

26/02/2024

# GRAPHS CLASS - 5

## 1. Course Schedule (Leetcode-207)

### Problem Statement:

There are a total of numCourses courses you have to take, labeled from 0 to numCourses - 1. You are given an array prerequisites where prerequisites[i] = [ai, bi] indicates that you must take course bi first if you want to take course ai.

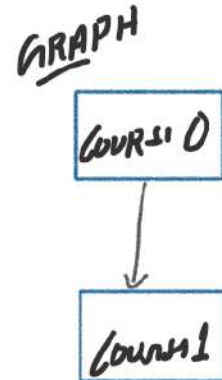
For example, the pair [0, 1], indicates that to take course 0 you have to first take course 1.

Return **true** if you can finish all courses. Otherwise, return **false**.

### Example 1:

Input: numCourses = 2, prerequisites = [[1,0]]

Output: true



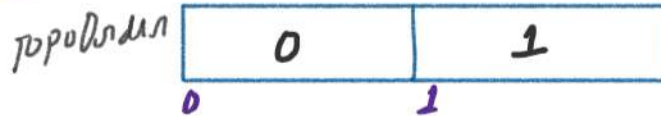
COURSE 1 ke liye  
pahle COURSE 0 lena ho.  
it means COURSE 1 depends on  
COURSE 0 so we can apply topological  
sorting to store the dependency  
order of course.

Example 1:

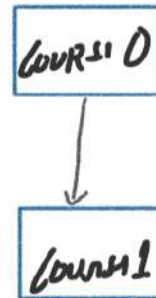
Input: numCourses = 2, prerequisites = [[1,0]]

Output: true

**STEP 1**



GRAPH



**STEP 2**

if (topoOrder.size() == numCourses) {

    // All courses are finished  
    → return true

}

else {

    // All courses are not finished

}

STEP 1

create topoOrder  
using topological sort

STEP 2

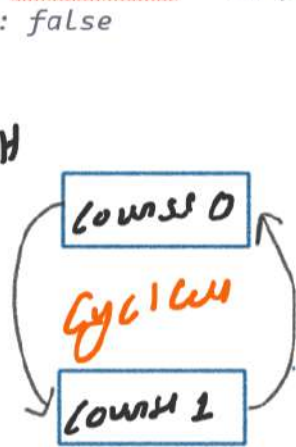
check All courses are  
finished or not

Example 2:

Input: numCourses = 2, prerequisites = [[1,0],[0,1]]

Output: false

GRAPH



if you want to take course 1. you have to finish course 0 first.  
if you want to take course 0. you have to finish course 1 first.

STEP 1

course 0's Indegree = 1  
course 1's Indegree = 1

STEP 2

topoList.size != numCourses  
↳ return false

topoOrder is always empty because topological sorting algo. push a node jiski indegree 0 ho. It means wo node independent hai. But yaha par dono course 0 and course 1 ek dusre par he depend kar rahi hai. jiska matlab courses number finished.

```
// 1. Course Schedule (Leetcode-207)

class Solution {
public:
    void topoSortUsingBFS(int n, vector<int> &topoOrder, unordered_map<int, list<int>> &adjList){...}

    bool canFinish(int numCourses, vector<vector<int>>& prerequisites) {
        // Create adjList first
        unordered_map<int, list<int>> adjList;
        for(auto i: prerequisites){
            int u = i[0];
            int v = i[1];
            adjList[v].push_back(u);
        }

        // Step 1: Create the topo order for checking all courses are finished or not
        vector<int> topoOrder;
        topoSortUsingBFS(numCourses, topoOrder, adjList);

        // Step 2: Check all courses are finished or not
        if(topoOrder.size() == numCourses){
            // Cycle does not present: means all course are finished
            return true;
        }
        else{
            // Cycle does present: means all course are never finished
            return false;
        }
    }
};
```

STEP 2

```
void topoSortUsingBFS(int n, vector<int> &topoOrder, unordered_map<int, list<int>> &adjList){
    queue<int> q;
    unordered_map<int, int> indegree;

    // Step 1: Initialize the indegree
    for(auto i: adjList){
        for(auto neighbour: i.second){
            indegree[neighbour]++;
        }
    }

    // Step 2: push all nodes jinki indegree zero hai
    for(int node = 0; node < n; node++){
        if(indegree[node] == 0){
            q.push(node);
        }
    }

    // Step 3: BFS on queue to print the order dependency wise
    while(!q.empty()){
        auto frontNode = q.front();
        q.pop();
        topoOrder.push_back(frontNode);

        for(auto neighbour: adjList[frontNode]){
            indegree[neighbour]--;

            // check neighbour node indegree is zero or not
            if(indegree[neighbour] == 0){
                q.push(neighbour);
            }
        }
    }
}
```

STEP 1

## 2. Course Schedule II (Leetcode-210)

### Problem Statement:

There are a total of numCourses courses you have to take, labeled from 0 to numCourses - 1. You are given an array prerequisites where prerequisites[i] = [ai, bi] indicates that you must take course bi first if you want to take course ai.

For example, the pair [0, 1], indicates that to take course 0 you have to first take course 1.



Return the **ordering of courses** you should take to finish all courses. If there are many valid answers, return any of them. If it is impossible to finish all courses, return an empty array.

Example 1:

Input: numCourses = 2, prerequisites = [[1,0]]

Output: [0,1]

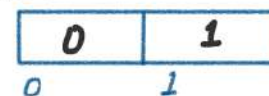
GRAPH



0's Indegree = 0

1's Indegree = 1 → It means 1 depends on 0.

STEP 1 Create topological order



STEP 2 Return topological order when all courses are finished otherwise return empty array

→ return 

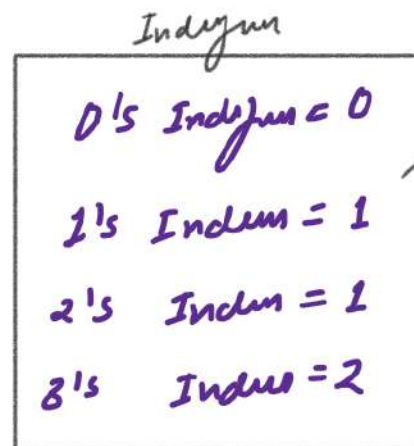
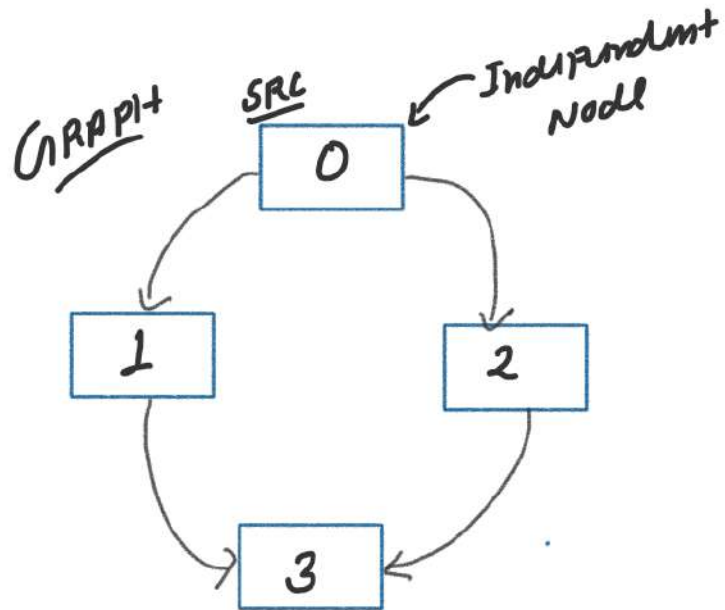
0	1
---	---



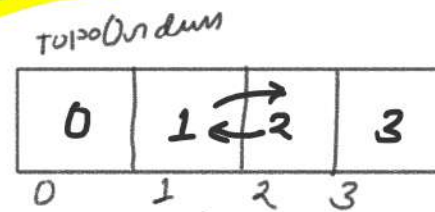
Example 2:

Input: numCourses = 4, prerequisites =  $[[1,0],[2,0],[3,1],[3,2]]$

Output:  $[0,2,1,3]$



**STEP 1**



↑ This is a dependency order

**STEP 2**

if (topoOrder.size == numCourses)  
    ↳ return topoOrder;

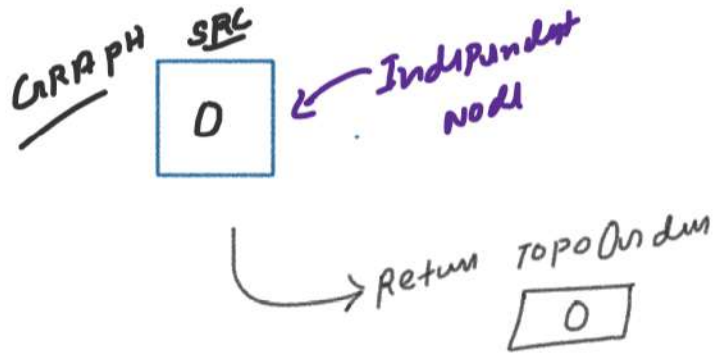
else

    ↳ return {};

Example 3:

Input: numCourses = 1, prerequisites = []

Output: [0]



STEP 2

if (topoOrder.size == numCourses)

↳ return topoOrder;

else

↳ return {};



```
// 2. Course Schedule II (Leetcode-210)
```

```
class Solution {
public:
    void topoSortUsingBFS(int n, vector<int> &topoOrder, unordered_map<int, list<int>> &adjList){...}
    vector<int> findOrder(int numCourses, vector<vector<int>>& prerequisites) {
        // Create adjList first
        unordered_map<int, list<int>> adjList;
        for(auto i: prerequisites){
            int u = i[0];
            int v = i[1];
            adjList[v].push_back(u);
        }

        // Create the topo order for checking all courses are finished or not
        vector<int> topoOrder;
        topoSortUsingBFS(numCourses, topoOrder, adjList);

        if(topoOrder.size() == numCourses){
            // Cycle does not present: means all course are finished
            return topoOrder;
        }
        else{
            // Cycle does present: means all course are never finished
            return {};
        }
    }
};
```

STEP 2

```
void topoSortUsingBFS(int n, vector<int> &topoOrder, unordered_map<int, list<int>> &adjList){
    queue<int> q;
    unordered_map<int, int> indegree;

    // Step 1: Initialize the indegree
    for(auto i: adjList){
        for(auto neighbour: i.second){
            indegree[neighbour]++;
        }
    }

    // Step 2: push all nodes jinki indegree zero hai
    for(int node = 0; node < n; node++){
        if(indegree[node] == 0){
            q.push(node);
        }
    }

    // Step 3: BFS on queue to print the order dependency wise
    while(!q.empty()){
        auto frontNode = q.front();
        q.pop();
        topoOrder.push_back(frontNode);

        for(auto neighbour: adjList[frontNode]){
            indegree[neighbour]--;

            // check neighbour node indegree is zero or not
            if(indegree[neighbour] == 0){
                q.push(neighbour);
            }
        }
    }
}
```

STEP 1

### 3. Path with Minimum Effort (Leetcode-1631)

#### Problem Statement:

You are a hiker preparing for an upcoming hike. You are given **heights**, a 2D array of size **rows** x **columns**, where **heights[row][col]** represents the height of cell (**row**, **col**). You are situated in the top-left cell, (**0**, **0**), and you hope to travel to the bottom-right cell, (**rows-1**, **columns-1**) (i.e., **0-indexed**). You can move **up**, **down**, **left**, or **right**, and you wish to find a route that requires the minimum **effort**.

A route's **effort** is the **maximum absolute difference** in heights between two consecutive **cells** of the route.

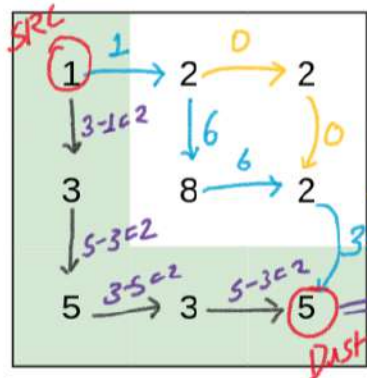
Return the **minimum effort** required to travel from the **top-left cell** to the **bottom-right cell**.

SRC DST

#### Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2



$$\min(3, 6, 2) = 2$$

$$\Rightarrow \max(1, 0, 0, 3) = 3$$

$$\Rightarrow \max(1, 6, 6, 3) = 6$$

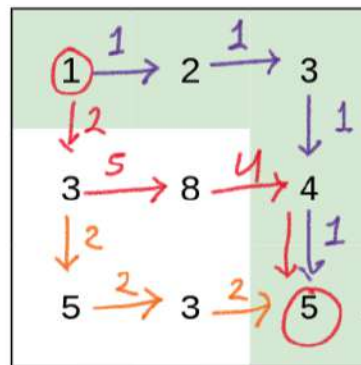
$$\Rightarrow \max(2, 2, 2, 2) = 2$$

$$\Rightarrow 2$$

#### Example 2:

Input: heights = [[1,2,3],[3,8,4],[5,3,5]]

Output: 1



$$\Rightarrow \max(2, 2, 2, 2) = 2$$

$$\Rightarrow \max(2, 5, 4, 2) = 5$$

$$\Rightarrow \max(1, 1, 1, 1) = 1$$

$$\min(2, 5, 1) = 1$$

maximum Effort

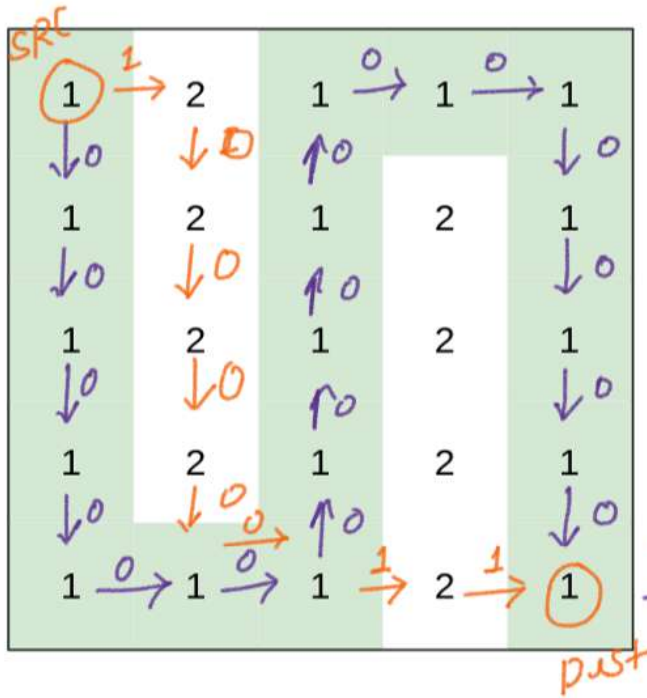
minimum Effort

**Example 3:**

Input: heights =

[[1,2,1,1,1],[1,2,1,2,1],[1,2,1,2,1],[1,2,1,2,1],[1,1,1,2,1]]

Output: 0



maximum Efforts

⇒ maxEffort = 1  
⇒ maxEffort = 0

minimum Efforts

↳ min(1, 0)

= 0

output

Example 1:

Input: heights =  $[[1, 2, 2], [3, 8, 2], [5, 3, 5]]$

Output: 2

*H*

<i>src</i> 1	2	2
3	8	2
5	3	<i>dest</i> 5

Heights

*D*

0	$\infty$	$\infty$
$\infty$	$\infty$	$\infty$
$\infty$	$\infty$	$\infty$

Difference

Initial state

$$\Rightarrow D[0][0] = 0$$

$$\Rightarrow \text{mini-push} \{0, \infty, \infty\}$$

*curDiff* *curX* *curY*



minHeap (mini)

$$H[\text{curX}][\text{curY}] = \text{CELL}$$

*curNode*

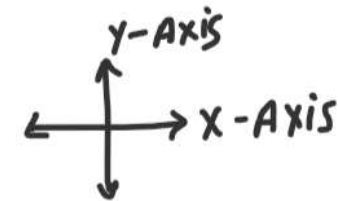
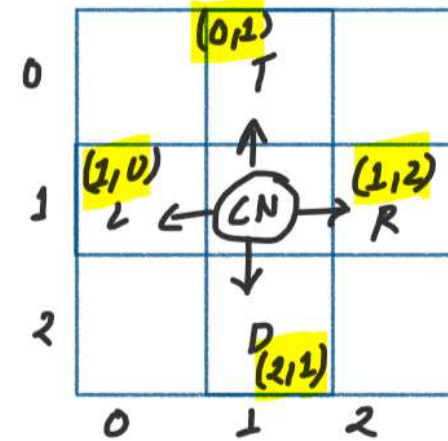
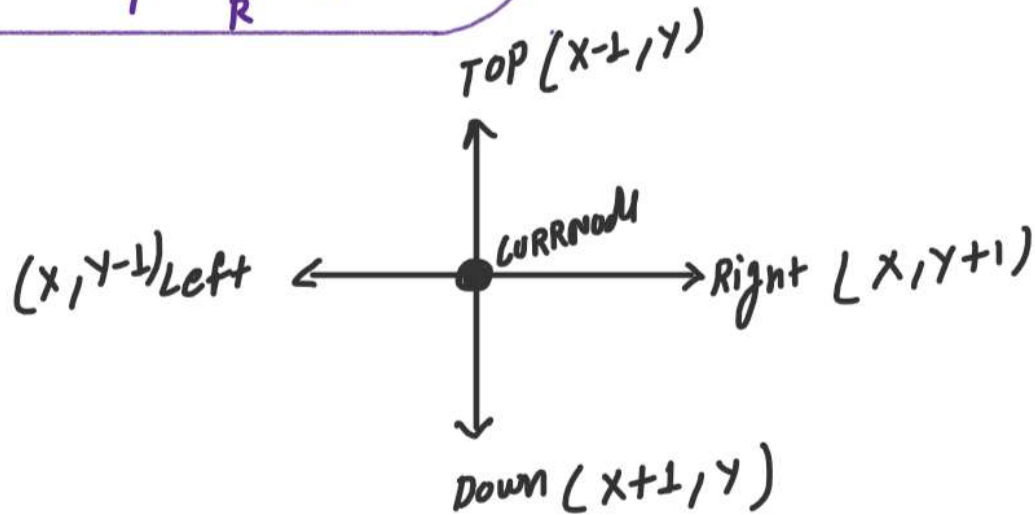
*Source* =  $H[0][0]$   
*Dest* =  $H[\text{row}-1][\text{col}-1]$

Now we can travel to all NBRNode  
i.e. Top, Right, Down, Left

```
int dx[] = {-1, 0, 1, 0};
int dy[] = {0, 1, 0, -1};
```



$-1 \Rightarrow$  X and Y decrement by 1  
 $0 \Rightarrow$  X and Y remains same  
 $1 \Rightarrow$  X and Y Increment by 1





Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

Item 1

	0	1	2
H	1	2	2
1	3	8	2
2	5	3	5

Heights


STEP 3

	0	1	2
D	0	$\infty$	$\infty$
	$\infty$	$\infty$	$\infty$
	$\infty$	$\infty$	$\infty$

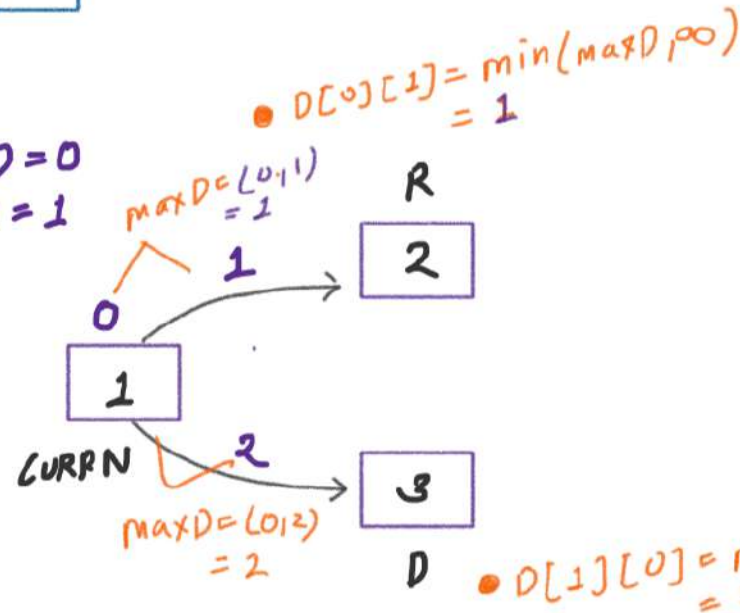
Differences

$D[0][0] < D[0][1]$   
 $0 < \infty \Rightarrow \text{isSafe} = \text{True}$   
 $D[0][0] < D[1][0]$   
 $0 < \infty \Rightarrow \text{True}$

STEP 1 Get TOP pair and pop it  $\Rightarrow \text{currD} = 0$   
 $\Rightarrow \text{currN} = 1$

STEP 2 We can travel to all nbrs  


STEP 3 Check the isSafe to update the diff when isSafe is True



{2, 3}
{1, 2}
<del>{0, 1}</del>

MinHeap (mini)



Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

Item 2

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	$\infty$
1	2	$\infty$	$\infty$
2	$\infty$	$\infty$	$\infty$

Differences

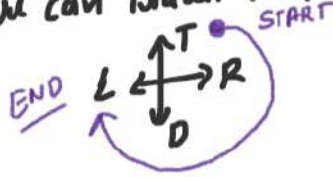
STEP 3

$D[0][1] < D[0][2]$   
 $1 < \infty \Rightarrow \text{isSafe} = \text{True}$

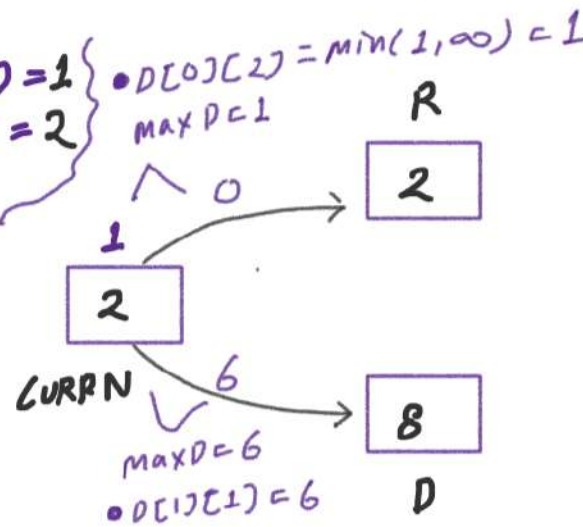
$D[0][1] < D[1][1]$   
 $1 < \infty \Rightarrow \text{True}$

STEP 1 Get TOP pair and pop it  $\Rightarrow \text{currD} = 1$   
 $\Rightarrow \text{currN} = 2$

STEP 2 We can travel to all Nbrs



STEP 3 Check the isSafe to update the diff when isSafe is True



$\{6, 8\}$
$\{1, 2\}$
$\{2, 3\}$
<del><math>\{1, 2\}</math></del>

MinHeap (mini)

Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

**Item 3**

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

**STEP 3**

	0	1	2
0	0	1	1
1	2	6	$\infty$
2	$\infty$	$\infty$	$\infty$

Differences

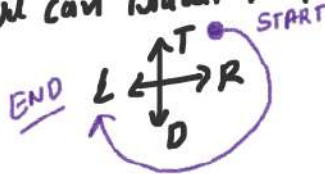
Down  $\Rightarrow (1 < \infty)$  ✓

{1,2}
{6,8}
<del>{1,2}</del>
{2,3}

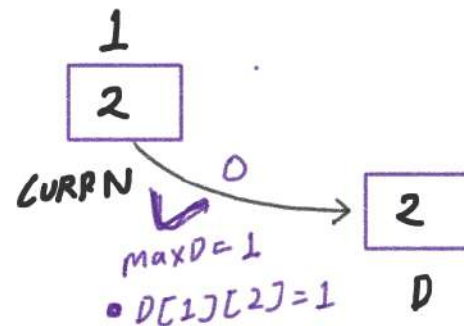
minHeap (mini)

**STEP 1** Get TOP pair and pop it  $\Rightarrow \text{currD} = 1$   
 $\Rightarrow \text{currN} = 2$

**STEP 2** We can travel to All Nbr



**STEP 3** Check the isSafe to update the diff when isSafe is True



Example 1:

Input: heights =  $[[1,2,2],[3,8,2],[5,3,5]]$

Output: 2

Iter 4

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	1
1	2	6	1
2	$\infty$	$\infty$	$\infty$

Differences

STEP3

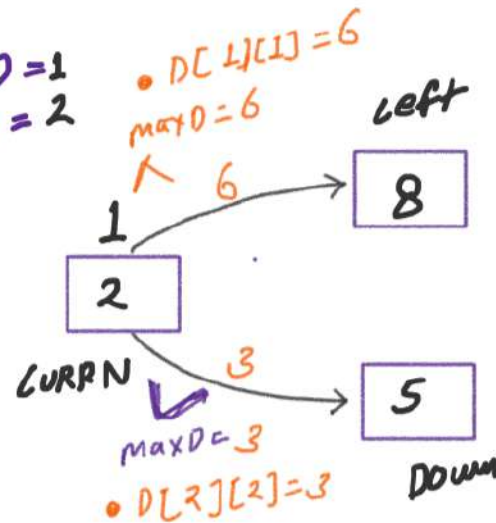
Down  $\Rightarrow 1 < \infty$  ✓  
Left  $\Rightarrow 1 < 6$  ✓

STEP1 Get Top pair and pop it  $\Rightarrow \text{currD} = 1$   
 $\Rightarrow \text{currN} = 2$

STEP2 We can travel to All nbrs



STEP3 Check the isSafe to update the diff when isSafe is true



<del>{1, 2}</del>
{6, 8}
{3, 5}
{2, 3}

MinHeap (mini)

Example 1:

Input: heights =  $[[1, 2, 2], [3, 8, 2], [5, 3, 5]]$

Output: 2

It's

	0	1	2
H	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
D	0	1	1
1	2	6	1
2	$\infty$	$\infty$	3

Differences

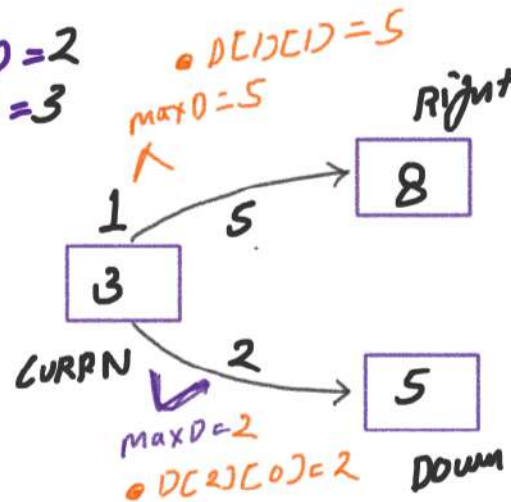
STEP3

Right  $\Rightarrow 2 < 6$  ✓  
Down  $\Rightarrow 2 < \infty$  ✓

STEP1 Get TOPRAIN and POP it  $\Rightarrow \text{CURRD} = 2$   
 $\Rightarrow \text{CURRN} = 3$

STEP2 We can Travel to All NBN  
END L T R D START

STEP3 Check the isSafe to update the diff when isSafe is True



{5, 8}
{6, 8}
{3, 5}
<del>{2, 3}</del>
{2, 5}

MinHeap (mini)

Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

It's

	0	1	2
H	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
D	0	1	1
1	2	5	1
2	2	$\infty$	3

Differences

STEP3

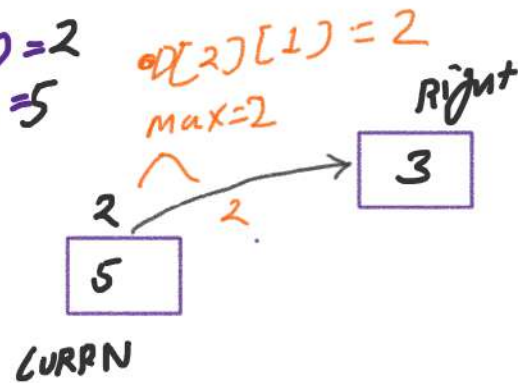
Right  $\Rightarrow 2 < \infty$  ✓

STEP1 Get TopPair and pop it  $\Rightarrow \text{currD} = 2$   
 $\Rightarrow \text{currN} = 5$

STEP2 We can travel to all Nbrs



STEP3 Check the isSafe to update the diff when isSafe is True



{5, 8}
{6, 8}
{3, 5}
{2, 3}
<del>{7, 5}</del>

MinHeap (mini)



Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

Item 7

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	3

Differences

STEP 3

TOP  $\Rightarrow (2 < 5)$  ✓

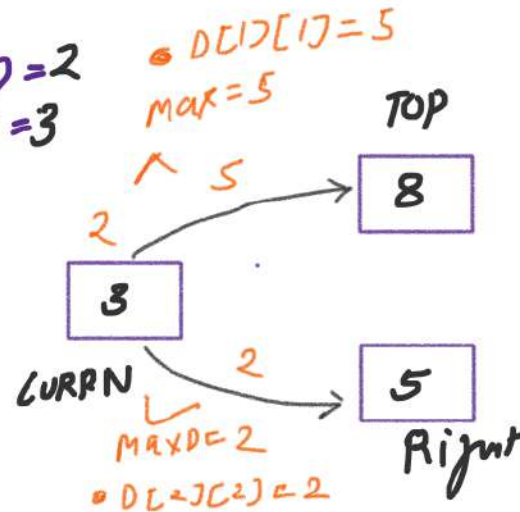
Right  $\Rightarrow (2 < 3)$  ✓

STEP 1 Get TOP pair and pop it  $\Rightarrow \text{currD} = 2$   
 $\Rightarrow \text{currN} = 3$

STEP 2 We can travel to All Nbrs



STEP 3 Check the isSafe to update the diff when isSafe is True



{5, 8}
{6, 8}
{3, 5}
<del>{2, 3}</del>
{2, 5}

MinHeap (mini)



Example 1:

Input: heights =  $[[1, 2, 2], [3, 8, 2], [5, 3, 5]]$

Output: 2

Item 8

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

STEP 3

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	2

Differences

Top  $\Rightarrow (2 < 1) \times$   
 Right  $\Rightarrow$  Array Index  $\times$   
 Down  $\Rightarrow$  Array Index  $\times$   
 Left  $\Rightarrow (2 < 2) \times$

STEP 1 Get top pair and pop it  $\Rightarrow$  currD = 2  
 $\Rightarrow$  currN = 5

STEP 2 We can travel to all Nbr



STEP 3 Check the isSafe to update the diff when isSafe is True

2  
 5  
 currN

{5, 8}
{6, 8}
{3, 5}
<del>{2, 5}</del>

minHeap (mini)

Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

Itum 9

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

STEP3

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	2

Differences

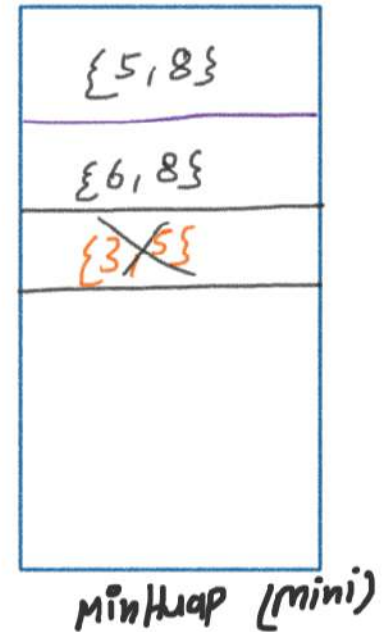
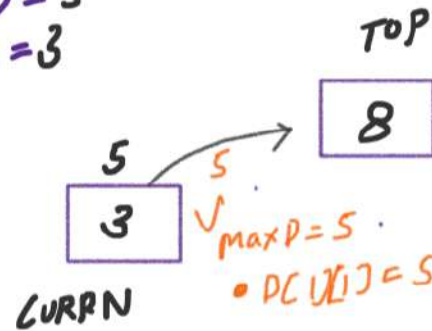
TOP  $\Rightarrow (2 < 5) \checkmark$

STEP1 Get TOP pair and pop it  $\Rightarrow \text{currD} = 5$   
 $\Rightarrow \text{currN} = 3$

STEP2 We can travel to all Nbrs



STEP3 Check the isSafe to update the diff when isSafe is True



Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

STEP3

Item 10

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	2

Differences

TOP  $\Rightarrow 5 < 2 \times$   
 Right  $\Rightarrow 5 < 1 \times$   
 Down  $\Rightarrow 5 < 2 \times$   
 Left  $\Rightarrow 5 < 2 \times$

<del>{5, 8}</del>
{6, 8}

minHeap (mini)

STEP1 Get TOP pair and pop it  $\Rightarrow \text{currD} = 5$   
 $\Rightarrow \text{currN} = 8$

STEP2 We can travel to All Nbr



STEP3 Check the isSafe to update the diff when isSafe is True

5  
8  
currN

Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

STEP3

Item 11

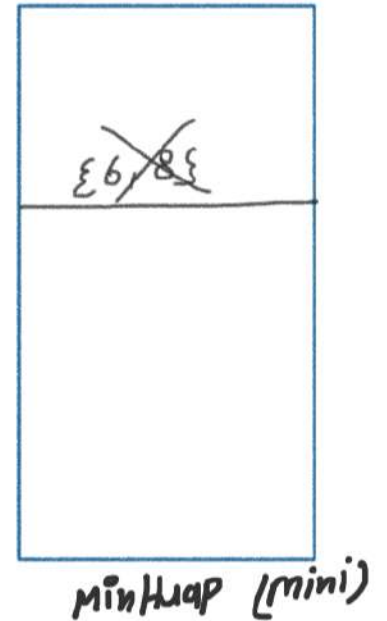
	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	2

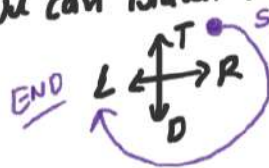
Differences

TOP  $\Rightarrow 5 < 1 \times$   
 Right  $\Rightarrow 5 < 1 \times$   
 Down  $\Rightarrow 5 < 2 \times$   
 Left  $\Rightarrow 5 < 2 \times$

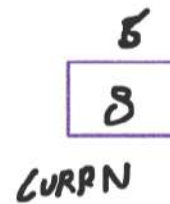


STEP1 Get TOPpair and pop it  $\Rightarrow \text{currD} = 6$   
 $\Rightarrow \text{currN} = 8$

STEP2 We can travel to All Node



STEP3 Check the isSafe to update the diff when isSafe is True



Example 1:

Input: heights = [[1,2,2],[3,8,2],[5,3,5]]

Output: 2

	0	1	2
0	1	2	2
1	3	8	2
2	5	3	5

Heights

	0	1	2
0	0	1	1
1	2	5	1
2	2	2	2

Differences

→ destination

return [2][2]  
⇒ (2)  
output

Empty  
now

minHeap (mini)

// 3. Path With Minimum Effort (Leetcode-1631)

```
class Solution {
public:
```

```
bool isSafe(int newX, int newY, int row, int col, vector<vector<int>>& diff, int currX, int currY) {
    if(newX >= 0 && newY >= 0 && newX < row && newY < col && diff[currX][currY] < diff[newX][newY]) {
        return true;
    }
    else {
        return false;
    }
}
```

```
int minimumEffortPath(vector<vector<int>>& heights) {
```

```
    priority_queue<pair<int, pair<int,int>>, vector<pair<int,pair<int,int>>>,
    greater<pair<int,pair<int,int>>>> mini;
```

```
    int row = heights.size();
    int col = heights[0].size();
    vector<vector<int>> diff(row, vector<int>(col, INT_MAX));
    int destX = row-1;
    int destY = col-1;
```

```
    // Initial state:
    // Set the src difference ==> 0
    diff[0][0] = 0;
    // Insert MinHeap Entry for src ==> {diff of src, (position cell of src)}
    mini.push({0, {0,0}});
```

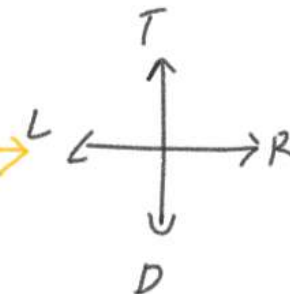
```
    while(!mini.empty()) {
        // Step 1: get the topPair and pop it
        pair<int, pair<int,int>> topPair = mini.top();
        mini.pop();
        int currDiff = topPair.first;
        pair<int,int> currNodeIndexPair = topPair.second;
        int currX = currNodeIndexPair.first;
        int currY = currNodeIndexPair.second;
```

```
    // Step 2: How we can travel to all nbr i.e. top, right, down, left
    int dx[] = {-1,0,1,0};
    int dy[] = {0,1,0,-1};
    for(int i=0; i<4; i++) {
        int newX = currX + dx[i];
        int newY = currY + dy[i];
```

```
    // Step 3: Check the isSafe to update the diff
    if(isSafe(newX, newY, row, col, diff, currX, currY)) {
        int maxDiff = max(currDiff, abs(heights[currX][currY]-heights[newX][newY]));
        // Update the minDiff in the diff array
        diff[newX][newY] = min(diff[newX][newY], maxDiff);
        // Create new entry for minHeap
        if(newX != destX || newY != destY) {
            mini.push({diff[newX][newY], {newX, newY}});
        }
    }
}
```

```
    // Return destination minimum difference effort
    return diff[destX][destY];
}
```

```
};
```



→ Github pan jahan code check karle  
Ajam visible nahin Raha Hai to  
push.