



Ordering the Display of Tuples

- List in alphabetic order the names of all instructors
select distinct *name*
from *instructor*
order by *name*
- We may specify **desc** for descending order or **asc** for ascending order, for each attribute; ascending order is the default.
 - Example: **order by** *name desc*
- Can sort on multiple attributes
 - Example: **order by** *dept_name, name*



Where Clause Predicates

- SQL includes a **between** comparison operator
- Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is, $\geq \$90,000$ and $\leq \$100,000$)
 - **select** *name*
 from *instructor*
 where *salary* **between** 90000 **and** 100000
- Tuple comparison
 - **select** *name, course_id*
 from *instructor, teaches*
 where (*instructor.ID, dept_name*) = (*teaches.ID, 'Biology'*);



Set Operations

- Find courses that ran in Fall 2009 or in Spring 2010

```
(select course_id from section where sem = 'Fall' and year = 2009)  
union  
(select course_id from section where sem = 'Spring' and year = 2010)
```

- Find courses that ran in Fall 2009 and in Spring 2010

```
(select course_id from section where sem = 'Fall' and year = 2009)  
intersect  
(select course_id from section where sem = 'Spring' and year = 2010)
```

- Find courses that ran in Fall 2009 but not in Spring 2010

```
(select course_id from section where sem = 'Fall' and year = 2009)  
except  
(select course_id from section where sem = 'Spring' and year = 2010)
```



Set Operations

- Set operations **union**, **intersect**, and **except**
 - Each of the above operations automatically eliminates duplicates
- To retain all duplicates use the corresponding multiset versions **union all**, **intersect all** and **except all**.

Suppose a tuple occurs m times in r and n times in s , then, it occurs:

- $m + n$ times in r **union all** s
- $\min(m, n)$ times in r **intersect all** s
- $\max(0, m - n)$ times in r **except all** s



Null Values

- It is possible for tuples to have a null value, denoted by *null*, for some of their attributes
- *null* signifies an unknown value or that a value does not exist.
- The result of any arithmetic expression involving *null* is *null*
 - Example: $5 + \text{null}$ returns null
- The predicate **is null** can be used to check for null values.
 - Example: Find all instructors whose salary is null.

```
select name
from instructor
where salary is null
```



Null Values and Three Valued Logic

- Any comparison with *null* returns *unknown*
 - Example: $5 < \text{null}$ or $\text{null} <> \text{null}$ or $\text{null} = \text{null}$
- Three-valued logic using the truth value *unknown*:
 - OR: $(\text{unknown} \text{ or } \text{true}) = \text{true}$,
 $(\text{unknown} \text{ or } \text{false}) = \text{unknown}$
 $(\text{unknown} \text{ or } \text{unknown}) = \text{unknown}$
 - AND: $(\text{true} \text{ and } \text{unknown}) = \text{unknown}$,
 $(\text{false} \text{ and } \text{unknown}) = \text{false}$,
 $(\text{unknown} \text{ and } \text{unknown}) = \text{unknown}$
 - NOT: $(\text{not } \text{unknown}) = \text{unknown}$
 - “*P* is unknown” evaluates to true if predicate *P* evaluates to *unknown*
- Result of **where** clause predicate is treated as *false* if it evaluates to *unknown*



Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

avg: average value

min: minimum value

max: maximum value

sum: sum of values

count: number of values



Aggregate Functions (Cont.)

- Find the average salary of instructors in the Computer Science department
 - **select avg** (*salary*)
from *instructor*
where *dept_name*= 'Comp. Sci.';
- Find the total number of instructors who teach a course in the Spring 2010 semester
 - **select count** (**distinct** *ID*)
from *teaches*
where *semester* = 'Spring' **and** *year* = 2010
- Find the number of tuples in the *course* relation
 - **select count** (*)
from *course*;



Aggregate Functions – Group By

- Find the average salary of instructors in each department
 - **select** *dept_name*, **avg** (*salary*)
from *instructor*
group by *dept_name*;
 - Note: departments with no instructor will not appear in result

<i>ID</i>	<i>name</i>	<i>dept_name</i>	<i>salary</i>
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

<i>dept_name</i>	<i>avg_salary</i>
Biology	72000
Comp. Sci.	77333
Elec. Eng.	80000
Finance	85000
History	61000
Music	40000
Physics	91000



Aggregation (Cont.)

- Attributes in **select** clause outside of aggregate functions must appear in **group by** list
 - */* erroneous query */*
select *dept_name, ID, avg (salary)*
from *instructor*
group by *dept_name;*



Aggregate Functions – Having Clause

- Find the names and average salaries of all departments whose average salary is greater than 42000

```
select dept_name, avg (salary)
from instructor
group by dept_name
having avg (salary) > 42000;
```

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups



Null Values and Aggregates

- Total all salaries

```
select sum (salary )  
from instructor
```

- Above statement ignores null amounts
- Result is *null* if there is no non-null amount
- All aggregate operations except **count(*)** ignore tuples with null values on the aggregated attributes
- What if collection has only null values?
 - count returns 0
 - all other aggregates return null



Nested Subqueries

- ❑ SQL provides a mechanism for the nesting of subqueries.
- ❑ A **subquery** is a **select-from-where** expression that is nested within another query.
- ❑ A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.



Example Query

- Find courses offered in Fall 2009 and in Spring 2010

```
select distinct course_id
from section
where semester = 'Fall' and year= 2009 and
       course_id in (select course_id
                      from section
                      where semester = 'Spring' and year= 2010);
```

- Find courses offered in Fall 2009 but not in Spring 2010

```
select distinct course_id
from section
where semester = 'Fall' and year= 2009 and
       course_id not in (select course_id
                          from section
                          where semester = 'Spring' and year= 2010);
```



Example Query

- Find the total number of (distinct) students who have taken course sections taught by the instructor with *ID* 10101

```
select count (distinct ID)  
from takes  
where (course_id, sec_id, semester, year) in  
        (select course_id, sec_id, semester, year  
         from teaches  
         where teaches.ID= 10101);
```

- **Note:** Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.



Set Comparison

- Find names of instructors with salary greater than that of some (at least one) instructor in the Biology department.

```
select distinct T.name  
from instructor as T, instructor as S  
where T.salary > S.salary and S.dept_name = 'Biology';
```

- Same query using > **some** clause

```
select name  
from instructor  
where salary > some (select salary  
                        from instructor  
                        where dept_name = 'Biology');
```




Definition of Some Clause

- $F \text{ <comp> some } r \Leftrightarrow \exists t \in r \text{ such that } (F \text{ <comp> } t)$
Where <comp> can be: $<$, \leq , $>$, $=$, \neq

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{true}$ (read: 5 < some tuple in the relation)

$(5 < \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 = \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true}$

$(5 \neq \text{some } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline \end{array}) = \text{true (since } 0 \neq 5)$

$(= \text{some}) \equiv \text{in}$

However, $(\neq \text{some}) \not\equiv \text{not in}$



Example Query

- Find the names of all instructors whose salary is greater than the salary of all instructors in the Biology department.

```
select name  
from instructor  
where salary > all (select salary  
                        from instructor  
                        where dept_name = 'Biology');
```



Definition of all Clause

□ $F \text{ <comp> all } r \Leftrightarrow \forall t \in r (F \text{ <comp> } t)$

$(5 < \text{all } \begin{array}{|c|} \hline 0 \\ \hline 5 \\ \hline 6 \\ \hline \end{array}) = \text{false}$

$(5 < \text{all } \begin{array}{|c|} \hline 6 \\ \hline 10 \\ \hline \end{array}) = \text{true}$

$(5 = \text{all } \begin{array}{|c|} \hline 4 \\ \hline 5 \\ \hline \end{array}) = \text{false}$

$(5 \neq \text{all } \begin{array}{|c|} \hline 4 \\ \hline 6 \\ \hline \end{array}) = \text{true (since } 5 \neq 4 \text{ and } 5 \neq 6)$

$(\neq \text{all}) \equiv \text{not in}$

However, $(= \text{all}) \not\equiv \text{in}$



Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists** $r \Leftrightarrow r \neq \emptyset$
- **not exists** $r \Leftrightarrow r = \emptyset$

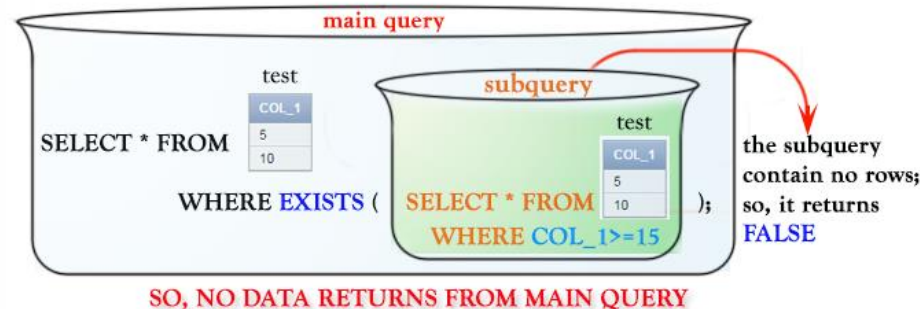
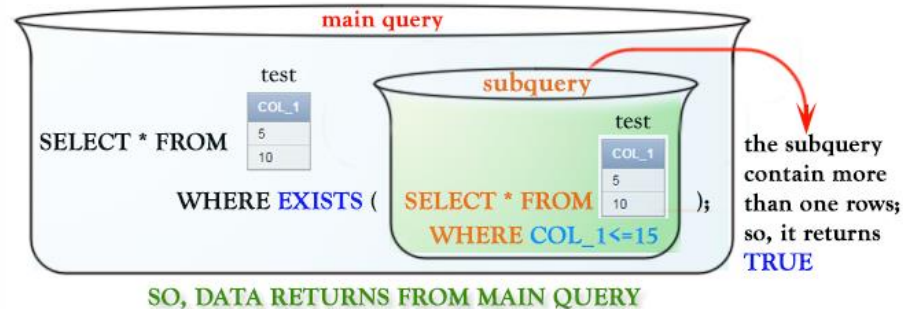


Test for Empty Relations

WHERE EXISTS (**subquery**);

EXISTS operator

- EXISTS is a Comparison operator
- Used in WHERE clause to validate an "IT EXISTS" condition.
- EXISTS will tell you whether a query returned any results.
- Returns a BOOLEAN, (TRUE or FALSE).
- Returns TRUE if a subquery contains any rows.



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Correlation Variables

- Yet another way of specifying the query “Find all courses taught in both the Fall 2009 semester and in the Spring 2010 semester”

```
select course_id
from section as S
where semester = 'Fall' and year= 2009 and
      exists (select *
              from section as T
              where semester = 'Spring' and year= 2010
              and S.course_id= T.course_id);
```

- **Correlated subquery**
- **Correlation name** or **correlation variable**



Not Exists

- Find all students who have taken all courses offered in the Biology department.

```
select distinct S.ID, S.name  
from student as S  
where not exists ( (select course_id  
                    from course  
                    where dept_name = 'Biology')  
except  
                  (select T.course_id  
                   from takes as T  
                   where S.ID = T.ID));
```

- Note that $X - Y = \emptyset \Leftrightarrow X \subseteq Y$
- Note: Cannot write this query using = **all** and its variants



Subqueries in the From Clause

- ❑ SQL allows a subquery expression to be used in the **from** clause
- ❑ Find the average instructors' salaries of those departments where the average salary is greater than \$42,000.

```
select dept_name, avg_salary
from (select dept_name, avg (salary) as avg_salary
      from instructor
      group by dept_name)
where avg_salary > 42000;
```

- ❑ Note that we do not need to use the **having** clause
- ❑ Another way to write above query

```
select dept_name, avg_salary
from (select dept_name, avg (salary)
      from instructor
      group by dept_name)
      as dept_avg (dept_name, avg_salary)
where avg_salary > 42000;
```




With Clause

- The **with** clause provides a way of defining a temporary view whose definition is available only to the query in which the **with** clause occurs.
- Find all departments with the maximum budget

```
with max_budget (value) as  
    (select max(budget)  
     from department)  
select budget  
from department, max_budget  
where department.budget = max_budget.value;
```



Complex Queries using With Clause

- With clause is very useful for writing complex queries
- Supported by most database systems, with minor syntax variations
- Find all departments where the total salary is greater than the average of the total salary at all departments

```
with dept_total (dept_name, value) as
    (select dept_name, sum(salary)
     from instructor
     group by dept_name),
dept_total_avg(value) as
    (select avg(value)
     from dept_total)
select dept_name
from dept_total, dept_total_avg
where dept_total.value >= dept_total_avg.value;
```



Scalar Subquery

- Scalar subquery is one which is used where a single value is expected
- E.g.

```
select dept_name,  
      (select count(*)  
       from instructor  
       where department.dept_name = instructor.dept_name)  
      as num_instructors  
from department;
```
- E.g.

```
select name  
from instructor  
where salary * 10 >  
      (select budget from department  
       where department.dept_name = instructor.dept_name)
```
- Runtime error if subquery returns more than one result tuple



Modification of the Database

- Deletion of tuples from a given relation
- Insertion of new tuples into a given relation
- Updating values in some tuples in a given relation



Modification of the Database – Deletion

- Delete all instructors

delete from *instructor*

- Delete all instructors from the Finance department

delete from *instructor*
where *dept_name*= 'Finance';

- Delete all tuples in the *instructor* relation for those instructors associated with a department located in the Watson building.

delete from *instructor*
where *dept_name* in (**select** *dept_name*
from *department*
where *building* = 'Watson');



Deletion (Cont.)

- Delete all instructors whose salary is less than the average salary of instructors

delete from *instructor*
where *salary* < (**select avg** (*salary*) **from** *instructor*);

- Problem: as we delete tuples from deposit, the average salary changes
- Solution used in SQL:
 1. First, compute **avg** salary and find all tuples to delete
 2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)



Modification of the Database – Insertion

- Add a new tuple to *course*

```
insert into course  
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- or equivalently

```
insert into course (course_id, title, dept_name, credits)  
values ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
```

- Add a new tuple to *student* with *tot_creds* set to null

```
insert into student  
values ('3003', 'Green', 'Finance', null);
```



Insertion (Cont.)

- Add all instructors to the *student* relation with *tot_creds* set to 0

insert into *student*

select *ID, name, dept_name, 0*

from *instructor*

- The **select from where** statement is evaluated fully before any of its results are inserted into the relation (otherwise queries like

insert into *table1* **select * from** *table1*

would cause problems, if *table1* did not have any primary key defined.



Modification of the Database – Updates

- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others receive a 5% raise
 - Write two **update** statements:

```
update instructor
  set salary = salary * 1.03
  where salary > 100000;
update instructor
  set salary = salary * 1.05
  where salary <= 100000;
```
 - The order is important
 - Can be done better using the **case** statement (next slide)



Case Statement for Conditional Updates

- Same query as before but with case statement

update *instructor*

set *salary* = **case**

when *salary* <= 100000 **then** *salary* * 1.05

else *salary* * 1.03

end



Updates with Scalar Subqueries

- Recompute and update `tot_creds` value for all students

update *student S*

set *tot_cred* = (**select** **sum**(*credits*)
 from *takes* **natural join** *course*
 where *S.ID* = *takes.ID* **and**
 takes.grade <> 'F' **and**
 takes.grade **is not null**);

- Sets *tot_creds* to null for students who have not taken any course
- Instead of **sum**(*credits*), use:

case

when **sum**(*credits*) **is not null** **then** **sum**(*credits*)

else 0

end