

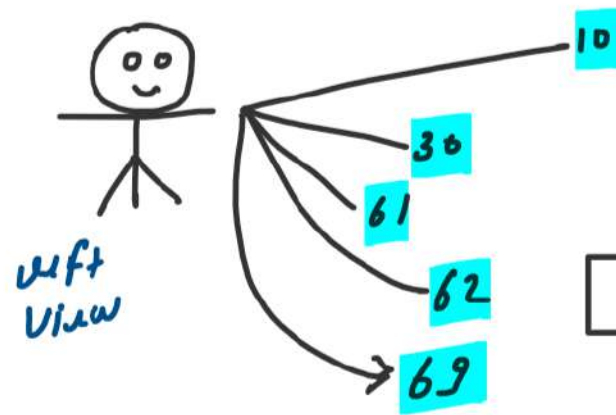
4/12/2023

BINARY TREE

CLASS - 3

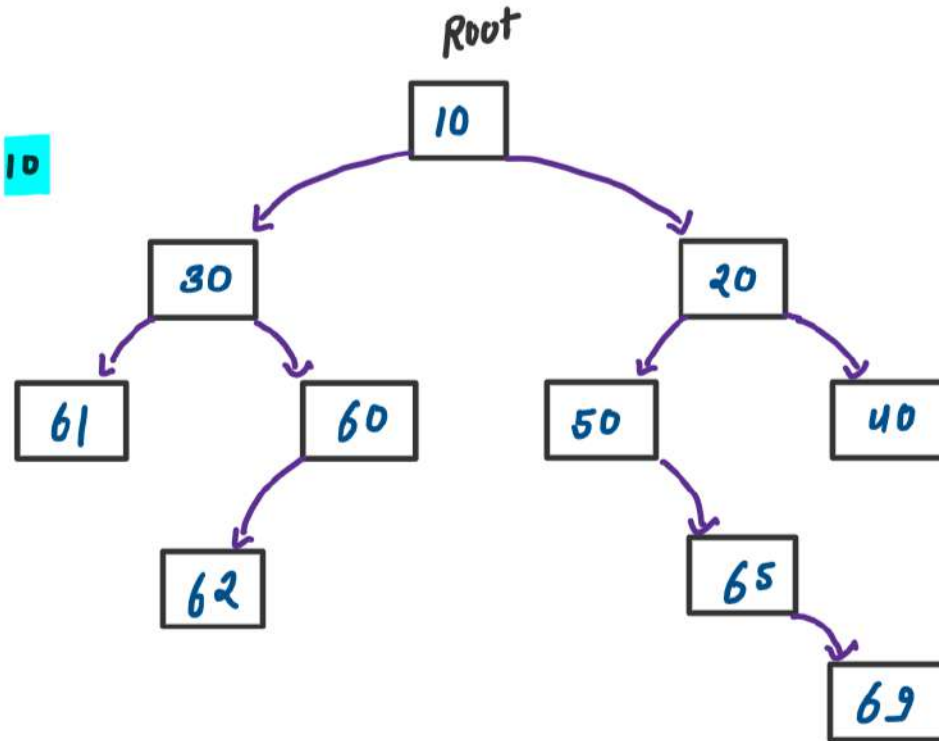
1. Left View of Binary Tree

Ex1



Output

10	30	61	62	69
----	----	----	----	----



LVL-0

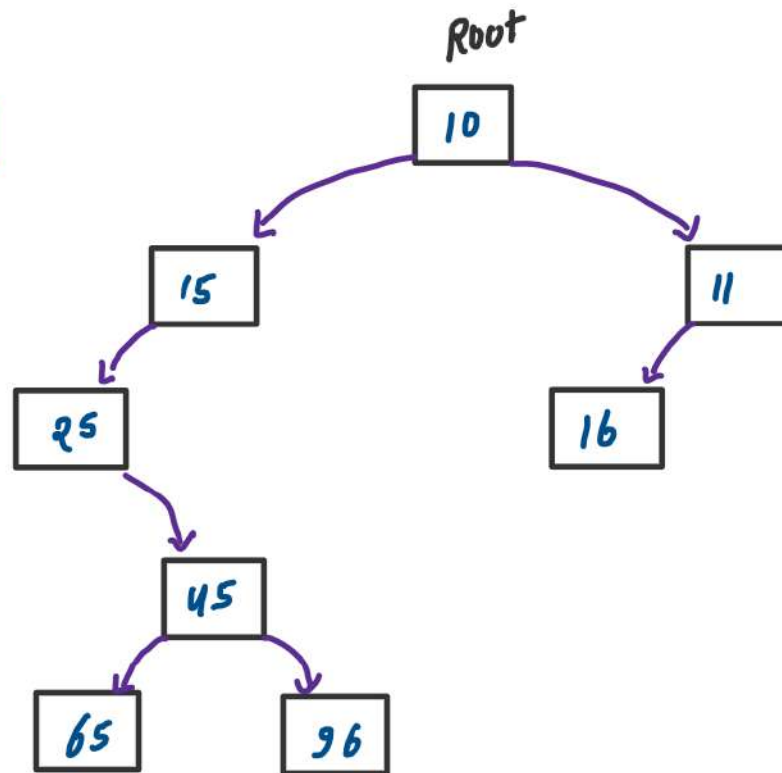
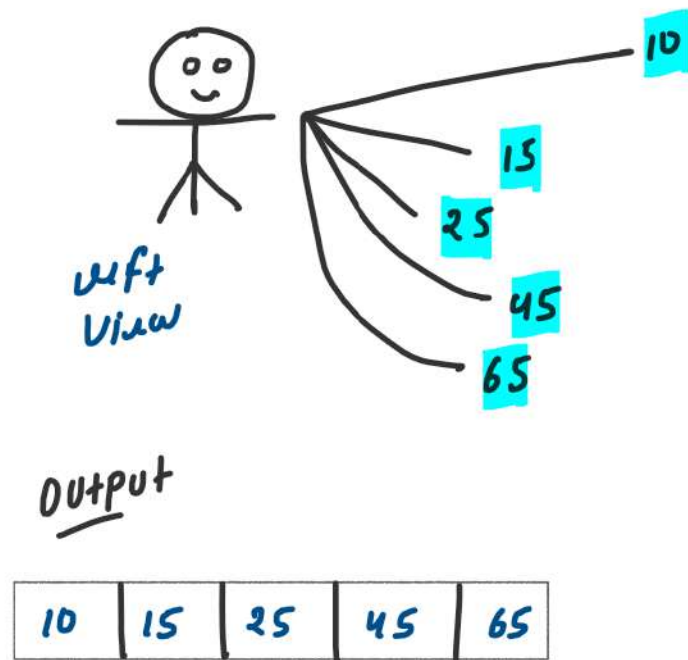
LVL-1

LVL-2

LVL-3

LVL-4

Ex2



LVL-0

LVL-1

LVL-2

LVL-3

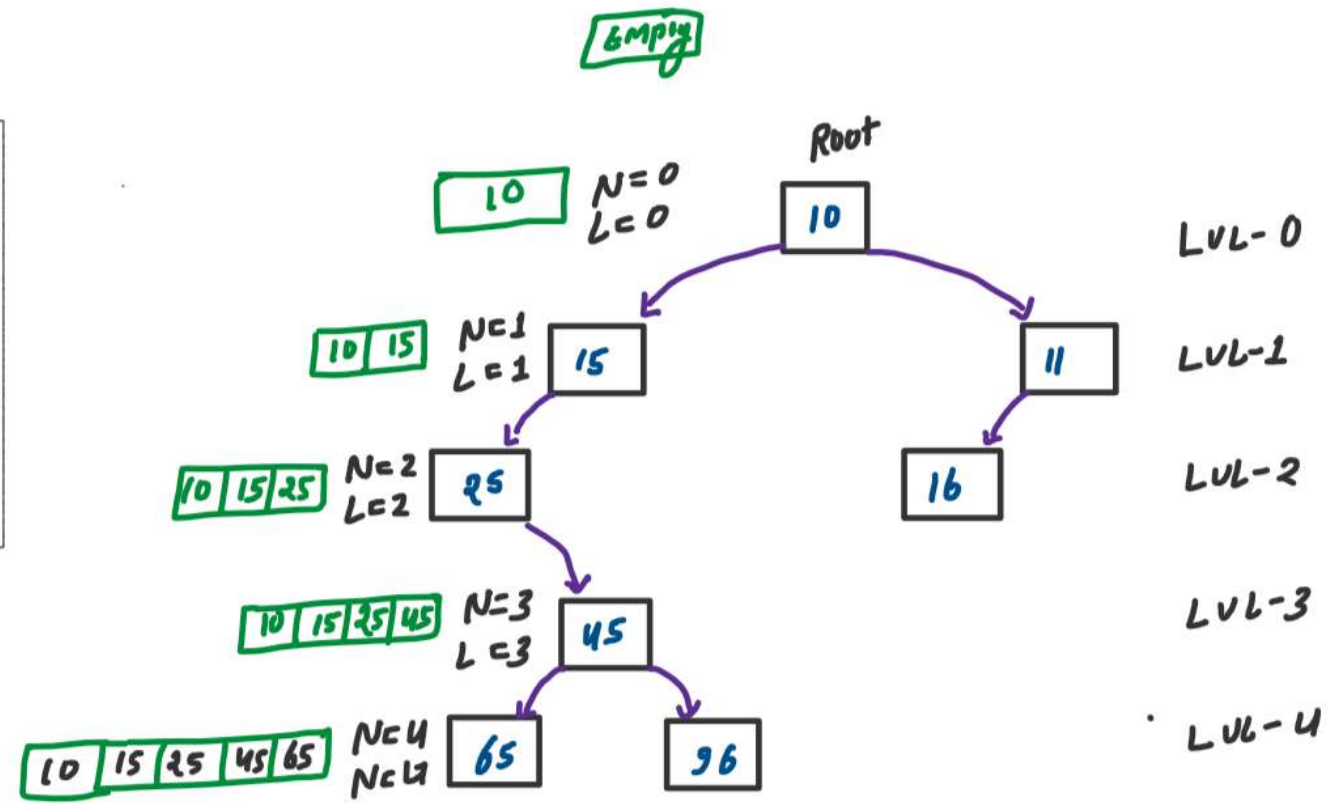
LVL-4

Logic Building

Vector<int> LeftView

10	15	25	45	65
0	1	2	3	4

$N = \text{LeftView.size}()$



```

// PROBLEM 01: Left view of binary tree
void printLeftView(Node* root, int level, vector<int> &leftView){
    // Base case
    if(root == NULL){
        return;
    }

    // 1 case hum solve kar lenge
    if(level == leftView.size()){
        leftView.push_back(root->data);
    }

    // Ab recursion solve kar lega
    printLeftView(root->left, level+1, leftView);
    printLeftView(root->right, level+1, leftView);
}

/*
Binary Tree Input: 10 15 25 -1 45 65 -1 -1 96 -1 -1 -1 11 16 -1 -1 -1

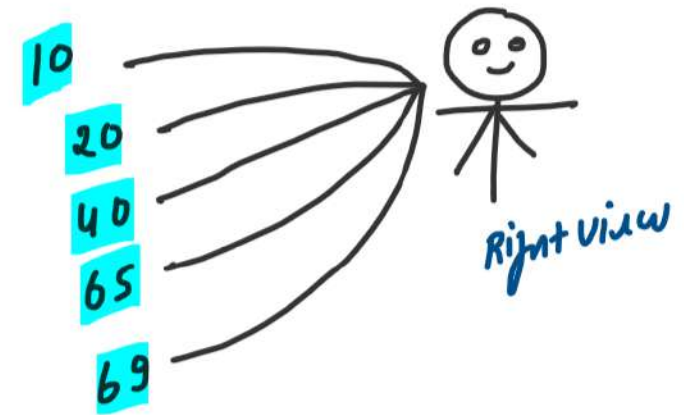
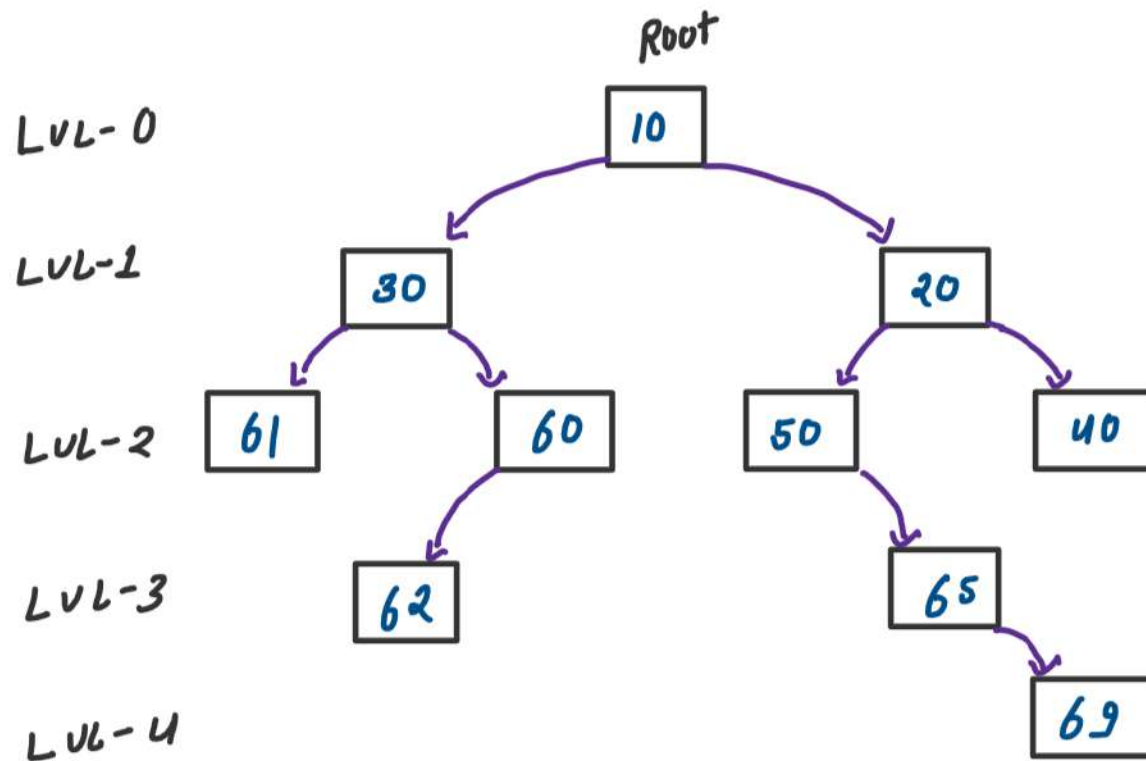
OUTPUT:
Left view:
10 15 25 45 65
*/

```

Time Complexity: $O(N)$,
where N is total number of nodes in binary tree

Space Complexity: $O(L)$,
where L is maximum number of nodes in the level of binary tree

2. Right View of Binary Tree



Output

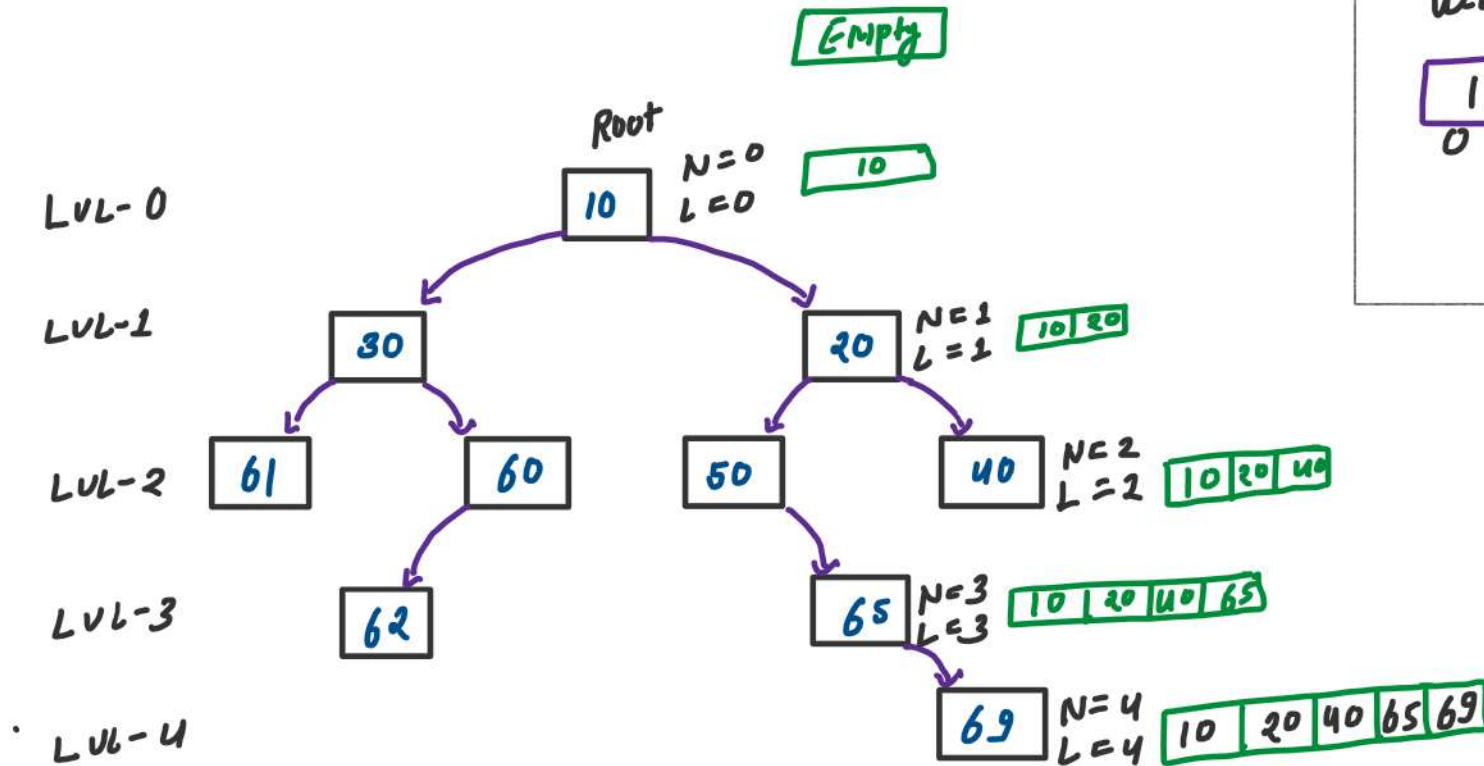
10	20	40	65	69
----	----	----	----	----

Logic Building

vector<int> RightView

10	20	40	65	69
0	1	2	3	4

$N = \text{RightView.size}()$



```

// PROBLEM 02: Right view of binary tree
void printRightView(Node* root, int level, vector<int> &rightView){
    // Base case
    if(root == NULL){
        return;
    }

    // 1 case hum solve kar lenge
    if(level == rightView.size()){
        rightView.push_back(root->data);
    }

    // Ab recursion solve kar lega
    printRightView(root->right, level+1, rightView);
    printRightView(root->left, level+1, rightView);
}

/*
Binary Tree Input:
10 30 61 -1 -1 60 62 -1 -1 -1 20 50 -1 65 -1 69 -1 -1 40 -1 -1

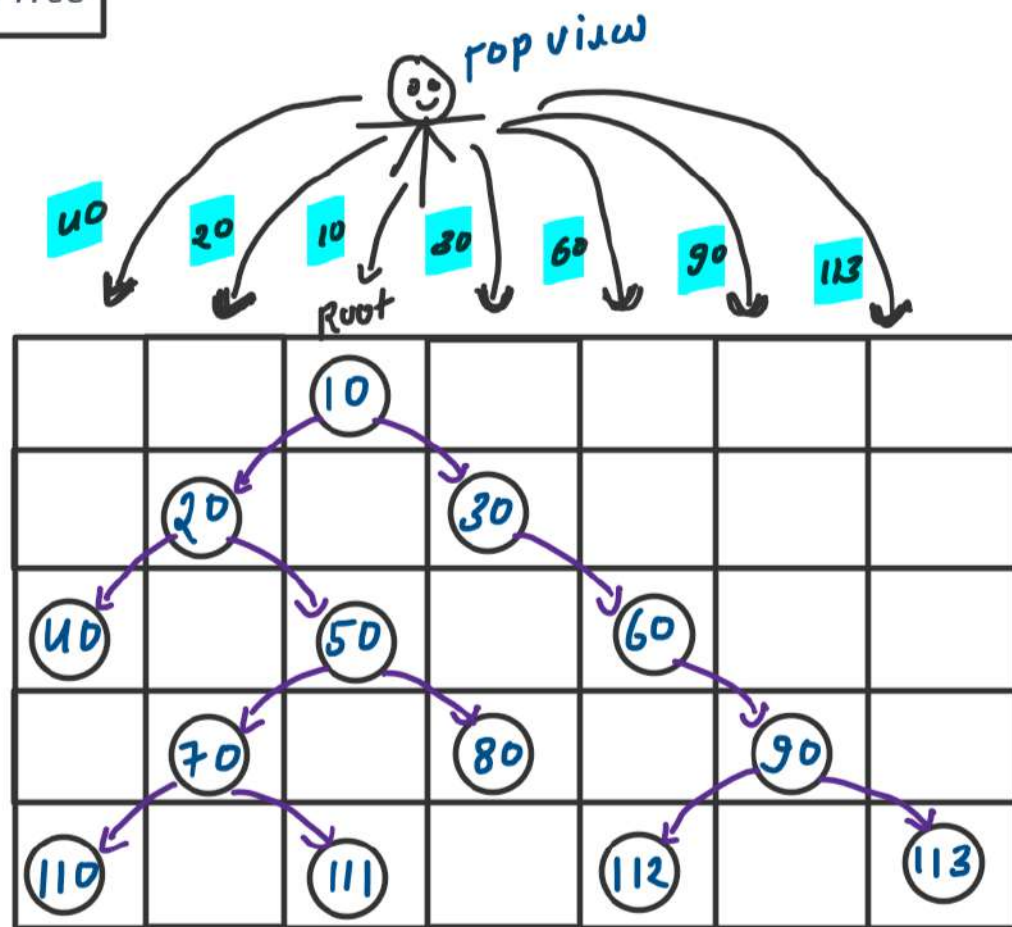
OUTPUT:
Right view:
10 20 40 65 69
*/

```

Time Complexity: $O(N)$,
where N is total number of nodes in binary tree

Space Complexity: $O(L)$,
where L is maximum number of nodes in the level of binary tree

3. Top View of Binary Tree



Output

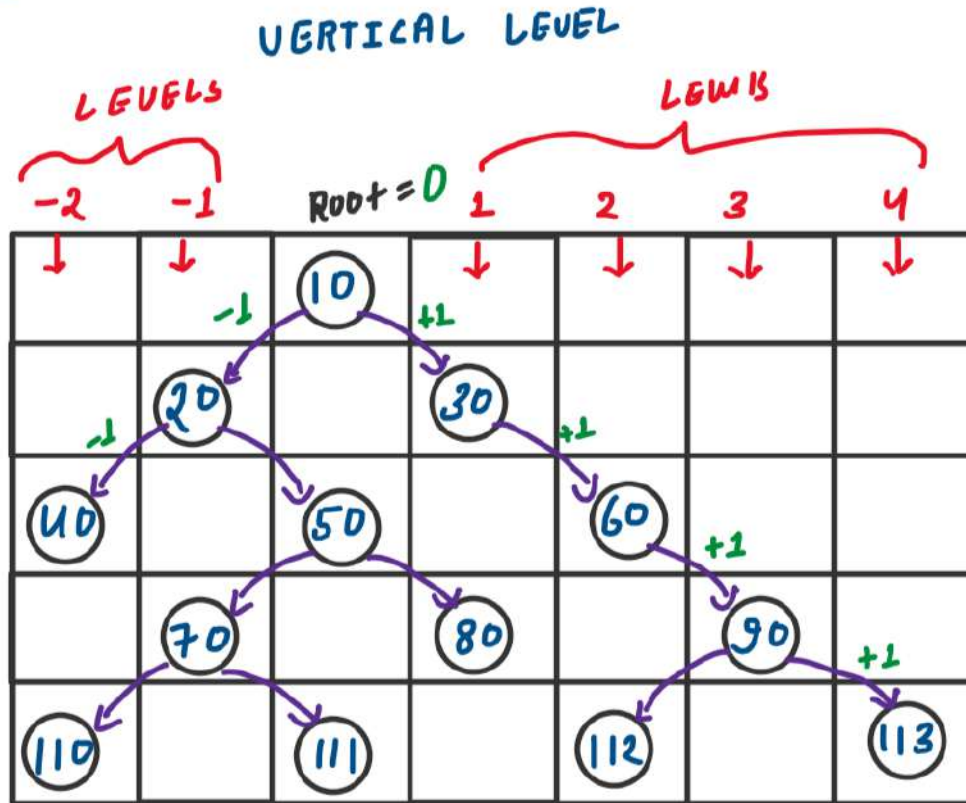
40	20	10	30	60	90	113
----	----	----	----	----	----	-----

Logic Building

Why use map
 ↳ Ans To Print
 The Top view
 Node data
 According to
 Horizontal distn
 wist of Node
 in order

LT4

-2	-1	0	1	2	3	4	LEVEL
40	20	10	30	60	90	113	Node Data
							map



MAP

Node → Data	Horizontal Distance TO Node
----------------	-----------------------------------

QUEUE

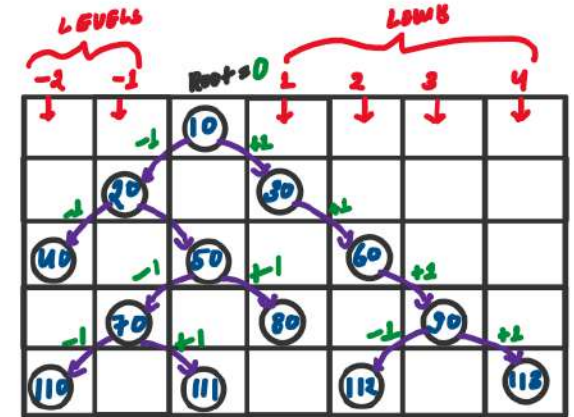
pair	
Node*	LEVEL

map <int, int> hdToNodeMP;
 queue<pair<Node*, int>> q;
 Initially
 ↳ q.push (max_pair (root, 0));

DRY RUN

map

Data	LEVEL
40	-2
20	-1
10	0
30	+1
60	2
90	3
113	4



NUCH

Queue

	Front	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT
	10	20	30	40	50	60	70	80	90	110	111	112	113
	0	-1	1	-2	0	2	-1	1	3	-2	0	2	4

LEVEL

0-1	0+1	-1-1	-1+1	1+1	0-1	0+1	2+1	-1-1	-1+1	3-1	3+1
-----	-----	------	------	-----	-----	-----	-----	------	------	-----	-----

STEP 1 Start null and level into queue and pop

STEP 2 Fetch level from queue and start data and level into map when unique level will occur.

now queue is empty - STOP

```
// PROBLEM 03: Top view of binary tree

void printTopView(Node* root){
    map<int, int> hdToNodeMap; // < level, data >
    queue<pair<Node*, int>> q;
    // Initially store the root node and level 0 into queue
    q.push(make_pair(root,0));

    while(!q.empty()){
        // Fetch front from queue and pop
        pair<Node*, int> front = q.front();
        q.pop();

        Node* frontNode = front.first;
        int level = front.second;

        // Store frontNode->data and level into the map when unique level will occur
        if(hdToNodeMap.find(level) == hdToNodeMap.end()){
            hdToNodeMap[level] = frontNode->data;
        }

        // Agar root ka left node exist krta hai to queue me push krdo with level-1
        if(frontNode->left != NULL){
            q.push(make_pair(frontNode->left, level+1));
        }
        // Agar root ka right node exist krta hai to queue me push krdo with level+1
        if(frontNode->right != NULL){
            q.push(make_pair(frontNode->right, level+1));
        }
    }

    cout<< "Printing Top View: " << endl;
    for(auto data: hdToNodeMap){
        cout<< data.second << " ";
    }
}

/*
Binary Tree Input:
10 20 40 -1 -1 50 70 110 -1 -1 111 -1 -1 80 -1 -1 30 -1 60 -1 9 112 -1 -1 113 -1 -1

OUTPUT:
10
20 30
40 50 60
70 80 9
110 111 112 113
Printing Top View:
40 20 10 30 60 9 113
*/
```

Time Complexity: $O(N)$, where N is total number of nodes in binary tree

Space Complexity: $O(N)$, where

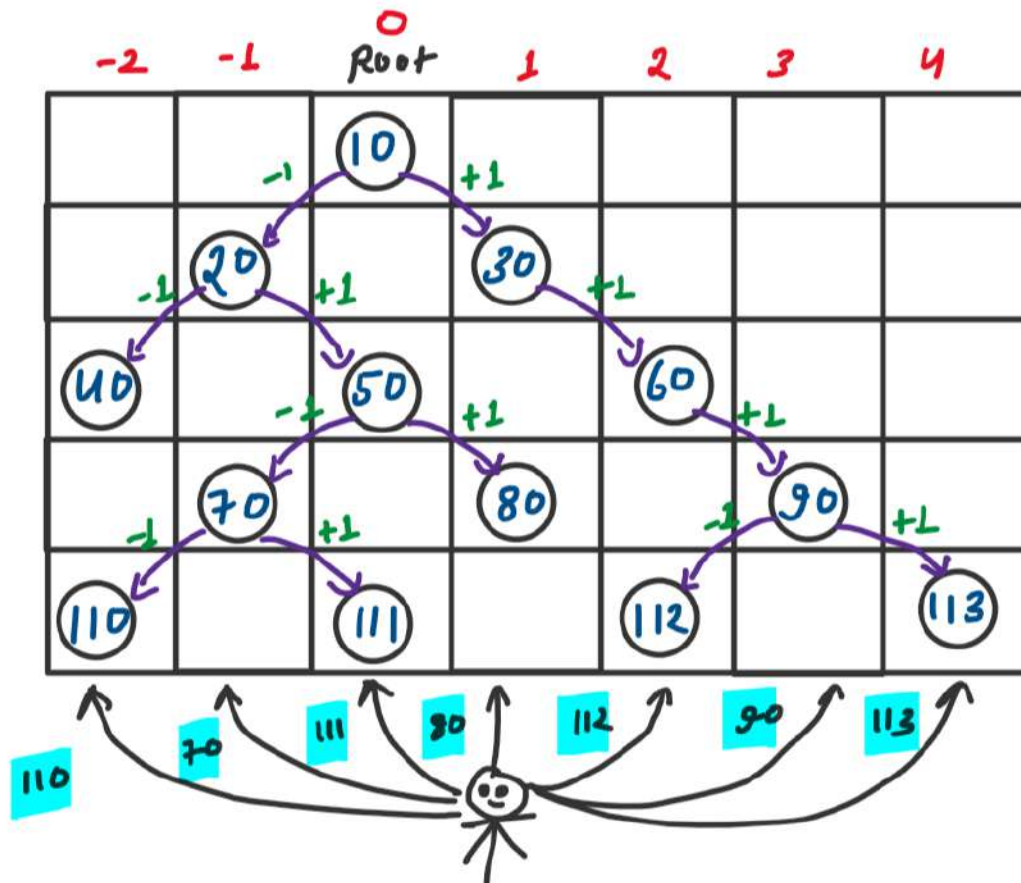
Case I - considering a skewed tree:

space complexity is $O(N)$

Case II - considering now skewed tree:

space complexity is $O(W)$, Where W is maximum width of the tree

4. Bottom View of Binary Tree



Output

110	70	111	80	112	90	113
-----	----	-----	----	-----	----	-----

STEP 1 Start Node and insert into queue and pop

STEP 2 Fetch data from queue and store data and insert into map

↪ Overwrite the exist entry of map
jab same entry occurs to Rahi Ho.

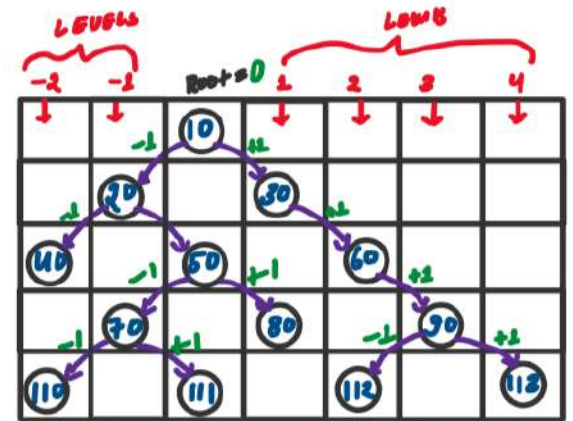
DRY RUN

map

Data	LEVEL
110	-2
70	-1
111	0
80	+1
112	2
90	3
113	4

second

first



	Front	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT	FRONT
Node	10	20	30	40	50	60	70	80	90	110	111	112	113
Level	0	-1	1	-2	0	2	-1	1	3	-2	0	2	4
		0-1	0+1	-1-1	-1+1	1+1	0-1	0+1	2+1	-1-1	-1+1	3-1	3+1

output [110 70 111 80 112 90 113]

now queue is empty - STOP

```
// PROBLEM 04: Bottom view of binary tree

void printBottomView(Node* root){
    map<int, int> hdToNodeMap; // < level, data >
    queue<pair<Node*, int>> q;
    // Initially store the root node and level 0 into queue
    q.push(make_pair(root,0));

    while(!q.empty()){
        // Fetch front from queue and pop
        pair<Node*, int> front = q.front();
        q.pop();

        Node* frontNode = front.first;
        int level = front.second;

        // OVERWRITE: Store frontNode->data and level into the map
        hdToNodeMap[level] = frontNode->data;

        // Agar root ka left node exist krta hai to queue me push krdo with level-1
        if(frontNode->left != NULL){
            q.push(make_pair(frontNode->left, level+1));
        }
        // Agar root ka right node exist krta hai to queue me push krdo with level+1
        if(frontNode->right != NULL){
            q.push(make_pair(frontNode->right, level+1));
        }
    }

    cout<< "Printing Bottom View: " << endl;
    for(auto data: hdToNodeMap){
        cout<< data.second << " ";
    }
}

/*
Binary Tree Input:
10 20 40 -1 -1 50 70 110 -1 -1 111 -1 -1 80 -1 -1 30 -1 60 -1 9 112 -1 -1 113 -1 -1

OUTPUT:
10
20 30
40 50 60
70 80 9
110 111 112 113
Printing Bottom View:
110 70 111 80 112 9 113
*/
```

Time Complexity: $O(N)$, where N is total number of nodes in binary tree

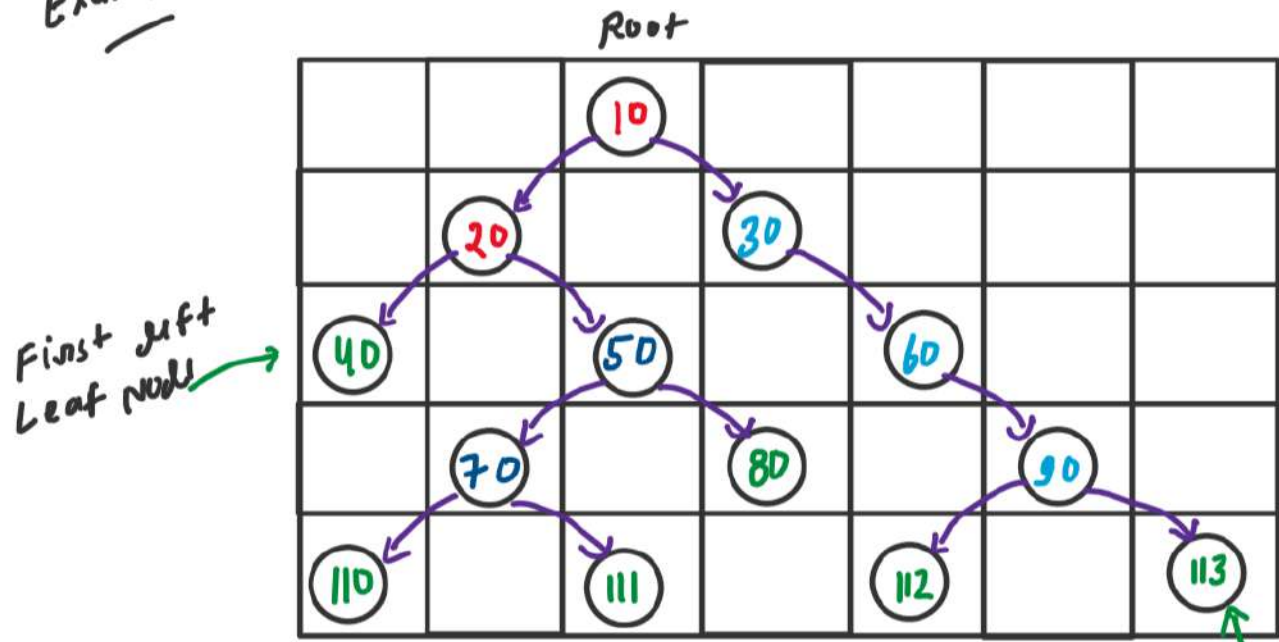
Space Complexity: $O(N)$, where

*Case I - considering a skewed tree:
space complexity is $O(N)$*

*Case II - considering now skewed tree:
space complexity is $O(W)$, Where W is maximum width of the tree*

5. Boundary Traversal of Binary Tree

Example: 01



Output

10 20 40 110 111 80
112 113 90 60 30

Print A \Rightarrow Left Node Boundary

Print B \Rightarrow Leaf Node Boundary

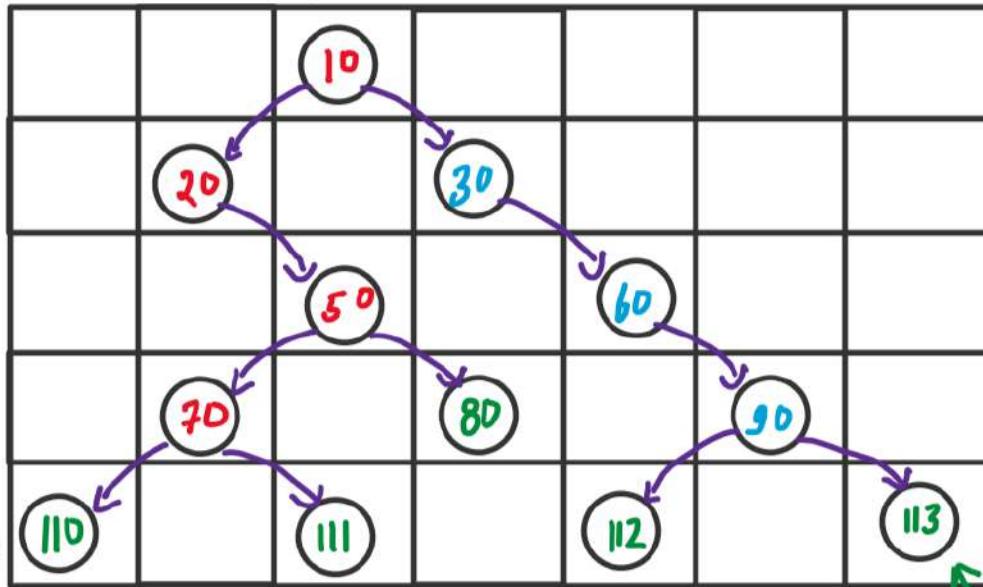
Print C \Rightarrow Right Node Boundary

Print (A) Print All node but for left node
jo EK leaf node hoga

First Right
Leaf Node

Example: 02

Root



First Left Leaf Node

Print (A) Print All nodes but for left nodes jo EK leaf node hoga

Output

10	20	50	70	110	111
80	112	113	90	60	30

Print A \Rightarrow Left Node Boundary

Print B \Rightarrow Leaf Node Boundary

Print C \Rightarrow Right Node Boundary

First Right Leaf Node

FUNCTION (A)

```
void printLeftBoundary(Node* Root) {  
    if (Root == NULL) return;  
    if (Root->Left == NULL &&  
        Root->Right == NULL)  
        return;  
    cout << Root->data << " ";  
}
```

First Left Leaf Node HANE par
Function se Bahar ho jayega ---

EX1 par DRY
RUN KARO

```
{  
    if (Root->Left != NULL)  
        printLeftBoundary(Root->Left);  
    else if (Root->Right != NULL)  
        printLeftBoundary(Root->Right);  
}
```

EX2 par DRY RUN
KARO

3

o/p 10 20 → EX1

o/p 10 20 50 20 → EX2

FUNCTION (B)

```
void printLeafBoundary (Node* root) {
```

```
    if (root == NULL) return;
```

```
    if (root->left == NULL &&  
        root->right == NULL) {
```

```
        cout << root->data << " ";
```

```
    }
```

```
    •
```

```
}
```

→ jab node leaf hai tabhi
print karna hai

```
{ printLeafBoundary (root->left);  
  printLeafBoundary (root->right);
```

o/p 40 110 111 80 112 113 → Ex 1

o/p 110 111 80 112 113 → Ex 2

FUNCTION (C)

```
void printRightBoundary ( Node* Root ) {
```

```
    if ( Root == NULL ) return;
```

```
    if ( Root->Left == NULL &&  
        Root->Right == NULL ) {  
        return;
```

```
    }
```

```
    if ( Root->Right != NULL )
```

```
        printRightBoundary ( Root->Right );
```

```
    else if ( Root->Left != NULL )
```

```
        printRightBoundary ( Root->Left );
```

```
}
```

First Right Leaf Node AANE par
Function se Bahar Ho jao....

cout << Root->data << " ";

o/p

90	60	30
----	----	----

 → Ex 1

o/p

90	60	30
----	----	----

 → Ex 2

}

Function Boundary Traversal

```
void boundaryTraversal(Node* Root) {  
    if (Root == NULL) return;
```

```
    Print Left Boundary (Root);
```

```
    Print Leaf Boundary (Root);
```

```
    Print Right Boundary (Root);  
}
```

WRONG

→ Output

Ex1

10	20	40	110	111	80
112	113	90	60	30	10

WRONG

→ Output

Ex2

10	20	50	70	110	111
80	112	113	90	60	30
10					


```

void boundaryTraversal(Node* Root) {
    if (Root == NULL) return;

    PrintLeftBoundary(Root);
    PrintLeafBoundary(Root);
    if (Root->right != NULL) {
        PrintRightBoundary(Root->right);
    }
    else {
        PrintRightBoundary(Root->left);
    }
}

```

connect
→ Output

Ex1

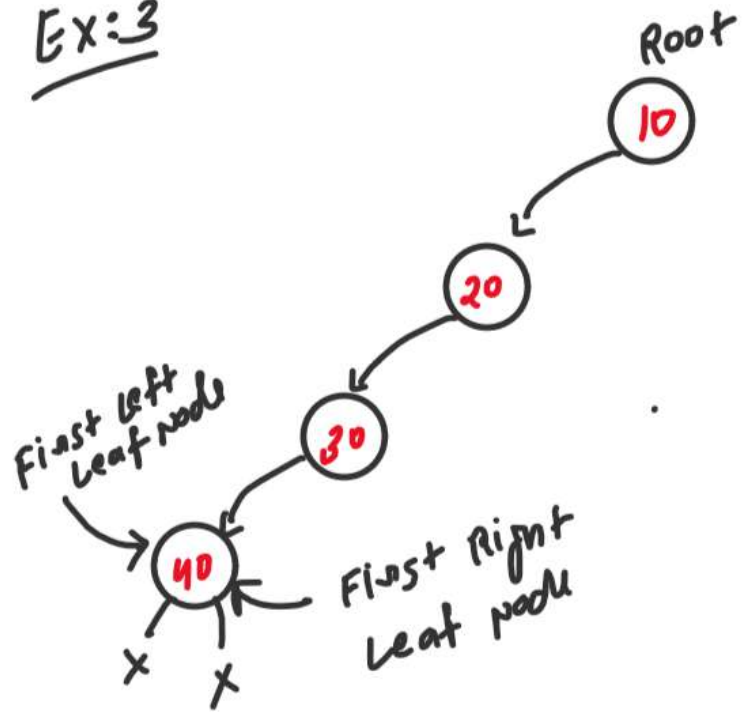
10	20	40	110	111	80
112	113	90	60	30	

connect
→ Output

Ex2

10	20	50	70	110	111
80	112	113	90	60	30

Ex: 3



Left \Rightarrow 10 20 30

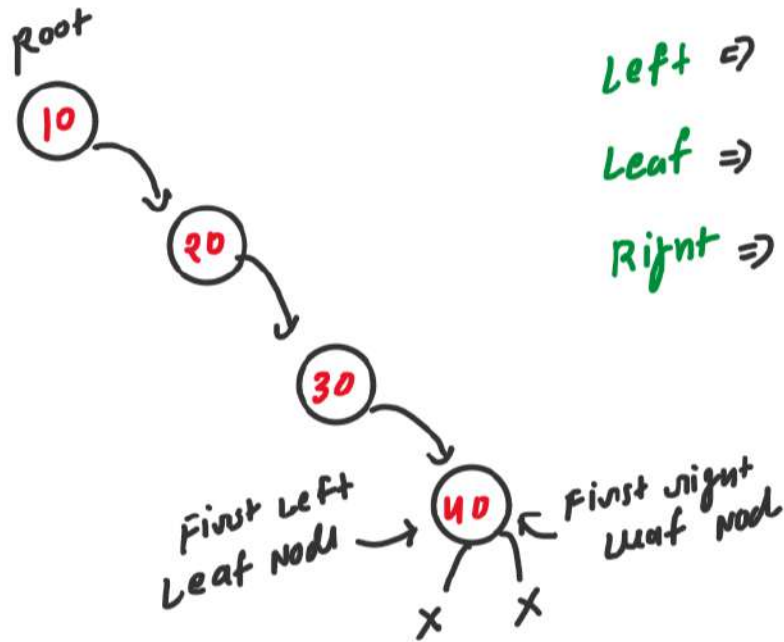
Leaf \Rightarrow 40

Right \Rightarrow 30 20

Output

10 20 30 40 30 20

Ex: 4



Left \Rightarrow 10 20 30

Leaf \Rightarrow 40

Right \Rightarrow 30 20

Output

10 20 30 40 30 20

A

```

void printLeftBoundary(Node* root){
    // Base Case
    if(root == NULL){
        return;
    }

    // First Left Leaf Node aane par function se bahar ho jao
    if(root->left == NULL && root->right == NULL){
        return;
    }

    cout<< root->data << " ";

    if(root->left != NULL){
        printLeftBoundary(root->left);
    }
    else if(root->right != NULL){
        printLeftBoundary(root->right);
    }
}

```

C

```

void printRightBoundary(Node* root){
    // Base case
    if(root == NULL){
        return;
    }

    // Jab first right leaf node aa jaye to function se bahar ho jao
    if(root->left == NULL && root->right == NULL){
        return;
    }

    if(root->right != NULL){
        printRightBoundary(root->right);
    }
    else if(root->left != NULL){
        printRightBoundary(root->left);
    }

    cout<< root->data << " ";
}

```

B

```

void printLeafBoundary(Node* root){
    // Base case
    if(root == NULL){
        return;
    }

    // Jab-2 leaf node ayega tabhi print karna hai
    if(root->left == NULL && root->right == NULL){
        cout<< root->data << " ";
    }

    printLeafBoundary(root->left);
    printLeafBoundary(root->right);
}

```

boundaryTraversal(Node* root){

```

    if(root == NULL){
        return;
    }

    printLeftBoundary(root);
    printLeafBoundary(root);
    if(root->right != NULL){
        printRightBoundary(root->right);
    }
    else if(root->left != NULL){
        printRightBoundary(root->left);
    }
}

```

Handwritten annotations: Arrows labeled A, B, and C point to the recursive calls in the code. Arrow A points to `printLeftBoundary(root);`, arrow B points to `printLeafBoundary(root);`, and arrow C points to the `if` block containing `printRightBoundary` calls.

Time Complexity and Space Complexity: $O(N)$,

Where N is number of nodes in binary tree