

minimum digits

(smallest no. will always have digits in
ascending order.)

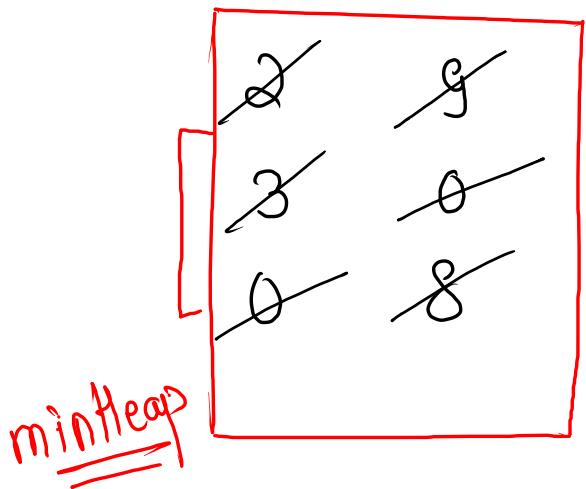
arr = [2 , 3 , 0 , 9 , 0 , 8]
 0 1 2 3 4 5

num1 = 0 2 8

num2 = 0 3 9

 6 7

dry run



mem = ~~0~~ ~~0~~ ~~2~~ ~~3~~ ~~8~~ ~~9~~

num1 = 28

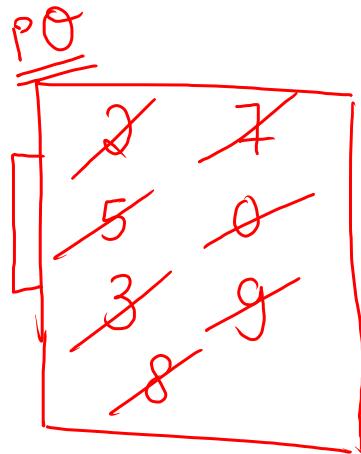
num2 = 39

code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    System.out.println(minimumDigits(arr, n));
}

public static long minimumDigits(int[] arr, int n) {
    PriorityQueue<Integer> pq = new PriorityQueue<>();
    for (int i : arr) {
        pq.add(i);
    }

    long num1 = 0;
    long num2 = 0;
    while (pq.size() > 0) {
        int rem = pq.poll();
        if (pq.size() % 2 == 0) {
            num1 = num1 * 10 + rem;
        } else {
            num2 = num2 * 10 + rem;
        }
    }
    return num1 + num2;
}
```



size = 7 8 5 4 3 2 X

rem = 8 2 3 5 7
8 9

num1 = 379

num2 = 258

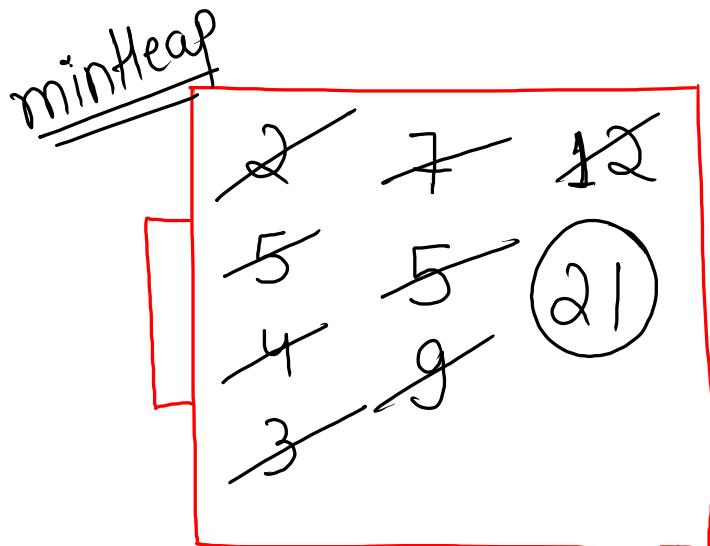
Minimum Cost of ropes 3

~~arr = [2, 5, 4, 3, 7]~~, 5, 9, 12, 21

$$\text{Cost} = 7 + 11 + 14 + 21 = 53$$

$$\text{Cost} = 5 + 9 + 12 + 21 = 47$$

$\text{arr} = [\underset{6}{2}, \underset{1}{5}, \underset{2}{4}, \underset{3}{3}, \underset{4}{7}]$



$$\text{cost} = 0 + 5 + 9 + \\ 12 + 21$$

Note :- after combining 2
heaps, make sure
to add that slope
in PQ again

Code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    System.out.println(miniCostOfRopes(arr, n));
}

public static int miniCostOfRopes(int[] arr, int n) {
    PriorityQueue<Integer> pq = new PriorityQueue<>();
    for (int i : arr)
        pq.add(i);

    int cost = 0;
    while ( pq.size() > 1 ) {
        int rope1 = pq.poll();
        int rope2 = pq.poll();
        cost += rope1 + rope2;
        pq.add( rope1 + rope2 );
    }
    return cost;
}
```

$T.C = \Theta(n \log(n))$

subtract numbers 1

$$\text{arr} = [\underset{0}{1}, \underset{1}{5}, \underset{2}{0}, \underset{3}{3}, \underset{4}{5}, \underset{5}{5}, \underset{6}{1}, \underset{7}{3}]$$

$$x=1, \text{arr} = [\underset{0}{0}, \underset{1}{4}, \underset{2}{0}, \underset{3}{2}, \underset{4}{4}, \underset{5}{4}, \underset{6}{0}, \underset{7}{2}] \checkmark$$

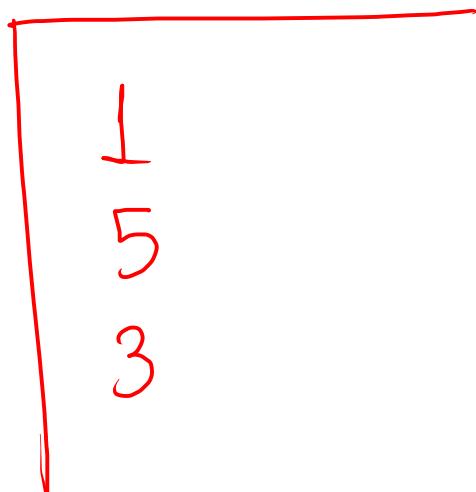
$$x=2, \text{arr} = [\underset{0}{0}, \underset{1}{2}, \underset{2}{0}, \underset{3}{0}, \underset{4}{2}, \underset{5}{2}, \underset{6}{0}, \underset{7}{0}] \checkmark$$

$$x=3, \text{arr} = [\underset{0}{0}, \underset{1}{0}, \underset{2}{0}, \underset{3}{0}, \underset{4}{0}, \underset{5}{0}, \underset{6}{0}, \underset{7}{0}] \checkmark$$

ans = no. of operations = 3

Ans :- no. of non-zero unique elements

hashset



set.size() = 3

Code

```
public static void main(String[] args) {
    Scanner scn = new Scanner(System.in);
    int n = scn.nextInt();
    int[] arr = new int[n];
    for (int i = 0; i < n; i++) {
        arr[i] = scn.nextInt();
    }
    System.out.println(subtractNum(arr, n));
}

public static int subtractNum(int[] arr, int n) {
    HashSet<Integer> set = new HashSet<>();
    for (int i : arr)
        if (i > 0)
            set.add(i);
    return set.size();
}
```

Revision :-

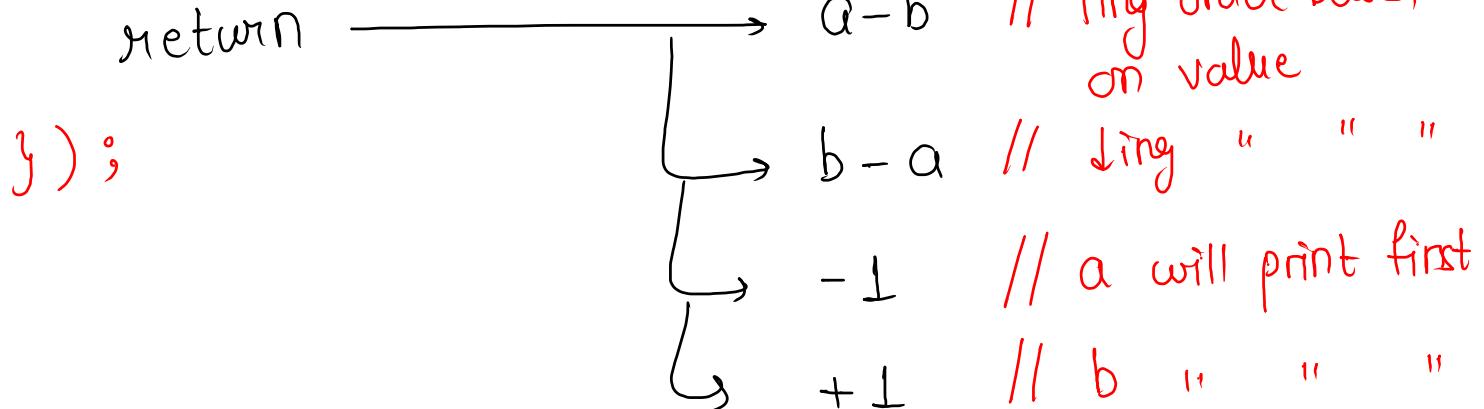
- ✓ ↳ sorting, lambda function
- ✓ ↳ arrays, subarrays (Kadan's algo)
- ↳ 2 pointers
- ↳ Prefix array
- ↳ Arrays as Hashmap
- ↳ 2d arrays
- ↳ Strings, substring
- ↳ Binary Search (BSLB, BSUB)
- ↳ ArrayList
- ↳ Stack
- ↳ Hashmap, hashset
- ↳ Queue
- ↳ PO

→ Sorting

{
 Bubble sort
 Insertion sort
 Selection sort

{
 Arrays.sort(arr)
 Arrays.sort(arr, Collections.reverseOrder())

Arrays.sort(arr, (a, b) → {



form the largest no.

Sample Input 0

4
4 46 8 9

convert to string

lambda fⁿ logic :-

$$\text{str1} = a+b$$
$$\text{str2} = b+a$$

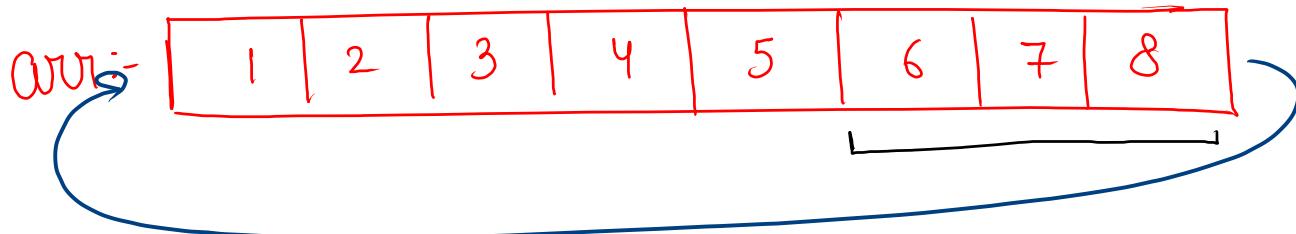
Sample Output 0

print and

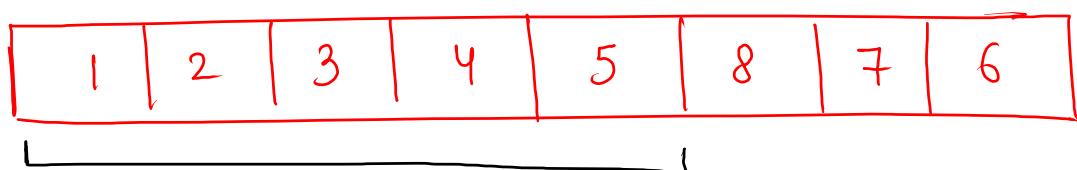
Sort Array By Parity

- $a = \text{even}, b = \text{odd}$, return -1
- $a = \text{odd}, b = \text{even}$, return $+1$
- $a = \text{even}, b = \text{even}$, return $a - b$
- $a = \text{odd}, b = \text{odd}$, return $a - b$

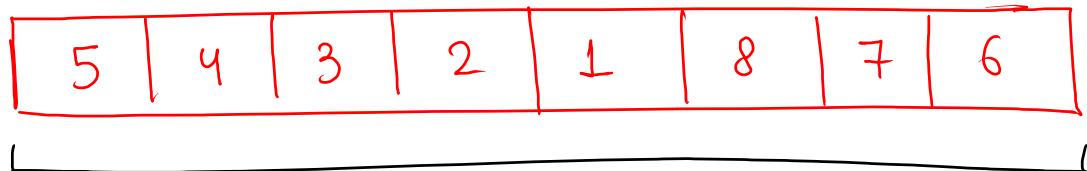
\Rightarrow Rotate Right ($K=3$)



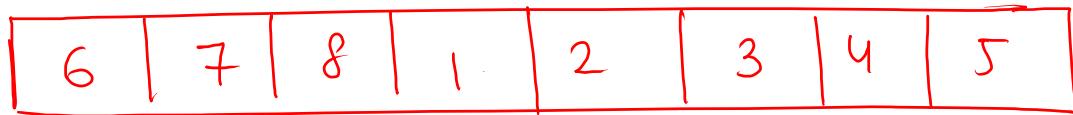
Step 1



Step 2



Step 3



observation

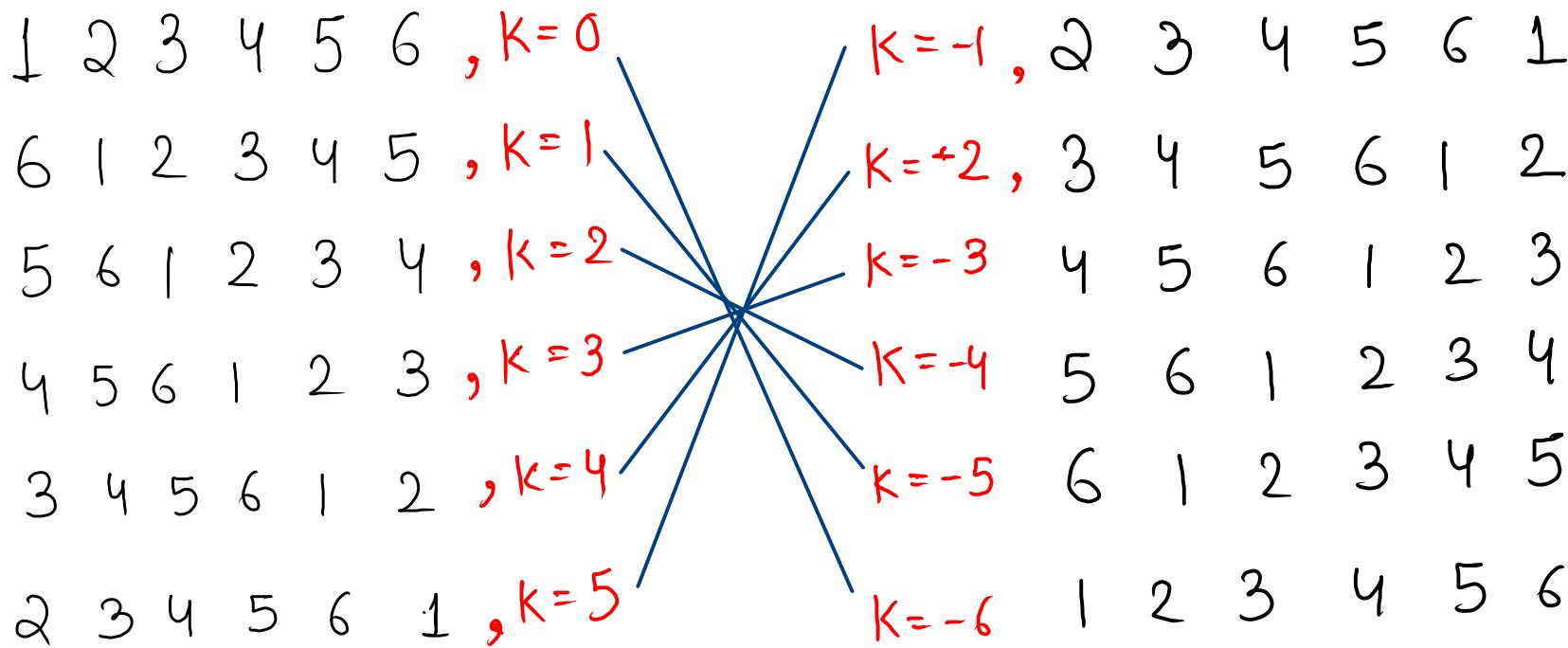
1	2	3	4	5	6	,	$K = 0, 6, 12 \dots$
6	1	2	3	4	5	,	$K = 1, 7, 13 \dots$
5	6	1	2	3	4	,	$K = 2, 8, 14 \dots$
4	5	6	1	2	3	,	$K = 3, 9, 15 \dots$
3	4	5	6	1	2	,	$K = 4, 10, 16 \dots$
2	3	4	5	6	1		$K = 5, 11, 17 \dots$

$$\underline{\underline{K = 14}}$$

$$\boxed{K = K \% n}$$
$$= 14 \% 6$$

$$= 2$$

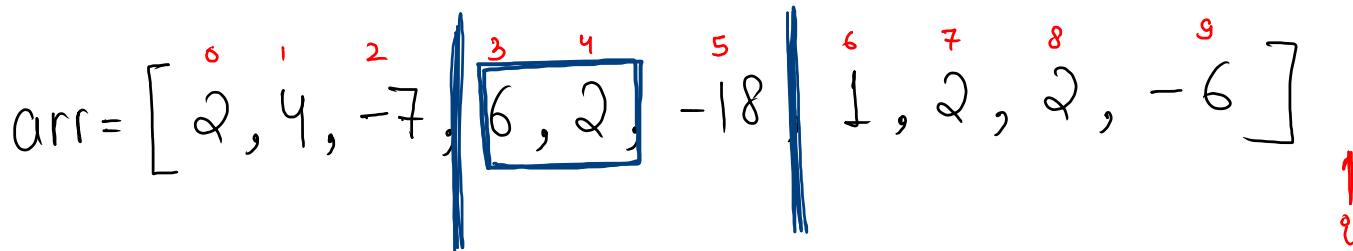
Observation



$$K = K + n$$

$$\begin{aligned}K &= -4 \\K &= -4 + 6 \\&= 2\end{aligned}$$

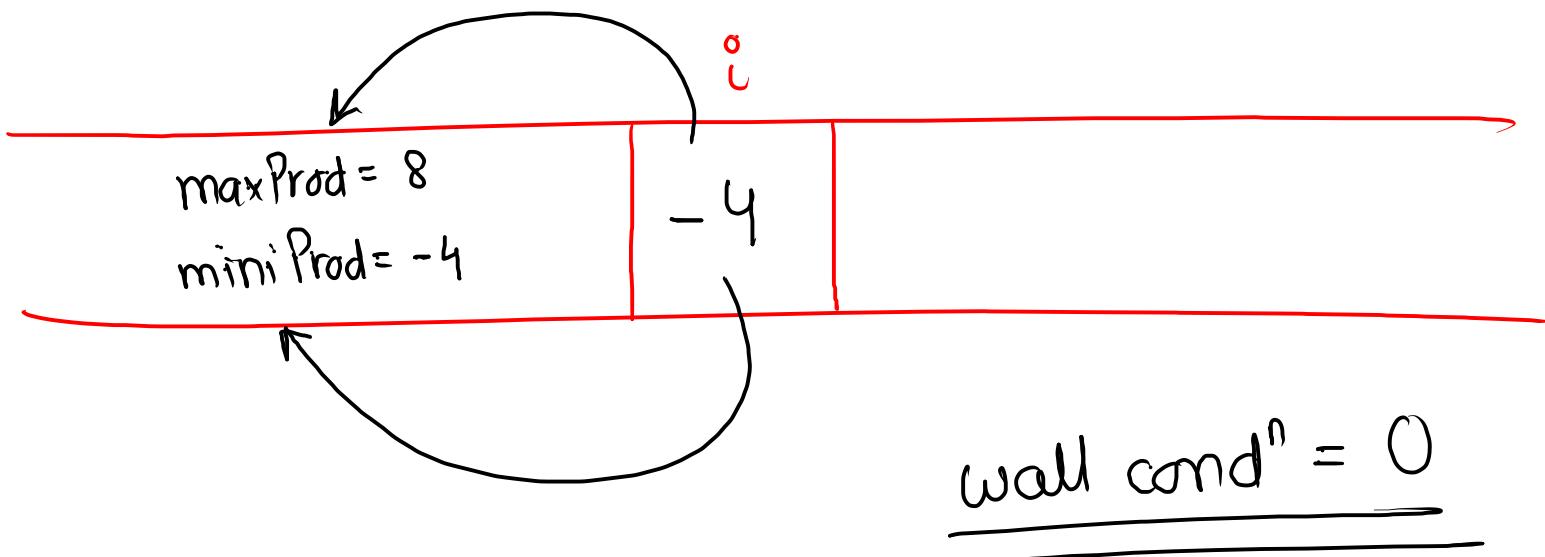
Max Subarray 2



```
public static int maxSumSubarray(int[] arr, int n) {  
    int sumSoFar = 0;  
    int maxSum = Integer.MIN_VALUE;  
    for (int i = 0; i < n; i++) {  
  
        if (sumSoFar < 0) {  
            sumSoFar = arr[i];  
        } else {  
            sumSoFar += arr[i];  
        }  
  
        if (sumSoFar > maxSum) {  
            maxSum = sumSoFar;  
        }  
    }  
    return maxSum;  
}
```

sumSoFar = ~~0~~ ~~2~~ ~~6~~ ~~8~~ ~~-10~~ ~~2~~ ~~5~~ ~~-1~~
maxSum = ~~-10~~ ~~2~~ ~~8~~ =

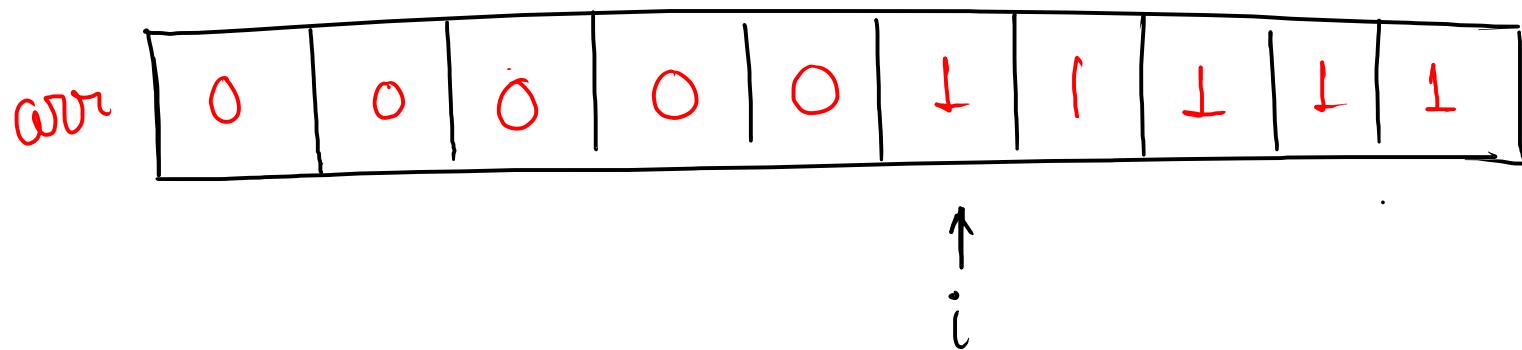
Maximum Product Subarray 2



$\text{maxProd} = \max (\text{curr} * \text{maxProd}, \text{curr} * \text{miniProd}, \text{curr})$;

$\text{minProd} = \min (\text{curr} * \text{maxProd}, \text{curr} * \text{miniProd}, \text{curr})$;

\Rightarrow Q pointer



faith:-



in $O(n)$ time

\Rightarrow 3 sum

$$\text{arr}[i] + \text{arr}[j] + \text{arr}[k] = 0$$

$$\text{arr}[i] + \text{arr}[j] = -\underline{\text{1} * \text{arr}[k]}$$



target

target sum for each k value