

Homework 5

Problem 1

1. The set of strings over the alphabet $\Sigma = \{a, b\}$ with more a's than b's.

$$R_1: S \rightarrow XaX$$

$$R_2: X \rightarrow aXb \mid bXa \mid XX \mid a \mid \epsilon$$

2. The complement of the language $\{a^n b^n \mid n \geq 0\}$

$$R_1: S \rightarrow Xa \mid bX \mid aSb$$

$$R_2: X \rightarrow aX \mid bX \mid \epsilon$$

Problem 2

$$A = \{a^i b^j c^k \mid i=j \text{ or } j=k \text{ where } i, j, k \geq 0\}$$

$$A_1 = \{a^i b^j c^k \mid i=j \text{ where } i, j, k \geq 0\}$$

$$A_2 = \{a^i b^j c^k \mid j=k \text{ where } i, j, k \geq 0\}$$

CFG of Language A, Union of languages A_1 and A_2

$$R_1: S \rightarrow MX \mid YN$$

$$R_2: M \rightarrow aMb \mid \epsilon$$

$$R_3: N \rightarrow bNc \mid \epsilon$$

$$R_4: X \rightarrow cX \mid \epsilon$$

$$R_5: Y \rightarrow aY \mid \epsilon$$

This grammar is ambiguous because it is the union of languages A_1 and A_2 which has an intersection that is of all strings $a^i b^i c^i$ where i, j, k are all equal. This means if MX is chosen it can lead to this as well as if YN is chosen in rule 1.

(Problem 2 continued)

For example two different left most derivations lead to the same string abc .

$$\textcircled{1} \quad S \Rightarrow MX \Rightarrow aMbX \Rightarrow aEbX \Rightarrow aEbCX \Rightarrow aEbCE = abc$$

$$\textcircled{2} \quad S \Rightarrow YN \Rightarrow aXN \Rightarrow aEN \Rightarrow aEbNC \Rightarrow aEbEN = abc$$

Problem 3

$$R_1: A \rightarrow BAB|B|E$$

$$R_2: B \rightarrow \emptyset|E$$

Step 0 create a new start variable

$$R_1: S \rightarrow A$$

$$R_2: A \rightarrow BAB|B|E$$

$$R_3: B \rightarrow \emptyset|E$$

Step 1 remove ϵ -productions

$$R_1: S \rightarrow A|E$$

$$R_2: A \rightarrow BAB|B|AB|BA|A|BB$$

$$R_3: B \rightarrow \emptyset$$

Step 2 remove Unit Production in R_2

$$R_2: A \rightarrow BAB|\emptyset|AB|BA|BB$$

Step 3 remove Unit Production in R_1

$$R_1: S \rightarrow BAB|\emptyset|AB|BA|BB|E$$

$$R_1: S \rightarrow BAB|\emptyset|AB|BA|BB|E$$

$$R_2: A \rightarrow BAB|\emptyset|AB|BA|BB$$

$$R_3: B \rightarrow \emptyset$$

Step 4 move all terminals to unit productions where RHS is one terminal

$$R_1: S \rightarrow BAB|xx|AB|BA|BB|\epsilon$$

$$R_2: A \rightarrow BAB|xx|AB|BA|BB$$

$$R_3: B \rightarrow xx$$

$$R_4: x \rightarrow 0$$

Step 5 create two variable RHS, by replacing long production rules

$$R_1: S \rightarrow YB|xx|AB|BA|BB|\epsilon$$

$$R_2: B \rightarrow YB|xx|AB|BA|BB$$

$$R_3: B \rightarrow xx$$

$$R_4: x \rightarrow 0$$

$$R_5: Y \rightarrow BA$$