

# CPSC 335 – Lecture 2

Efficiency Analysis Part 1  
(ADITA Ch 3.1-3.5)

# Measuring Resource Use (3.2)

Resource Types:

- **Time**
- *Space*
- I/O bandwidth
- Cache misses
- Energy

# Worst-Case Analysis (3.2.1)

**Complexity:** the amount of a **resource** consumed when an algorithm runs on a specific **instance** of a problem

**Worst-case complexity:** **maximum** amount of resource the algorithm may consume

- Related to **instance (= input) size**: running on bigger instances causes the algorithm to consume more resources

# Worst-Case Analysis (3.2.1)

Why worst case?

- The algorithm *might* use up only this much resource, but maybe more (e.g. take *at least* X seconds to run)

vs

- The algorithm will definitely not use more than this much resource (e.g. take at most X seconds to run)

Which is more useful?

# Complexity Functions (3.2.2)

**Size measure:** size of a problem's **instances** (e.g. number of elements in a list)

**Complexity function:** maps the **size measure** (= instance size, e.g.  $n$ ,  $m$ ) to the algorithm's **complexity** (= amt of resource consumed)

- Written in terms of the size measure, e.g.  $f(n)$
- Can never be negative
- Time complexity function:  $T(n)$

# Asymptotic Notation (3.3)

$O(f(n))$   $\approx$  *efficiency class*

$= \{g(n) \mid \exists c > 0, t \geq 0 \text{ such that } g(n) \leq c * f(n) \forall n \geq t\}$

- $g(n) \leq c * f(n)$

$g(n)$  = any given instance (of size  $n$ ) complexity

$f(n)$  = worst-case complexity

$c$  = constant factor, irrelevant to efficiency analysis  
(implementation code vs. algorithm pseudocode)

- $\forall n \geq t$

$t$  = threshold of  $n$  small enough ( $\geq 0$ ) to not have same complexity behavior as all  $n$  larger than threshold

# Definition of $O$ (3.3.1)

$O(f(n)) = O(n)$  ?

- $O(n)$ : set of all functions equivalent to  $f(n) = n$
- $O(n^2)$ : set of all functions equivalent to  $f(n) = n^2$
- $O(f(n)) = O(n)$ 
  - $7n \in O(n)$
  - $2n + 3 \in O(n)$
- $O(f(n)) = O(n^2)$ 
  - $3n^2 \in O(n^2)$

# Big Eight (3.4)

- $O(1)$ : constant
- $O(\log n)$ : logarithmic
- $O(n)$ : linear
- $O(n^2)$ : quadratic
- $O(n^3)$ : cubic
- $O(c^n)$ : exponential
- $O(n!)$ : factorial



# Experimental Analysis (3.5)

- **Experimental analysis:** deriving complexity functions by using scientific method to gather and analyze **observed data** (implementation running actual code)
- **Mathematical analysis:** deriving complexity functions via **models** (pseudocode) and **proofs**
- Projects = experimental analysis, homework = mathematical analysis

# Experimental Analysis (3.5)

1. **Question:** What is the time efficiency class of algorithm A?
2. **Hypothesis:** Algorithm A runs in  $O(n)$  time.
3. **Prediction:** Multiple runs of A on various instance (= input) sizes  $n$  will have runtimes that follow the trend of a straight line (linear)
4. **Testing:** implement A and run the code on various input sizes, plot the runtime results
5. **Analysis:** analyze scatterplot for linear fit

# Experimental Analysis (3.5)

## Practical considerations

- **Instrumental error:** inaccuracies due to the method of measurement
  - is the instrument sensitive enough? (seconds vs milliseconds vs nanoseconds)
  - is the measurement being affected by outside circumstances? (other programs running)
- **Benchmarks:** specific inputs and measurement criteria for meaningful real-world results
  - Hard to choose good ones
- **Random instances:** easier to generate, but real-world use may have non-random instance distribution