

7SENG010W Data Structures and Algorithms

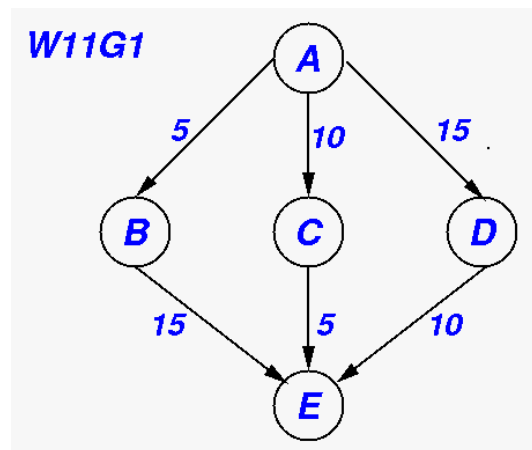
Week 10 Tutorial Exercises

Topics: Graph Shortest Paths using Dijkstra's Shortest Path Algorithm

These exercises cover the *Shortest Path* problem for weighted edge digraphs using Edsger W. Dijkstra's Shortest Path Algorithm. You can use either the Adjacency Matrix or Adjacency Lists representation of a graph developed in the Week 9 Tutorial. However, for this exercise it is recommended that you use Adjacency Lists to represent the graph.

Exercise 2.

Using the following graph **W11G1**



$V = \{A, B, C, D, E\}$

$E = \{ (A, B, 5), (A, C, 10), (A, D, 15), (B, E, 15), (C, E, 5), (D, E, 10) \}$

Using **A** as the *source* vertex **s** perform Dijkstra's Shortest Path algorithm on the graph.

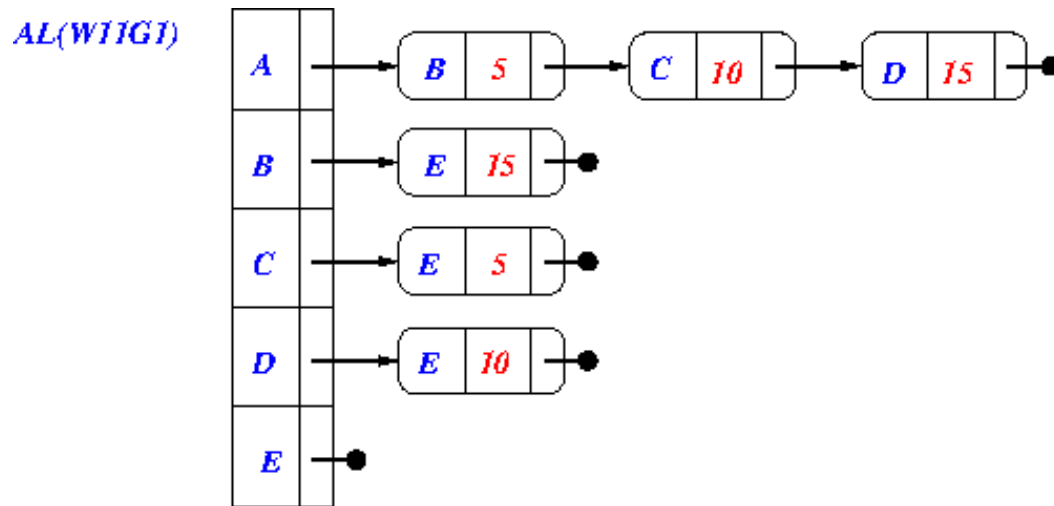
You should produce (pen and paper or an online tool) a "log" of the state of the search as it is performed, i.e. use a table similar to the one used in the lecture notes that records the: `edgeTo[v]`, `pathTo[v]`, `PriQueue` and `nearestVertex`.

After the algorithm has been completed reconstruct the shortest paths from **A** to each vertex from the `edgeTo[v]` edges for each vertex.

Compare your answers with other students in the tutorial. If there are any differences try to work why.

Exercise 2 Solution

(1) First construct the adjacency list representation for graph *W11G1*



(2) Now run *Dijkstra's Shortest Path Algorithm* with *A* as the source vertex *S*.

DSP Step 0: Call DijkstraShortestPath(A)

Initialise the data structures by:

- setting all $\text{edgeTo}[v]$ to none,
- setting all $\text{distTo}[v]$ to ∞ ,
- create an empty priority queue,

Start search from the start vertex A, by:

- setting its distance from itself to 0 &
- inserting it & its distance into the queue, i.e. (A, 0).

	A	B	C	D	E
edgeTo[v]	-	-	-	-	-
distTo[v]	0	∞	∞	∞	∞
PriQueue	< (A, 0) >				
nearestVertex	A				

DSP Step 1: Search from Nearest Vertex A

1) PriorityQueue equals $\langle \mathbf{A}, 0 \rangle$, so enter WHILE-loop.

2) Get: $\text{nearestVertex} \leftarrow \text{PriorityQueue.Dequeue}()$
 $\text{nearestVertex} = A;$

then $\text{PriorityQueue} = \langle \rangle$.

3) Find A's adjacent vertices: B, C & D, using edges: (A, B, 5), (A, C, 10), (A, D, 15).

4) **Found shorter path to B**, via (A, B, 5):

$$\text{distTo}(B) > \text{distTo}(A) + \text{weight}(A, B)$$
$$\infty > 0 + 5$$

Update:

$$\text{edgeTo}[B] \leftarrow (A, B, 5)$$
$$\text{distTo}[B] \leftarrow \text{distTo}(A) + \text{weight}(A, B) = 0 + 5 = 5$$
$$\text{PriorityQueue.insert}(\mathbf{B}, 5)$$

5) **Found shorter path to C**, via (A, C, 10):

$$\text{distTo}(C) > \text{distTo}(A) + \text{weight}(A, C)$$
$$\infty > 0 + 10$$

Update:

$$\text{edgeTo}[C] \leftarrow (A, C, 10)$$
$$\text{distTo}[C] \leftarrow \text{distTo}(A) + \text{weight}(A, C) = 0 + 10 = 10$$
$$\text{PriorityQueue.insert}(\mathbf{C}, 10)$$

6) **Found shorter path to D**, via (A, D, 15):

$$\text{distTo}(D) > \text{distTo}(A) + \text{weight}(A, D)$$
$$\infty > 0 + 15$$

Update:

$$\text{edgeTo}[D] \leftarrow (A, D, 15)$$
$$\text{distTo}[D] \leftarrow \text{distTo}(A) + \text{weight}(A, D) = 0 + 15 = 15$$
$$\text{PriorityQueue.insert}(\mathbf{D}, 15)$$

So found shorter paths for B, C & D.

	A	B	C	D	E
edgeTo[v]	-	(A, B, 5)	(A, C, 10)	(A, D, 15)	-
distTo[v]	0	5	10	15	∞
PriQueue	< (B, 5), (C, 10), (D, 15) >				
nearestVertex	A				

DSP Step 2: Search from Nearest Vertex B

1) PriQueue equals < (B, 5), (C, 10), (D, 15) >, so enter WHILE-loop.

2) Get: nearestVertex <-- PriQueue.Dequeue()
nearestVertex = B;

then PriQueue = < (C, 10), (D, 15) >

3) Find B's adjacent vertices: E, using edges: (B, E, 15).

4) *Found shorter path to E*, via (B, E, 15):

distTo(E) > distTo(B) + weight(B, E)
 $\infty > 5 + 15 = 20$

Update:

edgeTo[B] <-- (B, E, 15)
distTo[B] <-- distTo(B) + weight(B, E) = 5 + 15 = 20

PriQueue.insert((E, 20))

So found a shorter path for E of 20 via vertex B.

	A	B	C	D	E
edgeTo[v]	-	(A, B, 5)	(A, C, 10)	(A, D, 15)	(B, E, 15)
distTo[v]	0	5	10	15	20
PriQueue	< (C, 10), (D, 15), (E, 20) >				
nearestVertex	B				

DSP Step 3: Search from Nearest Vertex C

1) PriorityQueue equals $\langle (C, 10), (D, 15), (E, 20) \rangle$, so enter WHILE-loop.

2) Get: nearestVertex \leftarrow PriorityQueue.Dequeue()
nearestVertex = C;

then PriorityQueue = $\langle (D, 15), (E, 20) \rangle$

3) Find C's adjacent vertices: E, using edges: (C, E, 5).

4) **Found shorter path to E**, via (C, E, 5):

distTo(E) > distTo(C) + weight(C, E)
 $20 > 10 + 5 = 15$

Update:

edgeTo[E] \leftarrow (C, E, 5)
distTo[E] \leftarrow distTo(C) + weight(C, E) = $10 + 5 = 15$

PriorityQueue.insert(**(E, 15)**)

So found a new shorter path for E of 15 via C, which is shorter than the previous shortest path for E of 20 via B.

	A	B	C	D	E
edgeTo[v]	-	(A, B, 5)	(A, C, 10)	(A, D, 15)	(C, E, 5)
distTo[v]	0	5	10	15	15
PriorityQueue	$\langle (D, 15), (E, 15) \rangle$				
nearestVertex	C				

DSP Step 4: Search from Nearest Vertex D

1) PriorityQueue equals $\langle (D, 15), (E, 15) \rangle$, so enter WHILE-loop.

2) Get: nearestVertex \leftarrow PriorityQueue.Dequeue()
nearestVertex = D;

then PriorityQueue = $\langle (E, 15) \rangle$

3) Find D's adjacent vertices: E, using edge: (D, E, 10).

4) **Have not found a shorter path to E**, via (D, E, 10):

distTo(E) > distTo(D) + weight(D, E)
 $15 > 15 + 10 = 25$

So E's shortest path does not need to be updated or inserted into the queue.

So only found a new path for E of 25 via D, which is longer than the previous shortest path for E of 15 via C.

	A	B	C	D	E
edgeTo[v]	-	(A, B, 5)	(A, C, 10)	(A, D, 15)	(C, E, 5)
distTo[v]	0	5	10	15	15
PriQueue	$\langle (E, 15) \rangle$				
nearestVertex	D				

DSP Step 5: Search from Nearest Vertex E

1) PriorityQueue equals **< (E, 15) >**, so enter WHILE-loop.

2) Get: nearestVertex \leftarrow PriorityQueue.Dequeue()
nearestVertex = E;

then PriorityQueue = **< >**

3) Find E's adjacent vertices: **there are none.**

So nothing to update & the queue is empty so the while-loop terminates & the algorithm terminates. The final state is:

	A	B	C	D	E
edgeTo[v]	-	(A, B, 5)	(A, C, 10)	(A, D, 15)	(C, E, 5)
distTo[v]	0	5	10	15	15
PriQueue	< >				
nearestVertex	E				

DSP Step 6: Final Shortest Paths from A

The final state of the shortest path edges, distances & paths from A to all the other vertices is:

Vertex	edgeTo[v]	distTo[v]	Path Edges	Path
A	-	0	< >	< A >
B	(A, B, 5)	5	< (A, B, 5) >	< A, B >
C	(A, C, 10)	10	< (A, C, 10) >	< A, C >
D	(A, D, 15)	15	< (A, D, 15) >	< A, D >
E	(C, E, 5)	15	< (A, C, 10), (C, E, 5) >	< A, C, E >