

Q How to actually solve to get complexity?
Ans Few Few & ways :-

1) Plug & Chug

OR

2) Master's theorem.

OR

3) Akra Bazzi § 1996 § :- solve any

** Akra Bazzi :-

$$\text{formula: } T(n) = O\left(n^p + n^p \int_1^n \frac{g(u)}{u^{p+1}} du\right)$$

Words \hookrightarrow Time complexity = $T(n)$

$g(u)$:- this "g funt" is basically a time complexity
it-self. & we know that in time complexity
like constants & more dominating terms
are ignored.

we know that this going to be the simplest
form of the function because all the
less dominating terms and other things
will be removed.

P:

$$a_1 b_1^P + a_2 b_2^P + \dots = 1$$

$$\text{i.e. } \sum_{j=1}^k a_j b_j^P = 1 //$$

$$T(N) = T(N/2) + C$$

Q In constants why do we write $O(1)$?

Ans Because we don't really care about constants
for eg:- we can write $O(1)$ or $O(2k+c)$
all of these are constant, \therefore we can ignore the constants

\therefore Basically $O((2k+c) \times 1)$

\therefore Anything multiplied by anything, you can ignore the constant constants.

So, Hence anything in constants can be written as $O(1)$.

$$\text{Eg: 1) } T(N) = 2T\left(\frac{N}{2}\right) + (N-1)$$

$$a_1 = 2$$

$$b_1 = \frac{1}{2}$$

$$g(x) = x^{N-1}$$

$$\therefore 2 \times \left(\frac{1}{2}\right)^P = 1$$

$$\therefore P = 1$$

Once you've found the P then substitute it into the Akra-Bazzi formula:-

$$T(n) = O\left(n' + n' \int_1^n \frac{u-1}{u^2} du\right)$$

$$= O\left(n + n \int_1^n \frac{1}{u} - \frac{1}{u^2} du\right)$$

$$= O\left(n + n \left[\int_1^n \frac{du}{u} - \int_1^n \frac{du}{u^2} \right]\right) \quad \begin{matrix} = u^{-2} \\ = -\frac{1}{u} \end{matrix}$$

$$= O\left(n + n \left[(\log u) + \left[\frac{1}{u}\right] \right]\right)$$

↑
put ↑.

$$= O\left(n + n \left[\log n + \frac{1}{n} - 1 \right]\right)$$

$$= O\left(n + n \log n + \frac{1}{\cancel{n}} - \cancel{n}\right)$$

$$= O(n \log n + 1)$$

$$\therefore O(n \log n) \quad // \text{ Time complexity.}$$

So, for array of size N :-
Merge sort complexity = $O(N \log N)$.

$$2) \quad T(N) = 2T\left(\frac{N}{2}\right) + \frac{8}{9}T\left(\frac{3N}{4}\right) + N^2$$

$$\therefore 2 \times \left(\frac{1}{2}\right)^P + \frac{8^2 \times \frac{3}{4}}{9^3} = 1$$

$$1 + \frac{2}{3} \text{ so, put } P = 2$$

$$2 \times \frac{1}{4^2} + \frac{8}{9} \times \frac{69}{16} = 2$$

$$\therefore \frac{1}{2} + \frac{1}{2} = 1$$

$$\therefore P = 2 //$$

Substitute

$$T(n) = O\left(n^2 + n^2 \int_1^n \frac{x^2}{3x^3} dx\right)$$

$$= O\left(n^2 + n^2 \log n\right) \text{ ignore the less dominant term}$$

$$= O\left(n^2 \log n\right)$$

If you're unable to find value of P :

$$①. T(x) = 3T\left(\frac{x}{3}\right) + 4T\left(\frac{x}{4}\right) + x^2$$

② $P=1$ Case 1

$$= 3 \times \left(\frac{1}{3}\right) + 4 \times \left(\frac{1}{4}\right) = 1$$

$$\therefore 1 + 1 = 1$$

$$2 \neq 1$$

$\therefore 2 > 1$ This means

what? I need to increase the denominator

③ $P=2$ Case 2

$$3 \times \frac{1}{9} + 4 \times \frac{1}{16} = 1$$

$$\frac{1}{3} + \frac{1}{4} = \frac{7}{12} < 1$$

Hence, P is less than 2.

So, P actually lies in rd range 1 & 2

Note:

When $P < \text{power of } (g(x))$ then your
ans = $g(x)$.

Here, $g(x) = x^2$

$P < 2$ {i.e. power of $g(x)$ }

Hence, ans = $O(g(x))$.

$$T(n) = O \left(n^p + n^p \int_1^n \frac{u^2}{u^{p+1}} du \right)$$

$$= O \left(n^p + n^p \int_1^n u^{1-p} du \right)$$

$$= O(n^p + n^2)$$

$$\therefore p < 2$$

$\therefore n^p$ will become less dominating term n^2 Hence, ignore.

$$\therefore O(n^2) //$$

Hence proved.