

# COL106 - Data Structures and Algorithms

# Stack ADT

⊕ Stores Objects

⊗ Insertions & deletions are strictly  
"Last in - First out"

## Two Operations on stacks

push(x) : inserts Object x

pop() : removes the last inserted  
Object from the stack and  
returns it.

Additionally,

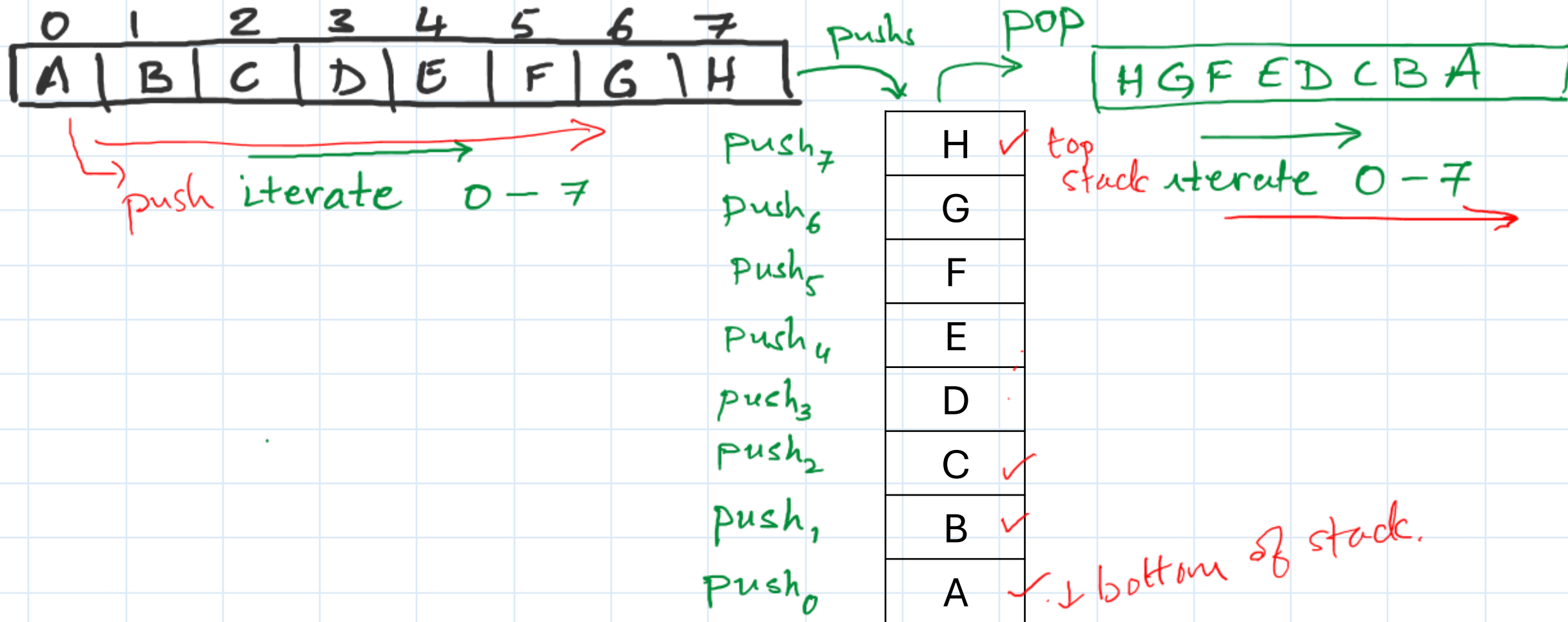
top() :

size() :

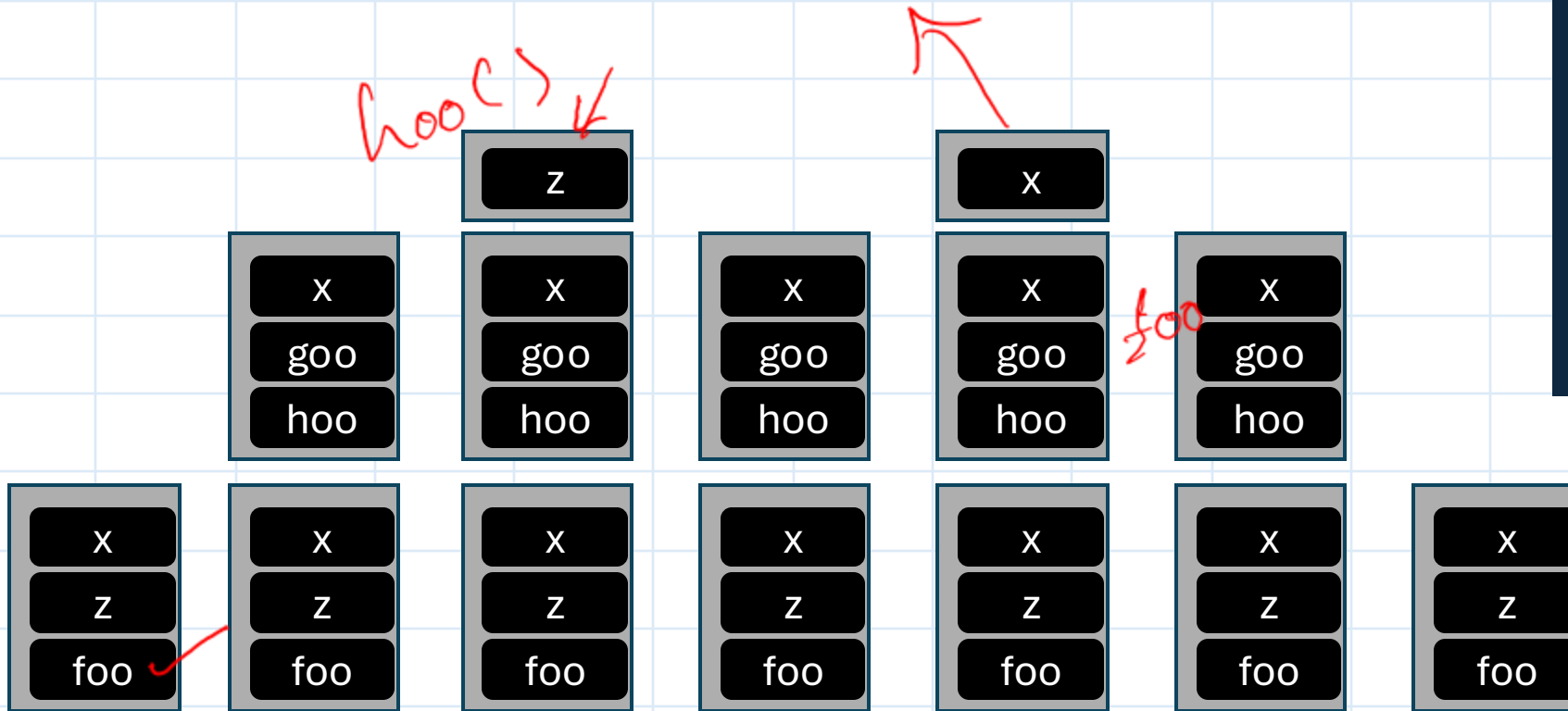
returns the last inserted object

# Applications of Stacks

## Reversing an array.



# Stacks for function calls and return values



```
def foo (x):  
    def goo ():  
        x = "abc"  
        print ("foo-goo x = ", x)  
    def hoo ():  
        z = x  
        print ("foo-hoo z = ", z)  
    x = x + 1  
    print ("foo x = ", x)  
    hoo ()  
    goo ()  
    print ("foo2 x = ", x)  
    return goo
```

```
x = 3  
z = foo(x)  
print ("x = ", x)  
print ("z = ", z)  
z()
```

# Parantheses matching

( [ { } ] )

match ✓  
match ✓  
match ✓

WHICH OF THE FOLLOWING IS VALID ?

1. ( ) ( ( ) ) { ( [ ( ) ] ) }

2. ( ( ( ) ( ( ) ) { ( [ ( ) ] ) }

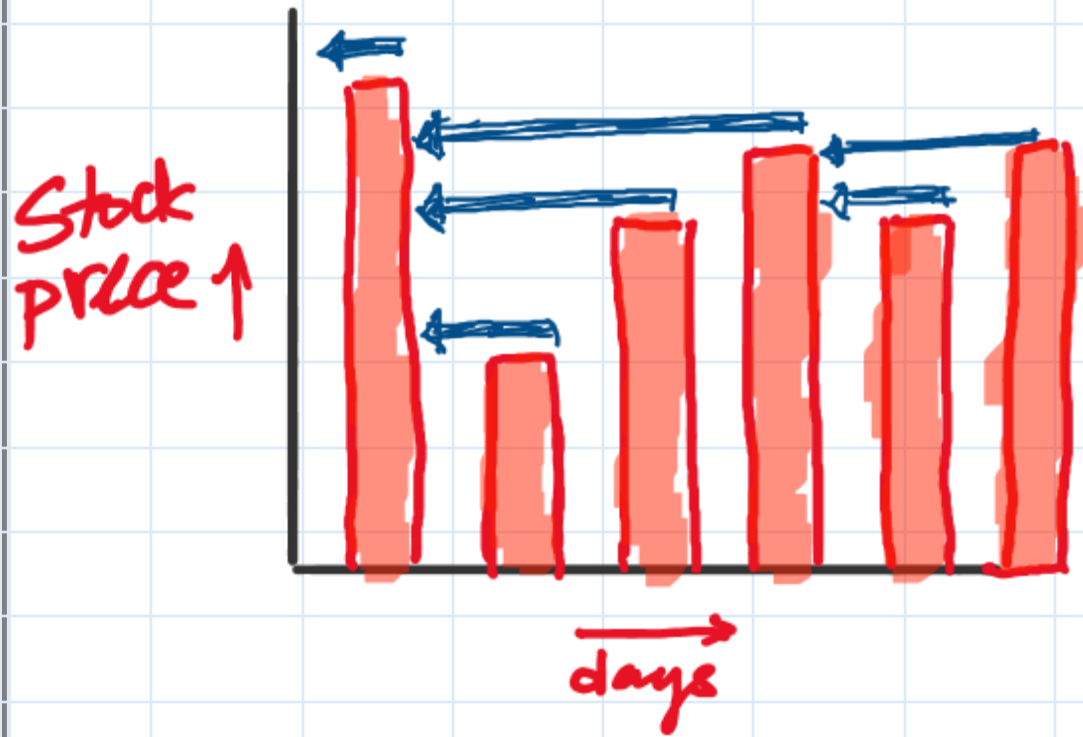
3. { [ ] ) }

```
public static boolean matching (String expression) {  
    final String opening = "( { [";  
    final String closing = ")}]";  
    Stack <Character> buffer = new LinkedList<>();  
  
    for (char c: expression.toCharArray()) {  
        if (opening.indexOf(c) != -1)  
            buffer.push(c);  
        else if (closing.indexOf(c) != -1) {  
            if (buffer.isEmpty()) return false;  
            if (closing.indexOf(c) !=  
opening.indexOf(buffer.pop())) return false;  
        }  
    }  
    return buffer.isEmpty();  
}
```

0, 1, 2

invalid  
expression

# COMPUTING THE SPAN OF EACH DAY.



For each day  $i$ , its span  $S[i]$  is the **maximum** number of **consecutive** elements  $X[j]$  immediately preceding  $X[i]$ , such that  $X[j] \leq X[i]$

We are on  $i$ -th day

→ Pop elements from the stack if  $\text{top}() \leq \text{stockPrice}[i]$ .

→  $\text{Span}[i] = \# \text{ elements popped.}$

→  $\text{push}(\text{stockPrice}[i])$

```

Stack_Span (X, n) {
    S ← empty array of n elements      1
    A ← empty stack                     1
    for i ← 0 to n-1 {                 n
        while (!A.isEmpty() && X[A.top()] ≤ X[i]) {  n
            A.pop()                     n
        }
        if (A.isEmpty()) then          n
            S[i] ← S[i]+1              n
        else
            S[i] ← i - A.top()          n
            A.push(i)                  n
    }
    return S;                          1
}

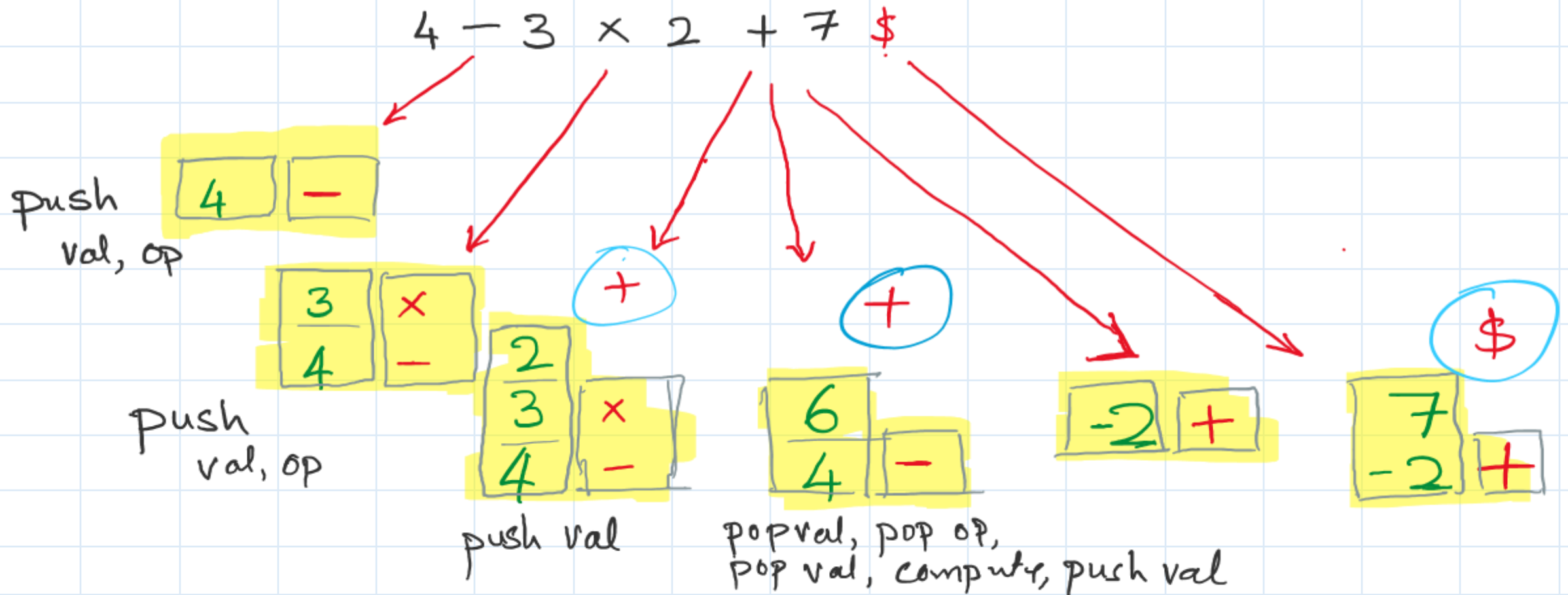
```

From



# Expression Evaluation

maintain two stacks : **opStack** **valStack**.



# Other Applications

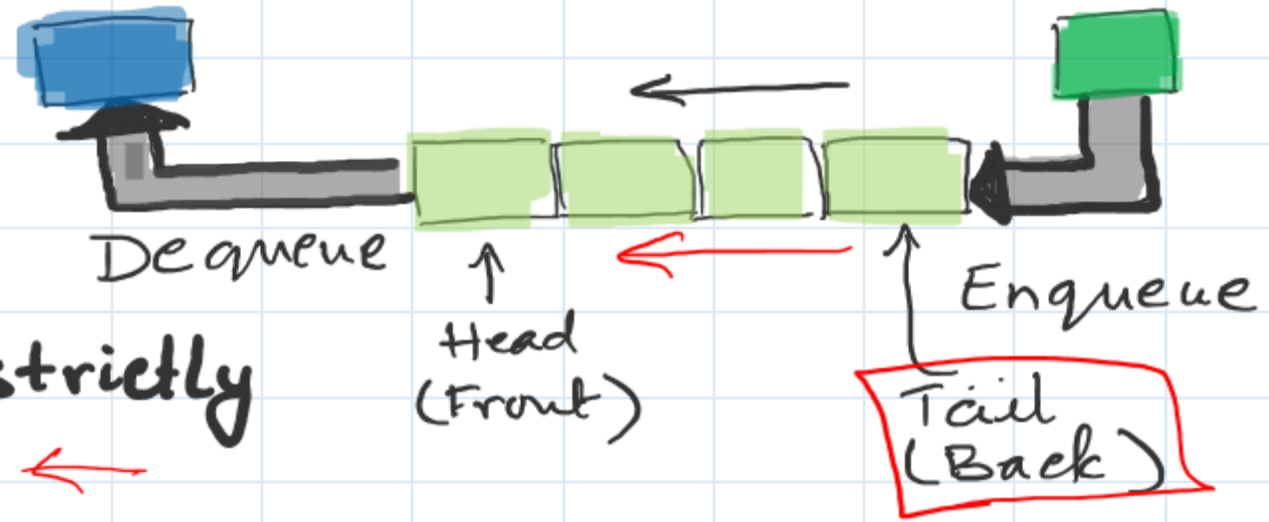
- Browser history  
separate stack for each tab
- HTML parsing
- Convex hull (Graham's Scan)

[https://en.wikipedia.org/wiki/Graham\\_scan](https://en.wikipedia.org/wiki/Graham_scan)



# Queue ADT

- ⊕ Stores Objects ✓
- ⊗ Insertions & deletions are strictly "First in - First out" ←



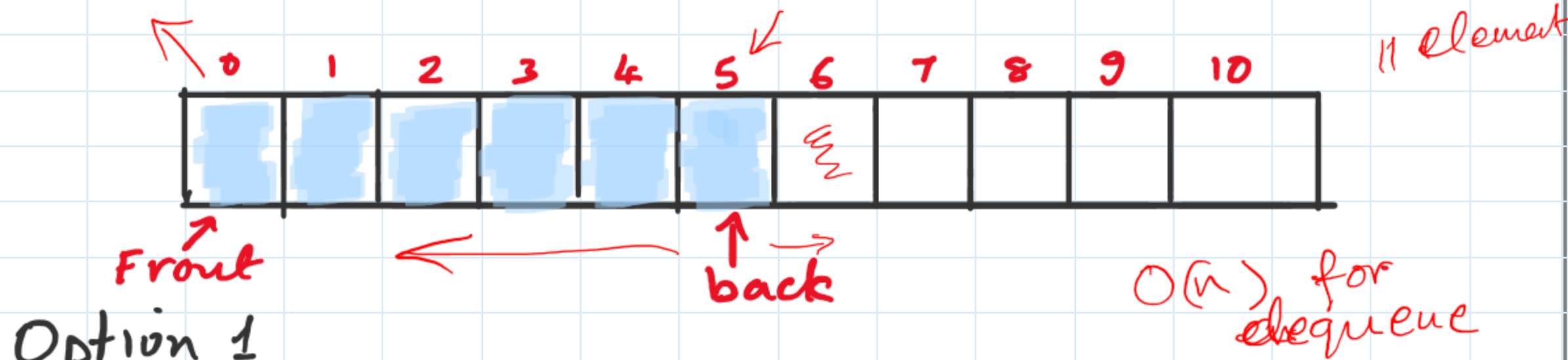
## Two Operations on Queues

OR  $\text{enqueue}(x)$   
 $\text{add}(x)$  : inserts object  $x$  at the tail.  $O(1)$

OR  $\text{dequeue}()$   
 $\text{remove}()$  : removes the objects from the head of the queue and returns it.  $O(1)$

OR  $\text{peek}()$   
 $\text{first}()$  : returns the object at the front of the queue without removing it.  $O(1)$

# Array based Implementation of a Queue.



Option 1

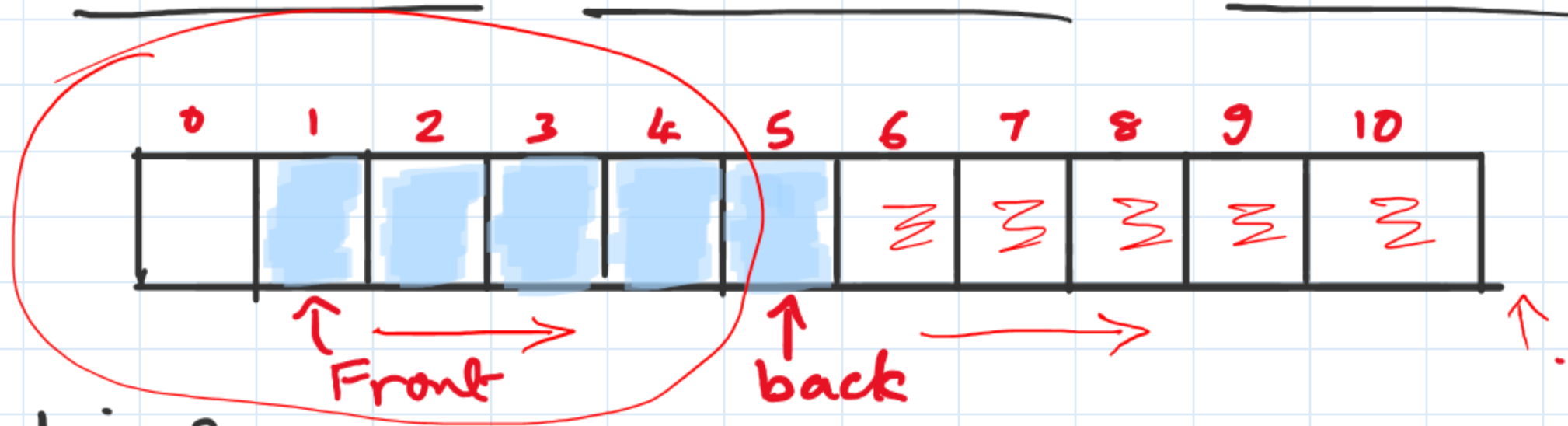
Store first element in index 0  
second — " — " — 1  
⋮

add(x) : increment back index, add x

remove() : remove the element at front  
and shift all elements right



# Array based Implementation of a Queue.



## Option 2

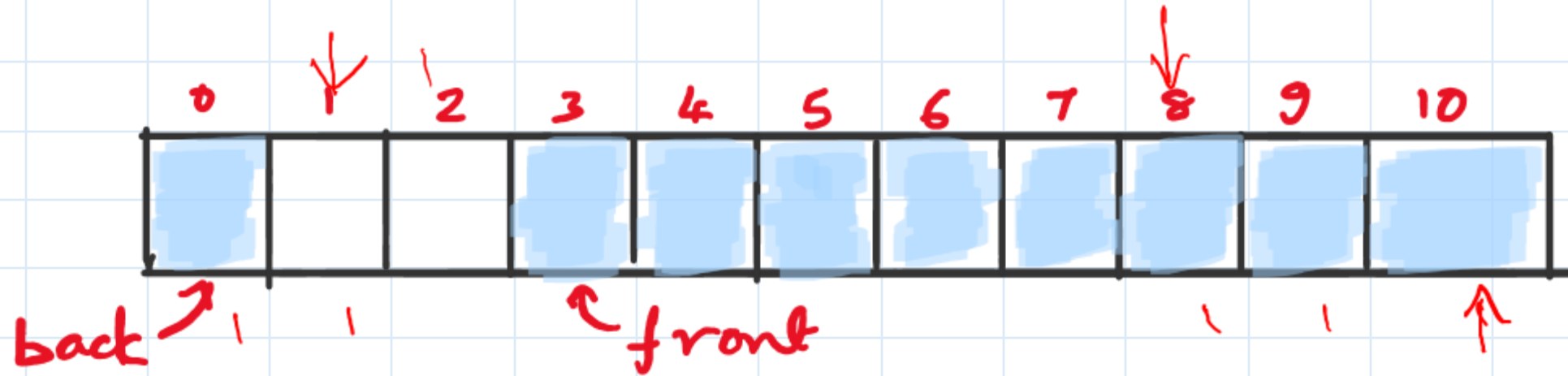
store first element in index 0  
second — " — " — " — 1

add(x) : increment back index, add x

remove() : remove the element at front.

increment the front index

# Array based Implementation of a Queue.



Option 3

Treat array as being "circular"  
Implementable using modulo arithmetic.

add(x): insert x at  $(\text{front} + \text{Size}) \% N$  Size of the queue  
remove(): remove the element at front.  $O(1)$

$\text{front} = (\text{front} + 1) \% N$   $O(1)$

# Linked List Implementation of Queue

- Often we need unbounded queue
- Linked List implementation is similar in approach to the implementation of stack.
- In general, linked list implementations are more expensive in implementation  
instantiation of new objects are expensive



# Double Ended Queue ADT (deque)

Combines the features of Stacks and Queue

add Last (x)

remove First()

add First (x)

remove Last()

first()

last()

size()

is Empty()

Implementation is left as an exercise