

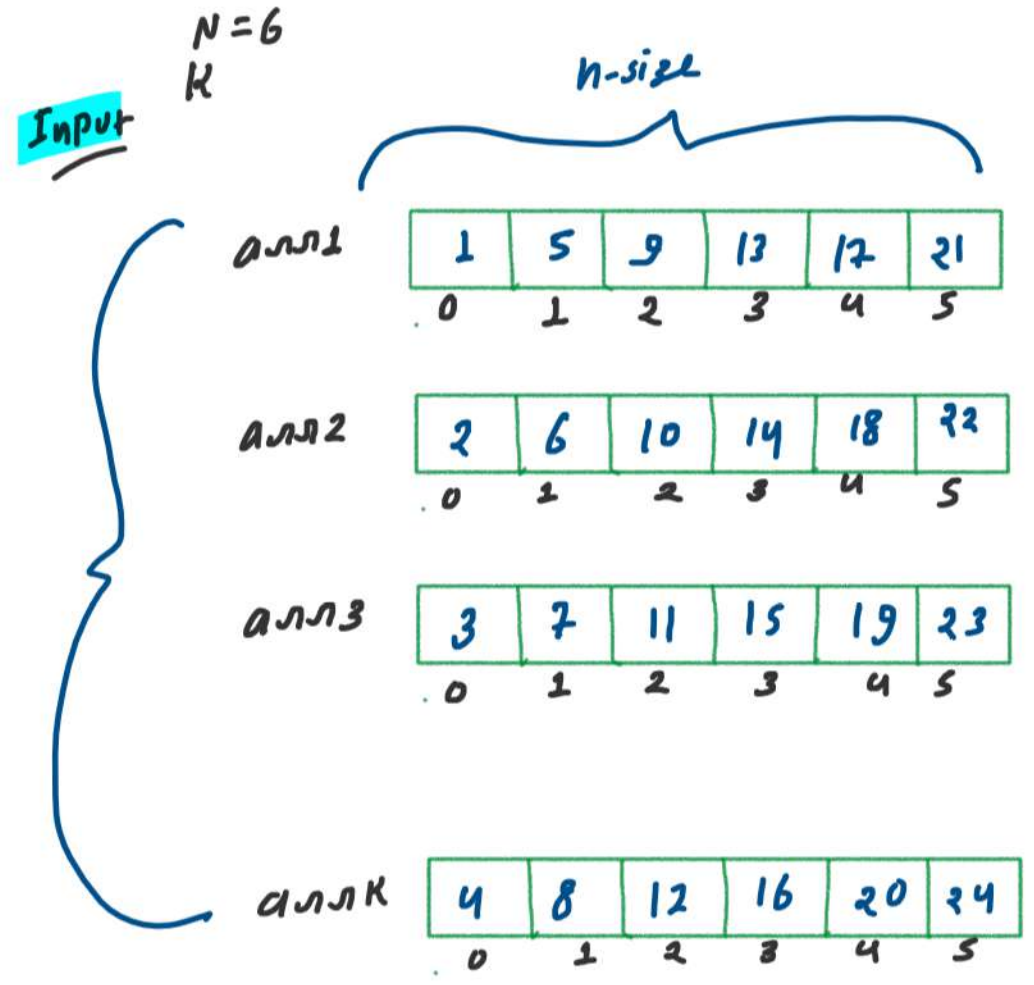
1. Merge K Sorted Arrays (GFG)

Output

single sorted Array

1	2	3	4	5	6	7	8	9	10	11	12
0	1	2	3	4	5	6	7	8	9	10	11

13	14	15	16	17	18	19	20	21	22	23	24
12	13	14	15	16	17	18	19	20	21	22	23



ALGORITHMS

STEP 1 Find First min element of K-Arrays

Note When we merge two sorted Arrays
→ to hum phle First min element
find karte hai OR ussi element ko
Ans Array me phle push karte hai
Right - YES

→ We find min-element using min-heap with
Time complexity $O(1)$

EX 1 Input $N=4$
 $K=3$

array 1

1	4	8	10
0	1	2	3

array 2

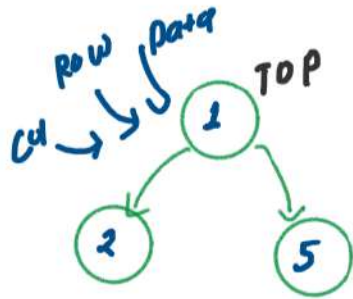
2	3	6	9
0	1	2	3

array 3

5	7	11	12
0	1	2	3

(I) CREATE MIN-HEAP USING FIRST ELEMENTS OF K-ARRAYS

Iteration 1



(II) FIND MIN ELEMENT

MIN = 1 RINDEX = 0 LINDEX = 0

(IV) UPDATE TOP and Heapify it
TOP = 4

array 0

1	4	8	10
0	1	2	3

array 1

2	3	6	9
0	1	2	3

array 2

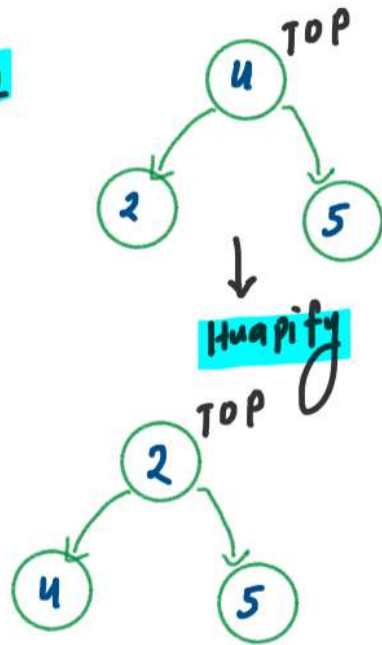
5	7	11	12
0	1	2	3

(III) Push MIN INTO ANSWER ARRAY

ANS

1											
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 2



- ② FIND MIN ELEMENT
 MIN = 2 RINDEX = 1 LINDEX = 0
- ④ UPDATE TOP and Heapify it
 TOP = 3

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

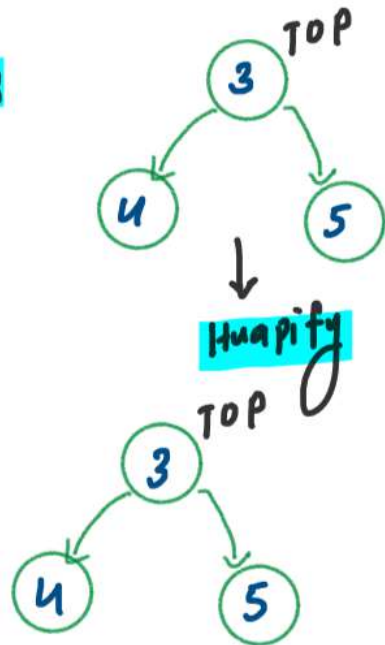
5	7	11	12
0	1	2	3

③ Push MIN INTO ANS ARRAY

ANS

1	2										
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 3



- ② FIND MIN ELEMENT
 MIN = 3 RINDEX = 1 LINDEX = 1
- ④ UPDATE TOP and Heapify it
 TOP = 6

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

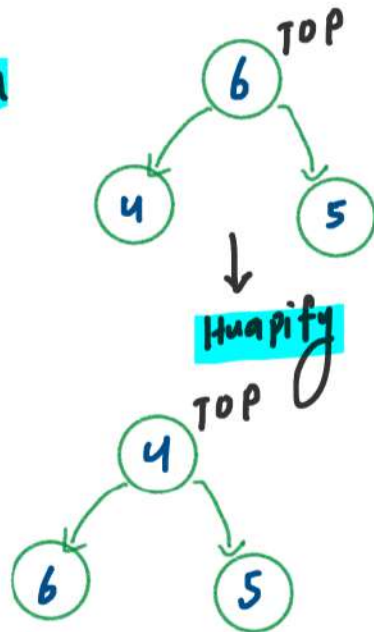
5	7	11	12
0	1	2	3

③ PUSH MIN INTO ANS ARRAY

ANS

1	2	3									
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 4



- ② FIND MIN ELEMENT
 MIN = 4 RINDEX = 0 LINDEX = 1
- ④ UPDATE TOP and Heapify it
 TOP = 8

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

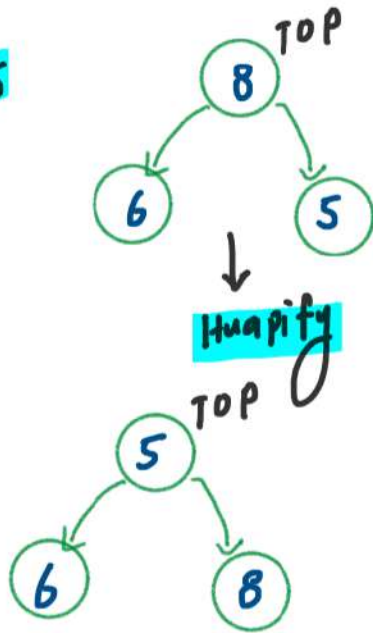
5	7	11	12
0	1	2	3

③ PUSH MIN INTO ANSW ARRAY

ANS

1	2	3	4								
0	1	2	3	4	5	6	7	8	9	10	11

Iterations



② FIND MIN ELEMENT

MIN = 5 RINDEX = 2 LINDEX = 0

④ UPDATE TOP and Heapify it
TOP = 7

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

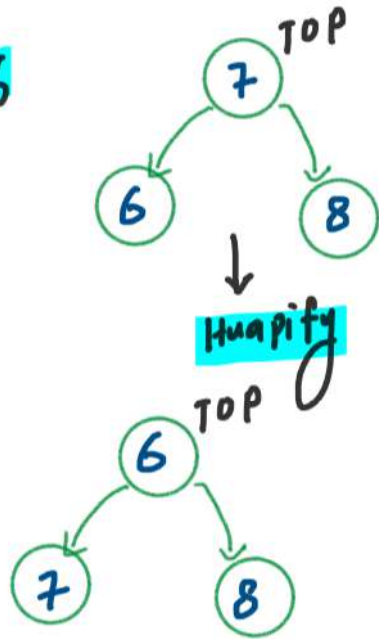
5	7	11	12
0	1	2	3

③ Push MIN INTO ANS ARRAY

ANS

1	2	3	4	5							
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 6



- ② FIND MIN ELEMENT
 MIN = 6 RINDEX = 1 LINDEX = 2
- ④ UPDATE TOP and Heapify it
 TOP = 9

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

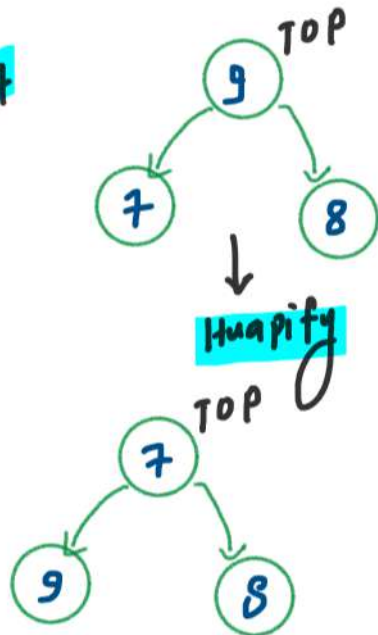
5	7	11	12
0	1	2	3

③ Push MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6						
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 7



- ② FIND MIN ELEMENT
 MIN = 7 RINDEX = 2 LINDEX = 1
- ④ UPDATE TOP and Heapify it
 TOP = 11

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

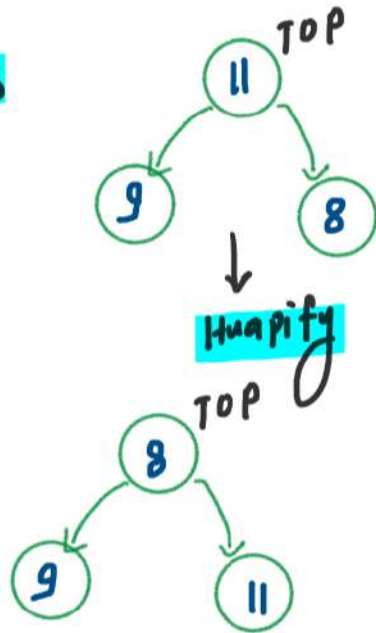
5	7	11	12
0	1	2	3

- ③ PUSH MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6	7					
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 8



- ② FIND MIN ELEMENT
 MIN = 8 RINDEX = 0 LINDEX = 2
- ④ UPDATE TOP and Heapify it
 TOP = 10

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

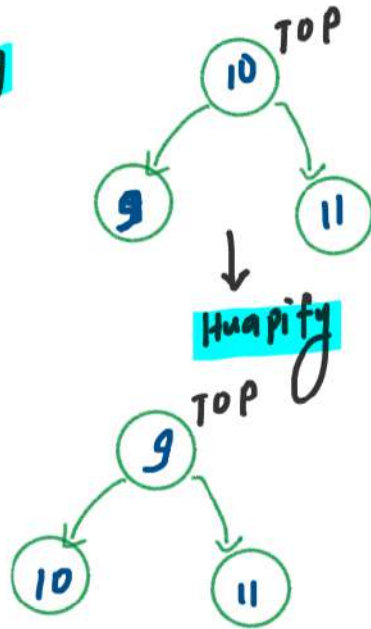
5	7	11	12
0	1	2	3

③ PUSH MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6	7	8				
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 9



② FIND MIN ELEMENT

MIN = 9 RINDEX = 1 LINDEX = 3

④ UPDATE TOP and Heapify it
TOP = 10

$CINDEX < N$
 $3 < 3 \times$

(GALTI YANHA
PAR HOTI HAI)

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

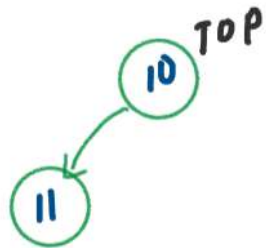
5	7	11	12
0	1	2	3

③ PUSH MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6	7	8	9			
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 10



(II) FIND MIN ELEMENT

MIN = 10 RINDEX = 0 LINDEX = 3

(IV) UPDATE TOP and Heapify it

TOP = 11

(LINDEX < N
3 < 3 X)

(GALTI YANHA
PAR HOTI HAI)

array 0

1	4	8	10
0	1	2	3

array 1

2	3	6	9
0	1	2	3

array 2

5	7	11	12
0	1	2	3

(III) PUSH MIN INTO ANSW ARRAY

ANS

1	2	3	4	5	6	7	8	9	10		
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 11

11 TOP

II FIND MIN ELEMENT

MIN = 11 RINDEX = 2 LINDEX = 2

IV Update TOP and Heapify it
TOP = 12

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

5	7	11	12
0	1	2	3

III Push MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6	7	8	9	10	11	
0	1	2	3	4	5	6	7	8	9	10	11

Iteration 12

12 ^{TOP} ~~X~~

II FIND MIN ELEMENT

MIN = 12 RINDEX = 2 LINDEX = 3

IV UPDATE TOP and Heapify it
TOP = TOP UPDATE

NAHI KAR
SARTE HAI

Heap is Empty
NOW

AND

COULD N
3 < 3 X

FINAL OUTPUT

ans1 0

1	4	8	10
0	1	2	3

ans2 1

2	3	6	9
0	1	2	3

ans3 2

5	7	11	12
0	1	2	3

III Push MIN INTO ANS ARRAY

ANS

1	2	3	4	5	6	7	8	9	10	11	12
0	1	2	3	4	5	6	7	8	9	10	11

(IV)

Array main jis min value ko push kar nam hai
uski next element k dubla har ki top element ko
update karne ke liye hume RowIndex, ColIndex and
data ki need hai Right → YES

→ TO HUM APNA KHUD KA EK NEW DATA TYPE
create karenge jisme yeh three properties
mikhni-

MIN
HEAP KA
NODE HAI
Info

data
rowIndex
colIndex

class Info {

public:

int data;

int rowIndex;

int colIndex;

Info(int data,
int rowIndex,
int colIndex) {

this->data = data;

this->rowIndex = rowIndex;

this->colIndex = colIndex;

}

}

```
// PROBLEM 1: Merge K Sorted Arrays (GFG)
#include<iostream>
#include<vector>
#include<queue>
using namespace std;

// OWN DATA TYPE
class Info
{
    ...
};

// OWN COMPARETOR TO RETURN THE MIN NODE FROM TWO DIFFERENT NODE -> true/false
class Compare
{
    ...
};

void mergeKSortedArrays(int arr[][4], int n, int k, vector<int> &ans){
    ....
}

int main(){
    int rowSize = 3;
    int colSize = 4;
    int arr[3][4] = {{1, 4, 8, 10},{2, 3, 6, 9},{5, 7, 11, 12}};

    int n = colSize;
    int k = rowSize;

    vector<int> ans;
    mergeKSortedArrays(arr, n, k, ans);

    cout<< " Printing Single Sorted Array: " << endl;
    for(int i = 0; i < ans.size(); i++){
        cout << ans[i] << " ";
    }

    return 0;
}

/*
Printing Single Sorted Array:
1 2 3 4 5 6 7 8 9 10 11 12
*/
```

```
// OWN DATA TYPE
class Info
{
    public:
        int data;
        int rowIndex;
        int colIndex;

        Info(int data, int rowIndex, int colIndex){
            this->data = data;
            this->rowIndex = rowIndex;
            this->colIndex = colIndex;
        }
};

// OWN COMPARETOR TO RETURN THE MIN NODE FROM TWO DIFFERENT NODE -> true/false
class Compare
{
    public:
        bool operator()(Info* first, Info* second){
            // Returns true if first = 1 comes before second=2 in the ordering
            return first->data > second->data; // Create Min Heap
        }
};
```

```

void mergeKSortedArrays(int arr[][4], int n, int k, vector<int> &ans){
    // Create MIN Heap
    priority_queue<Info*, vector<Info*>, Compare> pq;

    // I. process first k elements from k arrays
    for (int row = 0; row < k; row++)
    {
        int element = arr[row][0]; // arr[0][0], arr[1][0], arr[2][0]
        Info* tempNode = new Info(element, row, 0);
        pq.push(tempNode);
    }

    while (!pq.empty())
    {
        Info* topNode = pq.top();
        pq.pop();

        // II. Find topData (Min Value)
        int topData = topNode->data;
        int topRow = topNode->rowIndex;
        int topCol = topNode->colIndex;

        // III. Ab ans array me topData (Min Value) push kar do
        ans.push_back(topData);

        // IV. Ab next element kya hoga for the same row, jis row se element ko pop kiya hai
        // usse insert bhi to karna hai--> to topCol ko 1 se increment krdo
        if (topCol + 1 < n){
            // iska matlab present row me abhi or v elements baki hai
            Info* newNode = new Info(arr[topRow][topCol+1], topRow, topCol+1);
            pq.push(newNode);
        }
    }
}

```

Time Complexity

$$\left[\begin{array}{l} \text{Heap ki T.C.} = O(\log(K)) \\ \text{FOR LOOP ki T.C.} = O(K) \end{array} \right\} O(K * \log(K))$$

$K = \text{No. of Arrays}$

+

$$\left[\begin{array}{l} \text{WHILE LOOP ki T.C.} = O(N) \\ \text{Heap ki T.C.} = O(\log(K)) \end{array} \right\} O(N * \log(K))$$

$N = \text{Total Elements of All arrays}$

Overall T.C.

$$O(K * \log(K)) + O(N * \log(K))$$

SPACE complexity

(MIN Heap) Priority queue ki S.C. = $O(K)$
 K = no. of arrays

(Ans) Vector Array ki S.C. = $O(N)$
 N = Total elements
of All arrays

Overall S.C.

$$O(K) + O(N)$$

```
#include <iostream>
#include <queue>
```

```
// Custom comparison function for the min heap
```

```
struct Compare {
    bool operator()(int a, int b) {
        // Returns true if a comes before b in the ordering
        return a > b;
    }
};
```

```
int main() {
    // Creating a min heap of integers with the custom comparison function
    std::priority_queue<int, std::vector<int>, Compare> pq;
```

```
// Inserting elements into the min heap.
```

```
pq.push(5);
pq.push(2);
pq.push(8);
pq.push(1);
```

```
// Printing elements from the min heap
```

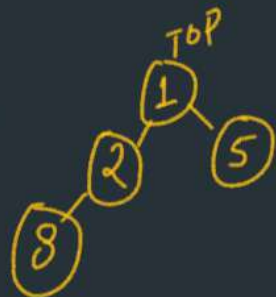
```
while (!pq.empty()) {
    std::cout << pq.top() << " ";
    pq.pop();
}
```

```
return 0;
```

```
}
```

```
/*
INPUT: 5 2 8 1
OUTPUT: 1 2 5 8 (MIN HEAP)
*/
```

MIN
HEAP



struct
re
place
class
u
use
ran
scape
to

```
#include <iostream>
#include <queue>
```

```
// Custom comparison function for the max heap
```

```
struct Compare {
    bool operator()(int a, int b) {
        // Returns false if a comes before b in the ordering
        return a < b;
    }
};
```

```
int main() {
    // Creating a max heap of integers with the custom comparison function
    std::priority_queue<int, std::vector<int>, Compare> pq;
```

```
// Inserting elements into the max heap
```

```
pq.push(5);
pq.push(2);
pq.push(8);
pq.push(1);
```

```
// Printing elements from the max heap
```

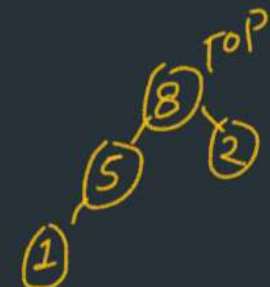
```
while (!pq.empty()) {
    std::cout << pq.top() << " ";
    pq.pop();
}
```

```
return 0;
```

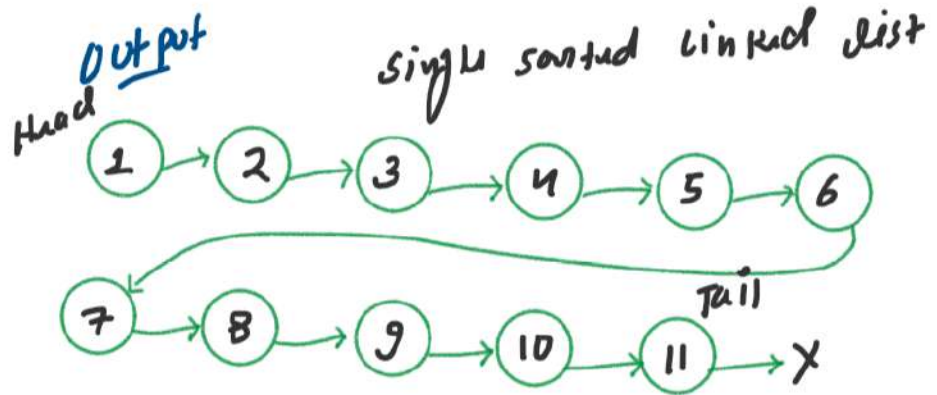
```
}
```

```
/*
INPUT: 5 2 8 1
OUTPUT: 8 5 2 1 (MAX HEAP)
*/
```

MAX HEAP

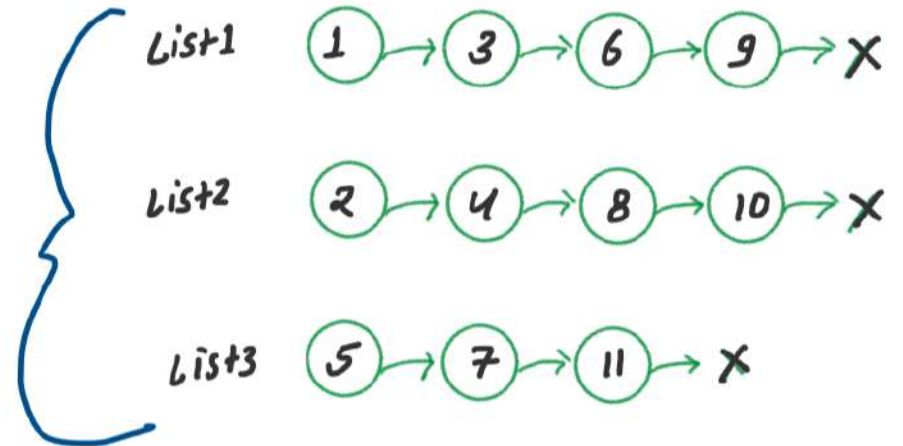


2. Merge K Sorted Linked List (Leetcode-23)



Input

lists[[1,3,6,9], [2,4,8,10], [5,7,11]]

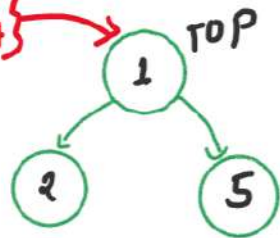


Algorithm

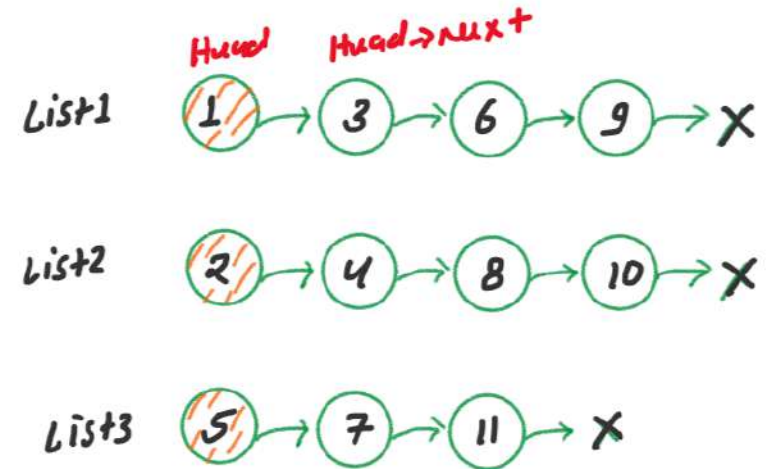
Iterate ↓

① Make min heap using First K - Nodes of all lists

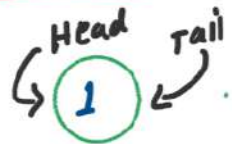
Head → data
Head → next



$$K = \text{lists.size}() = 3$$



New LL



② Create new linked list
Initiating → Head = X, Tail = X

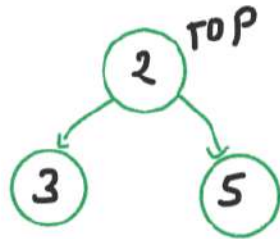
③ Fetch TOP and push pop to new linked list

TOP = 1

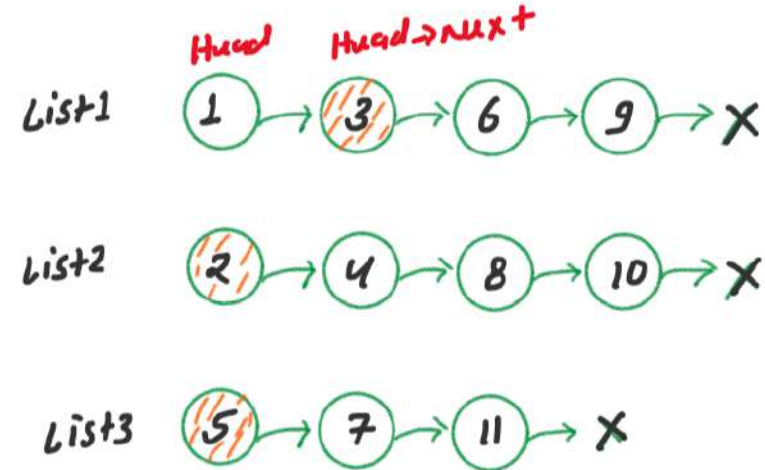
pop the TOP

Head = 1 and update the TOP if (tail → next != null)
tail = 1 → TOP = 3

Iterative 2



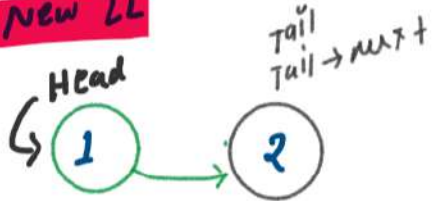
$$K = \text{lists.size}() \\ = 3$$



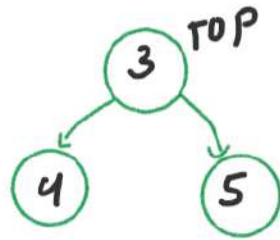
II Create new linked list
Head = 1 Tail = 1

III Fetch TOP and push pop to new linked list
TOP = 2
pop the TOP
tail -> next = 2 and update the TOP if (tail -> next != null)
tail = 2
↳ TOP = 4

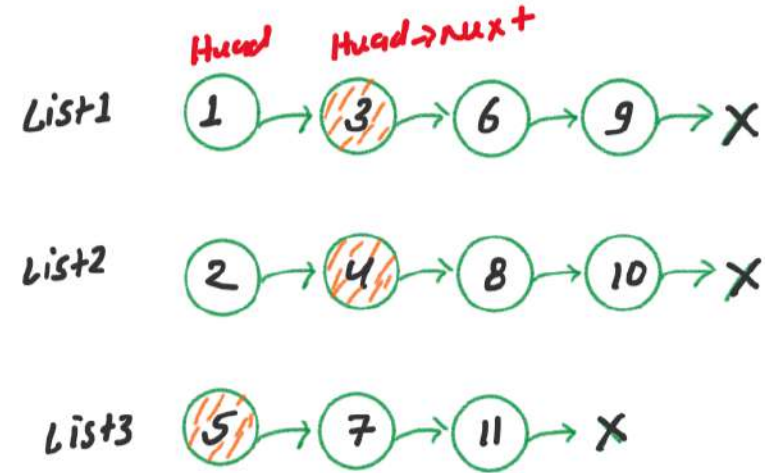
New LL



Iteration 3



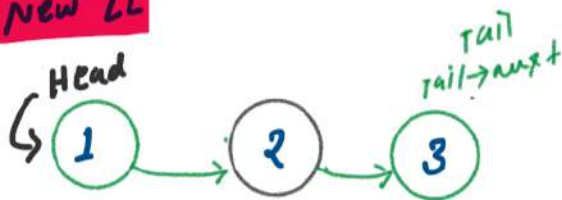
$$K = \text{Lists.size}() \\ = 3$$



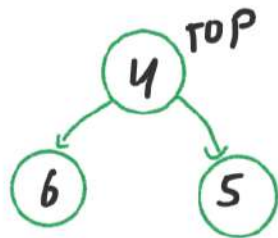
II Create New Linked List
Head = 1 Tail = 2

III Fetch TOP and push pop to New Linked List
TOP = 3
pop the TOP
tail -> next = 3 and update the TOP if (tail -> next != null)
tail = 3
↳ TOP = 6

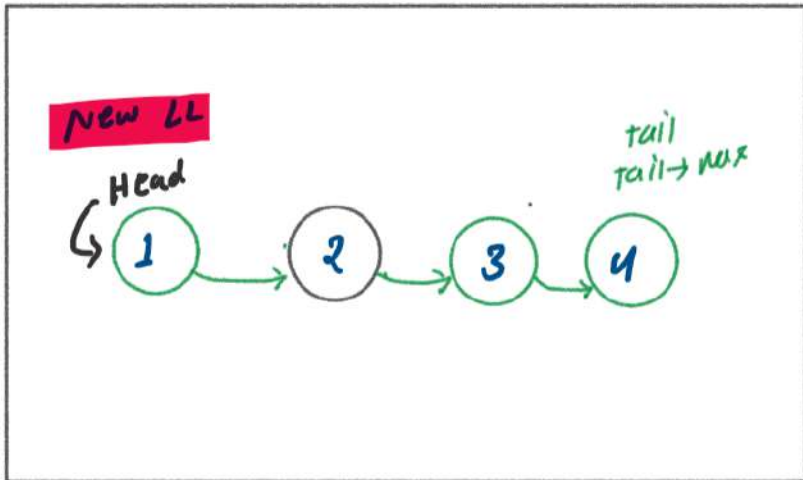
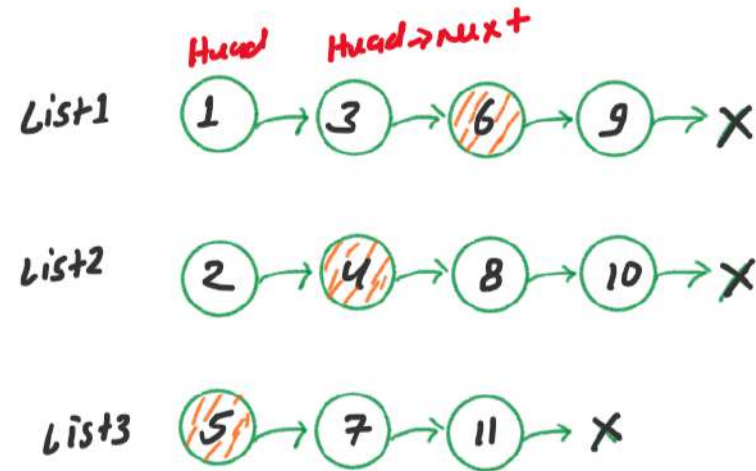
New LL



Iteration 4



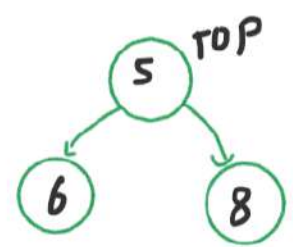
$$K = \text{lists.size}() \\ = 3$$



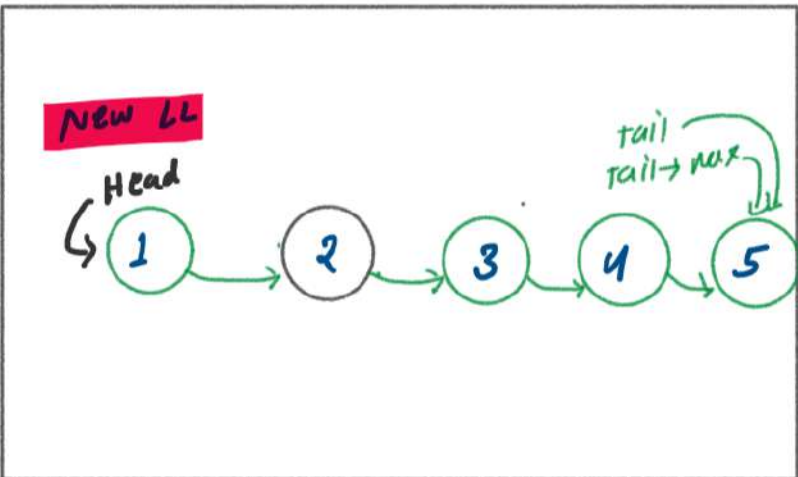
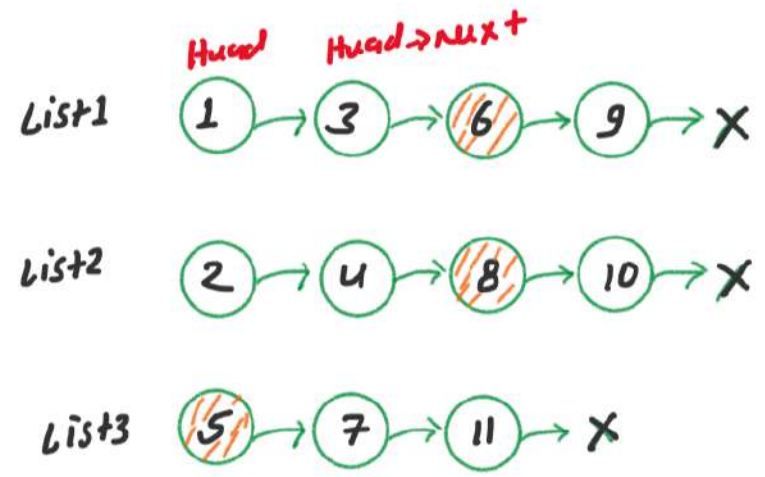
II Create new linked list
Head = 1 Tail = 3

III Fetch TOP and push pop to new linked list
 TOP = 4
 pop the TOP
 tail->next = 4 and update the TOP if (tail->next != null)
 tail = 4
 → TOP = 8

Iterations



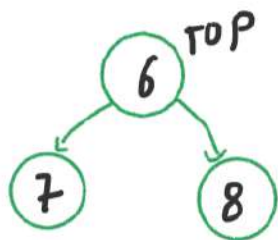
$$K = \text{lists.size}() \\ = 3$$



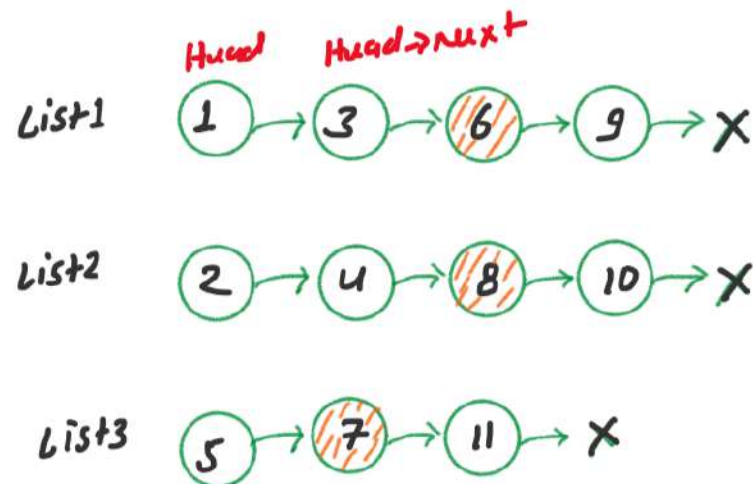
II Create new linked list
Head = 1 tail = 4

III fetch TOP and push pop to new linked list
TOP = 5
pop the TOP
tail -> next = 5 and update the TOP if (tail -> next != null)
tail = 5
 ↳ TOP = 7

Iteration 5



$$K = \text{Lists.size}() = 3$$

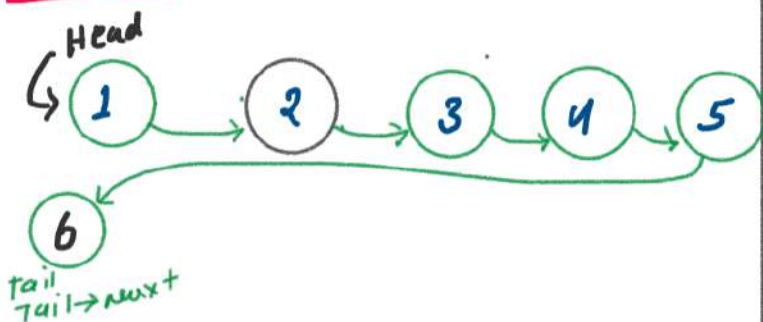


II Create new linked list
 Head = 1 Tail = 5

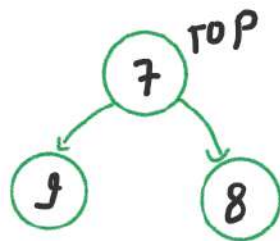
III Fetch TOP and push pop to new linked list
 TOP = 6

pop the TOP
 tail → next = 6 and update the TOP if (tail → next != null)
 tail = 6
 ↳ TOP = 9

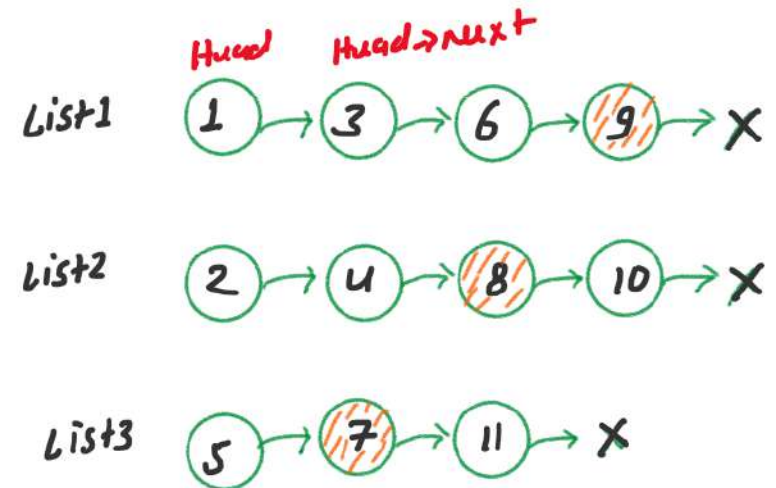
New LL



Iteration 7



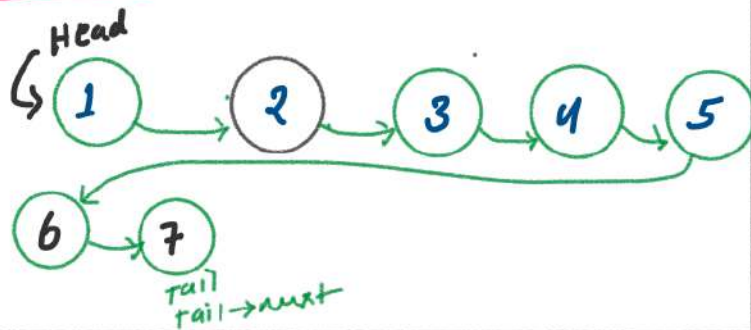
$$K = \text{lists.size}() = 3$$



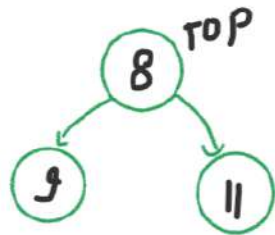
II Create new linked list
Head = 1 Tail = 6

III Fetch TOP and push pop to new linked list
TOP = 7
pop the TOP
tail → next = 7 and update the TOP if (tail → next != null)
tail = 7
→ TOP = 11

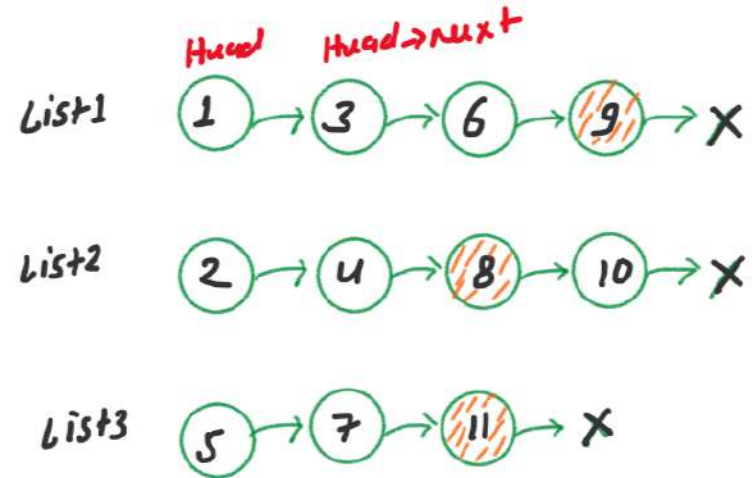
New LL



Iteration 8



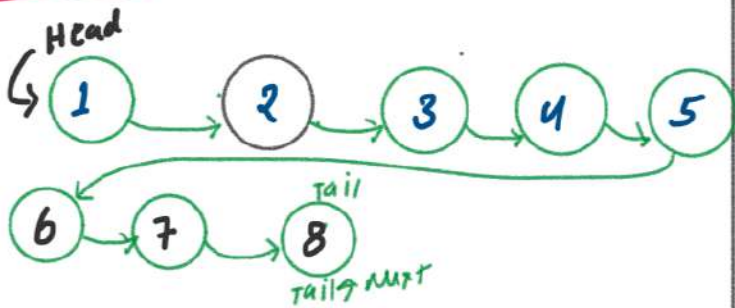
$$K = \text{lists.size}() \\ = 3$$



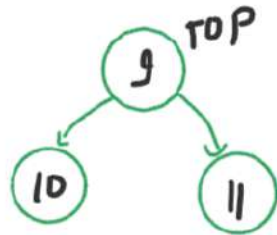
II Create new linked list
Head = 1 Tail = 7

III Fetch TOP and push pop to new linked list
TOP = 8
pop the TOP
tail -> next = 8 and update the TOP if (tail -> next != null)
tail = 8
↳ TOP = 10

New LL



Iteration 9



$$K = \text{lists.size}() \\ = 3$$



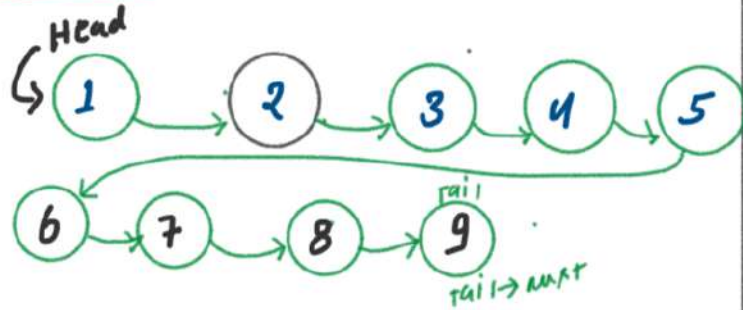
II Create new linked list
Head = 1 Tail = 8

III Fetch TOP and push pop to new linked list
TOP = 9

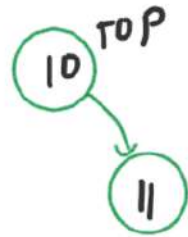
pop the TOP
tail->next = 9
tail = 9

and update the TOP if (tail->next != null)
TOP = TOP update
Nani Hoga

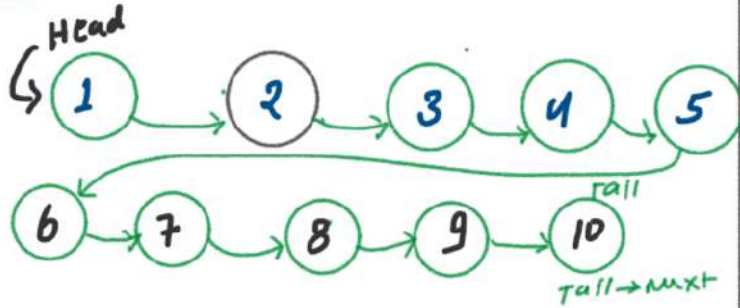
New LL



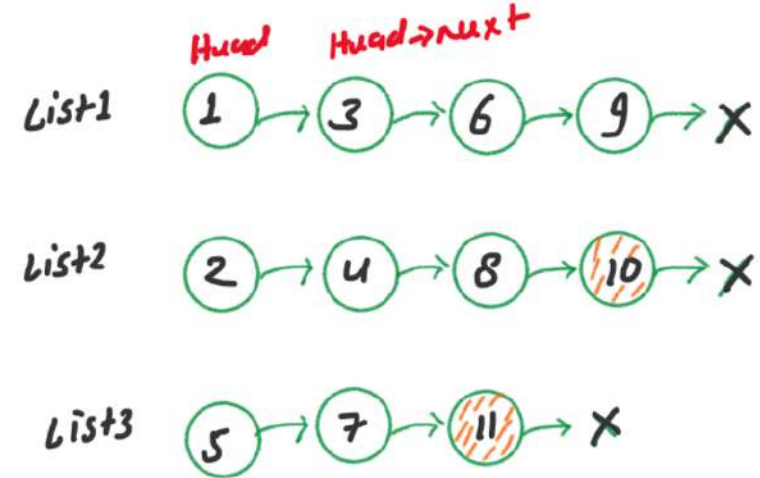
Iteration 10



New LL



$$K = \text{lists.size}() = 3$$



II Create new linked list
Head = 1 tail = 9

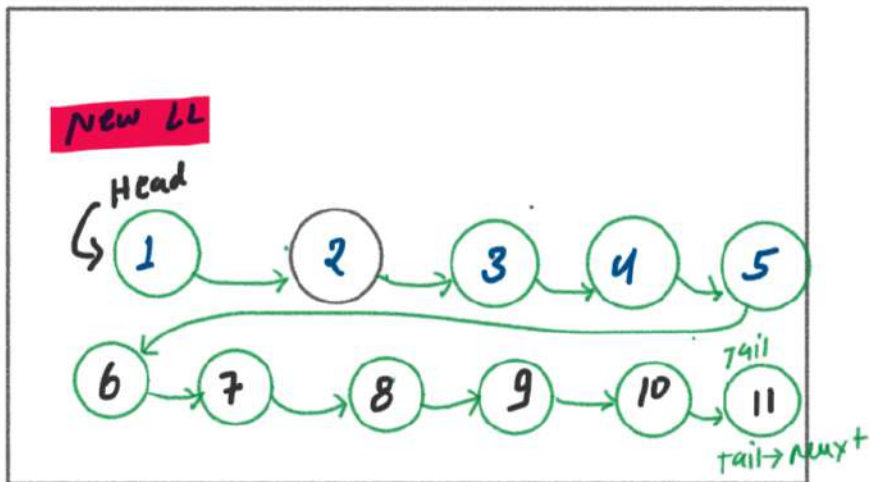
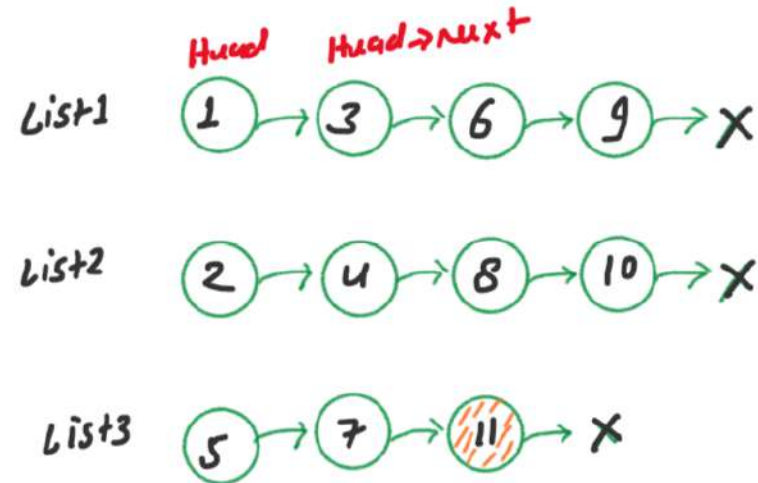
III Fetch TOP and push pop to new linked list

TOP = 10
pop the TOP
tail → next = 10 and update the TOP if (tail → next) = null
tail = 10
→ TOP = TOP update nahi Hoga

Iteration 1

11 TOP

$$K = \text{lists.size}() \\ = 3$$



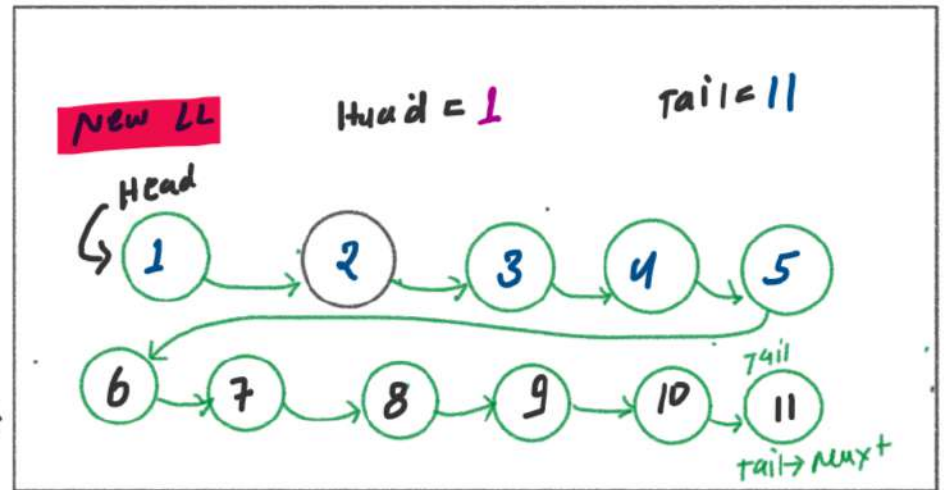
II Create new linked list
Head = 1 Tail = 10

III Fetch TOP and push pop to new linked list
 TOP = 11
 pop the TOP
 tail->next = 11 and update the TOP if (tail->next == null)
 tail = 11
 TOP = TOP update nahi hoga

Iteration 2

Heap is now Empty
→ stop

Final
Output ↗




```

class Solution {
public:
    // OWN COMPARATOR FOR MIN HEAP
    class Compare
    {
    public:
        bool operator()(ListNode* first, ListNode* second){
            return first->val > second->val;
        }
    };

    ListNode* mergeKLists(vector<ListNode*> &lists) {
        // Create MIN Heap
        priority_queue<ListNode*, vector<ListNode*>, Compare> pq;

        // I. process first k elements from k lists
        for (int i = 0; i < lists.size(); i++)
        {
            ListNode* listHead = lists[i];
            if(listHead != NULL){
                pq.push(listHead);
            }
        }

        // II. create new linked list
        ListNode* head = NULL;
        ListNode* tail = NULL;

        while(!pq.empty()){
            ListNode* topNode = pq.top();
            pq.pop();
            // III. Push first topNode in new linked list first time
            if(head == NULL && tail == NULL){
                head = topNode;
                tail = topNode;
                // IV. Update the topNode
                if(tail->next != NULL){
                    pq.push(tail->next);
                }
            }
            else{
                // III. Not pushing the first node now
                tail->next = topNode;
                tail = topNode;
                // IV. Update the topNode
                if(tail->next != NULL){
                    pq.push(tail->next);
                }
            }
        }
        return head;
    };
};

```

Time Complexity

$$\left\{ \begin{array}{l} \text{FOR LOOP HI T.C.} = O(K) \\ \text{MIN HEAP K T.C.} = O(\log K) \end{array} \right\}$$

where
K = Total linked lists

$$\rightarrow \text{T.C.} = O(K * \log K)$$

Merge Done K T.C.


$$\left\{ \begin{array}{l} \rightarrow \text{while loop T.C.} = O(N) \\ \text{MIN HEAP T.C.} = O(\log K) \end{array} \right\}$$

where
N = All nodes of
each linked lists

$$\rightarrow \text{T.C.} = O(N * \log K)$$

$$\left\{ \begin{array}{l} \text{Overall T.C.} = O(K * \log K) + O(N * \log K) \\ = O(N * \log K) \end{array} \right\}$$

SPACE complexity



min heap is taking space $O(K)$

where K is number of linked lists

3. Smallest Range in K Lists (Leetcode-632)

Input: NUMS = [
 [4, 10, 15, 24, 26],
 [0, 9, 12, 20],
 [5, 18, 22, 30]
]

Output: [20, 24]

Explanation:

List 1: [4, 10, 15, 24, 26], 24 is in range [20, 24].

List 2: [0, 9, 12, 20], 20 is in range [20, 24].

List 3: [5, 18, 22, 30], 22 is in range [20, 24].

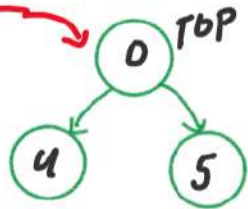
		Col0	Col1	Col2	Col3	Col4
Row0	arr1	4	10	15	24	26
		0	1	2	3	4
Row1	arr2	0	9	12	20	
		0	1	2	3	
Row2	arr3	5	18	22	30	
		0	1	2	3	

Algorithm

I Create MIN Heap of First K-Elements of All lists

Node

Data
Row I
Col I



Item 1

$K = \text{Total Row Size} = 3$

MIN = 0

MAX = 5

Range [MIN, MAX]
AnsStart = 0
AnsEnd = 5

II Insert New Element

New Element = 9

Col I = 1

Row I = 1

Update maxi = 9

II Fetch TOP and POP

TOP = 0

Update min = 0

Update the Ans (Range) when

AnsStart = 0

AnsEnd = 5

For smaller Range

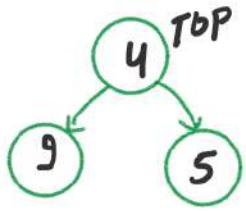
$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$5 - 0 < 5 - 0$$

$$5 < 5 \times$$

		Col 0	Col 1	Col 2	Col 3	Col 4
Row 0	Ans 0	4	10	15	24	26
		0	1	2	3	4
Row 1	Ans 1	0	9	12	20	
		0	1	2	3	
Row 2	Ans 2	6	18	22	30	
		0	1	2	3	

Item 2



newRange
 MIN = 4
 MAX = 9

oldRange
 Range [MIN, MAX]
 AnsStart = 0
 AnsEnd = 5

		Col0	Col1	Col2	Col3	Col4
Row0	Ans1	4 0	10 1	15 2	24 3	26 4
Row1	Ans2	0 0	9 1	12 2	20 3	
Row2	Ans3	5 0	18 1	22 2	30 3	

III Insert New Element
 NewElem = 10
 ColI = 1
 RowI = 0
 Update maxi = 10

II Fetch TOP and POP
 TOP = 4
 Update min = 4
 Update the Ans (Range) when
 AnsStart = 0
 AnsEnd = 5

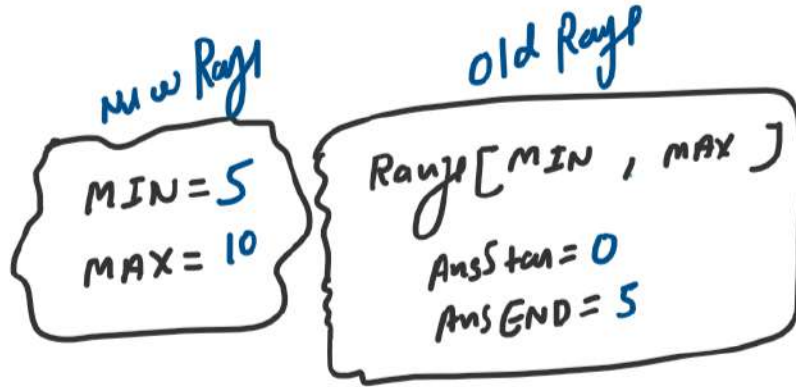
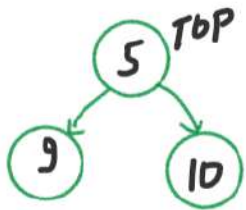
For smaller Range

$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$(9 - 4) < (5 - 0)$$

$$5 < 5 \times$$

Iter 3



		Col0	Col1	Col2	Col3	Col4
Row0	Ans1	4	10	15	24	26
		0	1	2	3	4
Row1	Ans2	0	9	12	20	
		0	1	2	3	
Row2	Ans3	5	18	22	30	
		0	1	2	3	

III Insert New Element
 NewElem = 18
 ColI = 1
 RowI = 2
 Update maxi = 18

II Fetch TOP and POP
 TOP = 5

Update min = 5
 Update the Ans (Range)
 AnsStart = 0
 AnsEnd = 5

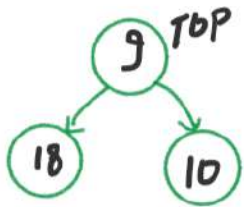
For smaller Range

$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$10 - 5 < 5 - 0$$

$$5 < 5 \times$$

Iter 4



new Range

MIN = 9

MAX = 18

old Range

Range [MIN, MAX]

AnsStart = 0

AnsEnd = 5

		Col0	Col1	Col2	Col3	Col4
Row0	ans1	4 0	10 1	15 2	24 3	26 4
Row1	ans2	0 0	9 1	12 2	20 3	
Row2	ans3	6 0	18 1	22 2	30 3	

Ⓐ Insert New Element
 NewElement = 12
 ColI = 2
 RowI = 1
 Update maxi = 18

Ⓑ Fetch TOP and POP
 TOP = 9
 Update min = 9
 Update the Ans (Range) when
 AnsStart = 0
 AnsEnd = 5

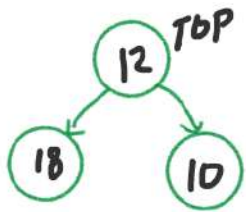
For smaller Range

$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$18 - 9 < 5 - 0$$

$$9 < 5 \quad \times$$

Item 5



new Range
 MIN = 12
 MAX = 18

old Range
 Range [MIN, MAX]
 AnsStart = 0
 AnsEnd = 5

		Col0	Col1	Col2	Col3	Col4
Row0	ans1	4 0	10 1	15 2	24 3	26 4
Row1	ans2	0 0	9 1	12 2	20 3	
Row2	ans3	6 0	18 1	22 2	30 3	

(I) Insert New Element

NewElem = 20

ColI = 3

RowI = 0

Update maxi = 20

(20, 18)

(II) Fetch TOP and POP

TOP = 12

Update min = 12

Update the Ans (Range)

AnsStart = 0

AnsEnd = 5

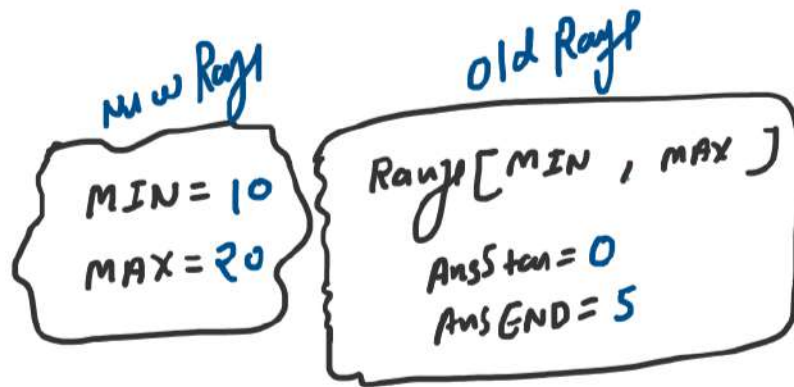
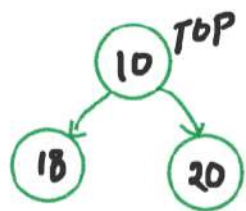
For smaller Range

$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$18 - 12 < 5 - 0$$

$$6 < 5 \quad \times$$

Item 6



		Col0	Col1	Col2	Col3	Col4
Row0	ans1	4	10	15	24	26
		0	1	2	3	4
Row1	ans2	0	9	12	20	
		0	1	2	3	
Row2	ans3	5	18	22	30	
		0	1	2	3	

(I) Insert New Element

New Element = 15

ColI = 2

RowI = 1

Update maxi = 20

(II) Fetch TOP and POP

TOP = 10

Update min = 10

Update the Ans (Range) when

AnsStart = 0

AnsEnd = 5

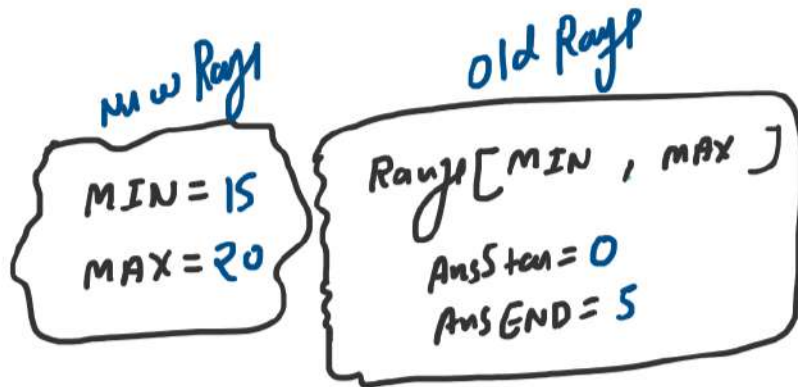
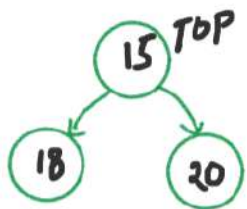
For smaller Range

$$[(\text{maxi} - \text{min}) < (\text{AnsEnd} - \text{AnsStart})]$$

$$20 - 10 < 5 - 0$$

$$10 < 5 \quad \times$$

Insert 7



		Col0	Col1	Col2	Col3	Col4
Row0	ans1	4	10	15	24	26
		0	1	2	3	4
Row1	ans2	0	9	12	20	
		0	1	2	3	
Row2	ans3	6	18	22	30	
		0	1	2	3	

(II) Insert New Element

New Element = 24

ColI = 3

RowI = 0

Update maxi = 24

(24 > 20)

(II) Fetch TOP and POP

TOP = 15

Update min = 15

Update the Ans (Range) when

AnsStart = 0

AnsEnd = 5

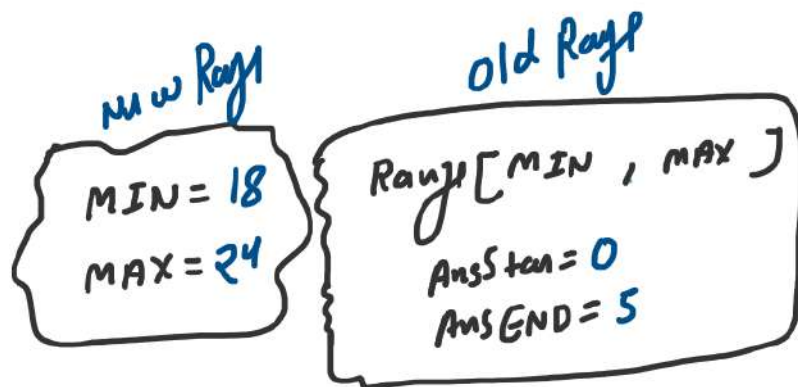
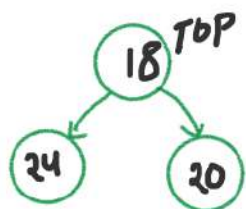
For smaller Range

$$[(\text{maxi} - \text{min}) < (\text{AnsEnd} - \text{AnsStart})]$$

$$20 - 15 < 5 - 0$$

$$5 < 5 \quad \times$$

Iter 8



		Col0	Col1	Col2	Col3	Col4
Row0	ans1	4	10	15	24	26
		0	1	2	3	4
Row1	ans2	0	9	12	20	
		0	1	2	3	
Row2	ans3	5	18	22	30	
		0	1	2	3	

⑦ Insert New Element

NewElement = 22

ColI = 2

RowI = 2

Update maxi = 24

(22 > 24) X

⑧ Fetch TOP and POP

TOP = 18

Update min = 18

Update the Ans (Range) when

AnsStart = 0

AnsEnd = 5

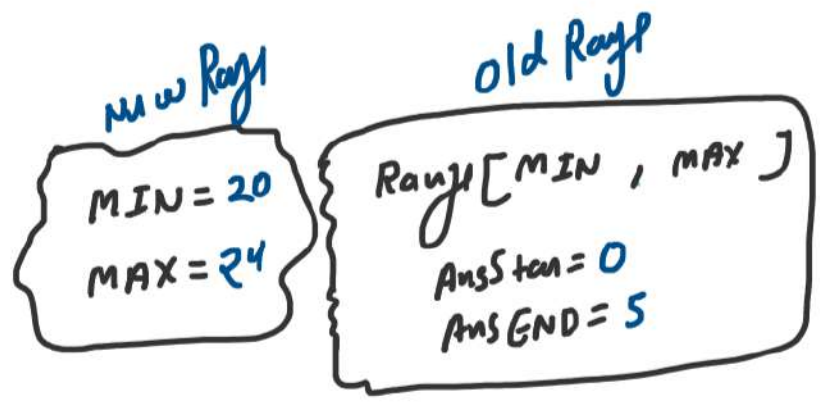
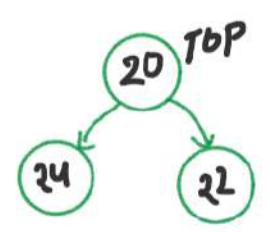
For smaller Range

$$[(\text{maxi} - \text{min}) < (\text{AnsEnd} - \text{AnsStart})]$$

$$24 - 18 < 5 - 0$$

$$6 < 5 \quad \text{X}$$

Item 9



		Col0	Col1	Col2	Col3	Col4
Row0	Ans1	4	10	15	24	26
		0	1	2	3	4
Row1	Ans2	0	9	12	20	
		0	1	2	3	
Row2	Ans3	6	18	22	30	
		0	1	2	3	

(II) Insert New Element

$R[1][4] \times$
 $ColI = 3 + 1 = 4$

STOP \Rightarrow loop ko Break kardo

(II) Fetch TOP and POP
 TOP = 20

update min = 20
 update the Ans (Range)
 AnsStart = 20
 AnsEnd = 24

For smaller Range

$$[(maxi - MIN) < (AnsEnd - AnsStart)]$$

$$24 - 20 < 5 - 0$$

$$4 < 5 \checkmark$$

Final output

Range [20, 24]
 Range [MIN, MAX]
 AnsStart = 20
 AnsEnd = 24

← Iss Range me jo Bhi Elements
 Hongi wo Har Ek
 List me Present Hongi

		Col0	Col1	Col2	Col3	Col4
Row0	Ans1	4	10	15	24	26
		0	1	2	3	4
Row1	Ans2	0	9	12	20	
		0	1	2	3	
Row2	Ans3	6	18	22	30	
		0	1	2	3	

20 21 22 23 24

20 21 22 23 24

20 21 22 23 24

```
class Solution {
public:
    // Own Data Type
    class Info
    {
        ...
    };
    // Own Comparator
    class Compare
    {
        ...
    };

    // Find smallest range
    vector<int> smallestRange(vector<vector<int>>& nums) {
        ...
    }
};
```

```
// Own Data Type
class Info
{
public:
    int data;
    int rowIndex;
    int colIndex;

    Info(int data, int rowIndex, int colIndex){
        this->data = data;
        this->rowIndex = rowIndex;
        this->colIndex = colIndex;
    }
};

// Own Comparator
class Compare
{
public:
    bool operator()(Info* first, Info* second){
        return first->data > second->data;
    }
};
```



```
// Find smallest range
vector<int> smallestRange(vector<vector<int>>& nums) {
    // Create MIN Heap
    priority_queue<Info*, vector<Info*>, Compare> pq;
    int maxi = INT_MIN;
    int mini = INT_MAX;

    // I. process first ke elements to crate the min heap [row=0,1,2,...][col=0]
    for(int i=0; i<nums.size(); i++){
        int element = nums[i][0];
        int row = i;
        int col = 0;
        Info* tempNode = new Info(element, row, col);
        pq.push(tempNode);
        maxi = max(maxi, element);
        mini = min(mini, element);
    }

    // Old Range
    int ansStart = mini;
    int ansEnd = maxi;

    // II. Fetch top and pop || update mini || update range
    while(!pq.empty()){
        ...
    }

    vector<int> ans;
    ans.push_back(ansStart);
    ans.push_back(ansEnd);
    return ans;
}
```

```
// II. Fetch top and pop || update mini || update range
while(!pq.empty()){
    Info* topNode = pq.top();
    int topData = topNode->data;
    int topRow = topNode->rowIndex;
    int topCol = topNode->colIndex;
    pq.pop();

    // Update mini
    mini = topNode->data;
    // maxi value pahle se ho updated hai
    // to ab smaller range ke liye compare kar lete hai
    int oldRangeDistance = ansEnd - ansStart;
    int newRangeDistance = maxi - mini;
    if(newRangeDistance < oldRangeDistance){
        // Update the old range
        ansStart = mini;
        ansEnd = maxi;
    }

    // III. Insert new element
    if(topCol + 1 < nums[topRow].size()){
        int newElement = nums[topRow][topCol+1];
        Info* newNode = new Info(newElement, topRow, topCol + 1);
        pq.push(newNode);

        // Update maxi
        maxi = max(maxi, newElement);
    }
    else{
        // agar koi bhi element nhi hai to Loop ko break krdo
        break;
    }
}
```

T.C.

$$\begin{cases} \text{For loop ki T.C.} = O(K) \\ \text{Heap ki T.C.} = O(\log K)^K \end{cases}$$

where
 $K = \text{Row / Numbers of lists}$

$$\rightarrow T.C. = O(K * \log K)$$

+

Range Find

$$\begin{cases} \text{while loop ki T.C.} = O(N) \\ \text{Heap ki T.C.} = O(\log K) \end{cases}$$

where
 $N = \text{Total Number of elements of each lists}$

$$\rightarrow T.C. = O(N * \log K)$$

$$\rightarrow \text{Overall T.C.} \Rightarrow O(K * \log K) + O(N * \log K)$$

$$\Rightarrow O(N * \log K)$$

Heap takes $O(K)$

space
 $\rightarrow S.C. = O(K)$