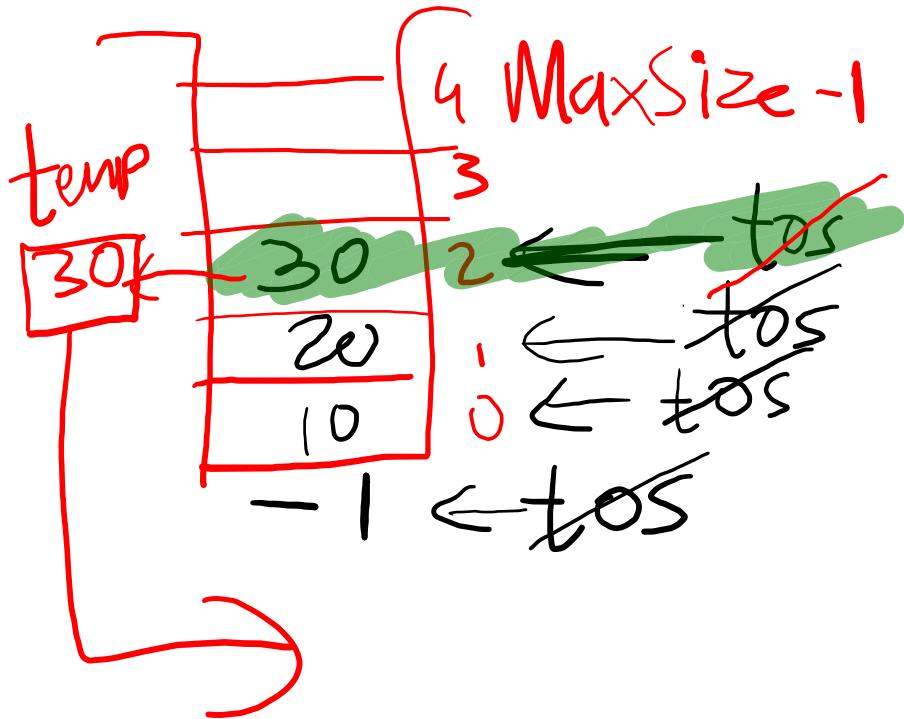


if  $tos == -1$   
 empty  
 if  $tos == MaxSize - 1$   
 Full

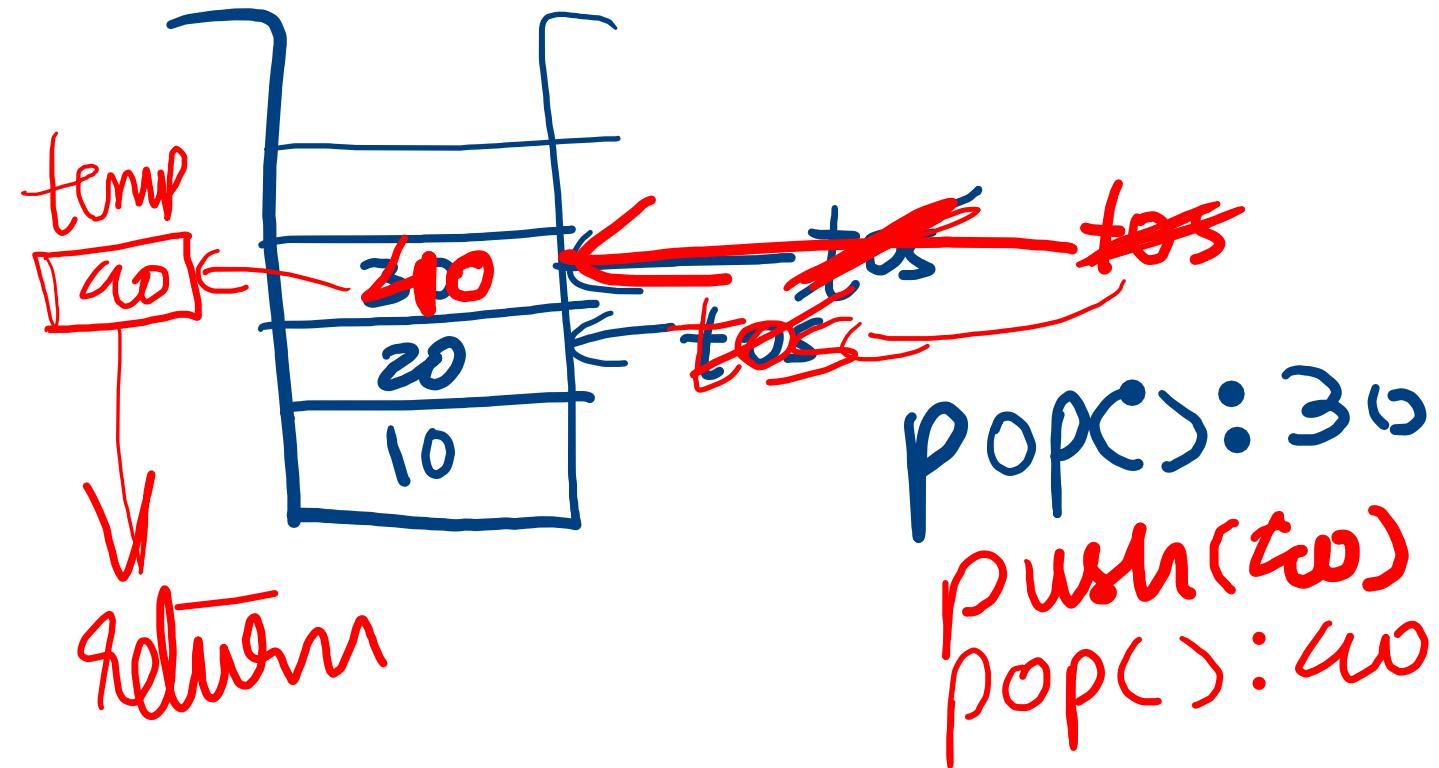
- Linear
- single ended
- mostly static

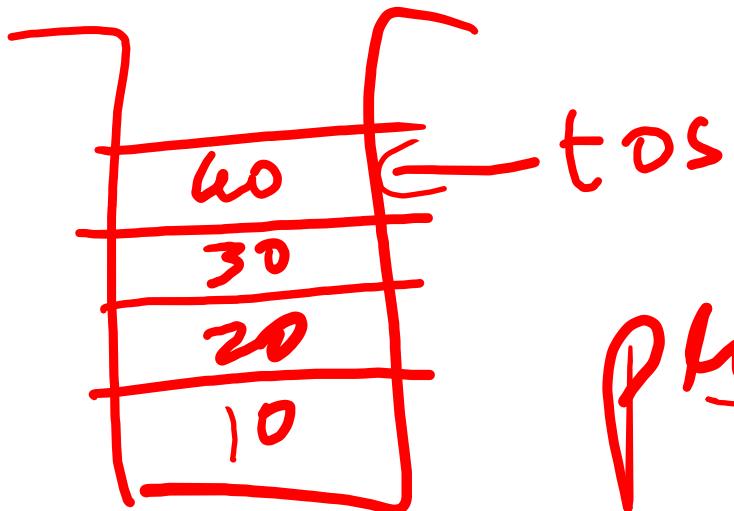
### ADT

↳ createStack(size)  
 ↳ push(e)  
 ↳ tos + 1  
 ↳ pop():e  
 ↳ peek():e  
 ↳ ~~tos~~

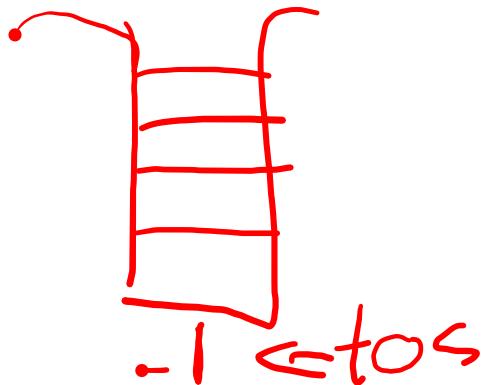


`push(10)`  
`push(20)`  
`push(30)`  
`peek(): 30`



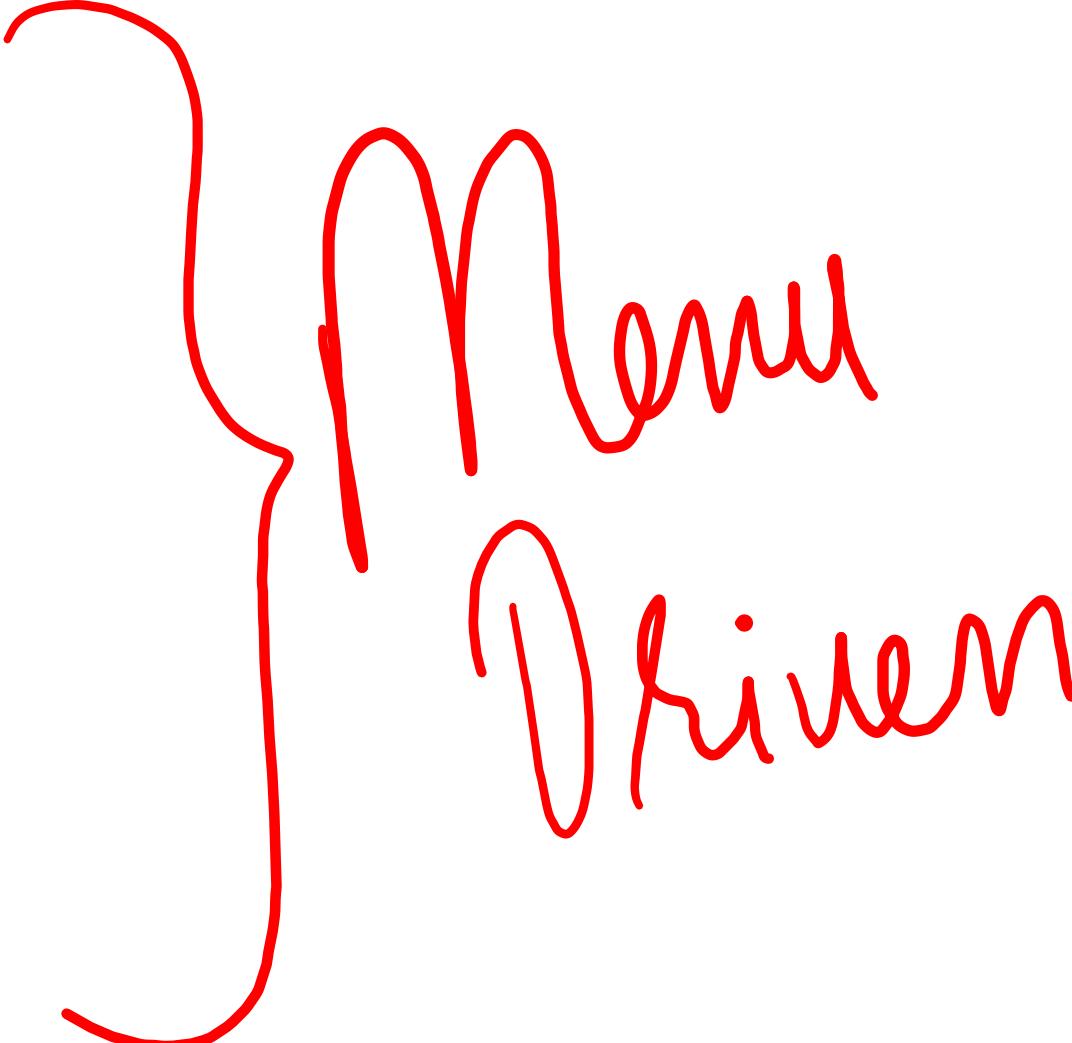


PrintStack()  
LIFO



40  
30  
20  
10  
↓ LIFO  
↓ order

```
do  
{  
    switch()  
    {  
    }  
} while
```



# App Stack

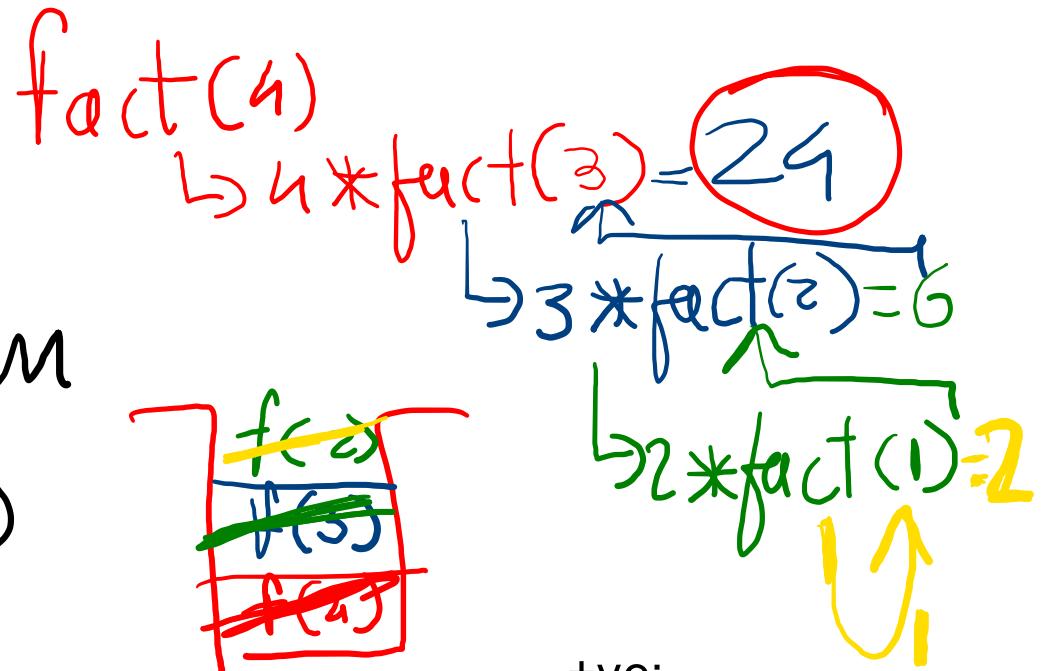
①  $\rightarrow$  Recursion

```
int fact(int no)
```

if(~~no == 1~~)  
return 1;

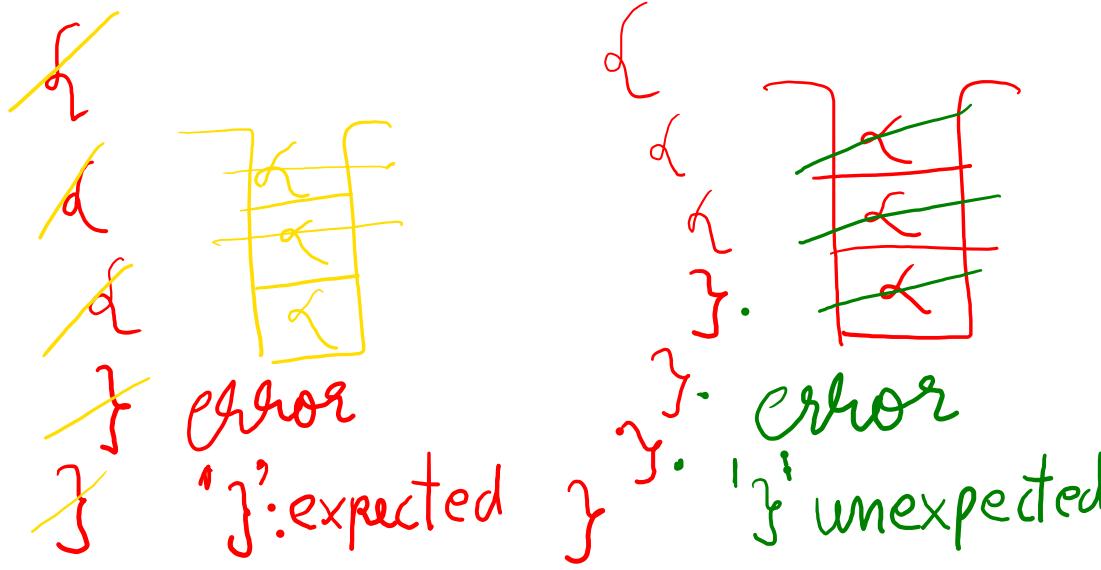
else return no \* fact(no - 1);

}



-ve:  
takes more memory, slow  
down, complex

## ② Wellness of Control



### ③ String reversal



$\Rightarrow \text{pop}()$   $\rightarrow$

The diagram shows the stack after popping the character 'o'. The stack is now empty, represented by a vertical rectangle with a single character 'o' at the top. Below the stack, the characters 'l', 'p', and 'a' are listed vertically, representing the current state of the stack.

④ Dec to binary

$$(12)_{10} \Rightarrow (0011)_2$$

The diagram illustrates the conversion of the decimal number 12 to binary using two methods. On the left, a circle encloses a division step where 6 is divided by 2, resulting in 3 with a remainder of 0. An arrow points from this result to the next division step. To the right of the circle, another division step shows 3 divided by 2, resulting in 1 with a remainder of 1. A curved arrow points from the 3 in the first step to the 1 in the second step. Below these, a third division step shows 1 divided by 2, resulting in 0 with a remainder of 1. This final remainder is placed in the least significant bit position of a vertical stack of four boxes. The boxes are arranged vertically, with the top one containing 1, the second containing 0, the third containing 0, and the bottom one containing 1, representing the binary number 0011.

# ⑤ Expression Conversion & Evaluation

Prefix:  $+ab(s/w)$   $\text{add}(a, b)$

infix:  $a + b$

Postfix:  $ab + (H/w)$

① Left to right if all equal

else 1<sup>st</sup> higher than lower

② (( )) then start from innermost

^ 3  
\* / % 2  
+ - 1

$$\begin{array}{r} a+b-c \\ \underline{+ab-c} \\ \hline A-c = -Ac \\ \underline{-tabc} \end{array}$$

$$\begin{array}{r} a+b-c \\ ab+ -c \\ \hline A-c = Ac- \\ \underline{ab+c-} \end{array}$$

$$\underline{(a+b)} * (c-d)$$

$$\begin{array}{r} \underline{+ab} * \underline{-cd} \\ \hline * +ab - cd \end{array}$$

$$(a+b\wedge c \cdot d - e) \quad | \quad (\underline{a \cdot b / c} - d \cdot e \% f)$$

$$\begin{array}{l} \cancel{(ab * /c - d * e \% f)} \\ \cancel{ab * c /} - \cancel{de * \% f} \end{array}$$

$$\cancel{ab * c /} - \cancel{de * f \%}$$

$$\cancel{ab * c /} \cancel{de * f \%} =$$

$$\begin{array}{l} \cancel{ab * /c - d * e \% f} \\ \cancel{(* abc} - d * e \% f \\ \cancel{(* abc} - \cancel{* de \% f} \\ \cancel{(* abc} - \cancel{\% * def} \\ = \cancel{(* abc \% * def} \end{array}$$

$$\frac{(a+b) * d - e}{\cancel{f * g / h}}$$

$$\frac{ab + \cancel{d} - e}{}$$

$$A * \cancel{d} = Ad *$$

$$\frac{ab + \cancel{d} * -e}{}$$

$$A - e = Ae -$$

$$\underline{\underline{ab + d * e - fg * h}}$$

$$\frac{fg * h}{}$$

$$\frac{A / h = Ah /}{}$$

$$\frac{fg * h /}{}$$

$$\underline{\underline{fg * h /}}$$

1.read infix from start to end

2.if read is

2.1 {[ push on stack

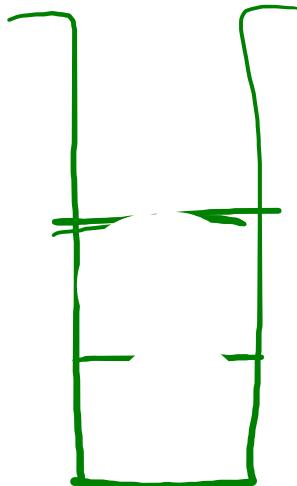
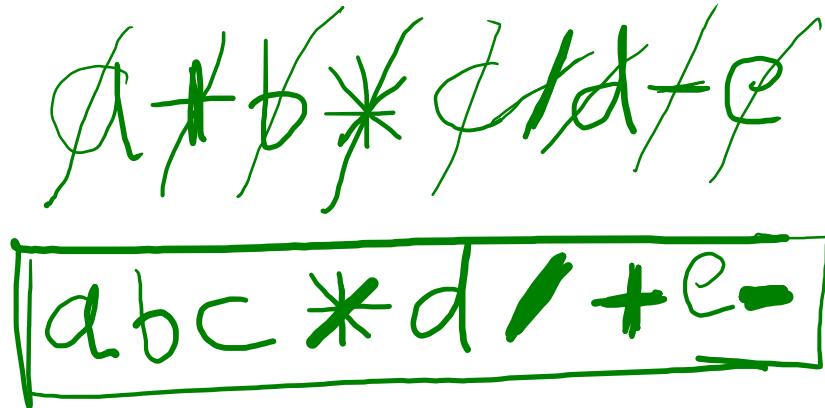
2.2 )]} then pop from stack till {[  
and copy popped to postfix

2.3 operand copy to postfix

2.4 if operator then compare it with  
tos if operator > then tos  
operator push else pop till either  
condition satisfies or stack  
becomes empty

3 continue till infix not over

4 copy left over of stack to postfix



1.read infix from end to start

2.if read is

2.1 )}] push on stack

2.2 {[ then pop from stack till )}] and copy  
poped to prefix

2.3 operand copy to prefix

2.4 if operator then compare it with tos if  
operator $\geq$  then tos operator push else pop till  
either condition satisfies or stack becomes  
empty

3 continue till infix not over

4 copy left over of stack to prefix

5 reverse prefix and use

a f b \* c

- +

+ a \* b c

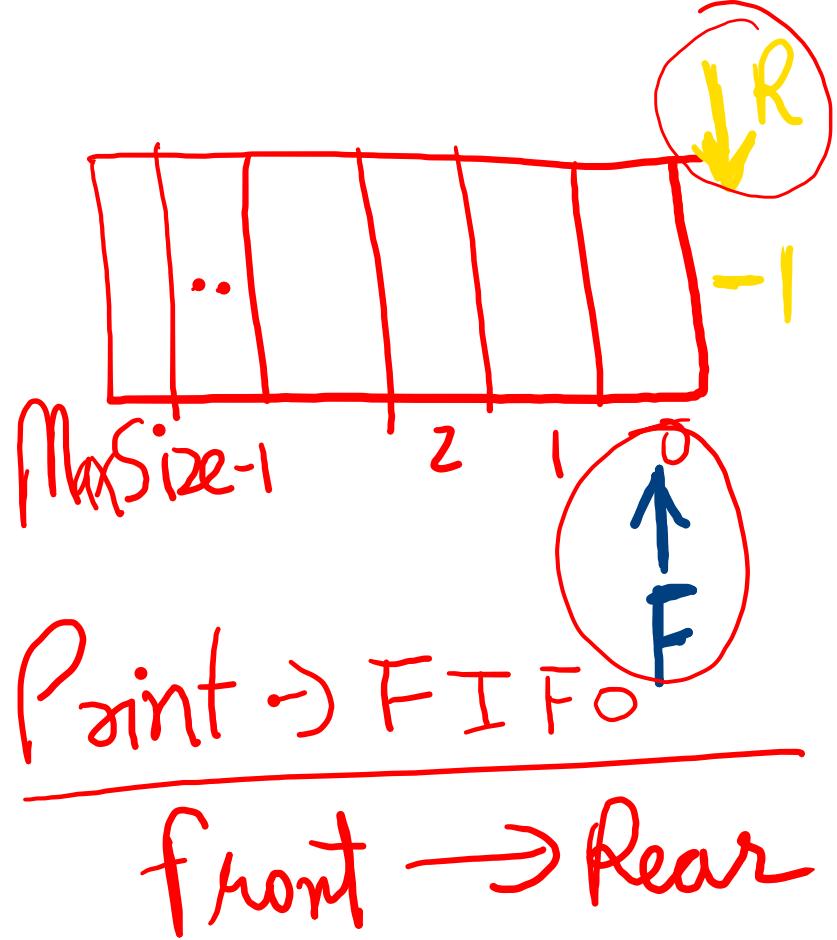
C b \* a +

Queue



rear

- FIFO
- Double ended
- Linear
- static



CreateQueue (size)

Enqueue(e): rear+1

Dequeue(): e front+1

if  $\text{rear} == \text{MaxSize}-1$   
then full

if  $\text{front} > \text{rear}$   
then empty

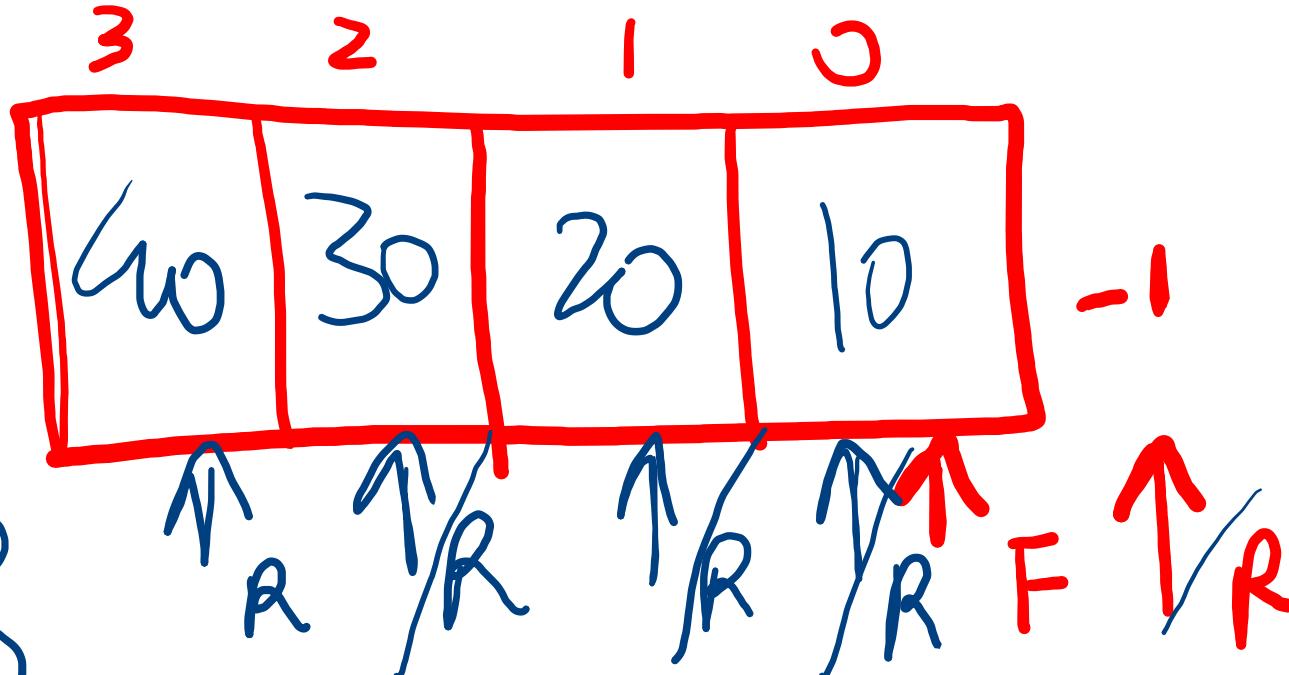


Enqueue(10)

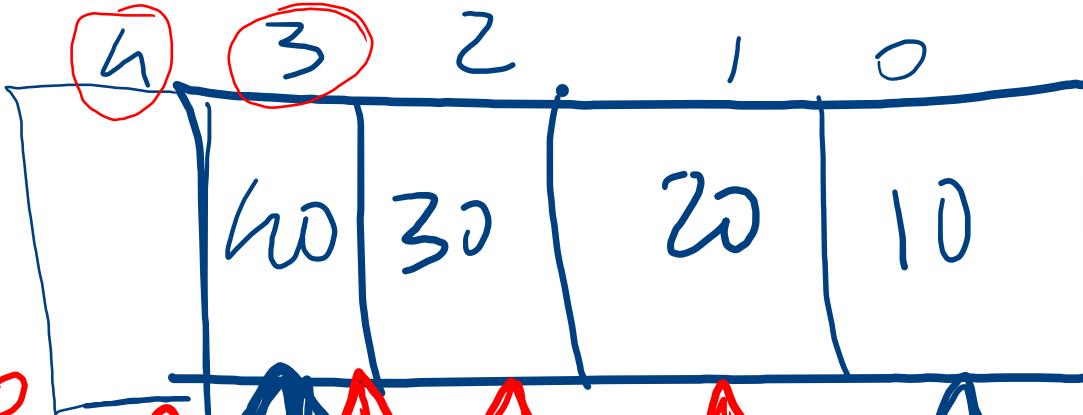
(20)

(30)

(40)



Dequeue():



() : 20

() : 30

() : 40

() : Empty

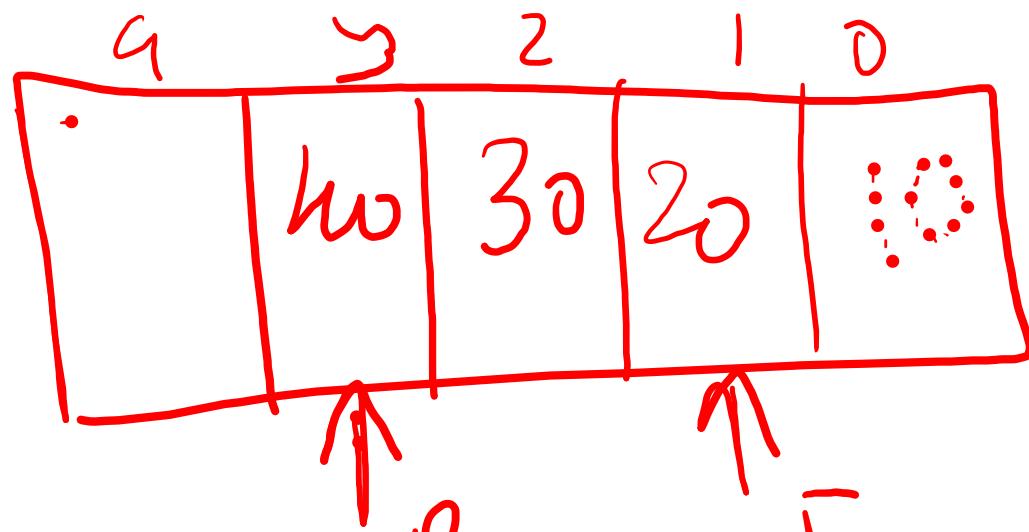
F R

F

/ F

/ F

/ F



Print from Front to rear (or)  
FIFO



...va ThreadJoin.java × ThreadPriority.java × ThreadSleep.java × ThreadSync.java × SortingDemo.java × SearchDemo.java × Treemain.java × GraphDemo.java × StackDemo.java × QueueLinearDemo.java

Source

History

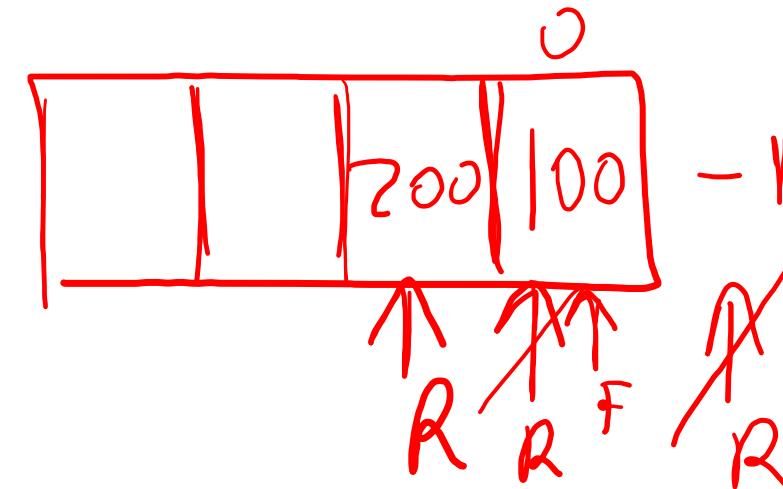
Projects

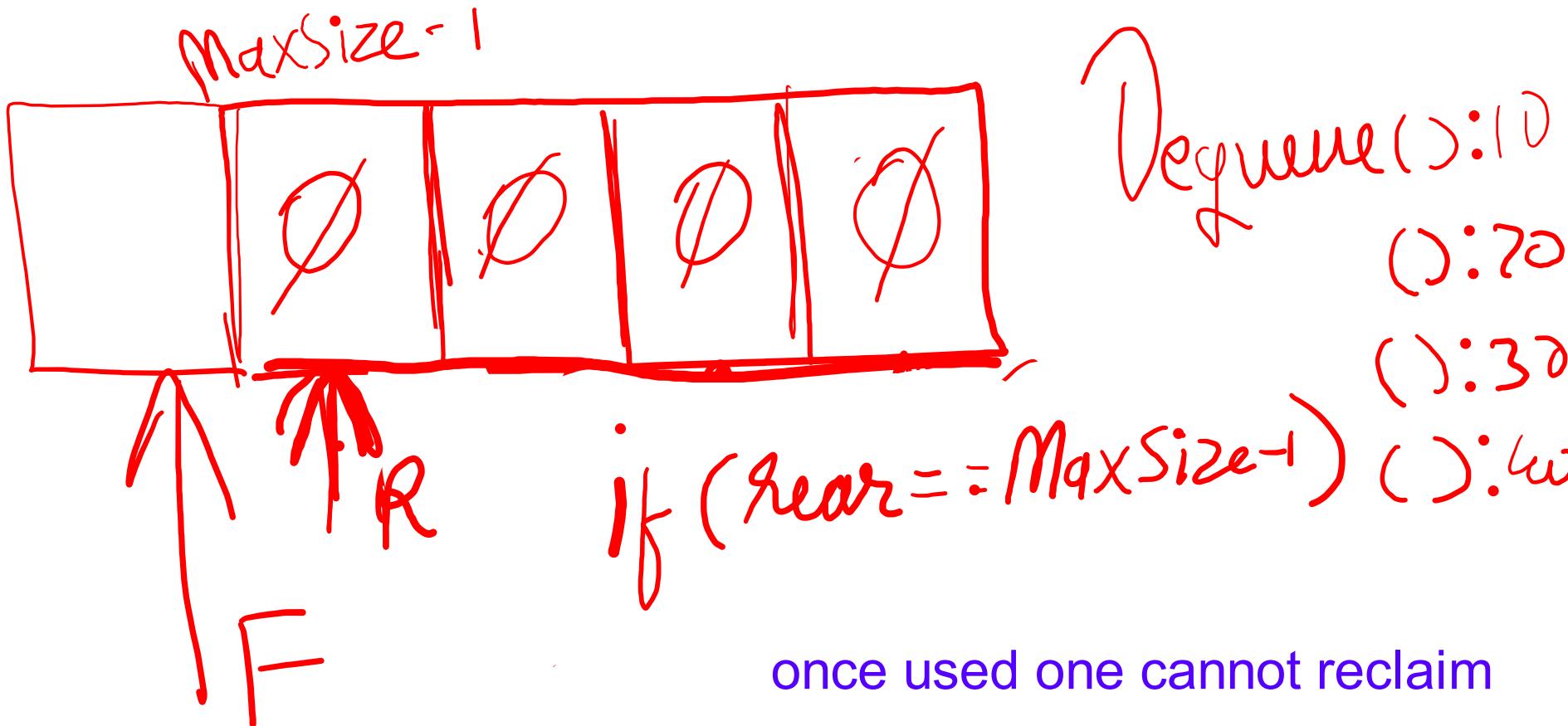
Files

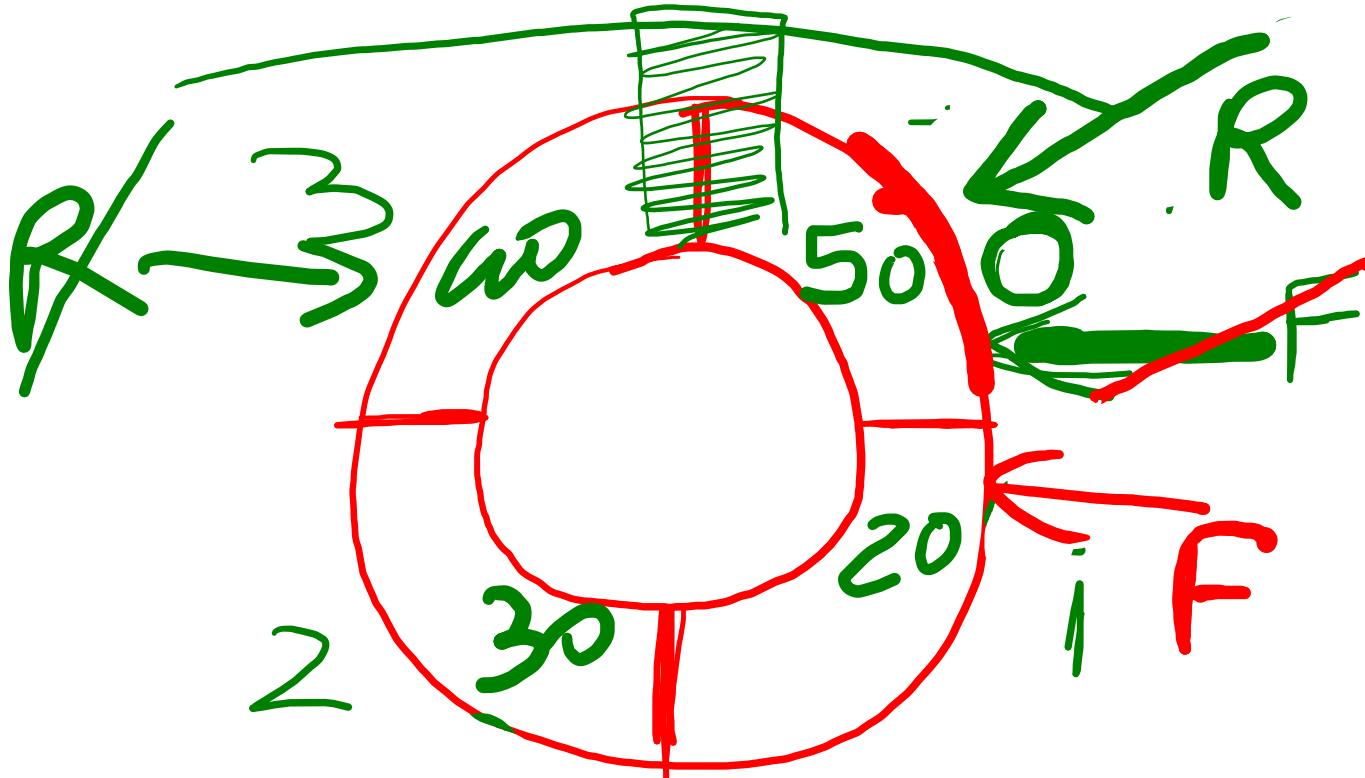
Services

Files

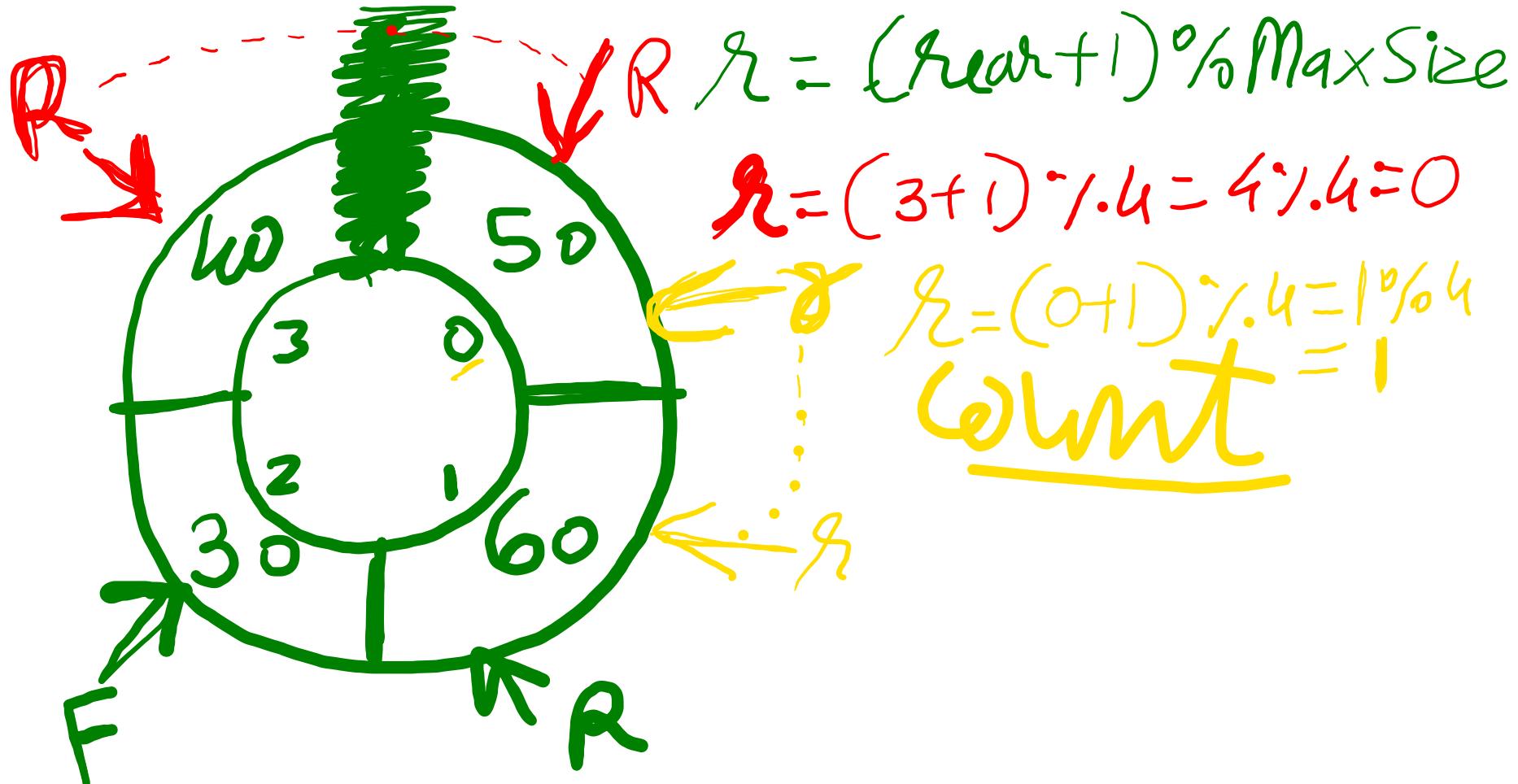
```
21 }  
22 void enqueue(int e)  
23 {  
24     rear++;  
25     q[rear]=e;  
26     //q[++rear]=e  
27 }  
28  
29 boolean isFull()  
30 {  
31     if(rear==MaxSize-1)  
32         return true;  
33     else  
34         return false;
```







$\ell(10)$   
 $\ell(20)$   
 $\ell(30)$   
 $\ell(40)$   
 $\ell(50)$   
 $\ell(10):10$



$$r = (\text{rear} + 1) \% \text{MaxSize}$$

$$r = (3 + 1) \% .4 = 4 \% .4 = 0$$

$$r = (0 + 1) \% .4 = 1 \% .4 = 1$$

Want



rear = -1  
Count = 0  
front = 0

if Count == 0  
Empty  
if Count == MaxSize  
Full

Enqueue

$r = (r+1) \% \text{MaxSize}$   
Count++

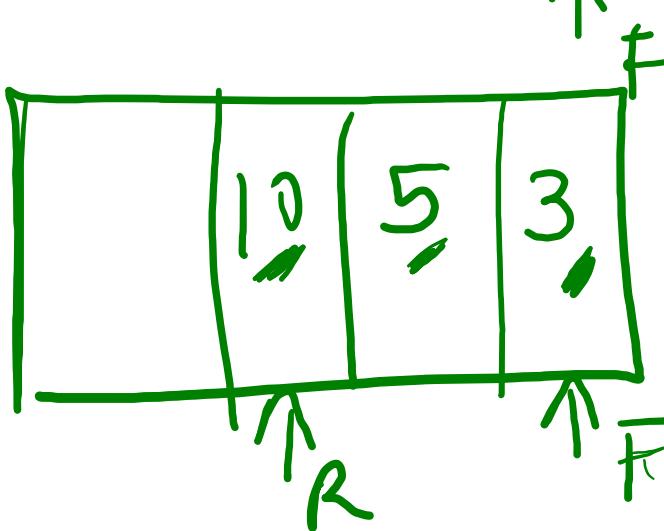
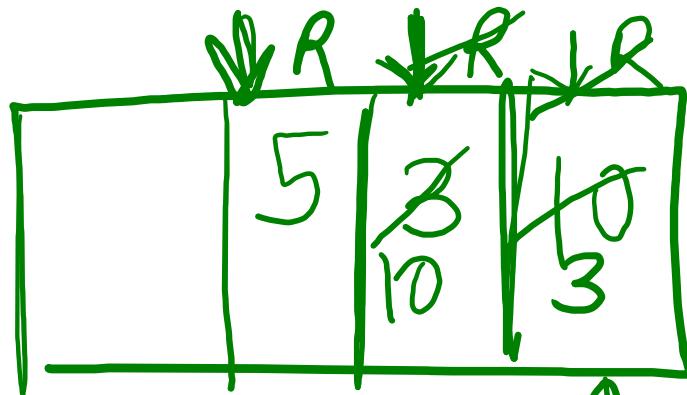
Dequeue

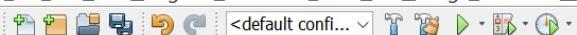
$f = (f+1) \% \text{MaxSize}$   
Count--



# Enqueue(c)

- 1. take element
  - 2. sort on ascending or descending as per need
- enqueue(10).  
enqueue(3)  
enqueue(5)



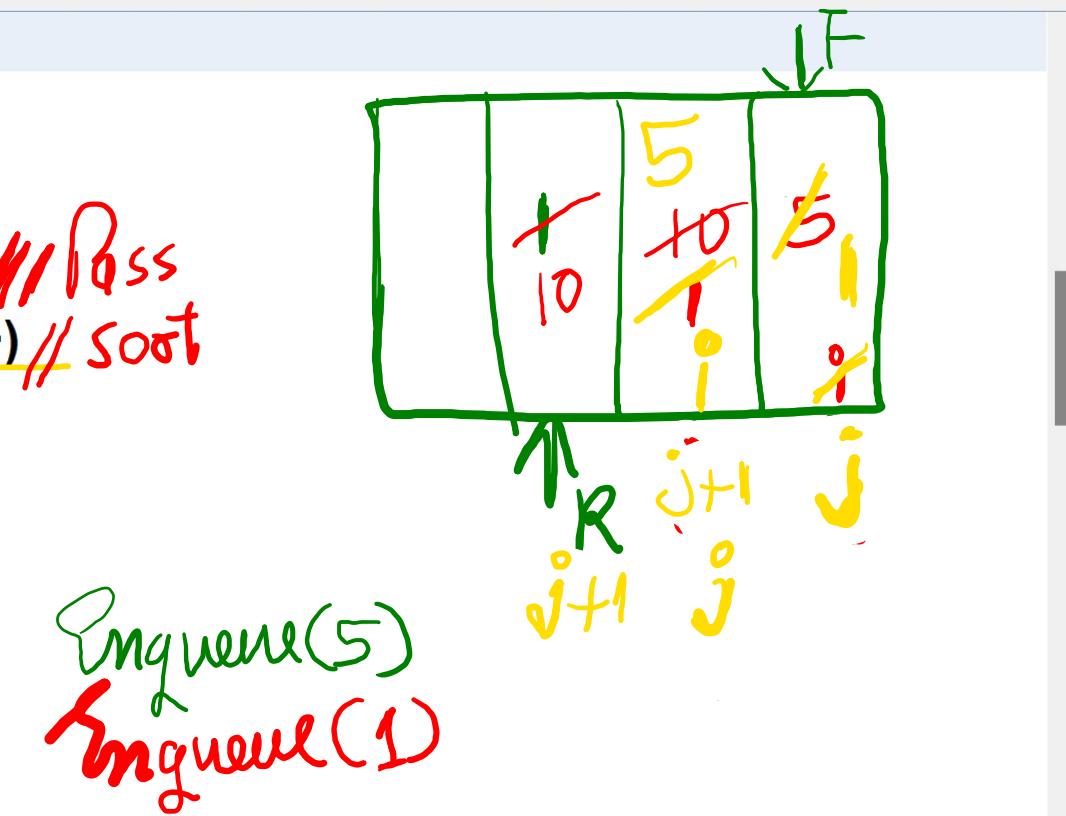


...va ThreadSync.java × SortingDemo.java × SearchDemo.java × Treemain.java × GraphDemo.java × StackDemo.java × QueueLinearDemo.java × CircularQueueDemo.java × PriorityQueueDemo.java

Source History

```

22 void enqueue(int e)
23 {
24     rear++;
25     q[rear]=e;
26     for(int i=front;i<rear;i++) // Pass
27     { for(int j=front,j<rear;j++)
28     { if(q[j]>q[j+1])
29     {
30         int t=q[j];
31         q[j]=q[j+1];
32         q[j+1]=t;
33     }
34 }
35 }
```



9821601163

Priority Queue: no LIFO no FIFO

↳ Dequeue():  $x \leftarrow \text{Max}$

↳ Dequeue(): 50      Dequeue(): 30

