

Escape the Matrix: Graphical Reasoning and Minimal Axioms for Quantum Circuits

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Outline

Equational Theories via Music

Completeness of Equational Theories

Minimality of Axioms

Minimal Axioms for Quantum Fragments

Example Axioms: Qubit Clifford

Conclusion

Motivation: The Matrix Problem

What happens if you try to verify a *20-gate* quantum circuit by hand?

- You multiply 20 matrices, each $2^n \times 2^n$, by hand
- Errors multiply and it becomes intractable for even 5 qubits

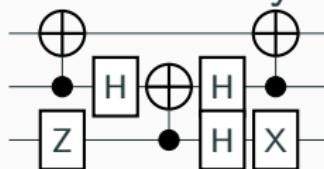
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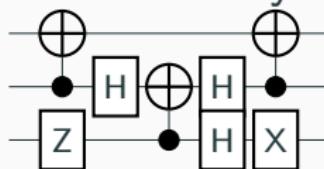
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- Rewrite rules on wires & nodes
- Local transformations, no full-blown matrix multiplication
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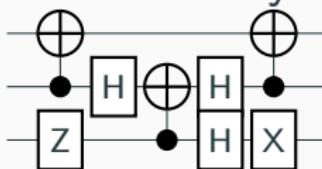
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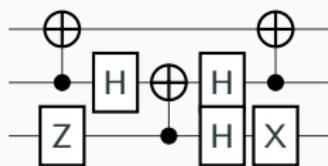


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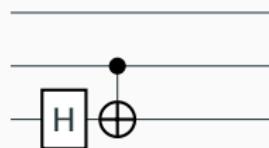
Circuit Optimization with Rewrites

Goal: Simplify a complex quantum circuit while preserving its function.

Before:



After:



Such rewrites come from a set of *equational rules* (axioms).

Equational Theory: Musical Analogy

- **Generators:** Musical note durations (whole, half, quarter, ...)
- **Equations:** Duration identities, e.g. $\text{half} + \text{half} = \text{whole}$
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Note Duration Example

Original

half note + half note

Rewritten

whole note

o	=	d d
j	=	• •
•	=	♪ ♪
♪	=	• ♩
♩	=	• ♩
♪ ♩	=	• ♩ ♩

A small, minimal rule delivers powerful rewriting capabilities.

What is Completeness?

A set of circuit equations is **complete** if:

- Any two circuits implementing the same unitary can be transformed into each other by these rules.
- Guarantees no missing identities: every true equality is derivable.

Why it matters: Ensures we can automate verification of equivalence.

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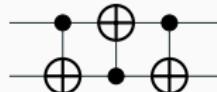
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Completeness Example: SWAP Circuits

Two implementations of a SWAP gate:



A complete theory provides a sequence of rewrites:



proving semantic equivalence without matrices.

What is Minimality?

- **Minimality:** No axiom is redundant.
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Proving Independence via Interpretations

- To show axiom X is necessary:
 - Construct an *alternate interpretation* where all other axioms hold,
 - But axiom X fails (gives different semantics).
- This counter-model demonstrates X cannot be derived from the rest.

Example: Interpretation which counts the parity of some generators.

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Qubit Clifford & Real Clifford

- **Qubit Clifford** : Core operations. Widely used for error correction and efficient classical simulation. It uses phases $\pm 1, \pm i$ (complex roots).
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Equational Theory for Qubit Clifford Circuits

$$\textcircled{w}^{\otimes 8} = \boxed{} \quad (\omega^8) \quad \boxed{H} \boxed{H} = \boxed{} \quad (H^2)$$

$$\boxed{S} \boxed{S} \boxed{S} \boxed{S} = \boxed{} \quad (S^4)$$

$$\boxed{H} \boxed{S} \boxed{H} = \textcircled{w} \quad \boxed{S^\dagger} \boxed{H} \boxed{S^\dagger} \quad (E)$$

$$\begin{array}{c} \bullet \\ \oplus \end{array} \boxed{S} \begin{array}{c} \bullet \\ \oplus \end{array} = \boxed{S} \quad (Cs) \quad \begin{array}{c} \bullet \\ \oplus \end{array} \begin{array}{c} \oplus \\ \bullet \end{array} = \begin{array}{c} \bullet \\ \oplus \end{array} \quad (B)$$

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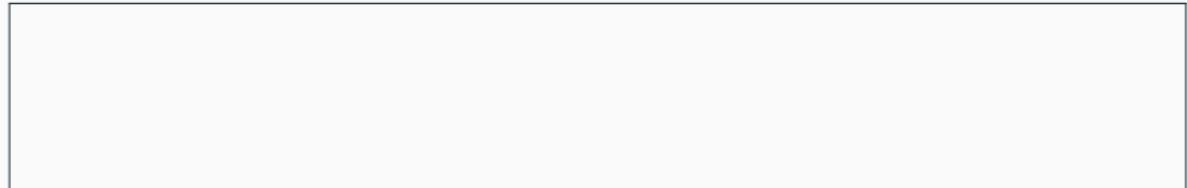
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Interpretation: Count the number of CNOT and SWAP gates in a circuit, mod 2.

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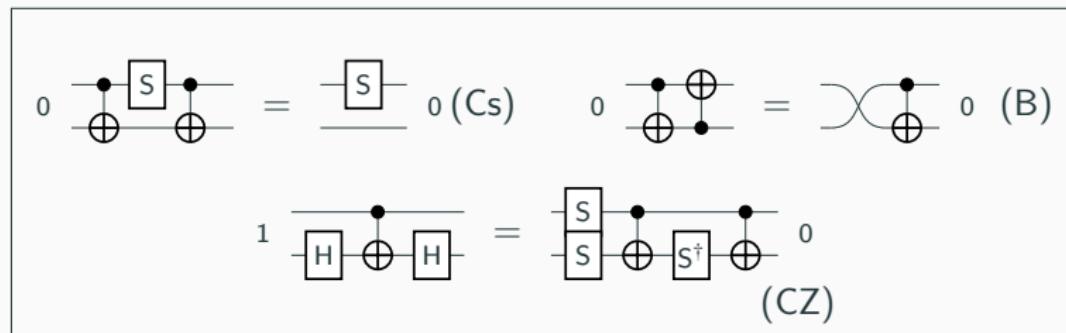
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Comparison of Axiom Set Sizes

Fragment	Previous Theory	Our Theory
Qubit Clifford	15	8 (minimal for n -qubits)
Real Clifford	16	10 (minimal for n -qubits)
Qutrit Clifford	17	11 (minimal for 2-qutrits, conjectured for n)
Clifford+T	18	11 (minimal for 1-qubit, simplified for 2, no result for n)

Summary

- Graphical equational theories let us *reason* and *optimize* circuits visually.
- Our rule sets are *complete* (all true equivalences provable) and are *minimal* in the sense that we can't derive some equations from the others.
- Enables scalable *automation* in quantum compilation and verification.

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