

Last time

Sign-implementation

- The Boolean function $F: \{0,1\}^n \rightarrow \{0,1\}$ is *sign-implemented* by Q_F^\pm if it maps $|x\rangle|00\dots 0\rangle$ to $(-1)^{F(x)}|x\rangle|0\dots 0\rangle$ for every $x \in \{0,1\}^n$
- Denote $f: \{0,1\}^n \rightarrow \{\pm 1\}$ by $f(x) = (-1)^{F(x)}$
- “Rotate”: Initialize n qubits $|0\rangle$ ’s. Put them through $H^{\otimes n}$. Obtain

$$|+\rangle \otimes \dots \otimes |+\rangle = \frac{1}{\sqrt{N}} \sum_{x \in \{0,1\}^n} |x\rangle, \text{ where } N = 2^n$$

- Definition: Uniform superposition.

Rotate, compute, rotate

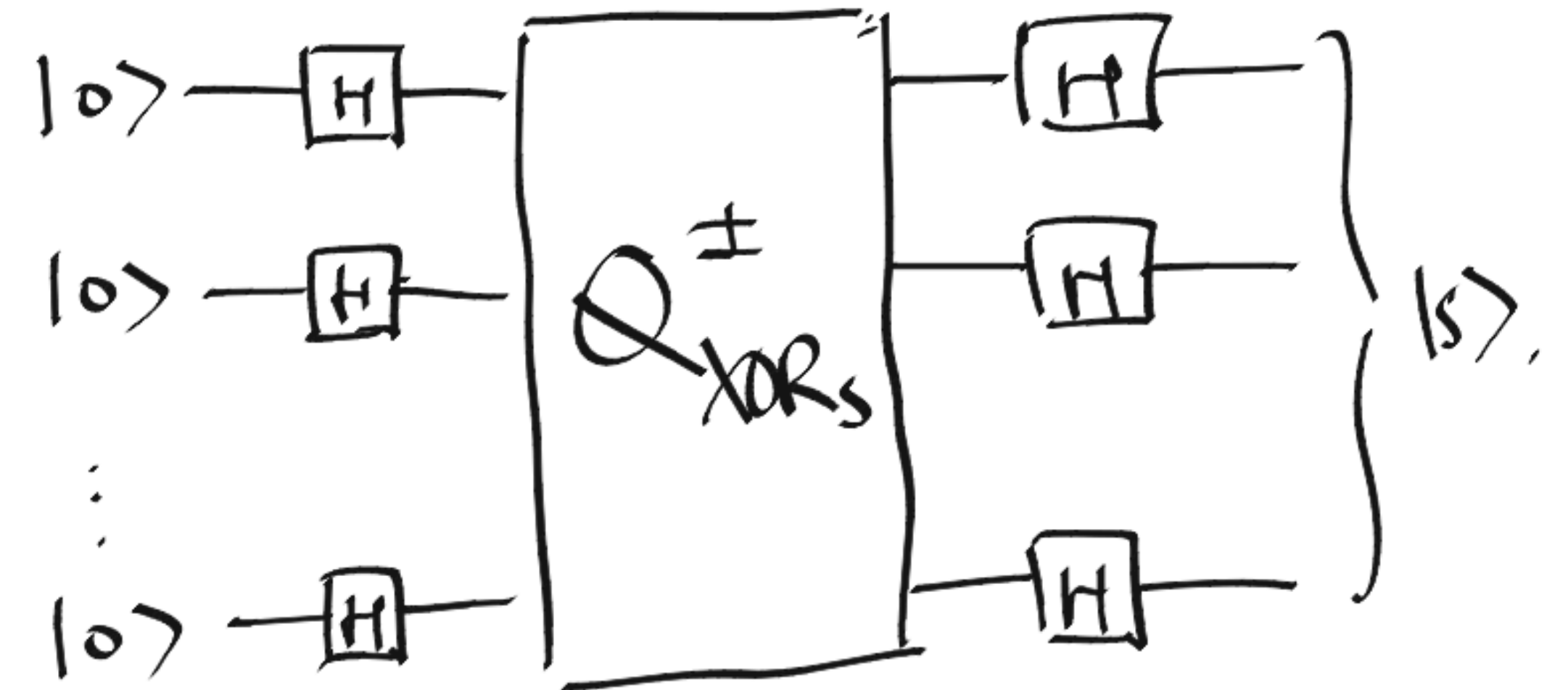
- Compute: Plug uniform superposition into Q_F^\pm and obtain ...
- Rotate: Apply Boolean Fourier Transform $H^{\otimes n}$ again, and obtain ...

- What is $H^{\otimes n}|x\rangle$ for $x \in \{0,1\}^n$?

- Example

- $x = 00\dots 0$

- $x = 011$



- $H^{\otimes n}|x\rangle = \frac{1}{\sqrt{N}} \sum_{s \in \{0,1\}^n} \pm |s\rangle$. What is the sign on $|s\rangle$?

What is the sign on $|s\rangle$?

- In general, the sign on $|s\rangle$ is a product of n signs
 - If $s_i = 0$, the i -th sign is always $+$
 - If $s_i = 1$, the i -th sign is $(-1)^{x_i}$
 - Upshot: the i -th sign is $(-1)^{s_i x_i}$
 - Overall, the sign is $(-1)^{\sum_i x_i s_i} = (-1)^{x \cdot s}$
- Theorem: For every $x \in \{0,1\}^n$,

$$H^{\otimes n}|x\rangle = \frac{1}{\sqrt{N}} \sum_{s \in \{0,1\}^n} (-1)^{x \cdot s} |s\rangle$$

- Theorem: For every $x \in \{0,1\}^n$,

$$H^{\otimes n}|x\rangle = \frac{1}{\sqrt{N}} \sum_{s \in \{0,1\}^n} (-1)^{x \cdot s} |s\rangle$$

- Corollary

$$H^{\otimes n} \left(\frac{1}{\sqrt{N}} \sum_{x \in \{0,1\}^n} (-1)^{x \cdot s} |x\rangle \right) = |s\rangle$$

- Application: Someone hands you a sign-implementation of XOR_s , defined by $\text{XOR}_s(x) = x \cdot s \bmod 2$ without telling you s . Only need to feed in 1 input to learn s with 100% accuracy!

Comparison

Classical inputs only

- $x \mapsto \text{XOR}_s(x) = x \cdot s \bmod 2$
- How many applications of XOR_s are needed?

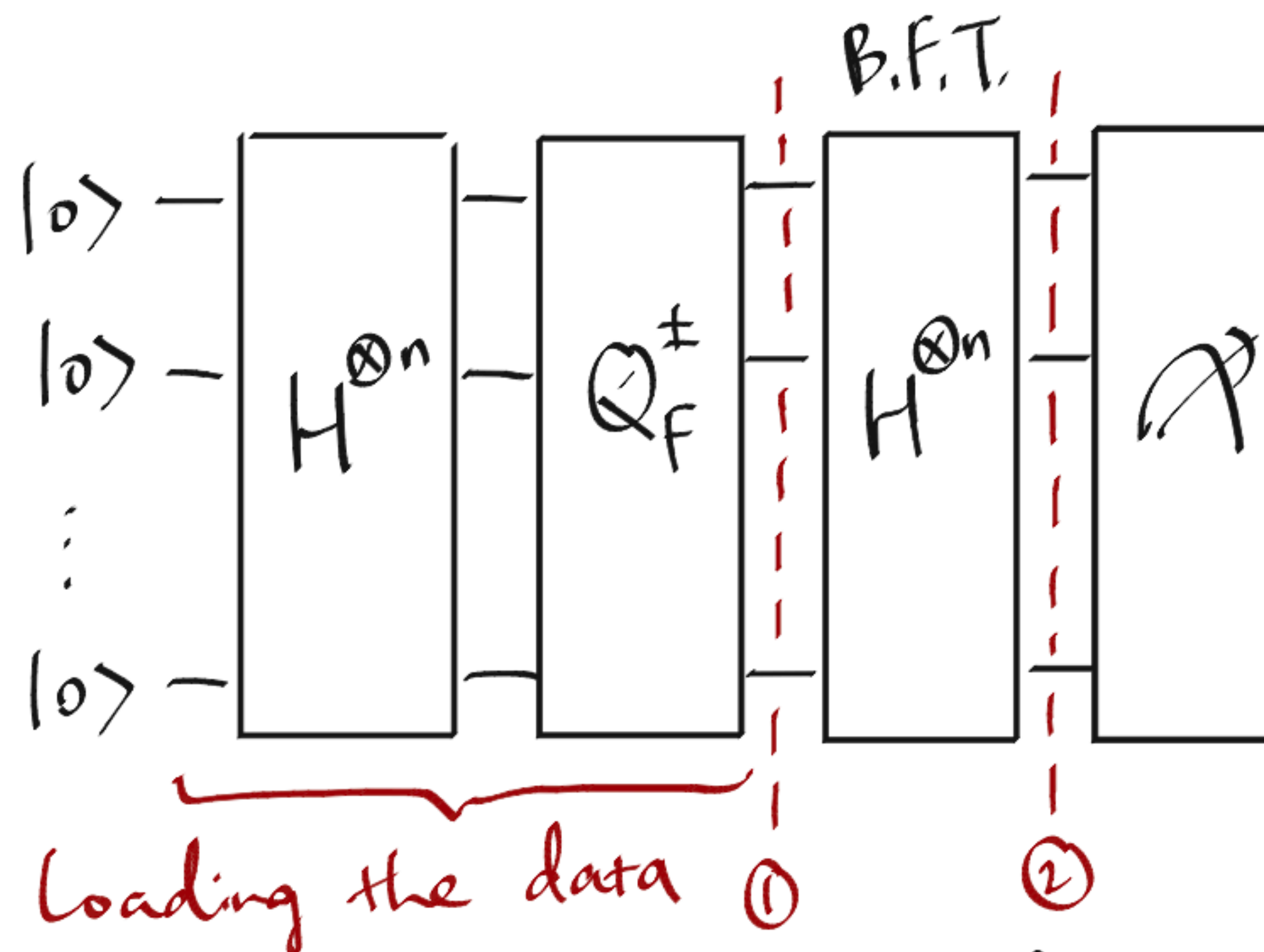
The general framework

to find patterns in implicitly-represented data

- Data vector $|g\rangle \in \mathbb{C}^N$ of length N
- Apply Fourier transform with “pattern vectors” $|\chi_0\rangle, \dots, |\chi_{N-1}\rangle$
- Obtain another vector of length N , where s -th entry is “strength of $|\chi_s\rangle$ pattern in the data vector”.
- Quantum:
 - $N = 2^n$ (say $n = 1000$)
 - Vectors implicitly represented by n -qubit state

- Pattern vectors $|\chi_0\rangle, \dots, |\chi_{N-1}\rangle$ can be any orthonormal basis of \mathbb{C}^N
- “Strength of patterns $|\chi_0\rangle, \dots, |\chi_{N-1}\rangle$ in $|g\rangle$ ” is just the coefficients when $|g\rangle$ is represented in $|\chi_0\rangle, \dots, |\chi_{N-1}\rangle$ basis, that is, $\langle \chi_s | g \rangle$.
- Let $U \in \mathbb{C}^{N \times N}$ be the matrix with columns $|\chi_0\rangle, \dots, |\chi_{N-1}\rangle$
 - U is unitary
 - U maps the standard basis to the χ -basis
 - U^{-1} or U^\dagger maps the χ -basis to the standard basis
 - U^\dagger maps $|g\rangle$ to ...
- Want:
 - Interesting / useful” pattern vectors
 - Associated change of basis U easy to implement by quantum gates

XOR pattern revisited



- $F: \{0,1\}^n \rightarrow \{0,1\}$
- The state at (1) is

$$\begin{array}{c}
 \begin{bmatrix} 0 \\ \vdots \\ 0 \\ \vdots \\ 1 \end{bmatrix} \xrightarrow{\quad} \begin{bmatrix} +1 \\ -1 \\ -1 \\ +1 \\ \vdots \end{bmatrix} \xrightarrow{\quad} \frac{1}{\sqrt{2}} \begin{bmatrix} +1 \\ -1 \\ -1 \\ +1 \\ \vdots \end{bmatrix} \\
 \begin{array}{c} F \uparrow \\ \text{length } 2^n \\ \text{(truth table)} \end{array} \quad \begin{array}{c} f \\ \text{unit vector} \\ \text{of length } N. \end{array}
 \end{array}$$

- What are the pattern vectors $|\chi_s\rangle$?
- $\chi_s: \{0,1\}^n \rightarrow \{\pm 1\}$ defined by $\chi_s(x) = (-1)^{s \cdot x}$
- $H^{\otimes n}|s\rangle = |\chi_s\rangle$

- $H^{\otimes n}$ maps the standard basis to the χ -basis
- $H^{\otimes n}$ maps the χ -basis to the standard basis (why?)
- At (1), the just state can be written as

$$|f\rangle = \sum_s \langle \chi_s | f \rangle |\chi_s\rangle$$

- At (2)

$$H^{\otimes n}|f\rangle = \sum_s \langle \chi_s | f \rangle |s\rangle$$

- Interpretation of the “strength” $\langle \chi_s \mid f \rangle$
 - Note $\chi_s, f: \{0,1\}^n \rightarrow \{\pm 1\}$
 - $\langle \chi_s \mid f \rangle = \dots$
- Another application [Deutsch–Jozsa '92]
 - Given Q_F implementing $F: \{0,1\}^n \rightarrow \{0,1\}$
 - Promised: either $F(x) = 0$ for all x or F is “balanced”
 - Find a way to decide which?