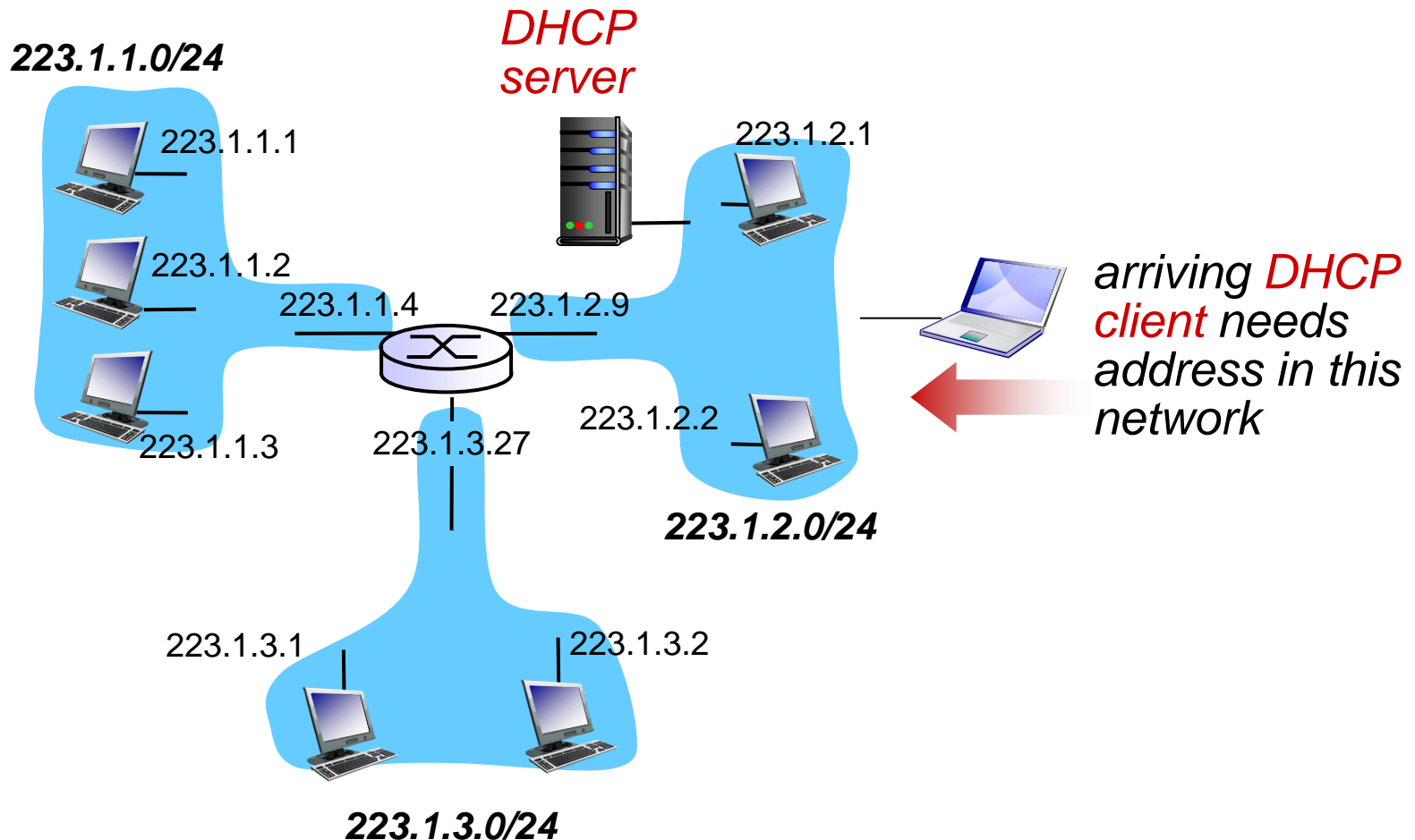
*goal:* allow host to *dynamically* obtain its IP address from network server when it joins network

- can renew its lease on address in use
- allows reuse of addresses (only hold address while connected/“on”)
- support for mobile users who want to join network (more shortly)

## *DHCP overview:*

- host broadcasts “**DHCP discover**” msg [optional]
- DHCP server responds with “**DHCP offer**” msg [optional]
- host requests IP address: “**DHCP request**” msg
- DHCP server sends address: “**DHCP ack**” msg

# DHCP client-server scenario

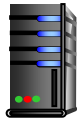


# DHCP client-server scenario

DHCP server: 223.1.2.5

DHCP discover

arriving  
client



Broadcast: is there a  
DHCP server out there?

DHCP offer

Broadcast: I'm a DHCP  
server! Here's an IP  
address you can use

DHCP request

Broadcast: OK. I'll take  
that IP address!

DHCP ACK

Broadcast: OK. You've  
got that IP address!

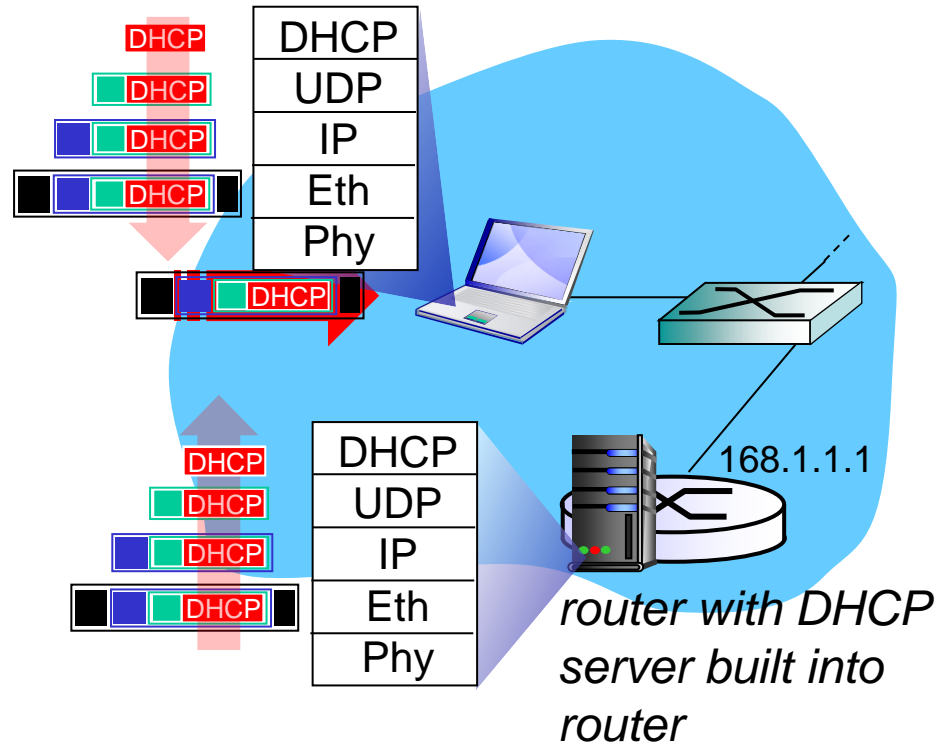


# DHCP: more than IP addresses

DHCP can return more than just allocated IP address on subnet:

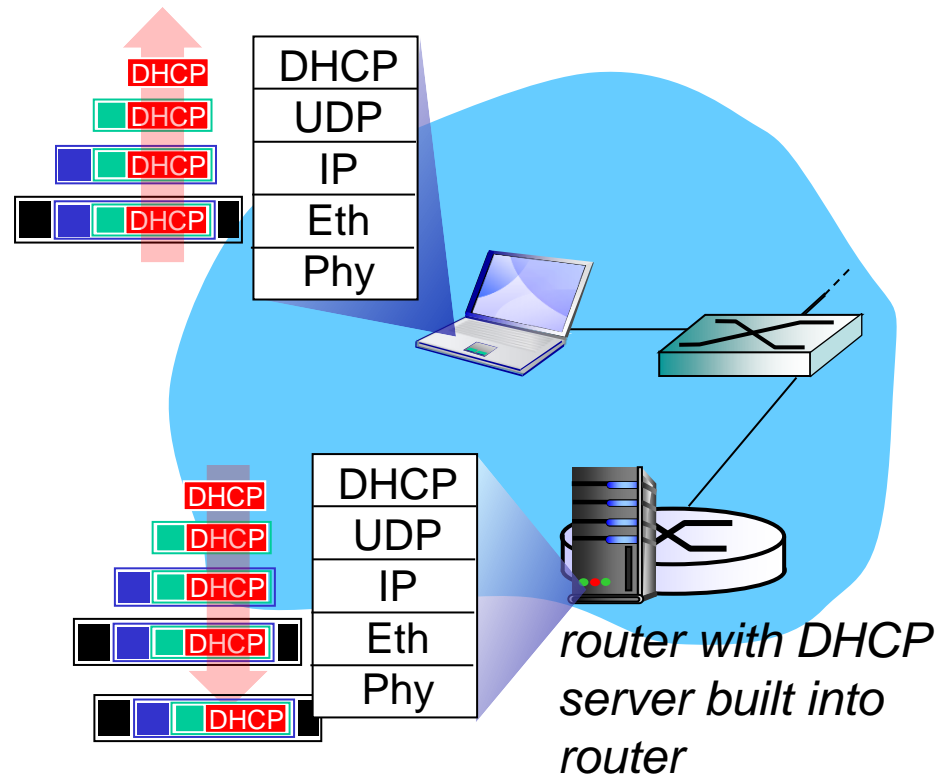
- address of first-hop router for client
- name and IP address of DNS sever
- network mask (indicating network versus host portion of address)

# DHCP: example



- ❖ connecting laptop needs its IP address, addr of first-hop router, addr of DNS server: use DHCP
- ❖ DHCP request encapsulated in UDP, encapsulated in IP, encapsulated in 802.1 Ethernet
- ❖ Ethernet frame broadcast (dest: FFFFFFFF) on LAN, received at router running DHCP server
- ❖ Ethernet demuxed to IP demuxed, UDP demuxed to DHCP

# DHCP: example



- ❖ DCP server formulates DHCP ACK containing client's IP address, IP address of first-hop router for client, name & IP address of DNS server
- ❖ encapsulation of DHCP server, frame forwarded to client, demuxing up to DHCP at client
- ❖ client now knows its IP address, name and IP address of DSN server, IP address of its first-hop router

# DHCP: Wireshark output (home LAN)

Message type: **Boot Request (1)**

Hardware type: Ethernet

Hardware address length: 6

Hops: 0

**Transaction ID: 0x6b3a11b7**

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

Client IP address: 0.0.0.0 (0.0.0.0)

Your (client) IP address: 0.0.0.0 (0.0.0.0)

Next server IP address: 0.0.0.0 (0.0.0.0)

Relay agent IP address: 0.0.0.0 (0.0.0.0)

**Client MAC address: Wistron\_23:68:8a (00:16:d3:23:68:8a)**

Server host name not given

Boot file name not given

Magic cookie: (OK)

Option: (t=53,l=1) **DHCP Message Type = DHCP Request**

Option: (61) Client identifier

Length: 7; Value: 010016D323688A;

Hardware type: Ethernet

Client MAC address: Wistron\_23:68:8a (00:16:d3:23:68:8a)

Option: (t=50,l=4) Requested IP Address = 192.168.1.101

Option: (t=12,l=5) Host Name = "nomad"

**Option: (55) Parameter Request List**

Length: 11; Value: 010F03062C2E2F1F21F92B

**1 = Subnet Mask; 15 = Domain Name**

**3 = Router; 6 = Domain Name Server**

44 = NetBIOS over TCP/IP Name Server

.....

request

Message type: **Boot Reply (2)**

Hardware type: Ethernet

Hardware address length: 6

Hops: 0

**Transaction ID: 0x6b3a11b7**

Seconds elapsed: 0

Bootp flags: 0x0000 (Unicast)

**Client IP address: 192.168.1.101 (192.168.1.101)**

Your (client) IP address: 0.0.0.0 (0.0.0.0)

**Next server IP address: 192.168.1.1 (192.168.1.1)**

Relay agent IP address: 0.0.0.0 (0.0.0.0)

Client MAC address: Wistron\_23:68:8a (00:16:d3:23:68:8a)

Server host name not given

Boot file name not given

Magic cookie: (OK)

**Option: (t=53,l=1) DHCP Message Type = DHCP ACK**

**Option: (t=54,l=4) Server Identifier = 192.168.1.1**

**Option: (t=1,l=4) Subnet Mask = 255.255.255.0**

**Option: (t=3,l=4) Router = 192.168.1.1**

**Option: (6) Domain Name Server**

Length: 12; Value: 445747E2445749F244574092;

IP Address: 68.87.71.226;

IP Address: 68.87.73.242;

IP Address: 68.87.64.146

**Option: (t=15,l=20) Domain Name = "hsd1.ma.comcast.net."**

reply

# IP addresses: how to get one?

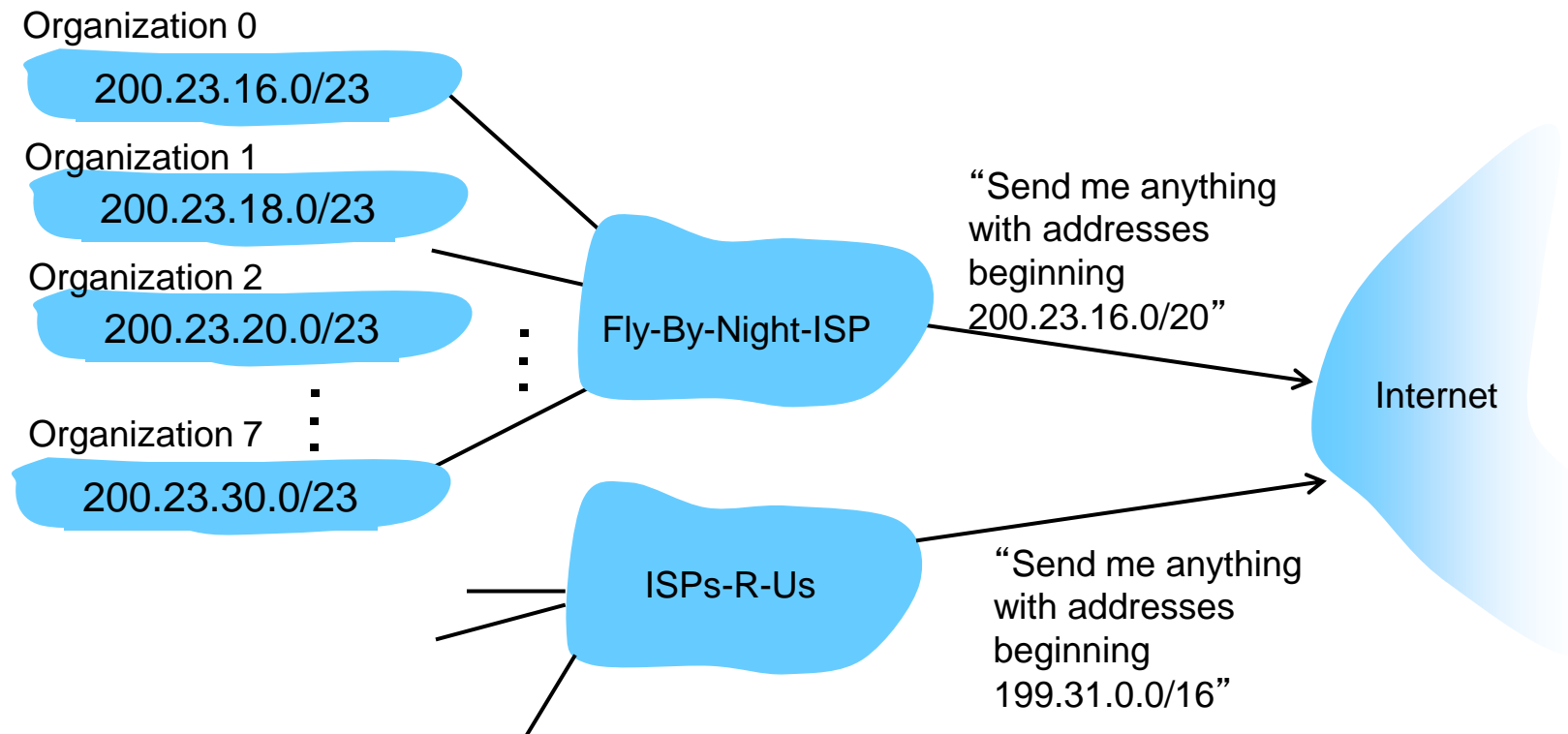
**Q:** how does *network* get subnet part of IP addr?

**A:** gets allocated portion of its provider ISP's address space

ISP's block	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/20
Organization 0	<u>11001000</u>	<u>00010111</u>	<u>00010000</u>	00000000	200.23.16.0/23
Organization 1	<u>11001000</u>	<u>00010111</u>	<u>00010010</u>	00000000	200.23.18.0/23
Organization 2	<u>11001000</u>	<u>00010111</u>	<u>00010100</u>	00000000	200.23.20.0/23
...	.....			....	....
Organization 7	<u>11001000</u>	<u>00010111</u>	<u>00011110</u>	00000000	200.23.30.0/23

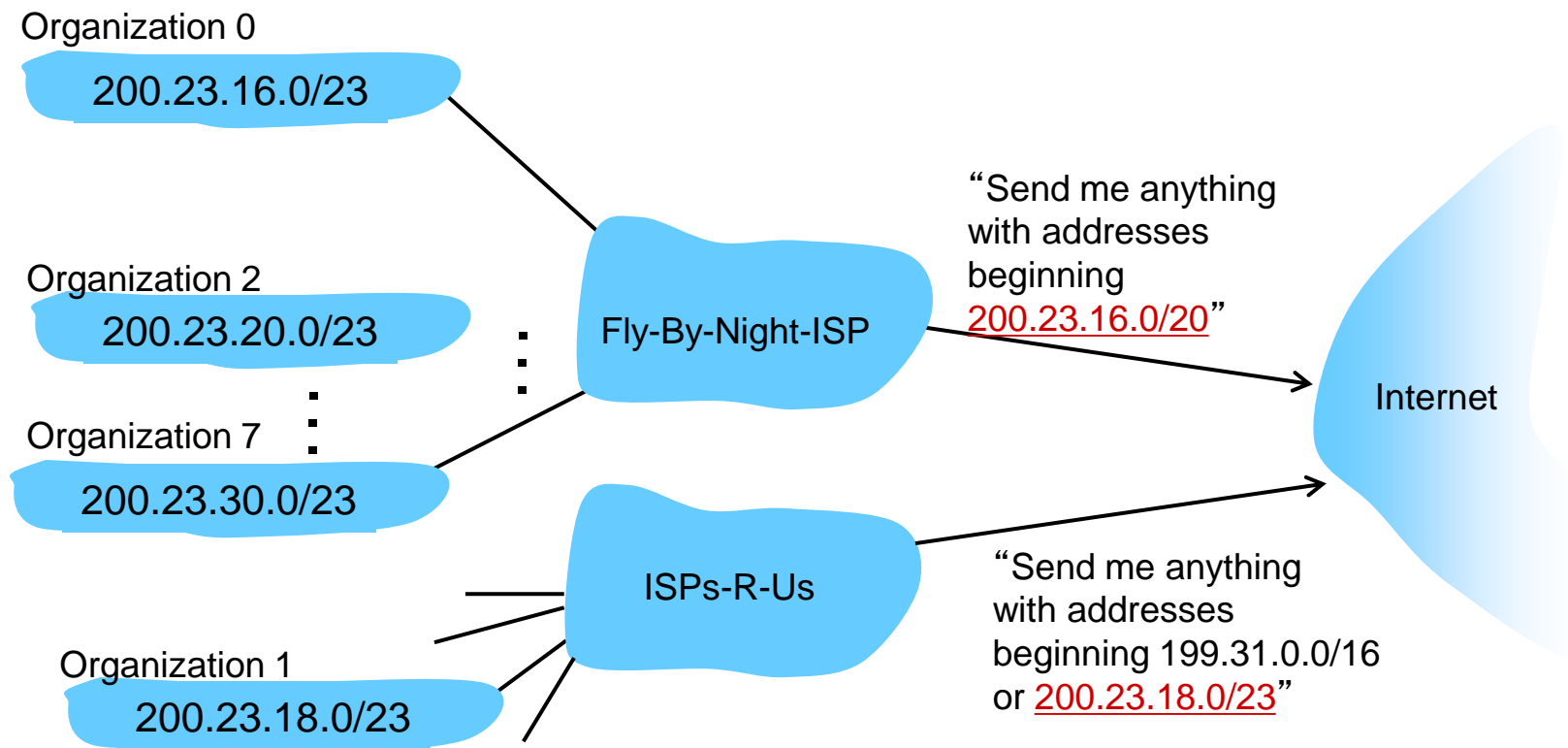
# Hierarchical addressing: route aggregation

hierarchical addressing allows efficient advertisement of routing information:



# Hierarchical addressing: more specific routes

ISPs-R-U has a more specific route to Organization 1



# IP addressing: the last word...

**Q:** how does an ISP get block of addresses?

**A: ICANN:** Internet Corporation for Assigned Names and Numbers <http://www.icann.org/>

- allocates addresses
- manages DNS
- assigns domain names, resolves disputes