



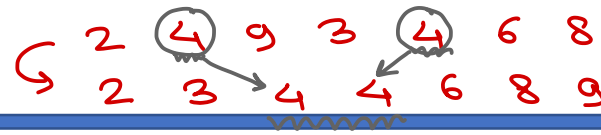
Data Structure & Algorithms

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Stack and Queue



- Stack & Queue are utility data structures.
- Can be implemented using array or linked lists.
- Usually time complexity of stack & queue operations is $O(1)$.
- Stack is Last-In-First-Out structure.
- Stack operations
 - ✓ push()
 - ✓ pop()
 - ✓ peek()
 - ✓ isEmpty()
 - ✓ isFull()*

- Simple queue is First-In-First-Out structure.

- Queue operations

- push() *enqueue()*
- pop() *dequeue()*
- peek()
- isEmpty() .
- isFull()*

- Queue types

- Linear queue
- Circular queue
- Deque
- Priority queue

ADT

- ① Array
 - ✓ get ele at pos
 - ✓ set ele at pos
 - ✓ sort, search
 - ✓ slice, ...

① Queue



Stack / Queue using Linked List

- Stack can be implemented using linked list.
 - add first
 - delete first
 - is empty
- Queue can be implemented using linked list.
 - add last
 - delete first
 - is empty



Linear Queue \rightarrow FIFO

deque

pop



front

enqueue

push



rear

init:

$f = -1$

$r = -1$

arr \rightarrow

-1

0

1

2

3

4

5

			44	55	66
--	--	--	----	----	----

\uparrow
 r

11 22 33

push:

$r++;$

$arr[r] = val;$

pop:

$f++;$

is full:

$r == size - 1$

peek:

return $arr[f + 1];$

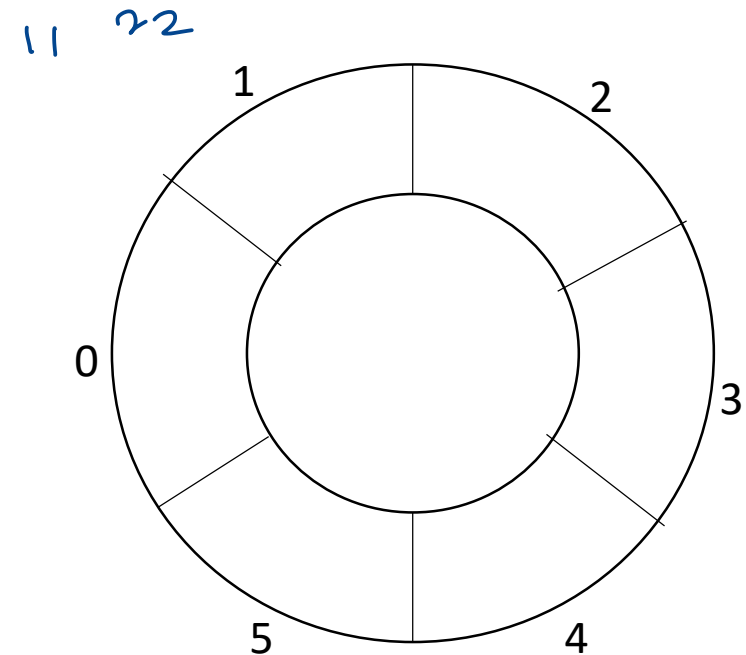
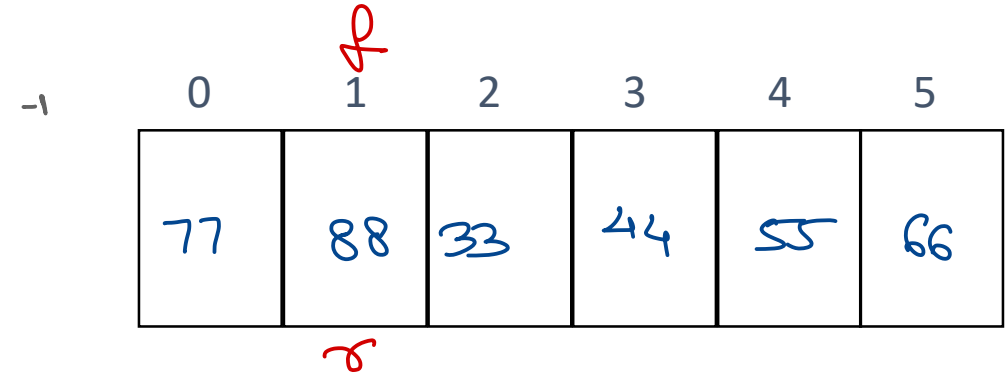
is empty:

$f == r$



Circular Queue

- In linear queue (using array) when rear reaches last index, further elements cannot be added, even if space is available due to deletion of elements from front. Thus space utilization is poor.
- Circular queue allows adding elements at the start of array if *rear* reaches last index and space is free at the start of the array.
- Thus *rear* and *front* can be incremented in circular fashion i.e. 0, 1, 2, 3, ..., n-1. So they are said to be circular queue.
- However queue full and empty conditions become tricky.



Circular Queue

q full $\rightarrow f == -1$ && $r == \text{size} - 1$

push

```
r++;  
if (r == size)  
    r = 0;  
arr[r] = val;
```

f

-1



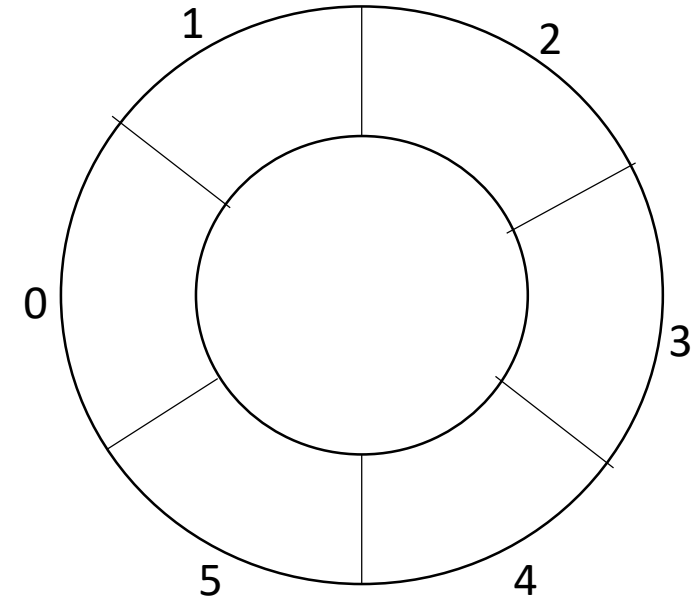
r

push

```
r = (r + 1) % size;  
arr[r] = val;
```

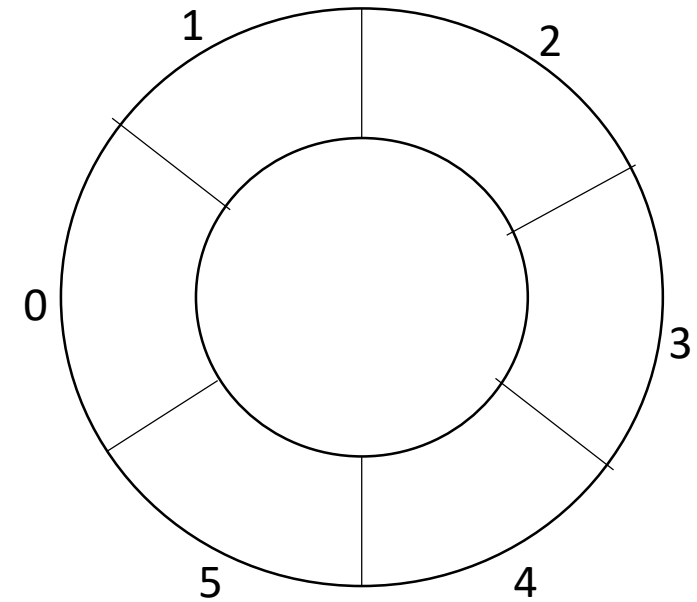
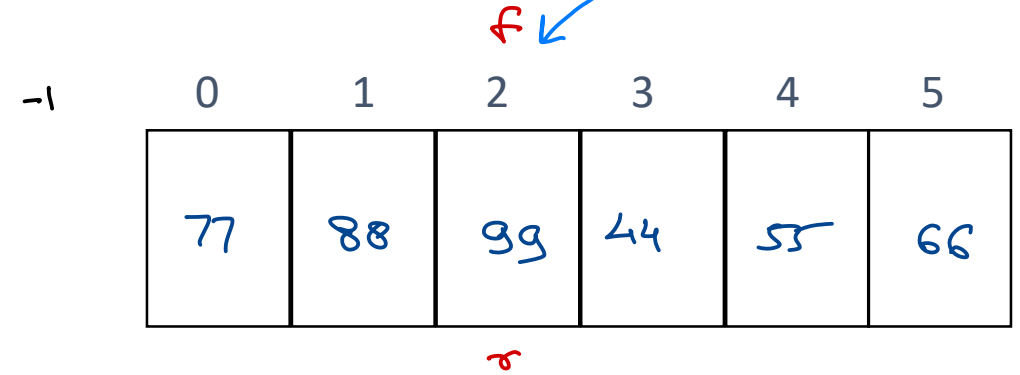
r/f

-1	$\xrightarrow{+1}$	0	$\xrightarrow{\%6}$	0
0	\rightarrow	1	\rightarrow	1
1	\rightarrow	2	\rightarrow	2
2	\rightarrow	3	\rightarrow	3
3	\rightarrow	4	\rightarrow	4
4	\rightarrow	5	\rightarrow	5
5	\rightarrow	6	\rightarrow	0



Circular Queue

q full : $f == r$ && $f != -1$



Circular Queue

q empty: $f == r$

POP:

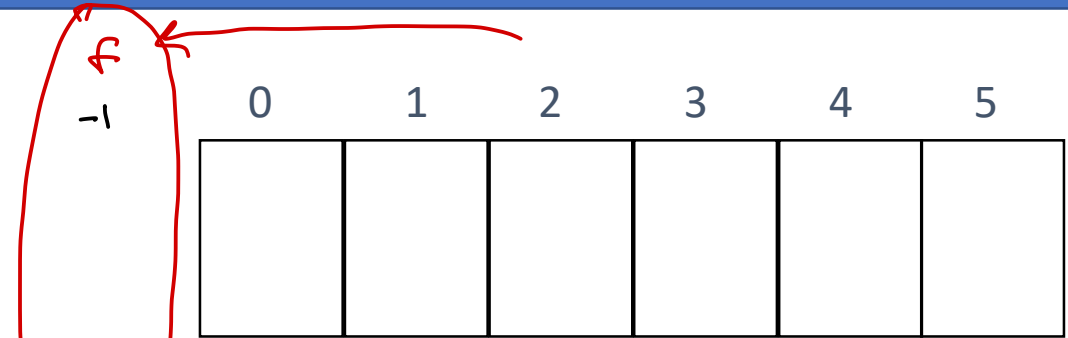
$f = (f + 1) \% \text{Size};$

if ($f == r$) \leftarrow check if ele
deleted was
last ele.

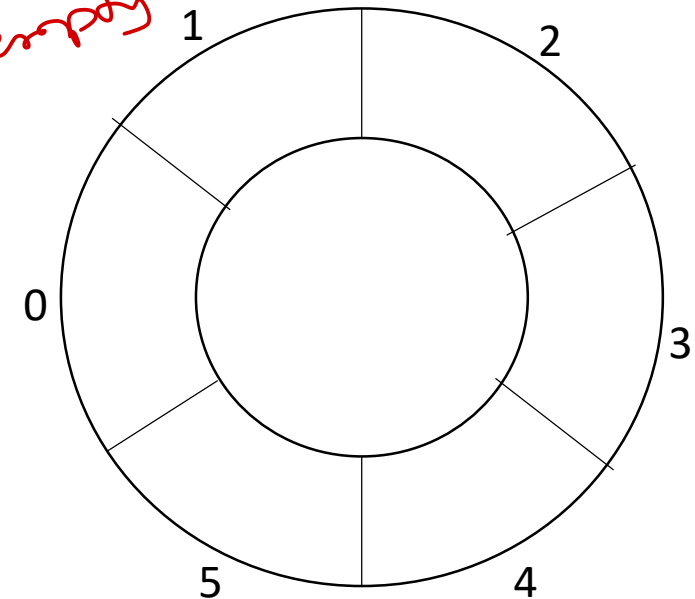
$\{$
 $f = -1;$
 $r = -1;$
 $\}$

3

\nearrow
trick
~~~~~



$\nwarrow$  q empty cond.





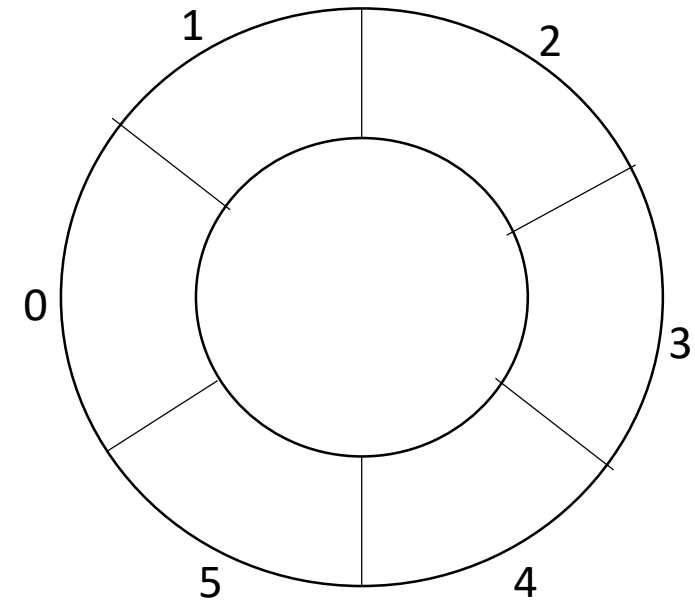
# Circular Queue

✓ q empty:  $f == r$  &&  $f == -1$

$f$   
-1



$r$



# Circular Queue

init:

$r = -1$   
 $f = -1$

push:

$r = (r + 1) \% \text{size};$   
 $\text{arr}[r] = \text{val};$

pop:

$f = (f + 1) \% \text{size};$   
 $\text{if}(f == r) \{$   
     $f = -1;$   
     $r = -1;$   
 $\}$

3

peek:

$i = (f + 1) \% \text{size};$   
 $\text{return arr}[i];$

is full:

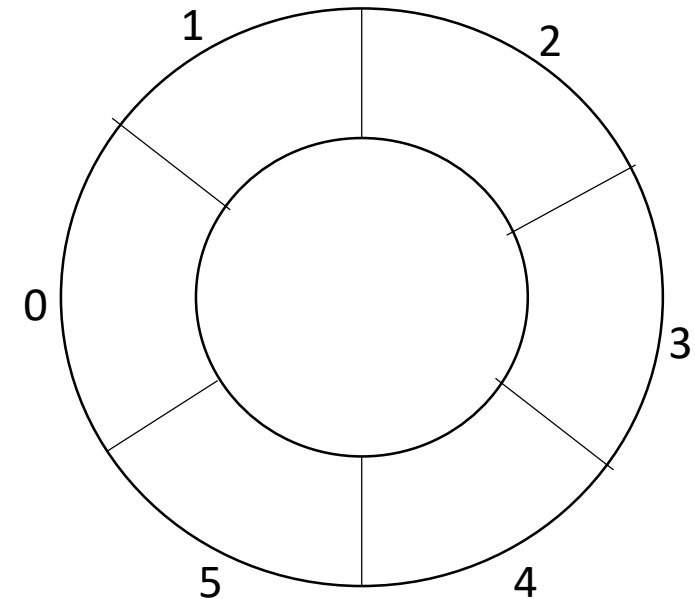
$(f == -1 \ \&\& \ r == \text{size} - 1)$

OR

$(f == r \ \&\& \ f != -1)$

is empty:

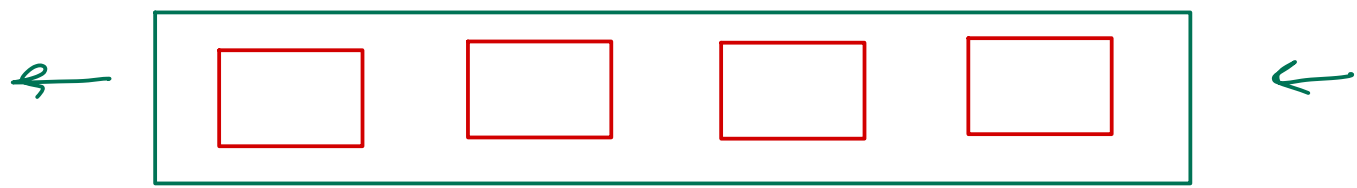
$(f == r \ \&\& \ f == -1)$



# DeQueue → java.util.Deque → 1.6

- In double ended queue, values can be added or deleted from front end or rear end.

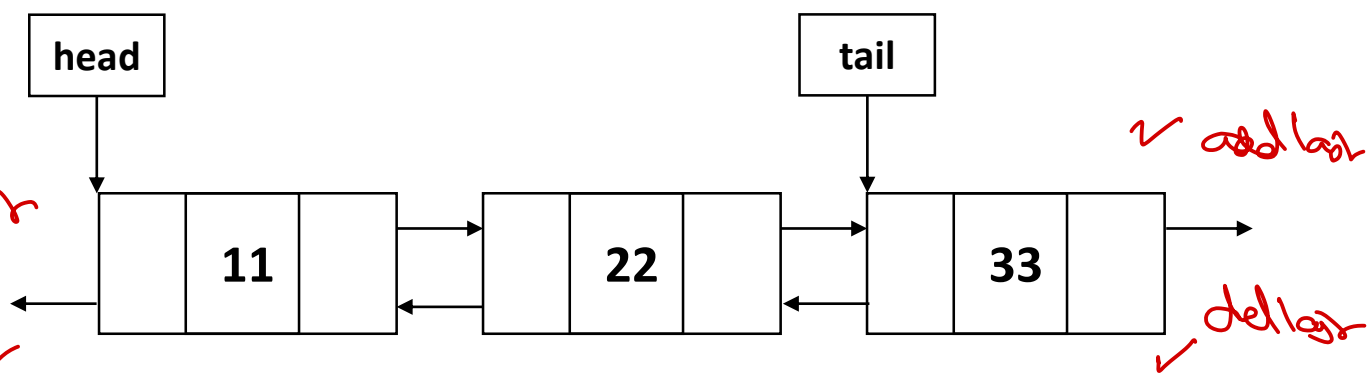
→ appn specific queue.



input restricted ← rear  
front ← output restricted



✓ add front  
✓ del from



# Priority queue

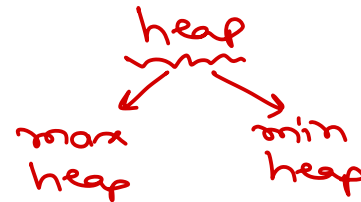
- In priority queue, element with highest priority is removed first.

push/pop  $\rightarrow$  may not be  $O(1)$ .

$\downarrow$   
elements need to be arranged  
as per their priorities

efficient impl  $\rightarrow$  push/pop  $\rightarrow$   $O(\log n)$

$\uparrow$   
implemented using



# applications

real life: queue at ticket counter.

① CPU scheduling

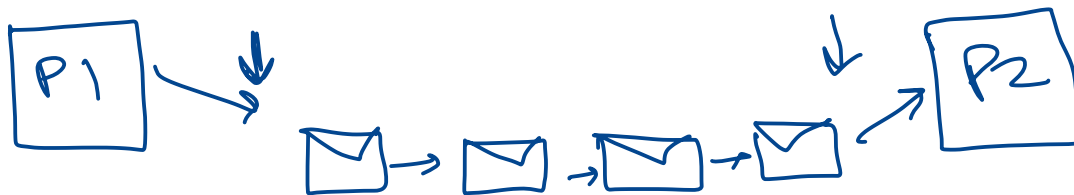
- FCFS, SJF, RR, ...  
    ↑          ↑  
  FIFO      Priority queue

④ processes wait for I/O devices in queue waiting queue.

② page replacement

- FIFO, LRU, ...

③ message queue (IPC)



Stack  $\rightarrow$  push/pop op from same end (top). - LIFO not

init:  
 $top = -1$

peek:  
 $\text{return arr[top]};$

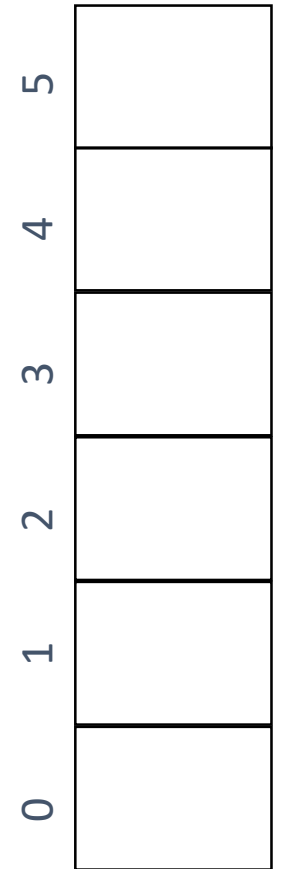
push:  
 $top++;$   
 $\text{arr}[top] = \text{val};$

pop:  
 $top--;$

isEmpty:  
 $top == -1$

isFull:  
 $top == \text{size} - 1$

22  
33  
44  
55  
11



top  $\rightarrow$

1



# Stack / Queue in Java collections

- class java.util.Stack<E>

- E push(E); ✓
- E pop(); ✓
- E peek(); ✓
- boolean isEmpty(); ✓

- interface java.util.Queue<E>

- boolean offer(E e); ✓ Push
- E poll(); ✓ POP
- E peek(); ✓
- boolean isEmpty(); ✓

LinkedList<>



infix, postfix, prefix

\* notations for math expression.

$a + b$  → infix ← human  
→ parenthesis  
 $+ a b$  → prefix  
 $a b +$  → postfix } Computer  
→ no parenthesis

precedence

① ( )

② \$

③ \* / L → R

④ + - L → R

$$5 + 9 - 4 * (8 - 6 / 2) + 1 $ (7 - 3)$$

Ⓐ  $5 + 9 - 4 * (8 - 6 / 2) + 1 $ (7 - 3)$

Ⓑ  $5 + 9 - 4 * 8 6 2 / - + 1 $ (7 - 3)$

Ⓒ  $5 + 9 - 4 * 8 6 2 / - + 1 $ 7 3 -$

Ⓓ  $5 + 9 - 4 * 8 6 2 / - + 1 7 3 - $$

Ⓔ  $5 + 9 - 4 8 6 2 / - * + 1 7 3 - $$

Ⓕ  $5 9 + - 4 8 6 2 / - * + 1 7 3 - $$

Ⓖ  $5 9 + 4 8 6 2 / - * - + 1 7 3 - $$

Ⓗ  $5 9 + 4 8 6 2 / - * - 1 7 3 - $ +$

$$5 + 9 - 4 * (8 - 6 / 2) + 1 $ (7 - 3)$$

$$+ - + 5 9 * 4 - 8 / 6 2 $ 1 - 7 3$$





# Infix to Postfix

$$5 + 9 - 4 * (8 - 6 / 2) + 1 \$ (7 - 3)$$

stack of  
operators

$$\bullet 5 + 9 - 4 * (8 - 6 / 2) + 1 \$ (7 - 3)$$

$$5 \ 9 \ + \ 4 \ 8 \ 6 \ 2 \ / \ - \ * \ - \ 1 \ 7 \ 3 \ - \ \$ \ +$$

- ① traverse infix expr left to right.
- ② if operand found, append to postfix result.
- ③ if operator found, push on stack.
  - \* if priority of topmost op in stack  $\geq$  priority of cur op, then pop & append to postfix.
- ④ when all space from infix completed,  
pop all ops from stack & append to postfix.
- ⑤ if opening ( found, push on stack.
- ⑥ if closing ) found, pop all ops from stack  
& append to postfix, until ( is found  
also pop & discard  $\_$ .

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# Infix to Prefix

•  $5 + 9 - 4 * (8 - 6 / 2) + 1 * (7 - 3) = -5$

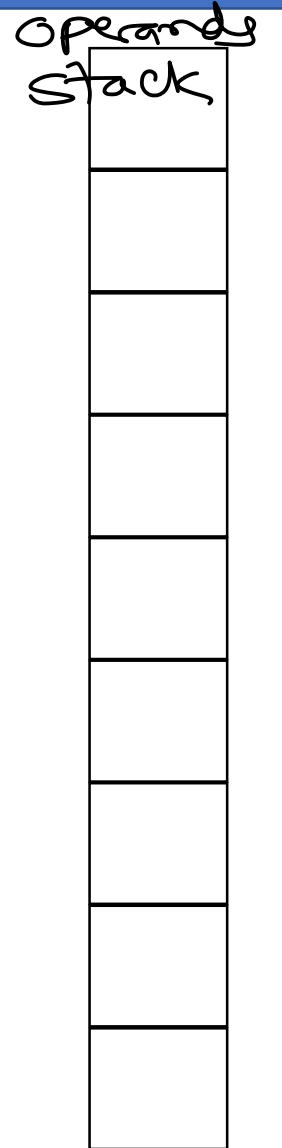
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# Postfix Evaluation

5 9 + 4 8 6 2 / - \* - 1 7 3 - \$ +

• 5 9 + 4 8 6 2 / - \* - 1 7 3 - \$ +



- ① parse postfix from left to right;
- ② if operand, push on stack.
- ③ if operator, pop two operands, calc result & push it on stack.
  - ↙ first popped is 2nd operand
  - ↘ second popped is 1st operand
- ④ when all symbols are done, pop final result from stack & return. -5



# Prefix Evaluation

• + - + 5 9 \* 4 - 8 / 6 2 \$ 1 - 7 3

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# Postfix to Infix

- While there are input symbol left
  - Read the next symbol from input.
  - If the symbol is an operand , Push it onto the stack.
  - Otherwise, the symbol is an operator.
  - If there are fewer than 2 values on the stack
    - Show Error
  - Else
    - Pop the top 2 values from the stack.
    - Put the operator, with the values as arguments and form a string.
    - Encapsulate the resulted string with parenthesis.
    - Push the resulted string back to stack.
  - If there is only one value in the stack
    - That value in the stack is the desired infix string.
  - If there are more values in the stack
    - Show Error

• a b c - + d e - f g - h + / \*

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# Prefix to Postfix

- Read the Prefix expression in reverse order (from right to left)
  - If the symbol is an operand, then push it onto the Stack
  - If the symbol is an operator, then pop two operands from the Stack
  - Create a string by concatenating the two operands and the operator after them.
  - $\text{string} = \text{operand1} + \text{operand2} + \text{operator}$
  - And push the resultant string back to Stack
  - Repeat the above steps until end of Prefix expression.
- $* + A B - C D$



# Parenthesis Balancing

•  $5 + ([9 - 4] * (8 - \{6 / 2\}))$

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*Thank you!*

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