# **COMPUTER SCIENCE**

**SUBJECT NAME: DBMS** 

**CHAPTER NO.: 1** 

**CHAPTER NAME:** FUNCTIONAL DEPENDENCY & NORMAL FORMS

**LECTURE NO.: 1** 

#### **SYLLABUS**

ER-model. Relational model: relational algebra, tuple calculus, SQL. Integrity constraints, normal forms. File organization, indexing (e.g., B and B+ trees). Transactions and concurrency control.

#### REFERENCE TEXTBOOKS:

- Fundamentals of Database System

  Authors: Elmasri Ramez, Navathe Shamkant
- Database System Concepts

  Authors: Abraham Silberschatz, Henry F. Korth
- Database Management Systems

  Authors: Raghu Ramakrishnan, Johannes Gehrke

RECORD: A collection of individual data items.

Dr. E. F CODD

DATABASE: An organized collection of records.

DBMS: A software that facilitates the end user to manage the underlying data

base.

### **FUNCTIONAL DEPENDENCY**

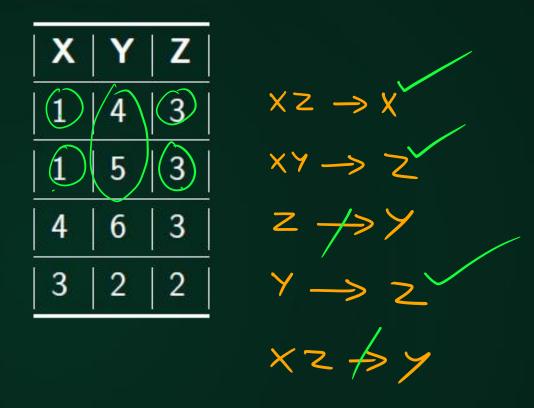


EACH VALUE OF A SHOULD BE ASSOCIATED
WITH THE SAME VALUE B'

A	В	C
a1	b1	c1
a2	b1	c1
a3	b2	c2
a1	b2	c1

$$A \neq B$$
 $B \neq C$ 
 $A \rightarrow C$ 
 $C \neq A$ 
 $AB \rightarrow C$ 
 $BC \neq A$ 
 $AC \neq B$ 
 $A \rightarrow A$ 
 $B \rightarrow B$ 
 $C \rightarrow C$ 
 $AB \rightarrow B$ 
 $C \rightarrow C$ 
 $C \rightarrow C$ 

# **FUNCTIONAL DEPENDENCY**



#### **FUNCTIONAL DEPENDENCY**

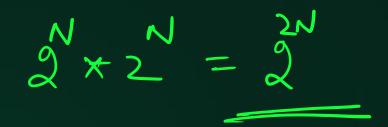
- (1) TRIVIAL FOR X->Y IF X2Y FJ:- AB->A AB->AB
- (2) NON TRIVIAL FOR X->Y IF X NY = \$\phi \text{ Fg = AB -> CD (AB N CD = \$\phi\$)
- (3) SEMI NONTRIVIAL FOR X->Y IF XNY # AND X ZY

# How many FDs are possible for a relation with "N" attributes?









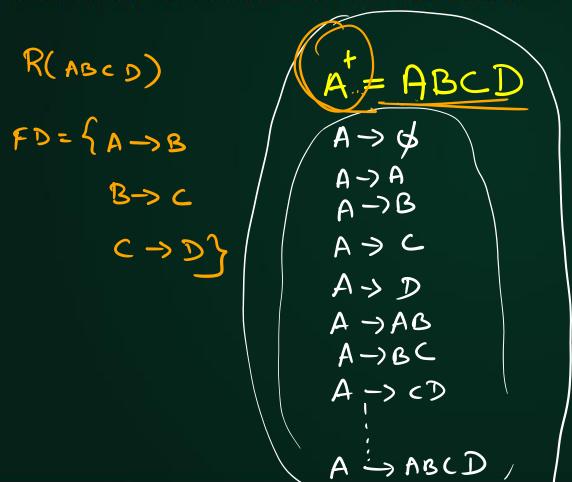
### **AMSTRONG'S AXIOMS:**

X > X

- (1) REFLEXIVITY IF X 2 Y THEN X →Y
- (2) AUGMENTATION IF X ->Y THEN XW->Y AND/OR XW ->YW
- (3) TRANSITIVITY IF X -> Y AND Y -> Z THEN X -> Z
- (4) PSEUDOTRANSITIVITY IF X->Y AND YW-Z THEN XW->Z
- (5) UNION IF X -> Y AND X -> Z THEN X -> YZ
- (6) DECOMPOSITION IF  $X \rightarrow YZ$  THEN  $X \rightarrow Y$  AND  $X \rightarrow Z$   $2C \rightarrow y$   $2C \rightarrow y$

#### **CLOSURE OF ATTRIBUTE SET:**

The set of all the attributes that are functionally determined by an attribute set S is called as Closure of Attribute Set S.



$$B^{\dagger} = BCD \qquad C = CD$$

$$C \rightarrow \phi$$

$$C \rightarrow C$$

$$C \rightarrow D$$

$$C \rightarrow C$$

#### KEY:

Attributes or set of attributes that determines the tuples uniquely in a relation.

# SUPER KEY:

(eid) = eid, ename, Salary

Attributes or set of attributes that possesses Uniqueness property

SUPER KEY

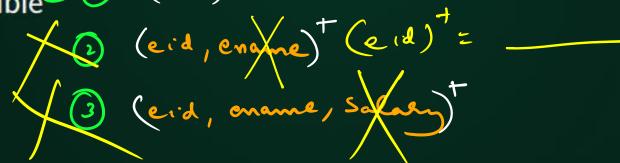
#### **CANDIDATE KEY:**

Attributes or set of attributes that possesses two properties

(1) Uniqueness

(e.d)

(2) Irreducible



EMPLOYEE



#### **PRIMARY KEY:**

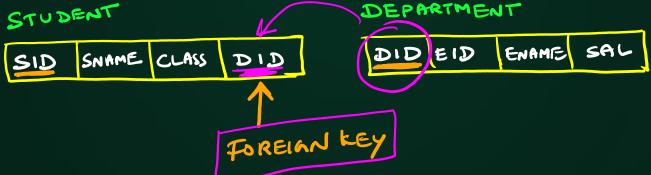
One of the Candidate Key selected by the designer of the database.

#### ALTERNATE KEY:

All Candidate Keys except the Primary Key

#### **FOREIGN KEY:**

Attributes or set of attributes in a relation that acts as Primary Key in another relation.



## FINDING ALL CANDIDATE KEYS

R(ABCD)

$$C \rightarrow A$$

$$(B)^{\dagger} = B$$

$$(2)(BCDA)^{\dagger} = BCDA$$

TOTAL 2 Candidak Key

AB & BC

RIABCRE)
F={AB>CD
C->D}

(AB) = ABCD

ONLY ONE CK

ABEL

BOTH LETE & RIGHT LEFT ONLY RIGHT ONLY AB (RB) = ABCD (ABE) = ABCDE CK as well as sle

R(ABC) F= {AB->C C->A} TOTAL CLY D(AB) = ABC CBD

ABX

< □ > < □ > < □ > < □ </p>

- (2) DB
  - 3 AC
- (1) D C
- (FB
- FC

Total 8 cks

FIND ALL THE CANDIDATE KEYS ?

CD->EF

D>A

C -> B

E>F

F > D}