

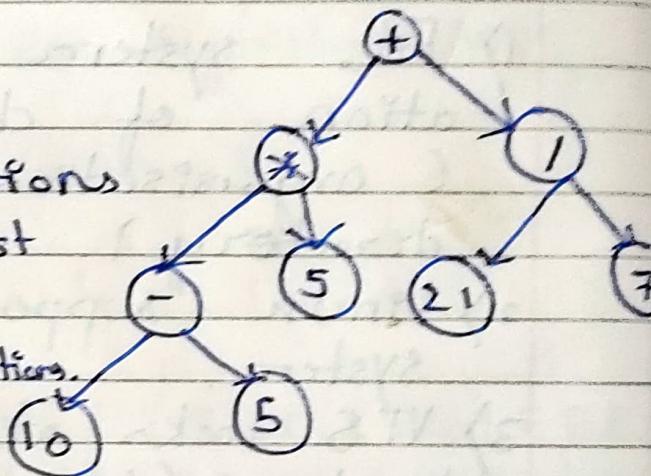
May 14

## ADS Unit 6.

## 9.1 Evaluation of expression tree.

logic :

do the operations  
in bracket, first  
preference to  
leaf nodes operations.



$$\rightarrow ((10 - 5) * 5) + (21 / 7)$$

$$\rightarrow 25 + 3$$

$$\rightarrow 28$$

Let's get into the processing part

```
public static double PROCESS
  (string op, double x, double y)
{
```

```
  if (op.equals("+"))
    return x + y;
```

basic struct q3

class Node {

    String val

    Node left

    Node right

    Node (String val) { this.val = val;

        this.left = this.right = null; }

} else if (op.equals("-")) {

    return x - y;

} else if (op.equals("\*")) {

    return x \* y;

} else if (op.equals("/")) {

    return x / y;

return 0;

}

public static double evu (Node root)

{ if (root == null) {

    return 0;

}

} if (isLeaf (root)) {

    return Double.parseDouble (root.val);

}

// recursively checking

double x = evu (root.left);

double y = evu (root.right);

return PROCESS (root.val, x, y);

}

## Hard coded method.

```
public int evu(Node node) {  
    if (node.value.equals("/")) {  
        return evu(node.left) / evu(node.  
            Right);  
    }  
    else if (node.value.equals("*")) {  
        return evu(node.left) * evu(node.  
            Right);  
    }  
    else if (node.value.equals("-")) {  
        return evu(node.left) - evu(node.  
            Right);  
    }  
    else if (node.value.equals("+")) {  
        return evu(node.left) + evu(node.  
            Right);  
    }  
    else {  
        return Integer.valueOf(node.val)  
    }  
}
```

Trick : make a complete binary tree (heap) and traverse traversal to get Max / Min heap.

## Q2 Heap and priority queue.

### a] Heap :-

→ A heap is a nearly complete binary tree with the following two properties:

1) All levels are full except the last one.

2) For Max heap :  $\text{parent}(n) \geq x$   
For Min heap :  $\text{parent}(n) \leq x$

→ In An Array

root :  $A[1]$

left child :  $A[i] = A[2i]$

Right child :  $A[i] = A[2i+1]$

Parent of  $A[i]$  :  $A[i] : A[i/2]$

→ Three operations can be done on heaps

→ Max - heapify (Restore max heap property)

→ build - Max - heap (building Max Heap)

→ Heapsort (sort an array)

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↳ Algo Max-heapify ( $A, i, n$ )

$l \leftarrow \text{left}(i)$

$r \leftarrow \text{Right}(i)$

if ( $l \leq n$  and  $A[l] > A[i]$ )

then largest  $\leftarrow l$

else largest  $\leftarrow i$

if ( $r \leq n$  and  $A[r] > A[\text{largest}]$ )

then largest  $\leftarrow r$

if largest  $\neq r$

then exchange  $A[i] \leftrightarrow A[\text{largest}]$

Max-heapify ( $A, \text{largest}, n$ )

Algo Build Max-heap ( $A$ )

$n = \text{length}[A]$

for  $i = n/2, i--$

do Max-heapify ( $A, i, n$ )

Priority queue

→ It is an ADT for maintaining a set  $S$  of elements with each element have a priority with it.

→ An element with high priority is served before an element with low priority and vice versa.

- If elements have same priority, they are served as their order in the queue.
- It has operations
  - 1) push(x) : inserting element x
  - 2) top or peek() : returns the element of S with highest or lowest priority.
  - 3) pop() : returns and pops the top element.

### 1.3 Maps :-

- Algorithm used here is Dijkstra's
- It is used in G Maps to find the shortest, easiest path to travel
- Algo : we will be having vertices and their weight  $G = (V, E)$  and we have to find shortest path to destination from source.

- 1> All paths are assumed to be  $\infty$  except the  $d[s]$ . which will be zero.
  - 2> We will be finding new paths and updating  $d[v]$  for all  $v \in \text{Adj}[u]$ . This is called as relaxation.
  - 3> Then we need to change the old path to new shorter path if  $d[u] + w(u, v) < d[v]$   
then  
$$d[v] = d[u] + w[u, v].$$
- 4> This is how it works.

## Unit 5.

### Q.1 File class with example

- 1> Acts like a wrapper class for file name.
- 2> A file name like "num.txt" has only string properties.
- 3> File has some very useful methods, they are as follow

- exists : tests if a file exist.
- canRead : tests if OS will let you read a file.
- canWrite : tests if the OS will let you write a file.
- delete : deletes the file.
- length : returns the no of bytes in that file.
- getName : returns the name of file.

↳ Eg.

```
File numFile = new File ("num.txt");
if (numFile.exists ()) {
    System.out.println (numFile.length ());
}
```

## 2. Serialization in file system.

→ To serialize an object means to convert its state to a byte stream so that the byte stream can be reverted back into a copy of

- As I/O stream accepts only binary it doesn't understand text.
- In order to tackle this problem whatever we want to accept and display we put it in a class that implements serializable.
- Java assigns a serial number to each object written out.
- The serializable class has class instance variable then they should also be serializable.

### Q.3 File parameter

- We can give the path of the file as parameter
- ```
File f = new File("Path");
        ^ name of file
parse(f.getAbsolutePath());
```
- This can be passed to methods

- Using PrintWriter we use write method we provide it string as parameter

File Writer myWriter = new

FileWriter ("Assi.txt")

myWriter.write ("Write to a file");  
myWriter.close();

→ During reading we just have to provide the path of file to file reader

File Reader reader = new FileReader

("My XD.txt")  
File name ↑

while ((character = reader.read()) != -1) {

System.out.print ((char) character)

; } // reads the program

}

reader.close(); // closing

Fil. IO stream and buffered reader

→ Java does provide strong, flexible support for I/O as it relates to files and network

- Java's I/O system is cohesive and consistent.
- Program Data is retrieved from input source.
- Program Results are sent to an o/p destination.
- Input stream : a stream that provides I/O to program. `System.in` is an input stream.  
Output stream : accept o/p from program. `System.out` is an o/p stream.

→ Code :-

```
import java.util.Scanner;  
class Stream {  
    public static void main (String  
        args []) {  
        Scanner sc = new Scanner (System.in);  
        System.out.println ("Enter a letter");  
        String s = sc.nextLine();  
    }  
    String s = sc.nextLine();
```

o/p  
stream

## b] Buffer Readers

→ To read using buffer means like reading in chunks.

→ There is a little delay but it's better than non buffered methods.

→ Disk operations per buffer of byte -- lower overhead.

→ Code :

```
import java.io.*;  
  
class mHead {  
    public static void main (String  
        args []) throws IOException {  
        char c;  
        BufferedReader br =  
            new BufferedReader (new  
                InputStreamReader (System.in));  
        System.out.println (" Enter  
            characters. In 'q' to quit ");  
        do {  
            c = br.read ();  
            System.out.println (c);  
        } while (c != 'q');  
    }  
}
```

## Q.5 File reader and File writer class

### a] File reader

1) FileReader is a subclass of Reader.

2) We have to import java.io.Reader. This is about how we create a FileReader.

Reader input = new FileReader();

3) It can take a txt file as parameter.

4) [read( char [] arr, int st, int len )] - reads the no of char equal to length and stores it into array.

### b] Print Writer (File Writer)

1) FileWriter is a subclass of OutputStreamWriter.

2) Syntax :

FileWriter fw = new FileWriter  
( );

3) In its constructor we can give it a path or an object of file class.

↳ Main method

`write (String txt)` → It is used to write the string into FileWriter.

`flush ()` → flushes data of FileWriter.

`close ()` → used to close FileWriter.

## Unit 4.

### Q.1 Hashing

→ Key, value pairs are stored in fixed size table called a hash table.

→ hash table is partitioned into many buckets.

→ Each bucket has many slots. Each slot holds record.

→ A hash  $f^n(x)$  transforms the identifier (key) into an address in the hash table.

Algo :-

a] Inserting :-

→ For inserting a new record, key must somehow be converted to array index.

→ Suppose  $n = 10$  (size of table)

we have to insert 5

$$5 \bmod 10 = 5$$

→ 5 will be inserted at 5<sup>th</sup> position.

b] Searching

→ We calculate the hash value first.

→ Later we check that location for the key, if the value is not there we continue to check array till end.

c) deleting :-

→ Records may also be deleted from hash table

parallel to each other

→ The only one thing is that the records should not be left empty.

### Collision resolution methods :

→ The spot where we try to insert an element onto an occupied spot

Not for 5% of rest ) H = 18

→ It can be resolved by collision resolution techniques

1) Quadratic Probing :-

$$\rightarrow f(q) = q^2$$

→ In quadratic probing, if  $i$ -th position is occupied we check the  $i+1$ st next we check the  $i+4$ th next  $i+9$ th etc.

2) Quad linear probing :-

$$\rightarrow f(i) = i$$

*When there is a collision we just probe the next slot in the table.*

3) Double hashing :-

*Firstly we need to find the first key.*

$$k_1 = H(\text{Key} \% \text{size of table}(M))$$

The second hash function we need to find value in range from  $k_1$  to capacity - 1

$$H(k_2) \text{hash}_2(\text{obj}) = 1 + (k_1 \% M)$$

$\therefore 59 = 1027 \leftarrow \frac{\text{size of table}}{1}$

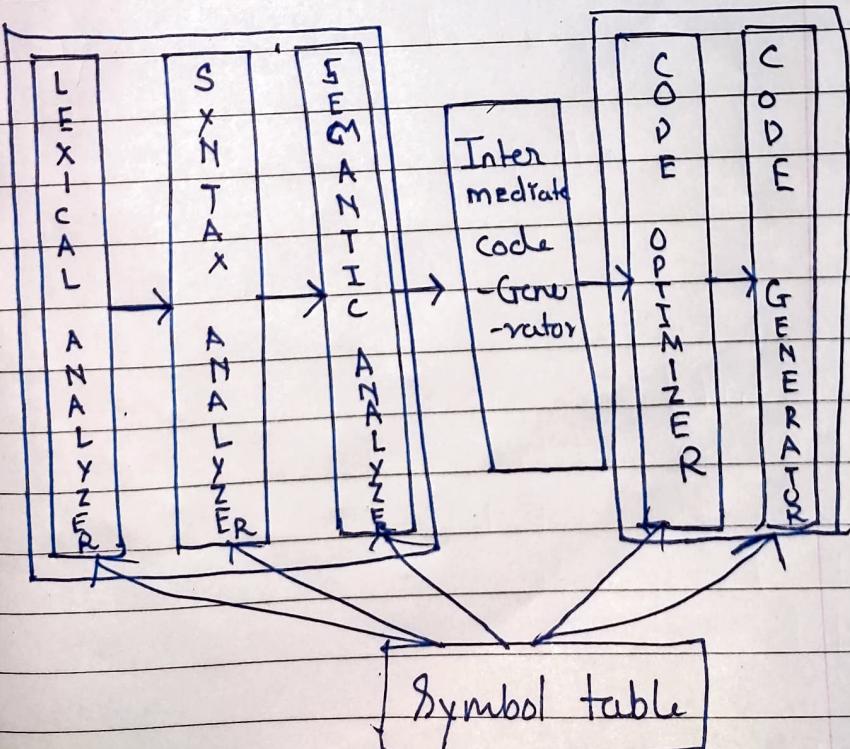
4) Open hashing (chaining)

*The key concept here is all elements that hash to a particular bucket are placed on that bucket's linked list.*

## 2 Symbol table

- Symbol table is a data structure used by a language translator such as a compiler or interpreter.
- It stores information about various source language.
- The information is collected by analysis phase of the compiler and used by the synthesis phases to generate the output which is in machine language.

→



→ Symbol table has three operations

- 1) insert()
- 2) Scope Management
- 3) lookup()

→ Each table is list of names and their associated attribute and the table are organized into a stack.