

28/11/2023

QUEUE CLASS - 3

1. First Unique Character in a String (Leetcode-387)

Ex1

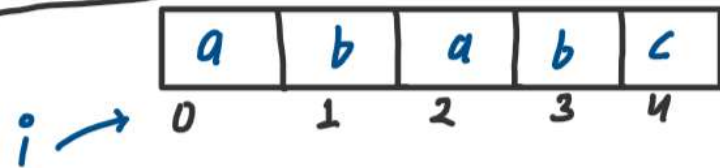
string $s = "ababc"$

Output = $aab\#c$

Ans

$aab\#c$

Explanation



$i=0$

$a \rightarrow a$

$i=1$

$a b \rightarrow a$

$i=2$

$a b a \rightarrow b$

$i=3$

$a b a b \rightarrow \#$

$i=4$

$a b a b c \rightarrow c$

First unique char

NO unique char

Ex 2

string $s = \text{"leetcode"}$

output = 1

First unique character is
l in whole string
Ans 1

Explanation

s	l	e	e	t	c	o	d	e
	0	1	2	3	4	5	6	7

$i=0$ l → 1
 $i=1$ l e → 1
 $i=2$ l e e → 1
 $i=3$ l e e t → 1
 $i=4$ l e e t c → 1
 $i=5$ l e e t c o → 1
 $i=6$ l e e t c o d → 1
 $i=7$ l e e t c o d e → 1

DRY RUN with building logic

$L \rightarrow R$

Ex1

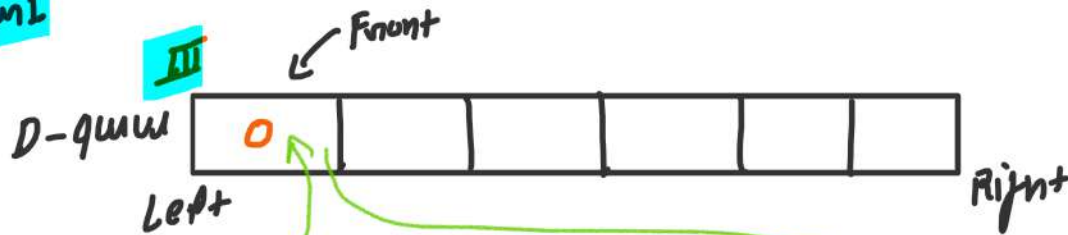
s	a	b	a	b	c
	0	1	2	3	4

I $ch = [a]$

II FREQUENCY TABLE

Char	count
a	1

Iteration 1



IV Ans \Rightarrow Unique character hai ya nahi check karengi

(I) pick $q.front() \Rightarrow 0$ index

(II) Go to the freq. Table to check the I^{th} Index char is unique or not
print = a

Item 2

S

a	b	a	b	c
-----	-----	-----	-----	-----

I $ch = [a][b]$

II FREQUENCY TABLE

Char	count
a	1
b	1

Diagram illustrating a queue implemented as an array. The array has 6 cells. The first cell contains '0' and is labeled 'Front' with an arrow. The second cell contains '1'. The first cell is also labeled 'D-queue' and 'Left' with an arrow. The last cell is labeled 'Right'.

IV Ans \Rightarrow Unique character hai ya nahi check karne hai

① 00

II aa

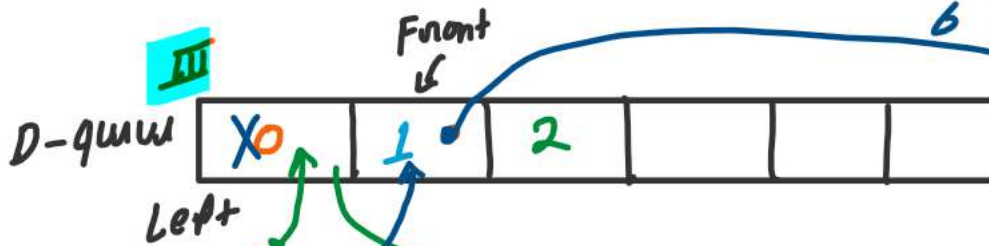
Example 3

S	a	b	a	b	c
	0	1	2	3	4

I ch = [a] [b] [a]

II FREQUENCY TABLE

Char	count
a	1 2
b	1



b is not repeating

a char repeat kar rha hai to pop index

IV Ans \Rightarrow Unique character hai ya nahi check karlenge

(I) 0 0 0 1

(II) a a b

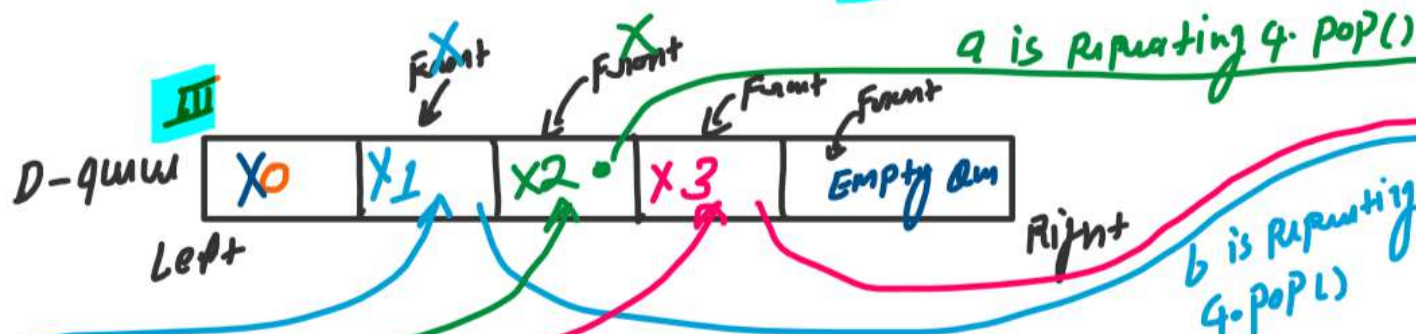
Ques 4

S a b a b c
 0 1 2 3 4

I ch = [a] [b] [a] [b] **II**

FREQUENCY TABLE

Char	Count
a	2
b	1 2



IV Ans \Rightarrow Unique character hai ya nahi check karlenge

I 0 0 0 1 1 2 3
II a a b #

Ques 5

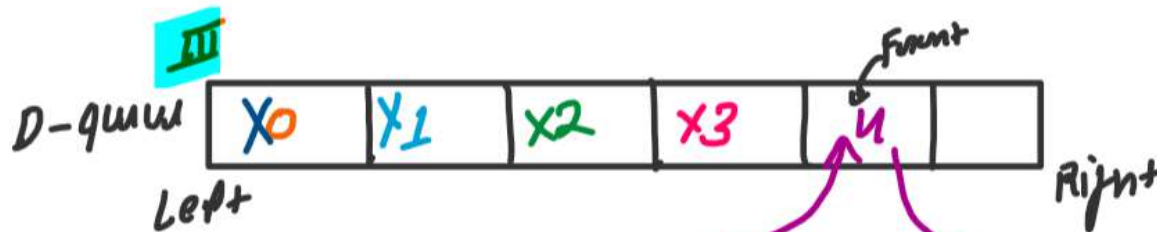
S

a	b	a	b	c
0	1	2	3	4

I ch = a b a b c

II FREQUENCY TABLE

Char	Count
a	2
b	2
c	1



IV Ans \Rightarrow Unique character hai ya nahi check karne hai

(I) 0 0 0 1 1 2 3 4
(II) a a b # c

Final output a a b # c


```

// 1. First Non Repeating/Unique Character in a String (Leetcode-387)
void firstUniqueChar(string s){
    deque<int> q;
    int freq[26] = {0};

    for (int i = 0; i < s.length(); i++)
    {
        // I -> picking current character from string
        char ch = s[i];

        // II -> storing the frequency of current character to array
        freq[ch - 'a']++;

        // III -> inserting the index of current character to queue
        q.push_back(i);

        // IV -> Ans find kro: Unique character hai ya nhi
        while(!q.empty()){
            // I -> picking the front value from queue
            char frontChar = s[q.front()];

            // II -> going to freq table to check the frontChar is unique or not
            if(freq[frontChar - 'a'] > 1){
                // Character is repeating-> no ans
                q.pop_front();
            }
            else{
                // Character is not repeating-> ans
                cout<< frontChar << " ";
                break;
            }
        }

        // Agar koi bhi ans nhi mila hai
        // iska mtlb queue empty ho chuka hai
        if(q.empty()){
            cout<< '#' << " ";
        }
    }
}

```

```

// LEETCODE PROBLEM SOLUTION CODE
class Solution {
public:
    int firstUniqChar(string s){
        deque<int> q;
        int freq[26] = {0};
        int ans = -1;

        for (int i = 0; i < s.length(); i++)
        {
            // I -> picking current character from string
            char ch = s[i];

            // II -> storing the frequency of current character to array
            freq[ch - 'a']++;

            // III -> inserting the index of current character to queue
            q.push_back(i);

            // IV -> Ans find kro: Unique character hai ya nhi
            while(!q.empty()){
                // I -> picking the front value from queue
                char frontChar = s[q.front()];

                // II -> going to freq table to check the frontChar is unique or not
                if(freq[frontChar - 'a'] > 1){
                    // Character is repeating-> no ans
                    q.pop_front();
                }
                else{
                    // Character is not repeating-> ans
                    ans = q.front();
                    break;
                }
            }

            // Agar koi bhi ans nhi mila hai
            // iska mtlb queue empty ho chuka hai
            if(q.empty()){
                ans = -1;
            }
        }

        return ans;
    }
};

```

Time Complexity: $O(N^2)$, Where N is Length of string

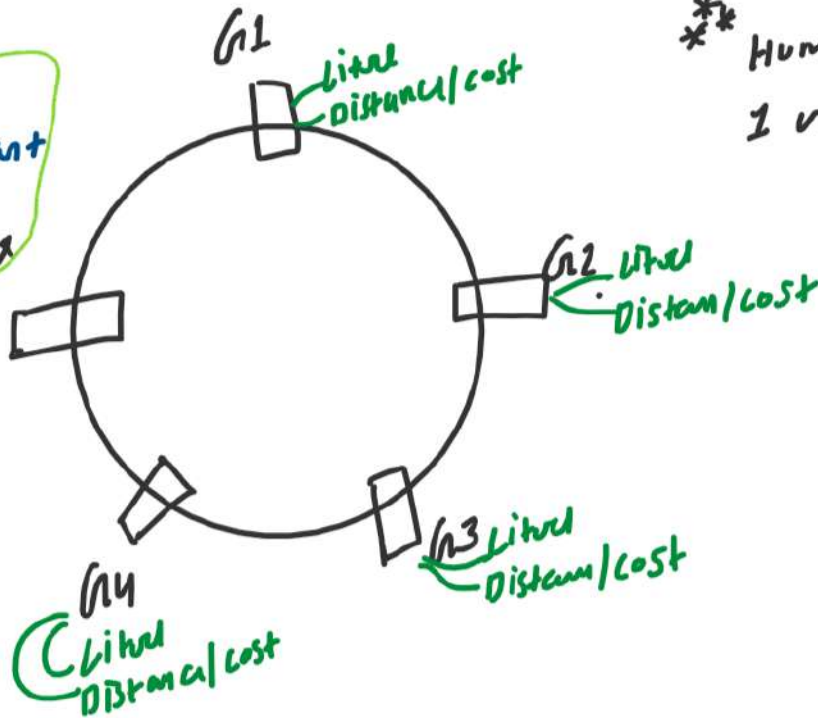
Space Complexity: $O(N)$, Where queue stores N number of characters

2. Gas Station (Leetcode-134) ***

find kya kaha hai?

man ko car n1 se chalna start
kanti hai aur wo n1 par AAKH
RUK jati hai TO Ans n1 ka Index
return kare -1 hog

Litral
Distance/cost n5



** Humari car 1 liter me
1 unit distance cover karugi only.

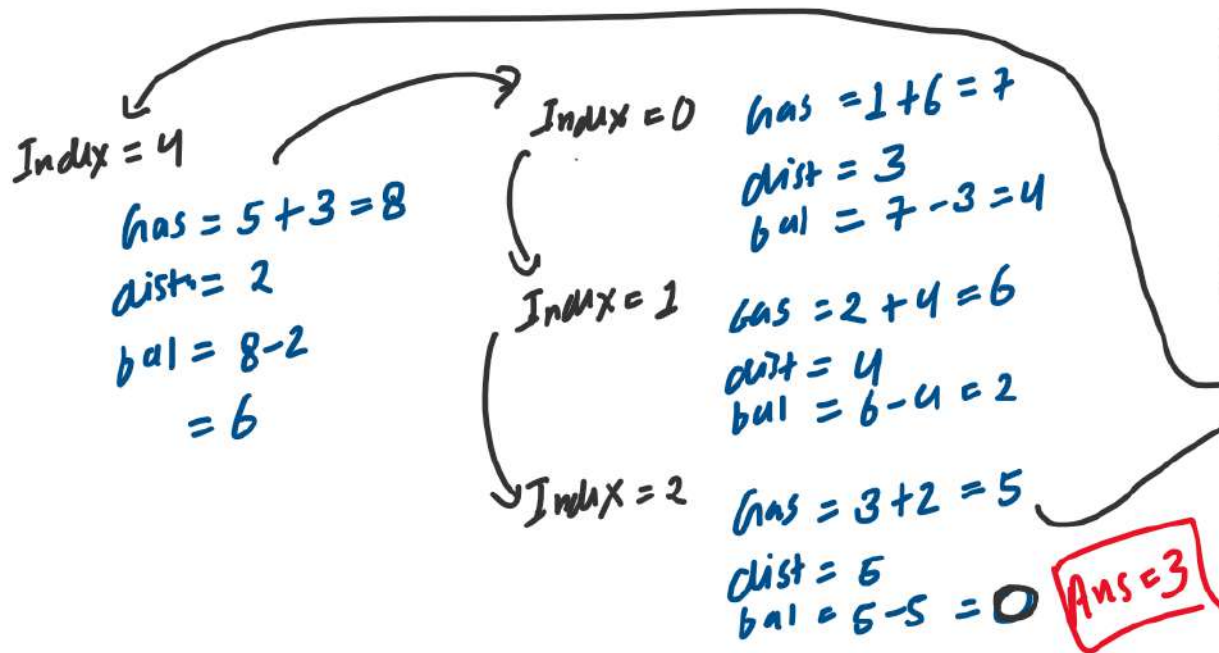
Ex 1

Input
gas

1	2	3	4	5
0	1	2	3	4

distance

3	4	5	1	2
0	1	2	3	4



BRUTE FORCE Approach

Index = 0 gas = 1
dist = 3
can not start from Index 0

Index = 1 gas = 2
dist = 4

Index = 2 gas = 3
dist = 5

Index = 3 gas = 4
dist = 1
bal = 4 - 1 = 3

$T.C = O(N^2)$

can start from Index 3 and Ans can be equal to 3

Ex2

Input
gas

2	3	4
0	1	2

distance

3	4	3
0	1	2

Output = -1

Can index 2 se start hve
but index 1 par ruk jati hai
iska matlab gas nahi rahi hogi
To Ans will be -1

BRUTE FORCE

Index = 0
X gas = 2
dist = 3

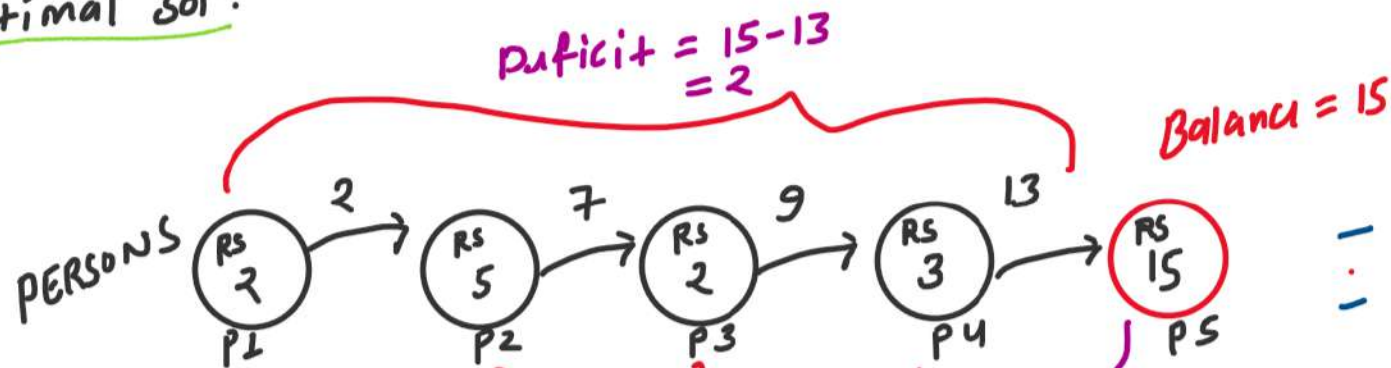
Index = 1
X gas = 3
dist = 4

Index = 2
gas = 4
dist = 3
bal = 4 - 3 = 1

Index = 0
gas = 2 + 1
dist = 3
bal = 0

Index = 1 loop stop at index 1

Optimal solⁿ.



Manlo P1 se RS na liye jaye
Tab gA PS ki help
Nahi ho sakti hai
OR le liye jaye
Tab bhi nahi ho
sakti hai

Manlo PS ki pass
RS 15 Alag hi hai
To wo khud se
Aphi help kar
sakta hai

T.C. = O(N)

- PS ko RS 15 ki need hai
- P1, P2, P3, and P4 RS collect karke PS ki help karne chahate hai

if (HELP)
 ↳ PS HAPPY
 ↳ Balance - deficit ≥ 0
 ↳ PS HAPPY

if (!HELP)
 ↳ PS NOT HAPPY
 ↳ Bal - Def < 0
 ↳ PS NOT HAPPY

Ex 1

Input	start				
	1	2	3	4	5
distanc					
	3	4	5	1	2

if (bal < 0) ← Absolutely
 ↳ deficit += bal;
 start = index + 1;
 bal = 0;

Iteration 1

deficit = 0
balance = 0
start = 0

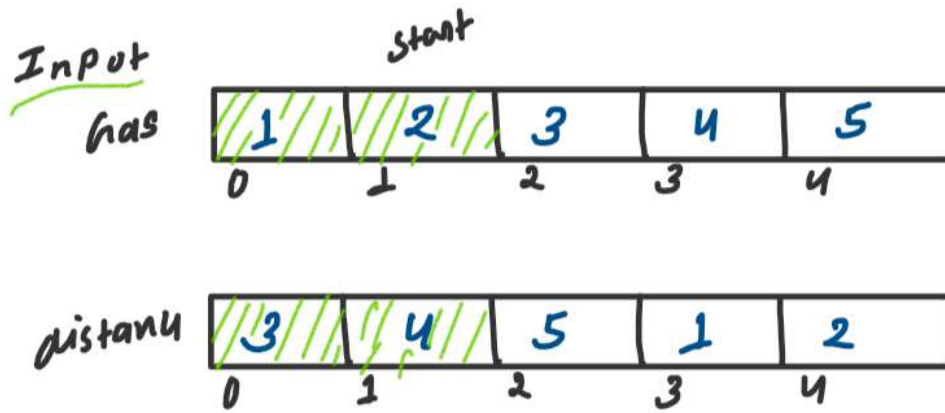
Index = 0

gas = 1

dist = 3

bal = 1 - 3 = -2

↳
$$\left(\begin{array}{l} \text{deficit} = 0 + 2 \\ \quad \quad = 2 \\ \text{start} = 1 \\ \text{bal} = 0 \end{array} \right)$$



if (bal < 0) ← Absolutely
 deficit += bal;
 start = index + 1;
 bal = 0;

Iteration 2

deficit = 2
 balance = 0
 start = 1

Index = 1

gas = 2
 dist = 4
 bal = 2 - 4 = -2

→ $\left(\begin{array}{l} \text{deficit} = 2 + 2 \\ \quad \quad = 4 \\ \text{start} = 2 \\ \text{bal} = 0 \end{array} \right)$

Input
has

			start		
1	2	3	4	5	
0	1	2	3	4	

distanc

3	4	5	1	2	
0	1	2	3	4	

if (bal < 0) ← Absolutung
 ↳ deficit += bal;
 start = index + 1;
 bal = 0;

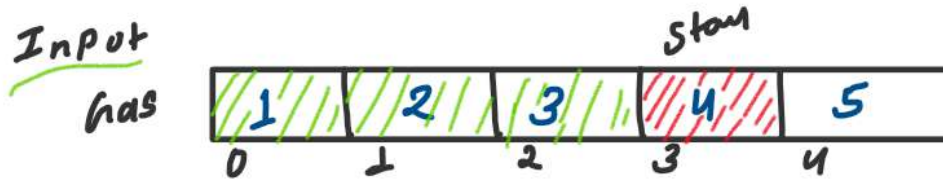
Iteration 3

deficit = 4
 balancer = 0
 start = 2

Index = 2

has = 3
 dist = 5
 bal = 3 - 5 = -2

↳ $\left(\begin{array}{l} \text{deficit} = 4 + 2 \\ \quad \quad = 6 \\ \text{start} = 3 \\ \text{bal} = 0 \end{array} \right)$



```

if(bal >= 0)
    bal += gas[index]
    index++
  
```

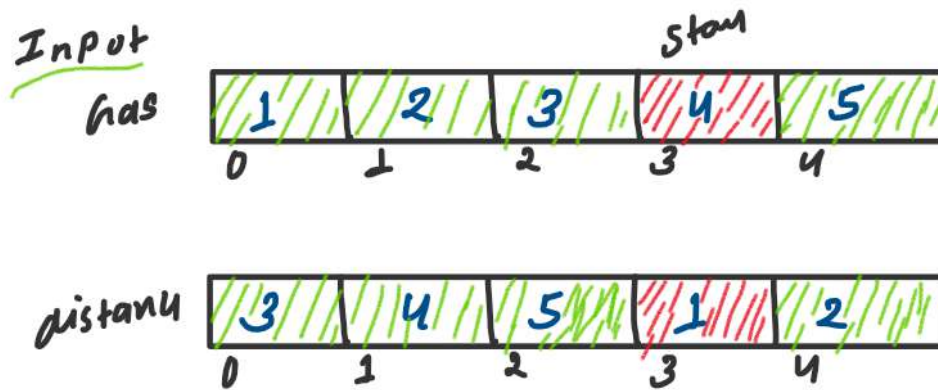
Iteration 4

deficit = 6
balance = 0
start = 3

Index = 3

gas = 4
dist = 1
bal = 4 - 1 = 3

bal += gas[index]
bal = 0 + 3
= 3



```

if (bal > 0)
    bal += bal
    index++
  
```

Iteration 5

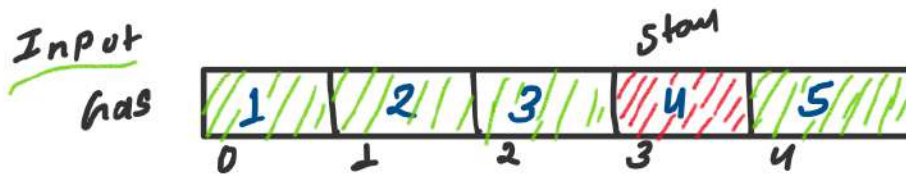
deficit = 6
balance = 3
start = 3

Index = 4

has = 5
dist = 2

bal = 5 - 2 = 3

→ $\left(\begin{array}{l} \text{bal} += \text{bal} \\ \text{bal} = 3 + 3 \\ = 6 \end{array} \right)$



Iteration 6

deficit = 6
balance = 6
start = 3

Index = 6
→ STOP THE LOOP

T.C. = $O(N)$
→ $N \Rightarrow$ length of array

if (bal - deficit == 0)
→ return start;
else if (bal - deficit < 0)
→ return -1;

Circle complete nahi
Hua Hai

CAR jaha zi chahq
start kahi thi AND
usi point par Ahar
Ruk gahi Hai
→ Iska milb circle
complete ho chuka
Hai

```

// 2. Gas Station (Leetcode-134)
// Optimal Solution Without Using Queue

class Solution {
public:
    int canCompleteCircuit(vector<int>& gas, vector<int>& cost) {
        int deficit = 0;
        int balance = 0;
        // Car start index se drive karna start karegi
        int start = 0;

        // Har ek index ko check karenge ki wo ans/starting index hai ya nahi
        for(int index = 0; index < gas.size(); index++){
            balance += gas[index] - cost[index];

            if(balance < 0){
                deficit += abs(balance);
                start = index + 1;
                balance = 0;
            }
        }

        // Car circle complete kar chuki hai
        if(balance - deficit >= 0){
            return start;
        }
        else{
            // Car circle complete nahi kar payi hai
            return -1;
        }
    }
};

```

Time Complexity: $O(N)$, Where N is size of array
Space Complexity: $O(1)$, no extra space used

3. Sliding Window Maximum (Leetcode-239)

Ex

nums

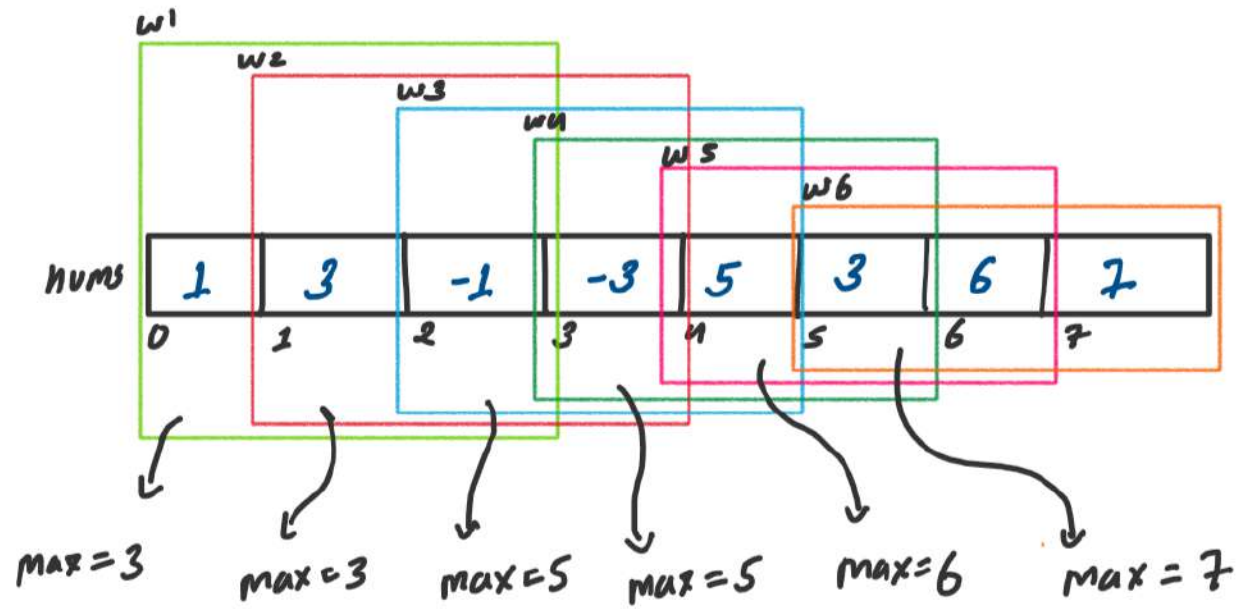
1	3	-1	-3	5	3	6	7
0	1	2	3	4	5	6	7

$K=3$

Output

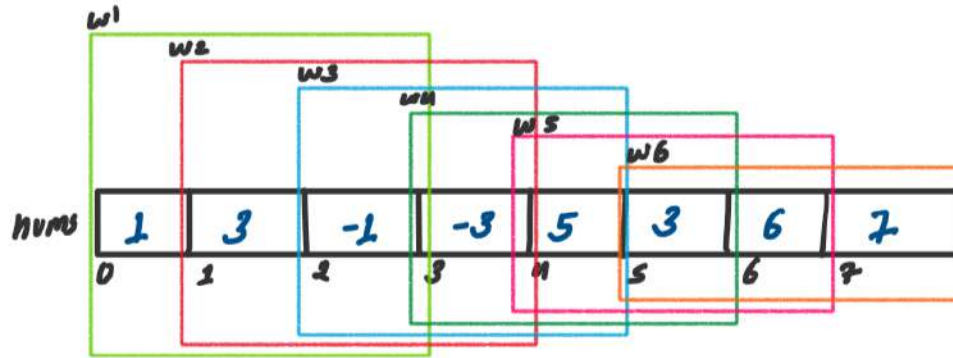
3	3	5	5	6	7
---	---	---	---	---	---

Explanation



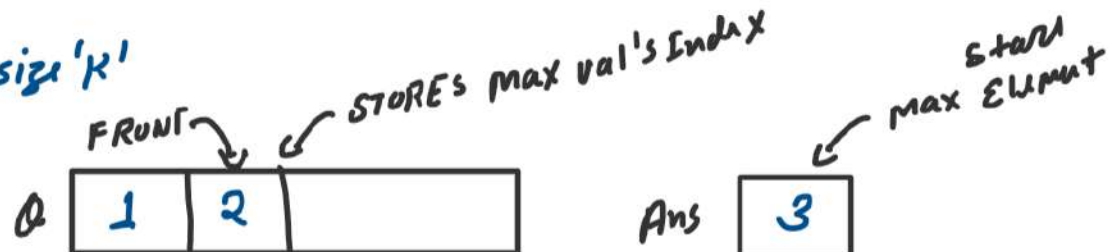
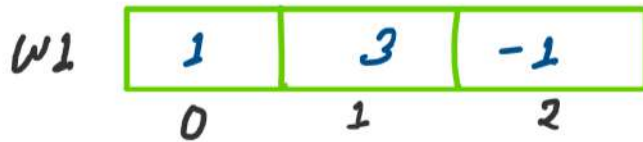
$K=3$
 $N = \text{size} = 8$

DRY RUN



$K=3$
 $N=size=8$

STEP 1 Process first window of size 'K'



YEH kaise fill ho gaya?
 ↓

STEP 1 Process first window of size 'K'

W1

1	3	-1
0	1	2

index = 0

REAR	FRONT
0	

Ans Empty

Q. pop-back()
 Q. push-back(1)

3 > 1 ← index = 1

REAR	FRONT
1	

Ans Empty

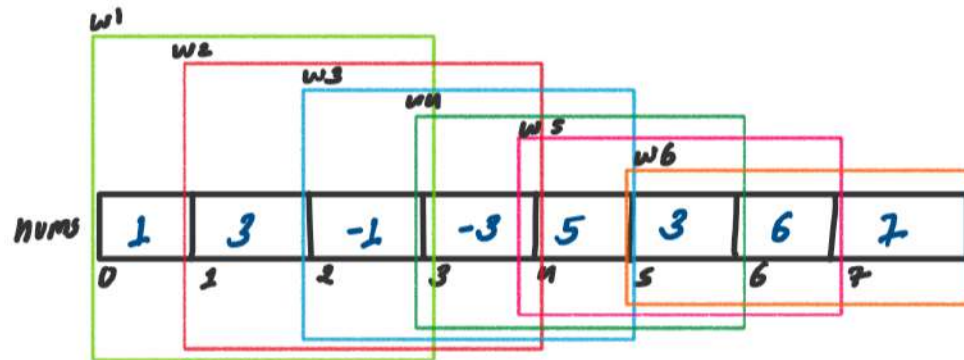
Q. push-back(2)

-1 > 3 X ← index = 2

FRONT	REAR
1	2

Ans 3 ← ans. push(3)

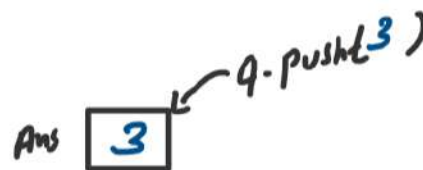
STEP 2 Process Remaining windows
REMOVE
Inclusion



$K=3$
 $N=size=8$



$\times \textcircled{I} 3 - 1 \neq 3$
 $\times \textcircled{II} -3 \neq -1$



REMOVE (in queue)
① Out of Range $[0, 1, 2]$
② Chhota Element removed
Index - $q.front + 1 \geq K$
 $q.pop_value()$ $\rightarrow q.pop_front()$

W3

-1	-3	5
2	3	4

✓ (I) $4-1 \geq 3$

✓ (II) $5-1$

Index = 4

	F		R	
X	2	3		

	FR			
X	2			

✓ (III) $5-3$

FR				
4				

q.push-back(4)

Ans

5

 ← q.push(5)

REMOVE (in queue)

- (I) Out of Range $[0, 1, 2]$
- (II) Chhota Element removed
 $\rightarrow \text{Index} - q.\text{front} + 1 \geq K$
 $\rightarrow q.\text{pop-front}();$
 $\rightarrow q.\text{pop-back}();$

W4

-3	5	3
3	4	5

X I 5-4 $\neq 3$

X II 3 $\neq 5$

Index = 5

F	R		
4	5		

Ans

5

 $\leftarrow q.push(5)$

REMOVE (in queue)

- ① Out of Range $[0, 1, 2]$
- ② Chhota Element removed
- $Index - q.front() \geq K$
- $q.pop_front();$
- $q.pop_back();$

WS

5	3	6
---	---	---

4 5 6

x (I) $6 - 4 \geq 3$

Index = 6 @

F	R		
4	5		

✓ (II) 6 > 4

FR
@

4			
---	--	--	--

✓ (II) 6 > 5

FR
@

6			
---	--	--	--

q.push-back(6)

Ans

6

 ← q.push(6)

REMOVE (in queue)

- (I) Out of Range [0, 1, 2]
- (II) Chhota Element removed

Index - q.front() + 1 >= K
 ↳ q.pop-front();
 ↳ q.pop-back();

W6

3	6	7
5	6	7

X (I) $7 - 6 > 3$

✓ (II) 7 7 6

Index = 7 @

FR	6			
----	---	--	--	--

@

FR	7			
----	---	--	--	--

q.push-back(7)

Ans

7

 ← q.push(7)

REMOVE (in queue)

- (I) Out of Range $[0, 1, 2]$
- (II) Chhoti Element removed

→ Index - q.front + 1 = K
 → q.pop-front();
 → q.pop-back();

```
// 3. Sliding Window Maximum (Leetcode-239)

class Solution {
public:
    vector<int> maxSlidingWindow(vector<int>& nums, int k) {
        deque<int> q; // store the max element's index
        vector<int> ans; // store the max element

        // Step 1: process the first window for "k time"
        for(int index = 0; index < k; index++){
            int element = nums[index];

            // Agar queue me element chotta hai
            while(!q.empty() && element > nums[q.back()]){
                q.pop_back();
            }

            // Yanha tabhi pahuch skta hu
            // Ya to queue me element chotta nhi hai
            // Ya queue empty ho chuka hai
            q.push_back(index);
        }

        // Step 2: process remaning windows
    }
};
```

```
// Step 2: process remaning windows
for(int index = k; index < nums.size(); index++){
    // Purani window ka ans store kardo
    ans.push_back(nums[q.front()]);

    // Remove
    // I -> remove the out of range index from queue
    if(!q.empty() && index - q.front() >= k){
        q.pop_front();
    }

    // II -> remove chotta index from queue
    // Agar queue me element chotta hai
    while(!q.empty() && nums[index] > nums[q.back()]){
        q.pop_back();
    }

    // Addition
    // Yanha tabhi pahuch skta hu
    // Ya to queue me element chotta nhi hai
    // Ya queue empty ho chuka hai
    q.push_back(index);
}

// Last window ka ans store karlo
ans.push_back(nums[q.front()]);

return ans;
```

Time Complexity: $O(N)$,
where N is size of array

Space Complexity: $O(K)$,
where K is the size of the window