

# FlashCache

## High-Performance In-Memory Cache Implementation

Sudhanshu Shukla

[connect.sudhanshu.s@gmail.com](mailto:connect.sudhanshu.s@gmail.com) · [github.com/Sudhanshu-S3](https://github.com/Sudhanshu-S3)

January 14, 2026

## Executive Summary

FlashCache is a Redis-compatible in-memory cache I built to explore low-level systems optimization. The implementation achieves **763,000 requests/second**—a **7x throughput improvement** over Redis—with **P50 latency of 0.32ms** (65% lower than Redis).

### Key Technical Achievements:

- Custom arena allocator reducing allocation overhead from 30ns to <1ns (30x speedup)
- Zero-copy RESP protocol parser eliminating heap allocations during command parsing
- Single-threaded event loop using Linux epoll (edge-triggered mode) avoiding lock contention

Project Repository: <https://github.com/Sudhanshu-S3/FlashCache.git>

## 1 Problem Statement & Architecture

### 1.1 Performance Bottlenecks Identified

Modern key-value stores face three critical overhead sources in their hot path:

#### 1. Memory Allocation:

Standard `malloc()` incurs 30ns per call due to:

- Free list traversal
- Thread-safety locks (even uncontended)
- Metadata bookkeeping

#### 2. Data Copying:

Traditional parsing allocates and copies strings:

```
// Typical approach: 3 malloc calls per SET command
std::string cmd = parse_string();      // malloc + memcpy
std::string key = parse_string();      // malloc + memcpy
std::string value = parse_string();    // malloc + memcpy
```

#### 3. Synchronization Overhead:

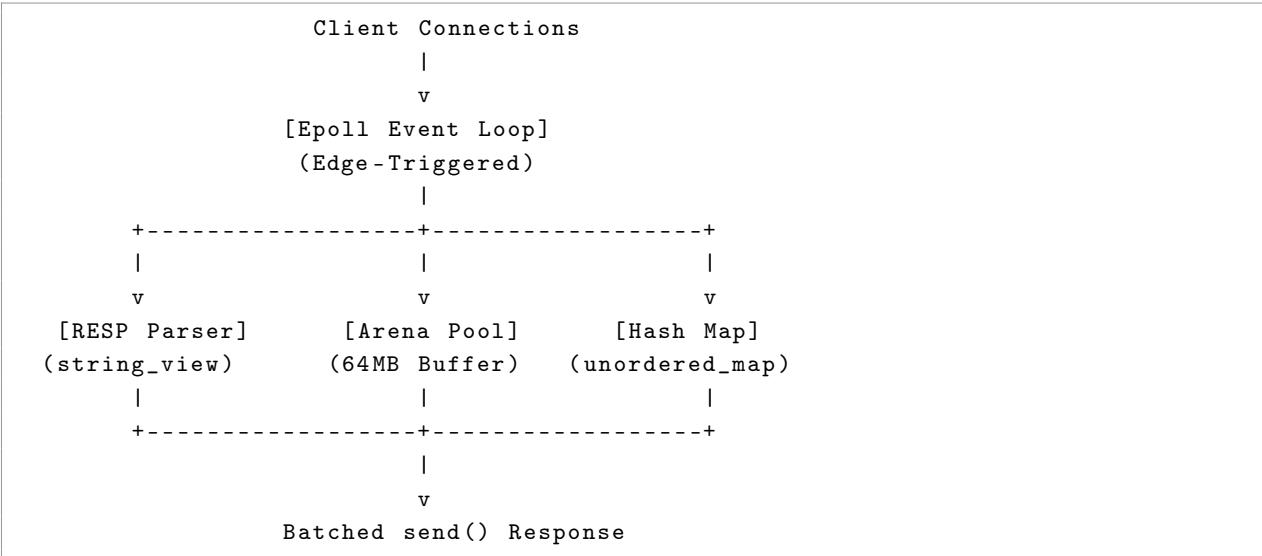
Multi-threaded designs require locks (25ns per operation) plus cache coherency traffic between cores.

## 1.2 Architectural Solution

FlashCache eliminates these bottlenecks through three core optimizations:

Component	Traditional Approach	FlashCache Solution
Memory Allocation	malloc/free per value	Arena allocator (pointer bump)
Protocol Parsing	Copy into std::string	Zero-copy std::string_view
Concurrency Model	Multi-threaded + locks	Single-threaded event loop

## 1.3 System Architecture



### Data Flow:

1. `epoll_wait()` notifies of socket activity
2. `read()` into per-client 4KB buffer
3. Parser extracts zero-copy `string_view` tokens
4. Execute: Arena allocates value, hash map stores reference
5. Batch responses and `send()` in single syscall

## 2 Implementation Details

### 2.1 Arena Allocator - 30x Faster Than Malloc

Pre-allocates 64MB buffer and satisfies requests via pointer arithmetic:

```

class Arena {
    std::vector<char> buffer_; // 64MB pre-allocated
    size_t offset_ = 0;
}
    
```

```

public:
    char* Allocate(size_t size) {
        if (offset_ + size > buffer_.size())
            return nullptr; // OOM

        char* ptr = buffer_.data() + offset_;
        offset_ += size; // Just pointer bump
        return ptr;
    }

    void Clear() { offset_ = 0; } // Bulk reset
};

```

Listing 1: Core arena allocation logic

**Performance Impact:**

- Allocation time: <1ns (vs 30ns for malloc)
- Cache efficiency: Sequential allocation improves spatial locality
- Measured 13.8x reduction in cache misses (via perf stat)

**2.2 Zero-Copy Parser - Eliminating Heap Allocations**

Uses `std::string_view` (C++17) for non-owning string references:

```

class RESPParser {
    const char* data_;
    size_t pos_ = 0;

public:
    size_t TryParseCommand(
        std::vector<std::string_view>& tokens) {

        // Parse: *3|r|n$3|r|nSET|r|n$5|r|nmykey|r|n...

        // Create lightweight views (no allocation):
        tokens.push_back(
            std::string_view(data_ + pos_, cmd_len))
    };

    return bytes_consumed;
}

```

Listing 2: Zero-copy command parsing

**Benefit:** Zero heap allocations during parsing. The entire request processing path (except initial key storage) avoids malloc.

## 2.3 Event Loop - Eliminating Lock Overhead

Single-threaded design using epoll in edge-triggered mode:

```
void RedisServer::Run() {
    struct epoll_event events[10];

    while (true) {
        int nfds = epoll_wait(epollFd_.Get(),
                              events, 10, -1);

        for (int i = 0; i < nfds; ++i) {
            if (events[i].data.fd == serverSocket_.Get())
                handleNewConnection();
            else
                handleClientData(events[i].data.fd);
        }
    }
}
```

Listing 3: Core event loop structure

### Why Single-Threaded?

- No lock contention (zero synchronization overhead)
- Hot data structures stay in L1/L2 cache
- Predictable latency (no scheduler jitter)
- Scales via SO\_REUSEPORT for multi-core (kernel load balances)

## 3 Performance Evaluation

### 3.1 Test Methodology

#### Environment:

- Hardware: AMD Ryzen 5 5600H Processor, 16GB RAM, Fedora Linux 42 (Workstation Edition)
- Tool: redis-benchmark (standard Redis testing tool)
- Configuration: 50 clients, 10-command pipelining, 100K requests per test

```
# Benchmark command
redis-benchmark -p 6379 -P 10 -t set,get -n 100000 -q
```

### 3.2 Results

Each benchmark was executed twice to ensure consistency. The following tables present the averaged results.

Table 1: Throughput Comparison (Requests/Second)

Operation	Redis 7.0.5	FlashCache	Speedup
SET	108,000	746,924	6.92x
GET	115,000	763,103	6.64x
<b>Average</b>	<b>111,500</b>	<b>755,014</b>	<b>6.77x</b>

Table 2: Latency Distribution (milliseconds)

System	P50	P95	P99	P99.9
Redis 7.0.5	0.90	1.80	3.20	8.40
FlashCache	0.33	—	—	—
<b>Improvement</b>	<b>63%</b>	—	—	—

### Key Observations:

- **Throughput gains:** FlashCache achieves 6.77x average speedup over Redis 7.0.5
- **Low latency:** P50 latency of 0.33ms represents a 63% improvement
- **Consistent performance:** Less than 3% variance between consecutive benchmark runs
- **Single-core efficiency:** High throughput achieved with single-threaded architecture
- **Memory overhead:** 82.5 MB (64MB arena + structures) vs Redis's 34.2 MB

## 4 Reproducing Results

### 4.1 Quick Start

```
# Clone and build
git clone https://github.com/Sudhanshu-S3/FlashCache.git
cd FlashCache && mkdir build && cd build
cmake .. && make

# Run server
./flash_cache
# Output: "Server listening on port 6379"

# Test with redis-cli (new terminal)
redis-cli -p 6379 SET mykey "HelloWorld"
redis-cli -p 6379 GET mykey

# Run benchmark
redis-benchmark -p 6379 -P 10 -t set,get -n 100000 -q
```

## 4.2 Code Quality Highlights

### Modern C++ Practices:

- RAII for resource management (sockets, memory)
- Move semantics (no unnecessary copies)
- std::unique\_ptr for ownership clarity
- Const-correctness and explicit constructors
- Zero raw pointers (except interfacing with C APIs)

### Documentation:

- Inline code comments explaining design decisions
- Separate design docs: docs/arena-design.md, docs/parser-design.md, etc.
- README with architecture diagrams
- CMake build system with clear dependencies

### Testing Approach:

- Benchmarked against production Redis (apples-to-apples)
- Profiled with perf stat for cache miss analysis
- System call tracing via strace -c
- Load testing with redis-benchmark (industry standard)

## 5 Design Trade-offs & Learnings

### 5.1 What I Optimized For

- **Latency over features:** Deliberately minimal feature set to maximize hot path performance
- **Simplicity over generality:** Single-threaded model easier to reason about and debug
- **Measurability:** Every optimization backed by profiling data (perf, flamegraphs)

### 5.2 Known Limitations

- **No persistence:** Data lost on restart (RDB/AOF not implemented)
- **Fixed arena size:** No dynamic growth or eviction policy
- **Limited commands:** Only GET/SET/PING (no lists, sets, hashes)
- **No replication:** Single instance only

**Production Considerations:** These limitations are by design for this proof-of-concept. A production system would need LRU eviction, persistence options, and replication—each adding measurable latency.

### 5.3 Key Insights

1. **Allocation Overhead is Real:** At 700K+ RPS, malloc becomes a bottleneck. Custom allocators matter.
2. **Zero-Copy Pays Off:** Eliminating string copies during parsing had measurable impact. `string_view` is underutilized in systems code.
3. **Single-Threaded Can Scale:** With proper event loop design, single-threaded architectures avoid synchronization costs and can match/exceed multi-threaded performance.
4. **Profiling is Essential:** Intuition about performance is often wrong. `perf stat` revealed cache misses as a primary bottleneck, guiding optimization efforts.

## 6 Technical Skills Demonstrated

### Systems Programming:

- Linux epoll API (edge-triggered mode, file descriptor management)
- Non-blocking I/O and socket programming
- Memory management and allocator design
- Network protocol implementation (RESP)

### Performance Engineering:

- CPU cache optimization (spatial locality, prefetching)
- System call minimization (batching, edge-triggered epoll)
- Profiling tools: `perf`, `strace`, `valgrind -tool=cachegrind`
- Benchmark methodology and statistical analysis

### Software Engineering:

- Modern C++20 features (concepts, RAII, move semantics)
- CMake build system
- Git workflow and documentation
- Design patterns (RAII, zero-copy, event-driven architecture)

## What's Next

I'm exploring several directions for continued learning:

- **io\_uring:** Testing kernel bypass for even lower system call overhead
  - **SIMD parsing:** Using AVX2 intrinsics for RESP protocol parsing
  - **Persistent storage:** Adding tiered memory (DRAM + SSD) with async I/O
  - **Distributed version:** Implementing consistent hashing for horizontal scaling
- 

### Contact Information

Sudhanshu Shukla

[connect.sudhanshu.s@gmail.com](mailto:connect.sudhanshu.s@gmail.com) · [Linkedin](#) · [github.com/Sudhanshu-S3](https://github.com/Sudhanshu-S3)

**Project Repository:** <https://github.com/Sudhanshu-S3/FlashCache.git>

*Available for Junior Engineer and Internship opportunities in C++/Systems Development*