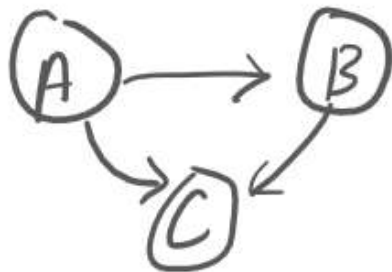


# BELLMAN FORD ALGORITHM (GRAPH)



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## 2. Bellman Ford Algorithm

 It is also used to find SSSP from source to all nodes

 **SSSP: SINGLE SOURCE SHORTEST PATH**

**What is the Bellman-Ford Algorithm:**


This Loop-based algorithm is used to find the shortest path from a single source node to all other destination nodes and detect the negative cycle in the graph.

**Note 1:** It is also used to find all shortest path from source to all other nodes in a negative and positive weighted, directed or undirected graph.

**Note 2:** It is also used to when a graph have the negative cycle.

**Step 1: to find the SSSP**

Relaxation N-1 times to get the all shortest path from , single node.

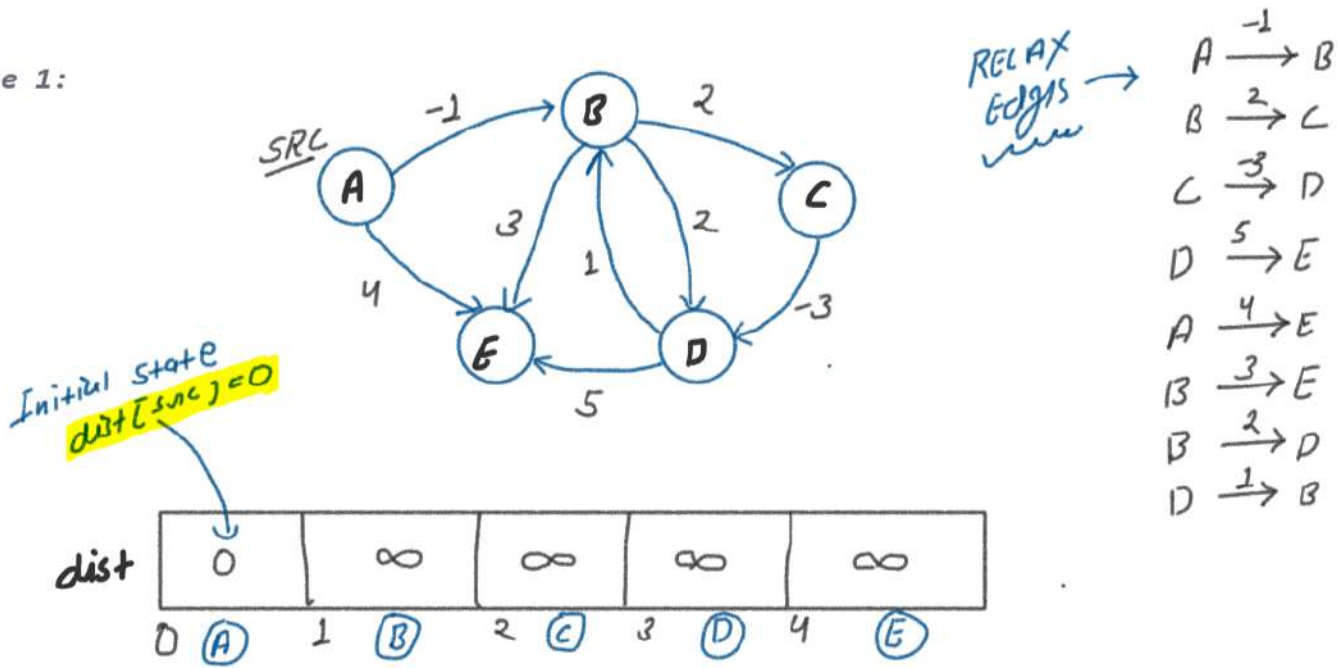
 
$$(\text{Dist}[\text{Node}] + \text{weight} < \text{Dist}[\text{Node}])$$

**Step 2: to detect the negative cycle**

One more relaxation after N-1 times to detect the negative cycle

$$\hookrightarrow \text{Flag} = \text{True} / \text{False}$$

Example 1:



Key	Value
A	$\{EB, -1\}, \{EE, 4\}$
B	$\{BE, 2\}, \{BD, 2\}, \{BE, 3\}$
C	$\{CD, -3\}$
D	$\{DE, 5\}, \{DB, 1\}$
E	x

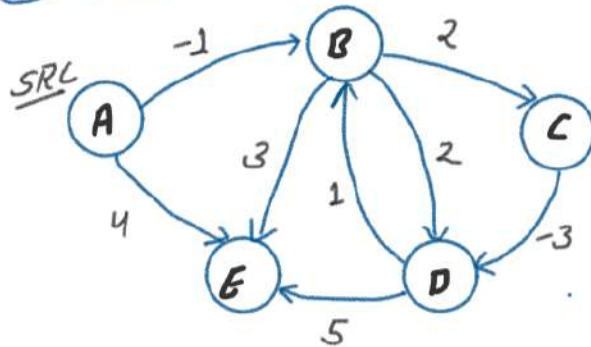
AdjList

DRY RUN

Relaxation 1:

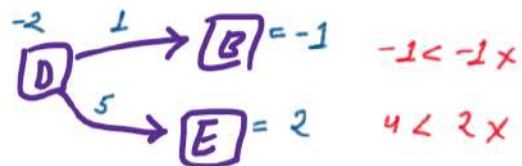
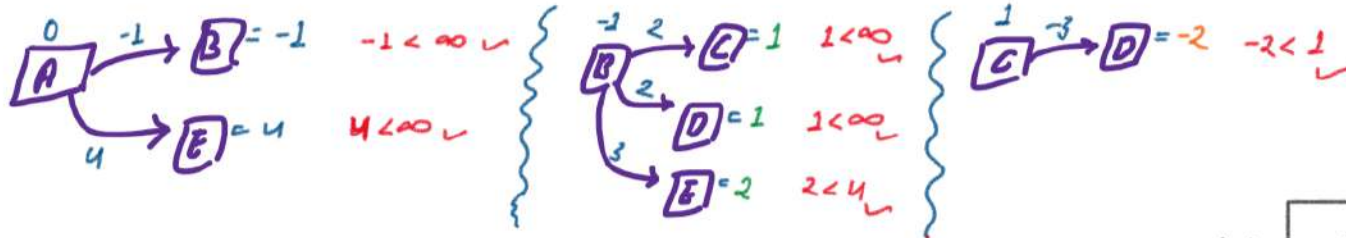
$N = \text{Total nodes} = 5$   
 $\text{Relaxation} = N - 1 = 4$

Step 1: to find the SSSP



Key	Value
A	$\{ \{B, -1\}, \{E, 4\} \}$
B	$\{ \{C, 2\}, \{D, 2\}, \{E, 3\} \}$
C	$\{ \{D, -3\} \}$
D	$\{ \{E, 5\}, \{B, 1\} \}$
E	X

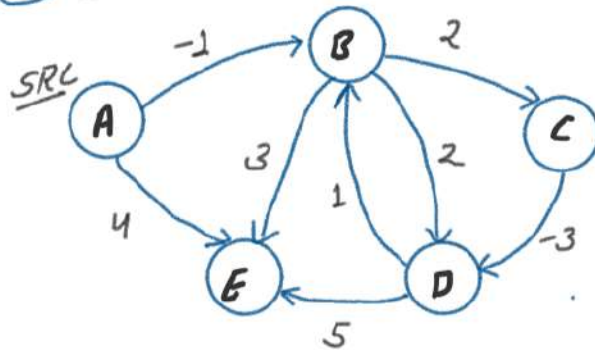
AdjList



dist	0	<del>-1</del>	<del>∞</del> 1	<del>∞</del> <del>2</del>	<del>∞</del> <del>2</del>
	0 (A)	1 (B)	2 (C)	3 (D)	4 (E)

Relaxation 2:

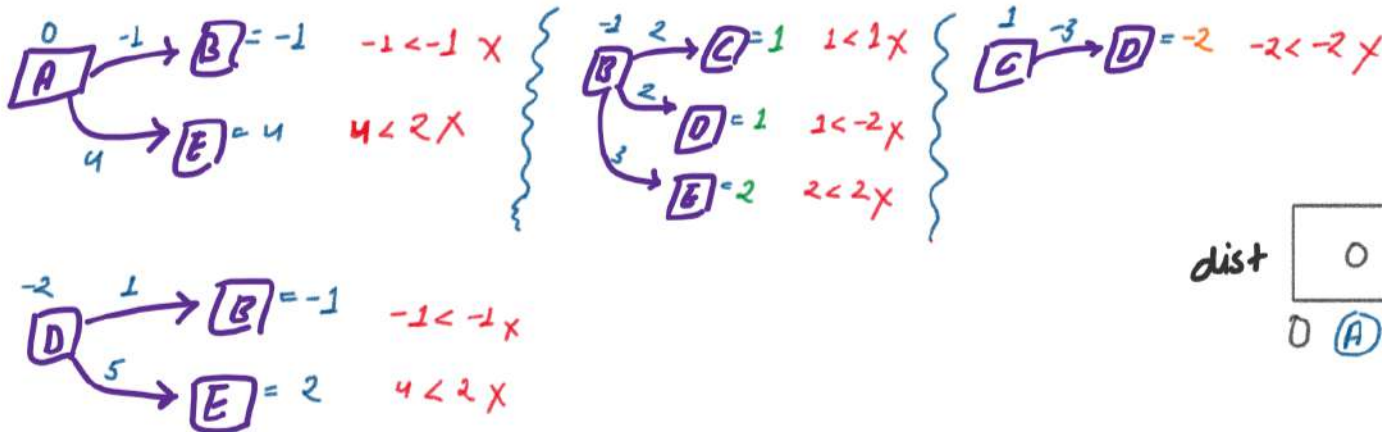
$N = \text{Total Nodes} = 5$   
 $\text{Relaxation} = N - 1 = 4$



No updation during Relaxation 2

Key	Value
A	{E, -1}, {E, 4}
B	{E, 2}, {D, 2}, {E, 3}
C	{E, D, -3}
D	{E, 5}, {B, 2}
E	x

AdjList

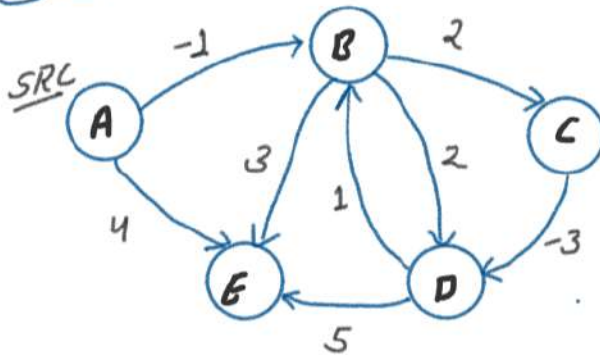


dist	0	-1	1	-2	2
	0 (A)	1 (B)	2 (C)	3 (D)	4 (E)



Relaxation 3:

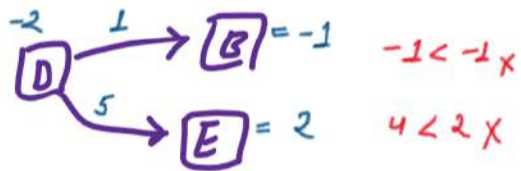
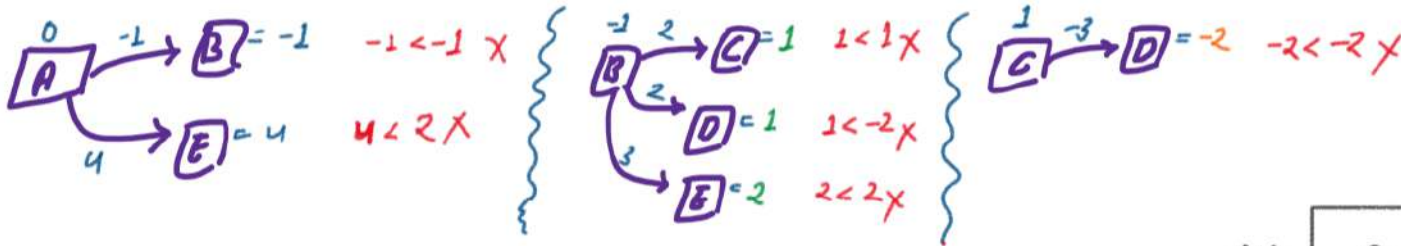
$N = \text{Total Nodes} = 5$   
 $\text{Relaxation} = N - 1 = 4$



No updation until  
 Relaxation 3

Key	Value
A	{E, -1}, {E, 4}
B	{E, 3}, {D, 2}, {C, 2}
C	{E, D, -3}
D	{E, 5}, {B, 2}
E	X

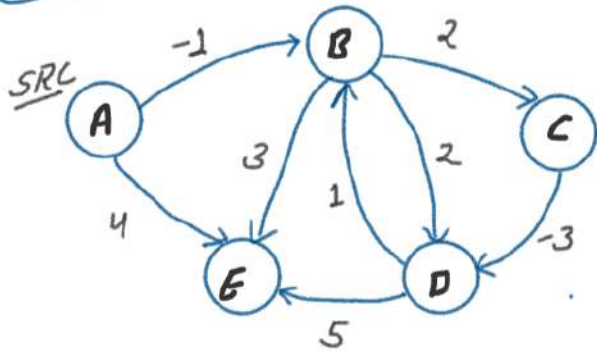
AdjList



dist	0	-1	1	-2	2
	0 (A)	1 (B)	2 (C)	3 (D)	4 (E)

Relaxation 4:

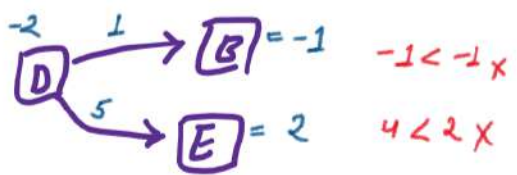
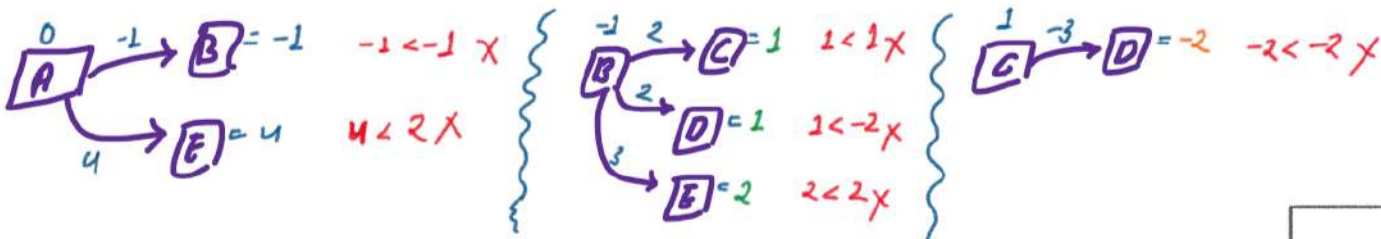
$N = \text{Total nodes} = 5$   
 $\text{Relaxation} = N - 1 = 4$



No updation while relaxation

Key	Value
A	{B, -1}, {E, 4}
B	{C, 2}, {D, 2}, {E, 3}
C	{D, -3}
D	{E, 5}, {B, 2}
E	x

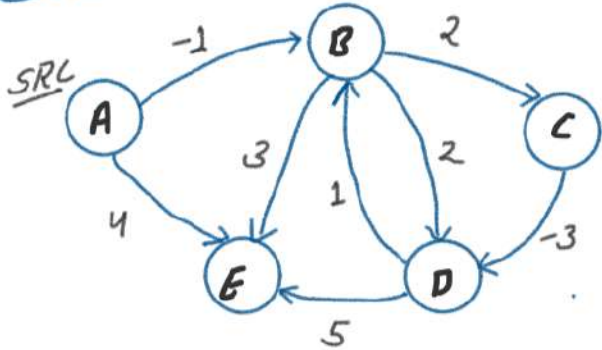
AdjList



dist	0	-1	1	-2	2
	0 (A)	1 (B)	2 (C)	3 (D)	4 (E)

Relaxation 5:

$N = \text{Total Nodes} = 5$   
 $\text{Relaxation} = N - 1 = 4$



STOP THE LOOP  
 due to Relaxation  
 is  $5 - 1 = 4$

Key	Value
A	{B, -1}, {E, 4}
B	{E, 2}, {D, 1}, {C, 3}
C	{D, -3}
D	{E, 5}, {B, 2}
E	x

AdjList

All shortest path from  
 SRC

dist

0	-1	1	-2	2
0 (A)	1 (B)	2 (C)	3 (D)	4 (E)



```

#include<iostream>
#include<vector>
#include<unordered_map>
#include<limits.h>
#include<list>

using namespace std;

class Graph
{
public:
    unordered_map<char, list<pair<char, int>>> adjList;

    void addEdges(char u, char v, int wt, int direction){
        if(direction == 1){
            // Directed Graph
            adjList[u].push_back({v,wt});
        }
        else{
            // Undirected Graph
            adjList[u].push_back({v,wt});
            adjList[v].push_back({u,wt});
        }
    }

    void bellManFord(int n, char src) {...}
};

int main(){
    Graph g;
    g.addEdges('A','B',-1,1);
    g.addEdges('A','E',4,1);
    g.addEdges('B','C',2,1);
    g.addEdges('B','D',2,1);
    g.addEdges('B','E',3,1);
    g.addEdges('C','D',-3,1);
    g.addEdges('D','E',5,1);
    g.addEdges('D','B',1,1);
    g.printAdjList();

    int n = 5;
    char src = 'A';
    g.bellManFord(n, src);

    return 0;
}

```

```

void bellManFord(int n, char src) {
    // Initial state
    vector<int> dist(n, INT_MAX);
    dist[src-'A'] = 0;

    // Step 1: Relaxation N-1 times to get the all shortest path from single node.
    for(int i=1; i<n; i++){
        // Traverse on entire edge list
        for(auto a: adjList){
            for(auto b: a.second){
                char u = a.first;
                char v = b.first;
                int wt = b.second;
                if(dist[u-'A'] != INT_MAX && dist[u-'A'] + wt < dist[v-'A']){
                    dist[v-'A'] = dist[u-'A'] + wt;
                }
            }
        }
    }

    // Step 2: One more relaxation after N-1 times to detect the negative cycle
    // 1 time relaxation
    bool anyUpdate = false;
    for(auto a: adjList){
        for(auto b: a.second){
            char u = a.first;
            char v = b.first;
            int wt = b.second;
            if(dist[u-'A'] != INT_MAX && dist[u-'A'] + wt < dist[v-'A']){
                anyUpdate = true;
                break;
            }
        }
    }

    // Expected Output
    if(anyUpdate == true){
        cout << "Negative Cycle Present" << endl;
    }
    else{
        cout << "No Negative Cycle Present in Graph" << endl;
        cout << "Printing Distance Array: ";
        for(auto i: dist){
            cout << i << " ";
        }
        cout << endl;
    }
}

```

T.C. = ?

S.C. = ?