Explanation CSE221 Assignment 06

Task 1: In this task, dijokana algorithm is used to find the shortest paths of all modes from the stant node of a directed weighted graph. In the dijkstra function, weighted graph. In the dijkstra function, we used a priority queue to stone the use used a priority queue to stone the nodes and its current distance from pq, we to eliminate the shortest distance from pq, we used heap data structure. Inside while loop, used heap data structure. Inside while loop, used heap data structure. Inside while loop, used heap data structure. We also stored a fictionary. We also stored a fictionary we distionary.

Task 2: This task is all most same as twok 1, but we used dijkstra method twice 1, but we used dijkstra method twice for the two nodes griven in import file. These two nodes are the start nodes. We stored the two nodes are the start nodes, we stored the for nodes in two different node 4 istar ces in two different node dictionaryes. Then we checked if there is dictionaryes.

cony common nodes in both distance dictionanies with presper values. Then we find out the total distance of that node for distance I and the total distance differences 2. We also find out the distance differences of both dictionanies for that node. Then we of both dictionanies for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance for all common nodes as this difference of distance of the selected node. In the given condition the formation of the selected node. In the given condition of the selected node. In the given condition of the selected node.

Tasks: This time, a molified dijustra function is used. This time, we find out the maximum between tistance from Stant to u (node) and fur local distance of v from u. Then we add it with priority queue with v (neighbour nede) and continue the loop. And finally we check if the distance is still infinity or not. It is infinity, then we return not of the distance value.