

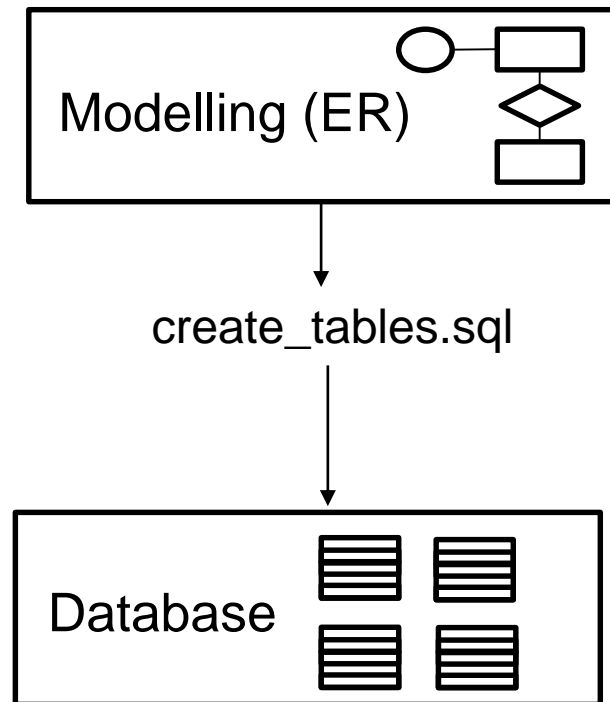


INFO20003 Database Systems

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Lecture 07
Relational Algebra

Week 4



SQL:

- Language for data manipulation
- Allow to create/delete tables, add/update/remove data, etc

Introduced next time

Relational algebra:

- The theory behind SQL
- Makes sure that SQL produces correct answers
- Inputs/outputs are relations

Today

How do we manipulate with relations?

1. **Selection** (σ): Selects a subset of *rows* from relation (horizontal filtering).
2. **Projection** (π): Retains only wanted *columns* from relation (vertical filtering).
3. **Cross-product** (\times): Allows us to combine two relations.
4. **Set-difference** ($-$): Tuples in one relation, but not in the other.
5. **Union** (\cup): Tuples in one relation and/or in the other.

Each operation returns a relation, operations can be composed

REEDUCATION

- Selection & Projection
- Union, Set Difference & Intersection
- Cross product & Joins
- Examples

Readings: Chapter 4, Ramakrishnan & Gehrke, Database Systems

Example Instances

**Reserves
(R1)**

<u>sid</u>	<u>bid</u>	<u>day</u>
22	101	10/10/96
58	103	11/12/96

Boats

<u>bid</u>	bname	color
101	Interlake	blue
102	Interlake	red
103	Clipper	green
104	Marine	red

**Sailors 1
(S1)**

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

**Sailors 2
(S2)**

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0



- Selection & Projection
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- Retains only attributes that are in the *projection list*
- *Schema* of result:
 - Only the fields in the projection list, with the same names that they had in the input relation
- Projection operator has to *eliminate duplicates*
 - How do they arise? Why remove them?
 - Note: real systems typically don't do duplicate elimination unless the user explicitly asks for it

1. Find ages of sailors :

$\pi_{age}(S2)$

2. Find names and rating of sailors :

$\pi_{sname, rating}(S2)$

sid	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

S2

sname	rating
yuppy	9
lubber	8
guppy	5
rusty	10

$\pi_{sname, rating}(S2)$

age
35.0
55.5

Removed duplicates

$\pi_{age}(S2)$

- Selects rows that satisfy *selection condition*
- Result is a relation. *Schema* of the result is same as that of the input relation.
- Do we need to do duplicate elimination?
- **Example:**

Find sailors whose rating is above 8

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

$\sigma_{rating > 8}(S2)$



<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
58	rusty	10	35.0

$\sigma_{rating > 8}(S2)$



- Conditions are standard arithmetic expressions

$>$, $<$, \geq , \leq , $=$, \neq

- Conditions are combined with AND/OR clauses

And: \wedge

Or: \vee

- **Example:**

Find sailors whose rating is above 8 and who are younger than 50

$$\sigma_{rating > 8 \wedge age < 50} (S2)$$

- Operations can be combined
- Select rows that satisfy *selection condition* & retain only *certain attributes (columns)*
- **Example:**

Find names and rating of sailors whose rating is above 8

sid	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0



sname	rating
yuppy	9
rusty	10

$\pi_{sname, rating}(\sigma_{rating > 8}(S2))$



- Selection & Projection
- Union, Set Difference & Intersection
- Cross product & Joins
- Examples

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- **Union:** Combines both relations together
- **Set-difference:** Retains rows of one relation that do not appear in the other relation
- These operations take two input relations, which must be *union-compatible*:
 - Same number of fields
 - Corresponding fields have the same type

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

S2

sid	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0
44	guppy	5	35.0
28	yuppy	9	35.0

$S1 \cup S2$

Duplicates are removed



<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

S2

sid	sname	rating	age
22	dustin	7	45.0

$S1 - S2$

Set Difference

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

S2

sid	sname	rating	age
22	dustin	7	45.0

$S1 - S2$

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
44	guppy	5	35.0

$S2 - S1$

Set-difference is not symmetrical

- In addition to the 5 basic operators, there are several additional “**Compound Operators**”
 - These add no computational power to the language, but are useful shorthands
 - Can be expressed solely with the basic operations
- **Intersection** retains rows that appear in *both* relations
- Intersection takes two input relations, which must be *union-compatible*
- Q: How to express it using basic operators?

$$R \cap S = R - (R - S)$$



Example:

Find sailors who appear in both relations S1 and S2

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	sname	rating	age
28	yuppy	9	35.0
31	lubber	8	55.5
44	guppy	5	35.0
58	rusty	10	35.0

S2

sid	sname	rating	age
31	lubber	8	55.5
58	rusty	10	35.0

$S1 \cap S2$



- Selection & Projection
- Union, Set Difference & Intersection
- Cross product & Joins
- Examples

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- **Cross product** combines two relations:
 - Each row of one input is merged with each row from another input
 - Output is a new relation with all attributes of *both* inputs
 - \times is used to denote cross-product
- Example: $S1 \times R1$
 - Each row of $S1$ paired with each row of $R1$
- Question: How many rows are in the result?
 - A: $\text{card}(S1) * \text{card}(R1)$



Cross Product Example

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	<u>bid</u>	<u>day</u>
22	101	10/10/96
58	103	11/12/96

R1

S1 X R1 =

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	22	101	10/10/96
22	dustin	7	45.0	58	103	11/12/96
31	lubber	8	55.5	22	101	10/10/96
31	lubber	8	55.5	58	103	11/12/96
58	rusty	10	35.0	22	101	10/10/96
58	rusty	10	35.0	58	103	11/12/96

- *Result schema* has one field per field of S1 and R1, with field names “inherited” if possible.
 - May have a *naming conflict*, i.e. both S1 and R1 have a field with the same name (e.g. *sid*).
 - In this case, can use the *renaming operator*.

$$\rho(C(1 \rightarrow \text{sid1}, 5 \rightarrow \text{sid2}), S1 \times R1)$$

Result relation name

C	sid1	sname	rating	age	sid2	bid	day
	22	dustin	7	45.0	22	101	10/10/96
	22	dustin	7	45.0	58	103	11/12/96
	31	lubber	8	55.5	22	101	10/10/96
	31	lubber	8	55.5	58	103	11/12/96
	58	rusty	10	35.0	22	101	10/10/96
	58	rusty	10	35.0	58	103	11/12/96

- Joins are compound operators involving cross product, selection, and (sometimes) projection.
- Most common type of join is a **natural join** (often just called **join**). $R \bowtie S$ conceptually is a cross product that matches rows where attributes that appear in both relations have equal values (and we omit duplicate attributes).
- To obtain cross product a DBMS must:
 1. Compute $R \times S$
 2. Select rows where attributes that appear in both relations have equal values
 3. Project all unique attributes and one copy of each of the common ones.

Example:

Find all sailors (from relation S1) who have reserved a boat

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

S1

<u>sid</u>	<u>bid</u>	<u>day</u>
22	101	10/10/96
58	103	11/12/96

R1

S1 ⋈ R1 =

<u>sid</u>	sname	rating	age	bid	day
22	dustin	7	45.0	101	10/10/96
58	rusty	10	35.0	103	11/12/96

Natural Join Example

1

S1 X R1 =

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	22	101	10/10/96
22	dustin	7	45.0	58	103	11/12/96
31	lubber	8	55.5	22	101	10/10/96
31	lubber	8	55.5	58	103	11/12/96
58	rusty	10	35.0	22	101	10/10/96
58	rusty	10	35.0	58	103	11/12/96

Natural Join Example

①
S1 X R1 =

②
σ

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	22	101	10/10/96
22	dustin	7	45.0	58	103	11/12/96
31	lubber	8	55.5	22	101	10/10/96
31	lubber	8	55.5	58	103	11/12/96
58	rusty	10	35.0	22	101	10/10/96
58	rusty	10	35.0	58	103	11/12/96

Natural Join Example

① $S1 \times R1 =$

② σ

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	22	101	10/10/96
22	dustin	7	45.0	58	103	11/12/96
31	lubber	8	55.5	22	101	10/10/96
31	lubber	8	55.5	58	103	11/12/96
58	rusty	10	35.0	22	101	10/10/96
58	rusty	10	35.0	58	103	11/12/96

π ③

$S1 \bowtie R1 =$

sid	sname	rating	age	bid	day
22	dustin	7	45.0	101	10/10/96
58	rusty	10	35.0	103	11/12/96

- **Condition Join (or theta-join)** is a cross product with a condition.

$$R \bowtie_c S = \sigma_c (R \times S)$$

(sid)	sname	rating	age	(sid)	bid	day
22	dustin	7	45.0	58	103	11/12/96
31	lubber	8	55.5	58	103	11/12/96

$$S1 \bowtie_{S1.sid < R1.sid} R1$$

–Result schema is the same as that of cross-product

- **Equi-Join** is a special case of condition join, where condition c contains only *equalities* (e.g. $S1.sid = R1.sid$)
 - Is this then a natural join? What is different?



- Selection & Projection
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Readings: Chapter 4, Ramakrishnan & Gehrke, Database Systems

Boats

<u>bid</u>	bname	color
101	Interlake	Blue
102	Interlake	Red
103	Clipper	Green
104	Marine	Red

Sailors

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
22	101	10/10/96
58	103	11/12/96

Find names of sailors who have reserved boat #103

Boats

<u>bid</u>	bname	color
101	Interlake	Blue
102	Interlake	Red
103	Clipper	Green
104	Marine	Red

Sailors

<u>sid</u>	sname	rating	age
22	dustin	7	45.0
31	lubber	8	55.5
58	rusty	10	35.0

Reserves

<u>sid</u>	<u>bid</u>	<u>day</u>
22	101	10/10/96
58	103	11/12/96

Solution 1:

$$\pi_{sname}((\sigma_{bid=103} \text{Reserves}) \bowtie \text{Sailors})$$

Solution 2:

$$\pi_{sname}(\sigma_{bid=103}(\text{Reserves} \bowtie \text{Sailors}))$$



Find all pairs of sailors in which the older sailor has a lower rating



- Relational Algebra Operations: Selection, Projection, Union, Set, Difference, Intersection, **JOINS**...
- Draw different queries with Relational Algebra operations



- Introducing SQL