

DAPA ASSIGNMENT

Generation of Combinations Lexicographically

Project done by GROUP 14:

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Problem Definition

Combinations refer to the combination of n things taken k at a time without repetition. To refer to combinations in which repetition is allowed, the terms k -selection, k -multiset, or k -combination with repetition are often used.

The combinations are generated in the groups of ascending nature in a lexicographical order, i.e. a general order of the way words are alphabetically ordered based on the alphabetical order of their component letters.

The corresponding algorithm is performed sequentially (normal execution) as well as parallelly (execution via threads and CUDA).

Assumptions

Several assumptions can be made via using the architecture used to make efficient processing such as, to enable a faster computation.

- The parallel algorithm uses an *EREW SM SIMD* model of computation.
- The time needed for thread creation is significant, which affects the time taken in parallel part of the algorithm, when the value of n is less.
- All the threads work independently with enough computations resources, so that there is never a case of deadlock, that can arise between them.

Algorithms

Sequential

procedure **sequentialCombination** (**n**, **m**)

Step 1:

(1.1) combinationArr \leftarrow [1, 2, 3, ..., m]

(1.2) produce combinationArr as output

Step 2:

(2.1) totalCombination \leftarrow nCr (n,m) - 1

(2.2) for i=1 to totalCombination do

 nextCombination (combinationArr,n,m)

endfor

procedure **nextCombination** (**combinationArr**, **n**, **m**)

Step 1: j \leftarrow m

Step 2: while (j>0) do

 if (combinationArr[j-1] < n-m+j) do

 (2.1) combinationArr[j-1] \leftarrow combinationArr[j-1] + 1

 (2.2) for i=j to m-1 do

 combinationArr[i] \leftarrow combinationArr[i-1] + 1

 (2.3) produce combinationArr as output

 (2.4) break

 else

 j \leftarrow j-1

 endif

endwhile

procedure **nCr** (**n**, **r**)

Step 1:

(1.1) p \leftarrow 1

(1.2) k \leftarrow 1

(1.2) if (n - r < r) do

 r \leftarrow n - r

endif

Step 2:

if (r != 0) do

 while (r) do

```

        (2.1)  $p \leftarrow p * n$ 
        (2.2)  $k \leftarrow k * r$ 
        (2.3)  $m \leftarrow \text{gcd}(p, k)$  // gcd - greatest common divisor
        (2.4)  $p \leftarrow p / m$ 
        (2.5)  $k \leftarrow k / m$ 
        (2.6)  $n \leftarrow n - 1$ 
        (2.7)  $r \leftarrow r - 1$ 
    endwhile
else
     $p \leftarrow 1$ 
endif
Step 3:
    return p

```

Parallel

procedure **parallelCombination** (n, m)

Step 1:

```

    (1.1)  $\text{MAX} \leftarrow 100$ 
    (1.2) declare integer variables k
    (1.2) declare integer arrays x[MAX], y[MAX], c[MAX]
    (1.3) declare boolean array z[MAX]

```

Step 2:

```

    (2.1) for i=1 to m do in parallel
        initialize(i)
    endfor
    (2.2) produce c as output

```

Step 3:

```

    while (z[0] == 0) do
        (3.1)  $k \leftarrow 0$ 
        (3.2) for i=2 to min(m,k) do in parallel
            task1(i)
        endfor
        (3.3) if (k==0) do
             $y[m-1] \leftarrow y[m-1] + 1$ 
        else
            for i=k to m do in parallel
                task2(i)
            endfor
        endif
    endwhile

```

```

(3.4)
for i=1 to m do in parallel
    task3(i)
endfor
(3.5) produce c as output
endwhile

```

```

procedure initialize (i)
    Step 1:  $x[i-1] \leftarrow n-m+i$ 
    Step 2:  $y[i-1] \leftarrow i$ 
    Step 3: if ( $y[i-1] == x[i-1]$ ) do
         $z[i-1] \leftarrow 1$ 
    else
         $z[i-1] \leftarrow 0$ 
    endif
    Step 4:  $c[i-1] \leftarrow i$ 

```

```

procedure task1 (i)
    if ( $z[i-2] == 0$  and  $z[i-1] == 1$ ) do
        (a)  $y[i-2] \leftarrow y[i-2] + 1$ 
        (b)  $k \leftarrow i$ 
    endif

```

```

procedure task2 (i)
     $y[i-1] \leftarrow y[k-2] + i - k + 1$ 

```

```

procedure task3 (i)
    Step 1:  $c[i-1] \leftarrow y[i-1]$ 
    Step 2:
        if ( $y[i-1] == x[i-1]$ ) do
             $z[i-1] \leftarrow 1$ 
        else
             $z[i-1] \leftarrow 0$ 
        endif

```

Implemented code (in C++)

Sequential code

```
#include <bits/stdc++.h>
using namespace std;

long long nCr(long long n, long long r)
{
    long long p = 1, k = 1;

    if (n - r < r)
        r = n - r;

    if (r != 0) {
        while (r) {
            p *= n;
            k *= r;

            long long m = __gcd(p, k);

            p /= m;
            k /= m;

            n--;
            r--;
        }
    }

    else
        p = 1;

    return p;
}

void printArray(long long arr[], long long n)
{
    for (long long i = 0; i < n; i++)
        cout << arr[i] << " ";
    cout << endl;
}
```

```

}
void nextCombination(long long combinationArr[], long long n, long long m)
{
    long long j = m;
    while (j > 0)
    {
        if (combinationArr[j - 1] < n - m + j)
        {
            combinationArr[j - 1]++;

            for (int i = j; i < m; i++)
                combinationArr[i] = combinationArr[i - 1] + 1;

            printArray(combinationArr, m);
            break;
        }
        else
            j--;
    }
}

void sequentialCombination(long long n, long long m)
{
    long long combinationArr[m];

    //initialize the array
    for (long long i = 0; i < m; i++)
        combinationArr[i] = i + 1;

    printArray(combinationArr, m);

    //number of iteration for the loop to run
    long long totalCombination = nCr(n, m) - 1;

    for (long long i = 0; i < totalCombination; i++)
        nextCombination(combinationArr, n, m);
}

int main()

```

```

{
    long long n, m;
    cout << "\nEnter the values of n and m (n>m) respectively: ";
    cin >> n >> m;

    clock_t start, end;
    start = clock();

    cout << "\nFollowing are the different combinations formed:\n";
    sequentialCombination(n, m);

    end = clock();
    double time_taken = double(end - start) / double(CLOCKS_PER_SEC);
    cout << "\nTime taken for sequential algorithm: " << time_taken << "secs" << "\n\n";
}

```

Parallel Code(using threads)

```

#include <bits/stdc++.h>
#include <thread>
#define MAX 100

using namespace std;

int N, M, k;
int x[MAX], y[MAX], c[MAX];
bool z[MAX];

void initialize(int i)
{
    int n = N;
    int m = M;
    x[i - 1] = n - m + i;
    y[i - 1] = i;
    if (y[i - 1] == x[i - 1])
        z[i - 1] = 1;
    else

```



```

        z[i - 1] = 0;
        c[i - 1] = i;
    }

void printArray(int arr[], int n)
{
    for (int i = 0; i < n; i++)
        cout << arr[i] << " ";
    cout << endl;
}

void printboolArray(bool arr[], int n)
{
    for (int i = 0; i < n; i++)
        cout << arr[i] << " ";
    cout << endl;
}

void task1(int i)
{
    if (z[i - 2] == 0 && z[i - 1] == 1)
    {
        y[i - 2] = y[i - 2] + 1;
        k = i;
    }
}

void task2(int i)
{
    y[i - 1] = y[k - 2] + (i - k + 1);
}

void task3(int i)
{
    c[i - 1] = y[i - 1];
    if (y[i - 1] == x[i - 1])
        z[i - 1] = 1;
    else
        z[i - 1] = 0;
}

```

```

}

int main()
{
    cout << "\nEnter the values of n and m (n>m) respectively: ";
    cin >> N >> M;

    cout << "\nFollowing are the combinations formed: \n";

    clock_t start, end;
    start = clock();

    vector<thread> threads;
    for (int i = 1; i <= M; i++)
    {
        threads.push_back(thread(initialize, i));
    }
    for (auto &th : threads)
    {
        th.join();
    }
    threads.clear();

    printArray(c, M);

    while (z[0] == 0)
    {
        threads.clear();
        k = 0;
        for (int i = 2; i <= M; i++)
        {
            threads.push_back(thread(task1, i));
            if (k == i)
                break;
        }
        for (auto &th1 : threads)
        {
            th1.join();
        }
    }
}

```

```

    threads.clear();
    if (k == 0)
    {
        y[M - 1] = y[M - 1] + 1;
    }
    else
    {
        for (int i = k; i <= M; i++)
        {
            threads.push_back(thread(task2, i));
        }
    }
    for (auto &th2 : threads)
    {
        th2.join();
    }
    threads.clear();
    for (int i = 1; i <= M; i++)
    {
        threads.push_back(thread(task3, i));
    }
    for (auto &th3 : threads)
    {
        th3.join();
    }
    threads.clear();
    printArray(c, M);
}

end = clock();
double time_taken = double(end - start) / double(CLOCKS_PER_SEC);
cout << "\nTime taken for parallel algorithm: " << time_taken << "secs" << "\n\n";
}

```

Parallel Code(using CUDA)

```
#include <cuda.h>
#include <cuda_runtime.h>
#include <stdio.h>
#include <stdlib.h>
#include <thrust/device_vector.h>
#include <thrust/host_vector.h>

#define int long long int

__device__ int Choose(int *fact, int n, int k) {
    return (n < k) ? 0 : (fact[n] / (fact[n - k] * fact[k]));
}

__global__ void PrintCombinations(int *fact, int n, int k, int *output) {
    int tid = blockIdx.x * blockDim.x + threadIdx.x;
    int m = tid;
    int nCr = fact[n] / (fact[n - k] * fact[k]);

    if (m >= nCr) {
        // printf("[thread %lld] Abort.\n", tid);
        return;
    }

    int idx = 0;

    int a = n, b = k;
    int x = (Choose(fact, n, k) - 1) - m;
    for (int i = 0; i < k; ++i) {
        a = a - 1;

        while (Choose(fact, a, b) > x) {
            a = a - 1;
        }
        output[m * k + idx++] = n - 1 - a;
        x = x - Choose(fact, a, b);
        b = b - 1;
    }
}
```

```

}

int32_t main() {
    int n = 5, k = 3;

    int *fact;
    fact = (int *)malloc(sizeof(int) * (n + 1));

    fact[0] = 1;
    for (int i = 1; i <= n; ++i) {
        fact[i] = fact[i - 1] * i;
    }

    printf("Factorial: ");
    for (int i = 0; i <= n; ++i) {
        printf("%lld ", fact[i]);
    }
    printf("\n");

    int nCk = fact[n] / (fact[n - k] * fact[k]);

    int *output = new int[nCk * k];
    int *d_fact, *d_output;
    cudaMalloc((void **)&d_fact, sizeof(int) * (n + 1));
    cudaMalloc((void **)&d_output, sizeof(int) * (nCk * k));
    cudaMemcpy(d_fact, fact, sizeof(int) * (n + 1), cudaMemcpyHostToDevice);
    cudaMemcpy(d_output, output, sizeof(int) * (nCk * k), cudaMemcpyHostToDevice);

    // Executing kernel
    int block_size = 256;
    int grid_size = ((nCk + block_size) / block_size);

    printf("Block size: %lld\n", block_size);
    printf("Grid size: %lld\n", grid_size);

    clock_t time_taken = clock();
    PrintCombinations<<<grid_size, block_size>>>(d_fact, n, k, d_output);
    cudaMemcpy(output, d_output, sizeof(int) * (nCk * k), cudaMemcpyDeviceToHost);

```

```
time_taken = clock() - time_taken;

printf("Time taken: %f s\n", ((float)time_taken / CLOCKS_PER_SEC) );

for (int i = 0; i < nCk; ++i) {
    printf("[Case #%lld] ", i);
    for (int j = 0; j < k; ++j) {
        printf("%lld ", output[i * k + j]);
    }
    printf("\n");
}

cudaFree(d_fact);
free(fact);
}
```

Analysis

Sequential

procedures **nCr** takes $O(n \cdot \log(r))$ and **nextCombination** takes $O(m)$ time in the worst case.

For the main procedure **sequentialCombination**,

Step 1 will take a worst case time-complexity of $O(m)$ for building and producing as output an array of size m .

Step 2.1 takes the time corresponding to procedure **nCr** i.e $O(n \cdot \log(m))$, and

Step 2.2 runs a sequential for-loop for $({}^nC_m - 1)$ times which executes the procedure **nextCombination** inside it, so Step 2.2 takes $({}^nC_m - 1) * O(m) = O(m * {}^nC_m)$ time.

Therefore the time-complexity of the whole procedure

sequentialCombination = $O(m) + O(n \cdot \log(m)) + O(m * {}^nC_m) = O(m * {}^nC_m)$

Parallel

procedures **initialize**, **task1**, **task2** and **task3** - all take $O(1)$ worst-case time-complexity since all the statements inside the procedures are assignment or branch statements which take constant time($O(1)$) for execution.

For the main procedure **parallelCombination**,

Step 1 will take $O(1)$ time, Step 2 will take $O(1)(\text{Step 2.1}) + O(m)(\text{Step 2.2}) = O(m)$ time.

Step 3.2, 3.3 and 3.4 consists of parallel loops which are executing **task1**, **task2**, **task3** respectively which requires **broadcasting** to be done before execution, that results in $O(\log(m))$ time for these steps.

The while-loop in Step 3 will take $O(nC_m)$ time for execution which leads to total execution time of the whole procedure **parallelCombination** =

$$O(m) + O({}^nC_m) * O(\log(m)) = O({}^nC_m * \log(m))$$

No. of processors used in **parallelCombination** = m .

Therefore,

the total cost of **parallelCombination** = $m * O({}^nC_m * \log(m))$

$$= O(m * {}^nC_m * \log(m)),$$

which is more than that of the sequential algorithm(**sequentialCombination**).

So the parallel algorithm is not optimal.

Results and Comparison

Data used for making Graph:

$$\mathbf{M = 5}$$

$N = [5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19]$

Parallel code time(in secs) = [0.000074, 0.000026, 0.000032, 0.000032, 0.000034, 0.000059, 0.000052, 0.000043, 0.000058, 0.000080, 0.000091, 0.000128, 0.000155, 0.000198, 0.000249]

Sequential code time(in secs) = [

0.000001, 0.000048, 0.000038, 0.000144, 0.000357, 0.000742, 0.001978, 0.003054, 0.007136, 0.009940, 0.012100, 0.026192, 0.030796, 0.043713, 0.054005]

$$\mathbf{M = 8}$$

$N = [8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19]$

Parallel code time(in secs) = [0.000140, 0.000033, 0.000034, 0.000037, 0.000058, 0.000076, 0.000115, 0.000228, 0.000409, 0.000572, 0.002734, 0.001474]

Sequential code time(in secs) = [0.000002, 0.000209, 0.000271, 0.001524, 0.003608, 0.011365, 0.013766, 0.025012, 0.070656, 0.186326, 0.374656, 0.532259]

M = 10

N = [10, 11, 12, 13, 14, 15, 16, 17, 18, 19, 20]

Parallel code time(in secs) = [0.000032, 0.000033, 0.000033, 0.000048, 0.000088, 0.000351, 0.000329, 0.000575, 0.001096, 0.005969, 0.013348]

Sequential code time(in secs) = [0.000001, 0.000061, 0.000350, 0.001613, 0.006815, 0.023217, 0.044254, 0.099836, 0.238743, 0.471609, 1.011403]

Output Screenshot:

Input: n = 5, m = 3

Sequential Code:

```
subhradeep@Subhradeep:/mnt/c/Users/HP/Desktop/Dapa-Assignment$ g++ sequential.cpp -o sequential
subhradeep@Subhradeep:/mnt/c/Users/HP/Desktop/Dapa-Assignment$ ./sequential

Enter the values of n and m (n>m) respectively: 5 3

Following are the different combinations formed:
1 2 3
1 2 4
1 2 5
1 3 4
1 3 5
1 4 5
2 3 4
2 3 5
2 4 5
3 4 5

Time taken for sequential algorithm: 0secs
```

CUDA Code:



The screenshot shows a Jupyter Notebook titled "Untitled1.ipynb". The interface includes a top menu bar with options like File, Edit, View, Insert, Runtime, Tools, and Help. Below the menu is a toolbar with icons for code, text, and other functions. The main area is divided into two sections: a code editor on the left and an output area on the right.

The code in the editor is a C++ program that calculates factorials using CUDA. It defines a block size of 256 and a grid size of 1. It prints the block and grid sizes, then calculates the time taken for the calculation. Finally, it prints the results for cases 0 through 9.

```
cudaMemcpy(d_output, output, sizeof(int) * (nCK * K), cudaMemcpyHostToDevice);

// Executing kernel
int block_size = 256;
int grid_size = (nCK + block_size) / block_size;

printf("Block size: %lld\n", block_size);
printf("Grid size: %lld\n", grid_size);

clock_t time_taken = clock();
PrintCombinations<<grid_size, block_size>>(d_fact, n, k, d_output);
cudaMemcpy(output, d_output, sizeof(int) * (nCK * k), cudaMemcpyDeviceToHost);

time_taken = clock() - time_taken;

printf("Time taken: %f s\n", ((float)time_taken / CLOCKS_PER_SEC) );

for (int i = 0; i < nCK; ++i) {
    printf("[Case #%lld] ", i);
    for (int j = 0; j < k; ++j) {
        printf("%lld ", output[i * k + j]);
    }
    printf("\n");
}

cudaFree(d_fact);
free(fact);
}
```

The output area shows the results of the program execution:

```
Factorial: 1 1 2 6 24 120
Block size: 256
Grid size: 1
Time taken: 0.000112 s
[Case #0] 0 1 2
[Case #1] 0 1 3
[Case #2] 0 1 4
[Case #3] 0 2 3
[Case #4] 0 2 4
[Case #5] 0 3 4
[Case #6] 1 2 3
[Case #7] 1 2 4
[Case #8] 1 3 4
[Case #9] 2 3 4
```

Code for Graph generation(in Python)

```
import math

import random

import time

import matplotlib.pyplot as plt

sarr = [5, 6, 7, 8, 9, 10, 11, 12, 13, 14, 15, 16, 17, 18, 19]

ptarr = [0.000074, 0.000026, 0.000032, 0.000032, 0.000034, 0.000059, 0.000052,
0.000043, 0.000058, 0.000080, 0.000091, 0.000128, 0.000155, 0.000198, 0.000249]

starr = [0.000001, 0.000048, 0.000038, 0.000144, 0.000357, 0.000742, 0.001978,
0.003054, 0.007136, 0.009940, 0.012100, 0.026192, 0.030796, 0.043713, 0.054005]

plt.plot(sarr, starr, color = 'green', marker='o', markerfacecolor='blue', label = 'Serial')

plt.plot(sarr, ptarr, color = 'red', marker='o', markerfacecolor='black', label = 'Parallel')

plt.ylim(0,0.01)

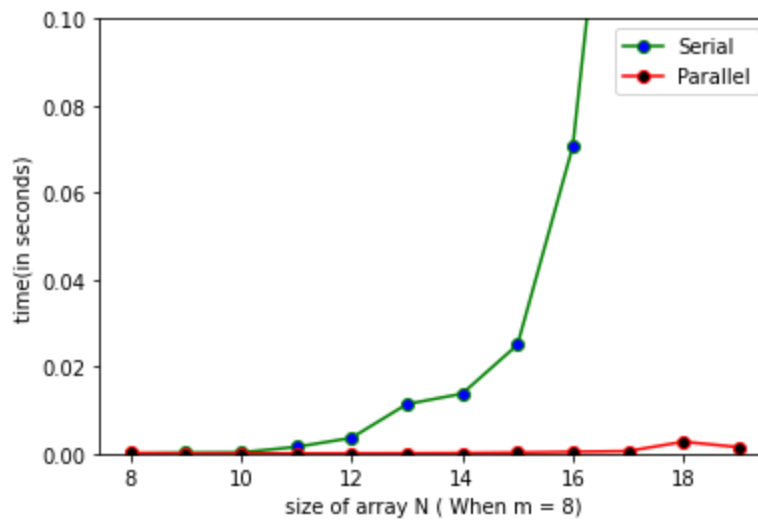
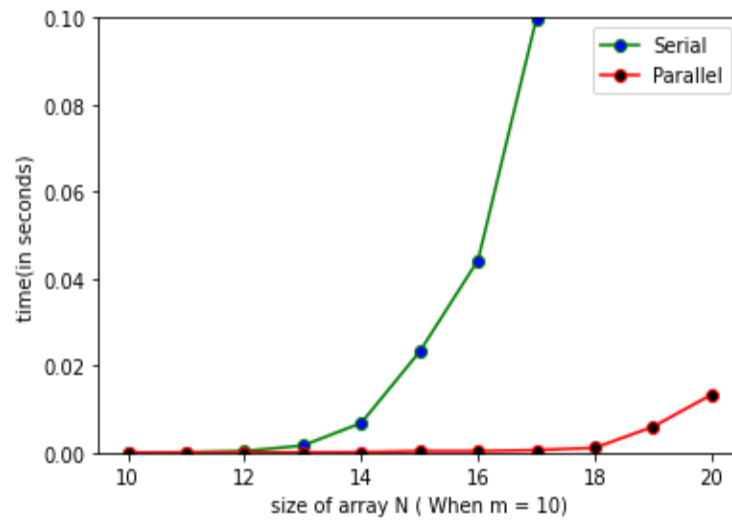
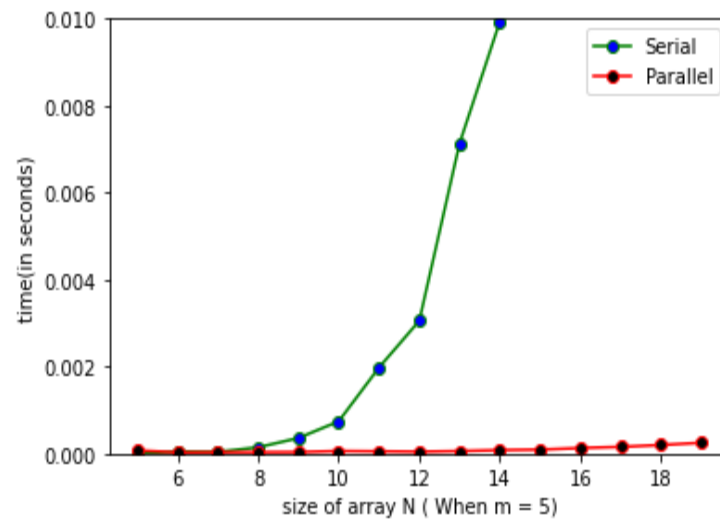
plt.xlabel('size of array N ( When m = 5)')

plt.ylabel('time(in seconds)')

plt.legend()

plt.show()
```

Graph



Conclusion

- When the input is small, the time taken by the parallel algorithm is more than the sequential algorithm. That's because for a small input, the time of splitting and assigning of elements to each thread is significant. This forms the extra overhead for the parallel algorithm.
- On increasing the number of elements, the difference between the time taken by both algorithms decreases.
- At one junction of time, we have no time difference between parallel and sequential algorithm.
- With increasing number of elements, we have significant time improvement for parallelly over the sequential one.

On a small number of elements, the sequential algorithm would be the best choice of algorithm since the process would be quicker than the parallel algorithm.

The parallel algorithm would be the best choice if the number of elements is large, because the thread creation time can be compensated for the number of threads which are operating in parallel, and thus, the overhead for thread creation can be saved.