Explanation (Assignment 4) ID: 22101489

Task 1 (a): To print the graph which is

directed and weight as adjacency matrix

we need a list and its elements will also

we need a list and its elements will be and

be lists. There initial elements will be armited that the

ust will be a O. After toking the inputs.

We will be and them in the output file,

we need to use a nested loop.

1(b) To point the directed and weighted graph,

We need a dictionary. After taking m, N grown

We need a dictionary. After take the other values

the input file, we will take the other in

the input file and store them in

times from the input file and dictionary [ui] = (VI, WI).

Then we will store them in the output file

then we will store them in the output file

with a nested, loop.

Task &: To perform bfs (breadth first search) we need a queue structure. After taking the inputs as directed, the graph will be stored as a diedionary. Then weitwill be sent to the bfs method with needs and with the result array, where the result will will array the method will implement the arguments along the method will implement the arguments along the method will implement the regult entray. Then it will be abled the regult entray. Then it will be abled the regult entray.

There was A) property X to

Tagk 3: DFS we is an recentariste algorithm where the depost ventex of a node will be discovered first. To implement the method, we use a recursive method. As before we store the graph from the imput of we create another a l'initial al. dietionary we store the initial color of the store the dfs method, we add it each node as a perform that, we add it we take a node as a Before that, we add it we take a not find the venter, and then it's venter, and then it's venter, and then we take the regult we will the venter as parameter. It is the regult we will the venter as parameter. It is the regult we will the venter as parameter. finally, we change resturn the rosm H.

Task 4: We use the dfg method to find a cycle in a greaph. We use to Hedionarry to storce the greath itself and to store the colon of nodes. After storing them the greaph and putting the initial color, we call the necursive If;
we have used a global Hag to
method. We used a will change
thore output. The method will change
thore output. the nodes ealon as 'Si' at the bogining. The nodes ealon as so not the bogining.

If any of it's vantex's there's a eycle.

That means there's a eycle.

The sold in flag and roturned. It is the end of the nodes colon will be charged to so.

Task 5: BFS method is used to check the

Task 5: BFS method is used to check the stored from the stored from the bfs method shortest path the graph will be stored from the input as di chromany then the input as di chromany then mode. While will be earlied with destination inode. While will be earlied with destination on not. If it is trea varising in the method will execute in guine. The node is the destination on not if it is the node is the path stored in guine. It will rectum Wone. If the node is not found, it will rectum Wone.

Tagk 6: To find the number of Liamonds, we will use of method. The greap will be Stoned as grid, in adjacency matrix style.

To find the first diamond as stand, we will

To find the first diamond as stand, we will use two nested loops to After & finding it, we will call the Hs method with the index position of it as parameter. At Jinst the method will check the conditions given led the method will check if the selected in question. Then it will check if the value in question. Then it will check if the related in question. Then it will recent will recent then it will recent value is it recently. Then it will recent then of counter it related for the 4 neighbours of counter call itself for the 4 neighbours. reconsively call itself for the 4 neighbourn the counter.

The wall return the wall of counter.

Task 7: To find the two nodes whe have the longest distance, we use the bis method. the longest distance, we so the faithest valued the faithest we find out the faithest want then we again can the stant. Then we again can the stant and find that no de as stant and find by with that no de as stant wo are the its farthest node. These two are the its farthest value. output value.