

Some more C functions for you:

sigprocmask pthread\_sigmask pthread\_self() atexit sigaction

The big problem: How to implement the mutex lock

### Hardware CPU instruction simplified solution

(‘Atomic\_Exchange’ swaps values at two addresses as an *uninterruptable* operation)

```
typedef p_mutex_t int;
pthread_mutex_init(p_mutex_t* m)      { *m = 1; }
pthread_mutex_lock(p_mutex_t* m)      { int local=0;
                                       do {
                                           ATOMIC_EXCHANGE(m, &local);
                                       } while(!value);
                                       }
```

```
pthread_mutex_unlock(p_mutex_t* m)    { *m = 1; }
```

**C-Code Candidate # 0** (Review) Protect our critical section with a mutex. But how should it work!?

```
pthread_mutex_lock(p_mutex_t* m)      { while(m->lock) {}; m->lock = 1;}
pthread_mutex_unlock(p_mutex_t* m)    { m->lock = 0; }
```

Problems?

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### Psuedo code Candidate # 1

```
wait until your flag is lowered
raise my flag
    // Do Critical Section stuff
lower my flag
```

```
wait until your flag is lowered
raise my flag
    // Do Critical Section stuff
lower my flag
```

// Threads do other stuff and then will repeat at sometime in the future

Problems with 1b? Fix?

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### Candidate #2

```
raise my flag
wait until your flag is lowered
    // Do Critical Section stuff
lower my flag
```

```
raise my flag
wait until your flag is lowered
    // Do Critical Section stuff
lower my flag
```

// Threads do other stuff and then will repeat at sometime in the future

Problems with 2?

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### Candidate #3

```
wait until my turn (turn==id?)
    // Do Critical Section stuff
turn = yourid
```

```
wait until my turn (turn==id?)
    // Do Critical Section stuff
turn = yourid
```

// Threads do other stuff and then will repeat at sometime in the future

Problems with 3?

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Three desirable properties of the solution to the Critical Section Problem?

//Ex 1 Make these functions thread-safe

// int thread\_mutex\_lock(pthread\_mutex\_t \*mutex); ("p\_m\_lock")

// int pthread\_mutex\_unlock(pthread\_mutex\_t \*mutex); ("p\_m\_unlock") will be useful

// Compile with -pthread

// Create a mutex that is ready to be locked!

pthread\_mutex\_t m = PTHREAD\_MUTEX\_INITIALIZER;

float sensor[2];

float old\_vals[2];

void on\_sensor\_update(float x, float y) {

// memcpy - fastest way to copy memory regions that do not overlap.

memcpy( old\_vals, sensor, sizeof( sensor ) );

sensor[0]=x;

sensor[1]=y;

}

float moved2() {

float dx = (sensor[0] - old\_vals[0]);

float dy = (sensor[1] - old\_vals[1]);

return dx\*dx + dy\*dy;

}

//#Ex 2

// Use a counting semaphore to ensure a maximum of 20 threads are running at a time. Threads that cannot acquire a music pass must wait (block) until one is released.

// hint: int sem\_init(sem\_t \*sem, int pshared, unsigned int value);

// int sem\_destroy(sem\_t \*sem); "The effect of destroying a semaphore upon which other threads are currently blocked is undefined."

// int sem\_post(sem\_t \*sem);

// int sem\_wait(sem\_t \*sem); will be useful

sem\_t s;

void init() { ?

void acquireMusicPass() { ?

void releaseMusicPass() { ?

//Ex3 Carefully explain when and how the following code can copy more than size+1 bytes to the target address

// when used with more than one thread. Your answer should include the interleaving of the two threads' actions.

char \*buffer = calloc(1,1);

size\_t size=1;

void append(char c) { buffer = realloc(buffer, ++size); buffer[size-2] = c; buffer[size-1]='\0'; }