Pimpri-Chinchwad Educational Trusts Pimpri-Chinchwad College of Engineering And Research, Ravet, Pune



DEPARTMENT OF COMPUTER ENGINEERING

LAB MANUAL

DATA STRUCTURES LABORATORY

Course Code: 210246

Vision – Mission of the Institute

Vision

To be a Premier institute of technical education & research to serve the need of society and all the stakeholders.

Mission

To establish state-of-the-art facilities to create an environment resulting in individuals who are technically sound having professionalism, research and innovative aptitude with high moral and ethical values.

Vision – Mission of the Computer Department

Vision

To strive for excellence in the field of Computer Engineering and Research through Creative Problem Solving related to societal needs

Mission:

- 1. Establish strong fundamentals, domain knowledge and skills among the students with analytical thinking, conceptual knowledge, social awareness and expertise in the latest tools & technologies to serve industrial demands
- 2. Establish leadership skills, team spirit and high ethical values among the students to serve industrial demands and societal needs
- 3. Guide students towards Research and Development, and a willingness to learn by connecting themselves to the global society

Program Educational Objectives (PEO)

- 1. To prepare globally competent graduates having strong fundamentals, domain knowledge, upd with modern technology to provide the effective solutions for engineering problems.
- 2. To prepare the graduates to work as a committed professional with strong professional ethics values, sense of responsibilities, understanding of legal, safety, health, societal, cultural environmental issues.
- 3. To prepare committed and motivated graduates with research attitude, lifelong learn investigative approach, and multidisciplinary thinking.
- 4. To prepare the graduates with strong managerial and communication skills to work effectivel individual as well as in teams.

Program Specific Outcomes (PSO)

A graduate of the Computer Engineering Program will demonstrate-

PSO1: Problem-Solving Skills- The ability to apply standard practices and strategies in software project development using open-ended programming environments to deliver a quality project.

PSO2: Professional Skills-The ability to understand, analyze and develop computer programs in the areas related to algorithms, software testing, application software, web design, data analytics, IOT and networking for efficient design of computer-based systems.

PSO3: Successful Career and Entrepreneurship- The ability to employ modern computer languages, environments, and platforms in creating innovative career paths to be an entrepreneur, and a zest for higher studies, and to generate IPR & Deliver a quality project.

Companion Course: Fundamentals of Data Structures (210242)

Course Objectives:

To understand basic techniques and strategies of algorithm analysis, the memory requirement for various data structures like array, linked list, stack, queue etc. using concepts of python and C++ programming language.

Course Outcomes:

On completion of the course, student will be able to-

CO	Statements
C206.1	Use algorithms on various linear data structure using sequential organization to solve real life problems.
C206.2	Analyze problems to apply suitable searching and sorting algorithm to various applications.
C206.3	Analyze problems to use variants of linked list and solve various real-life problems.
C206.4	Implement stack and queue data structures and algorithms for solving different kinds of problems.

Index

Group No.	Ass No	Experiment Name	CO	Page No.
Group	1	In second year, computer engineering class, group A student's play	C206.1	01
A		cricket, group B students play badminton and group C students play		
		football. Write a Python program using functions to compute		
		following: -		
		a) List of students who play both cricket and badminton		
		b) List of students who play either cricket or badminton but not both		
		c) Number of students who play neither cricket nor badminton		
		d) Number of students who play cricket and football but not		
		badminton.		
		(Note-While realizing the group, duplicate entries should be avoided,		
		do not use SET built-in functions)		
	2	Write a Python program to store marks scored in subject "Fundamental	C206.1	05
		of Data Structure" by N students in the class. Write functions to		
		compute following:		
		a) The average score of class		
		b) Highest score and lowest score of class		
		c) Count of students who were absent for the test		
		d) Display mark with highest frequency		
	_	Write a Python program to compute following computation on matrix:	C206.1	08
	3	a) Addition of two matrices		
		b) Subtraction of two matrices		
		c) Multiplication of two matrices		
		d) Transpose of a matrix		
Group	4	a) Write a Python program to store names and mobile numbers of your	C206.2	13
В		friends in sorted order on names. Search your friend from list using		
		binary search (recursive and non-recursive). Insert friend if not present		
		in phonebook		
		b) Write a Python program to store names and mobile numbers of your		

		friends in sorted order on names. Search your friend from list using Fibonacci search. Insert friend if not present in phonebook.		
	5	Write a Python program to store first year percentage of students in	C206.2	18
		array. Write function for sorting array of floating-point numbers in		
		ascending order using a) Selection Sort b) Bubble sort and display top		
		five scores		
	6	Write a Python program to store first year percentage of students in	C206.2	22
		array. Write function for sorting array of floating-point numbers in		
		ascending order using quick sort and display top five scores.		
Group	7	Department of Computer Engineering has student's club named	C206.3	26
C		'Pinnacle Club'. Students of Second, third and final year of department		
		can be granted membership on request. Similarly, one may cancel the		
		membership of club. First node is reserved for president of club and		
		last node is reserved for secretary of club. Write C++ program to		
		maintain club member's information using singly linked list. Store		
		student PRN and Name. Write functions to		
		a) Add and delete the members as well as president or even		
		secretary.		
		b) Compute total number of members of club		
		c) Display members		
		d) Display list in reverse order using recursion		
		e) Two linked lists exist for two divisions. Concatenate two lists.		
	8	Second year Computer Engineering class, set A of students like	C206.3	35
		Vanilla Ice-cream and set B of students like butterscotch ice-cream.		
		Write C/C++ program to store two sets using linked list. compute and		
		display i. Set of students who like either vanilla or butterscotch or both		
		ii. Set of students who like both vanilla and butterscotch iii. Set of		
		students who like only vanilla not butterscotch iv. Set of students who		
		like only butterscotch not vanilla v. Number of students who like		
		neither vanilla nor butterscotch		

Group	9	A palindrome is a string of character that's the same forward and		43	
D		backward. Typically, punctuation, capitalization, and spaces are			
		ignored. For example, Poor Dan is in a droop is a palindrome, as can			
		be seen by examining the characters —poor danisina droop and			
		observing that they are the same forward and backward. One way to			
		check for a palindrome is to reverse the characters in the string and			
		then compare with them the original-			
		in a palindrome, the sequence will be identical. Write C++			
		program with functions-			
		1. To check whether given string is palindrome or not that uses			
		a stack to determine whether a string is a palindrome.			
		2. To remove spaces and punctuation in string, convert all the			
		Characters to lowercase, and then call above Palindrome			
		checking function to check for a palindrome			
	3. To print string in reverse order using stack				
	In any language program mostly syntax error occurs due to		C206.4	47	
	10	unbalancing delimiter such as (),{},[]. Write C++ program using stack			
		to check whether given expression is well parenthesized or not.			
Group			C206.4	52	
E	11	example is the creation of a job queue by an operating system. If the			
		operating system does not use priorities, then the jobs are processed in			
		the order they enter the system. Write C++ program for simulating job			
		queue. Write functions to add job and delete job from queue.			
		A double-ended queue (deque) is a linear list in which additions and	C206.4	58	
	12	deletions may be made at either end. Obtain a data representation			
		mapping a deque into a one- dimensional array. Write C++ program to			
		simulate deque with functions to add and delete elements from			
		either end of the deque.			
	13	Pizza parlor accepting maximum M orders. Orders are served in first	C206.4	61	
		come first served basis. Order once placed cannot be cancelled. Write			
		C++ program to simulate the system using circular queue using array.			

References:

- 1. Steven S S. Skiena, "The Algorithm Design Manual", Springer, 2nd ed. 2008 Edition, ISBN-13: 978-1849967204, ISBN-10: 1849967202.
- 2. Allen Downey, Jeffery Elkner, Chris Meyers, "How to think like a Computer Scientist: Learning with Python", Dreamtech Press, ISBN: 9789351198147.
- 3. M. Weiss, "Data Structures and Algorithm Analysis in C++", 2nd edition, Pearson Education, 2002, ISBN-81-7808-670-0.
- 4. Brassard and Bratley, "Fundamentals of Algorithmic", Prentice Hall India/Pearson Education, ISBN 13-9788120311312.

Rubrics for Continuous Assessment in the laboratory

Parameters	Marks	Points		
		2 Marks	1 Marks	0 Marks
Timely		Students Completed	Students completed	Students completed
completion of	02	laboratory journal within	laboratory journal after a	laboratory journal at the
laboratory		time.	week.	end of Semester.

journal						
		4 Marks	3 Marks	2 Marks	1 Marks	0 Marks
Understanding of the assignment	04	Students performed the experiment/ executed the given program, extend the program by adding new features and answered all the questions.	Students performed the experiment/ executed the given program and answered all the questions.	Students performed the experiment / executed the given program with the help of peers and answered only a few questions.	Students performed the experiment/ executed the given program with the help of peers and did not answer any questions.	Students did not perform the experiment and did not answer any questions.
		4 Marks	3 Marks	2 Marks	1 Marks	0 Marks
		written in an extraordinary style and voice, Well organized interesting and voi Somew	Written in an	Whitton in little	Written in	Written in
Presentation /	0.4		interesting style	Written in little style, Gives	bad style, Gives some	bad style, no new
Clarity of journal writing	04		1	some new	new	information
Journal Writing			informative and	information but	information	and very
		Very informative.	organized.	poorly organized.	but poorly	poorly
			organized.		organized.	organized.

Abbreviation:

- T- Timely completion of laboratory journal
 U- Understanding of the assignment
 P- Presentation / Clarity of journal writing

- S- Sum

This is to certify that Mr./Miss/Mrs				
Roll No.:Exam.	Seat No.:	of SE/TE/BE Computer has carried out above		
practical /term work wi	thin PCCOER as	prescribed by Savitribai	Phule Pune University, Pune	
during the academic ye	ear 20 -20 .	His/Her performance	is satisfactory and attendance	
is%.				
Date:	Faculty I/C	HOD	Principal	

Assignment no: 01

Aim: To study and understand the concept of set theory using python.

Problem definition:

In second year, computer engineering class, group A student's play cricket, group B students play badminton and group C students play football. Write a Python program using functions to compute following: -

a) List of students who play both cricket and badminton

b) List of students who play either cricket or badminton but not both

c) Number of students who play neither cricket nor badminton

d) Number of students who play cricket and football but not badminton.

(Note-While realizing the group, duplicate entries should be avoided, do not use SET built-in functions)

Learning Objectives:

To understand basic techniques and strategies of algorithm using concepts of python.

Learning Outcome:

Students will be able to use algorithms on various linear data structure using sequential organization to solve real life problems.

Theory:

Python:

Python is a high-level, interpreted, interactive and object-oriented scripting language. Python is designed to be highly readable. It uses English keywords frequently where as other languages use punctuation, and it has fewer syntactical constructions than other languages.

Set Operations:

We have to perform here different set operations like Union, Intersections, Difference, Symmetric Difference.

Universal set U:

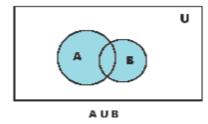
Often a discussion involves subsets of some particular set called the universe of discourse (or briefly universe), universal set or space. The elements of a space are often called the points of the space. We denote the universal set by U.

Example. The set of all even integers could be considered a subset of a universal set consisting of all the integers. Or they could be considered a subset of a universal set consisting of all the rational numbers. Or of all the real numbers.

Often the universal set may not be explicitly stated and it may be unclear as to just what it is. At other times it will be clear.

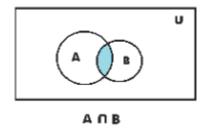
1. Union Operation: In set theory the union of collection of sets is the set of all distinct elements in the collection. The union of two sets A and B is the set consisting of all elements in A plus all elements in B and is denoted by $A \cup B$ or A + B.

Example. If $A = \{a, b, c, d\}$ and $B = \{b, c, e, f, g\}$ then $A \cup B = \{a, b, c, d, e, f, g\}$.



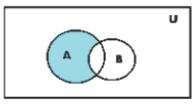
2. Intersection operation: In set theory intersection is a operation where we collect common elements of different sets. The intersection of two sets A and B is the set consisting of all elements that occur in both A and B (i.e. all elements common to both) and is denoted by $A \cap B$, A \cdot B or AB.

Example. If
$$A = \{a, b, c, d\}$$
 and $B = \{b, c, e, f, g\}$ then $A \cap B = \{b, c\}.$



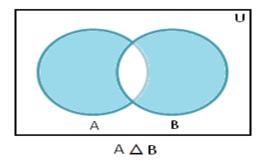
3. Difference Operation: It is a generalization of the idea of the compliment of a set and as a such is sometimes called the relative compliment of T with respect to S where T and s are two sets. The set consisting of all elements of a set A that do not belong to a set B is called the difference of A and B and denoted by A - B.

Example. If $A = \{a, b, c, d\}$ and $B = \{b, c, e, f, g\}$ then $A - B = \{a, d\}$.



A - B

4. Symmetric Difference: The symmetric difference between two sets S and t is the union of S-T and T-S. The symmetric difference using Venn diagram of two subsets A and B is a sub set of U, denoted by A \triangle B and is defined by A \triangle B = (A – B) \cup (B – A)



Input: Enter the total number of students in class, also enter the student who plays cricket, badminton and football.

Output: Union, intersection, set difference of entered students.

Algorithm/Pseudo code:

1. Function for union:

```
def find_union_set(A,B,C):
    for i in range(len(A)):
        C.append(A[i])
    for i in range(len(B)):
        flag = search_set(A,B[i]);
        if(flag == 0) :
            C.append(B[i])
```

2. Function for Intersection:

```
def find_intersection_set(A,B,C):
  for i in range(len(A)):
    flag = search_set(B,A[i]);
    if(flag == 1) :
        C.append(A[i])
```

3. Function for Difference:

```
def find_difference_set(A,B,C):
  for i in range(len(A)):
    flag = search_set(B,A[i]);
    if(flag == 0) :
        C.append(A[i])
```

Software required: Open Source Python, Programming tool like Jupyter Notebook, Pycharm, Spyder.

Conclusion: Thus, we have studied use of set operations using python.

Assignment No-02

Aim: To illustrate the various functions in python.

Problem Statement: Write a Python program to store marks scored in subject "Fundamental of

Data Structure" by N students in the class. Write functions to compute following:

a) The average score of class

b) Highest score and lowest score of class

c) Count of students who were absent for the test

d) Display mark with highest frequency

Learning Objectives:

To understand basic techniques and strategies of algorithm analysis, the memory requirement for

various data structures using the concepts of python.

Learning Outcome: Students will be able to use algorithms on various linear data structure using

sequential organization to solve real life problems

Theory: Lists are one of the most powerful tools in python. They are just like the arrays declared

in other languages. But the most powerful thing is that list need not be always homogenous. A

single list can contain strings, integers, as well as objects. Lists can also be used for implementing

stacks and queues. Lists are mutable, i.e., they can be altered once declared.

A tuple is a sequence of immutable Python objects. Tuples are just like lists with the exception that

tuples cannot be changed once declared. Tuples are usually faster than lists

Input: enter marks scored in subject "Fundamental of Data Structure" by N students in the class.

Output: Average score, highest score, lowest score, count of absentees, marks with highest

frequency.

Algorithm/Pseudocode:

• The average score of class

```
def find_average_score_of_class(A):
    sum = 0
    for i in range(len(A)):
        if(A[i] != -1):
            sum = sum + A[i]
        avg = sum / len(A)
        display_marks(A)
        print("\nAverage score of class is %.2f\n\n"%avg)
```

• Highest score and lowest score of class

```
def find_highest_and_lowest_score_of_class(A):
    max = -1
    min = 31
    for i in range(1,len(A)):
        if(max < A[i]):
        max = A[i]
        max_ind = i
        if(min > A[i] and A[i] != -1):
        min = A[i]
        min_ind = i
        display_marks(A)
        print("Highest Mark Score of class is %d scored by student %d"%(max,max_ind+1))
print("Lowest Mark Score of class is %d scored by student %d"%(min,min_ind+1)
```

• Count of students who were absent for the test:

```
def find_count_of_absent_students(A) :
  count = 0
  for i in range(len(A)):
```

```
if(A[i] == -1):
    count += 1
display_marks(A)
print("\tAbsent Student Count = %d"%count)
```

• Display mark with highest frequency:

```
def display_mark_with_highest_frequency(A):
    freq = 0
    for i in range(len(A)):
        count = 0
        if(A[i] != -1):
            for j in range(len(A)):
            if(A[i] == A[j]):
                 count += 1
        if(freq < count):
            Marks = A[i]
            freq = count
        display_marks(A)
        print("\nMarks with highest frequency is %d (%d)"%(Marks,freq))</pre>
```

Software required: Open Source Python, Programming tool like Jupyter Notebook, PyCharm, Spyder.

Conclusion: Thus, we have studied the implementation of various python operations.

Assignment no: 03

Aim: To study and understand the concept of matrix realization using Python.

Problem definition:

Write a Python program to compute following computation on matrix:

- a) Addition of two matrices
- b) Subtraction of two matrices
- c) Multiplication of two matrices
- d) Transpose of a matrix

Learning Objectives:

To understand concept matrix using python.

Learning Outcome:

Students will be able to use algorithms on various linear data structure using sequential organization to solve real life problems.

Theory:

Matrix:

We say a matrix is $m \times n$ if it has m rows and n columns. These values are sometimes called the dimensions of the matrix. Note that, in contrast to Cartesian coordinates, we specify the number of rows (the vertical dimension) and then the number of columns (the horizontal dimension). In most contexts, the rows and columns are numbered starting with 1. Several programming APIs, however, index rows and columns from 0. We use aij to refer to the entry in i th row and j th column of the matrix A.

The Order of a Matrix is its size or dimensions. The order is always given as the number of rows by the number of columns ($R \times C$).

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix}$$
 This matrix is a 3 x 2 matrix.
$$\begin{bmatrix} 1 & 5 & 6 & 7 & 8 \\ -2 & 3 & 5 & 4 & -1 \end{bmatrix}$$
 rows by columns
$$2 \times 5 \text{ matrix}$$

For two matrices to be added or subtracted, the dimensions must be the same. If they are the same, then the corresponding entries are added or subtracted whichever the operation.

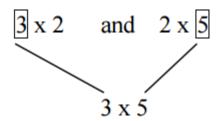
$$\begin{bmatrix} 2 & 5 \\ 3 & -6 \\ 4 & 0 \end{bmatrix} + \begin{bmatrix} -1 & 7 \\ 2 & 4 \\ -4 & 5 \end{bmatrix} = \begin{bmatrix} 1 & 12 \\ 5 & -2 \\ 0 & 5 \end{bmatrix}$$

For two matrices to be multiplied, their dimensions need to be analyzed to determine if it is possible. The number of columns of the first matrix MUST EQUAL the number of rows of the second matrix.

$$\begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix} \cdot \begin{bmatrix} 1 & 5 & 6 & 7 & 8 \\ -2 & 3 & 5 & 4 & -1 \end{bmatrix}$$
3 x2 and 2x 5

They are the same so we can multiply these two matrices.

The outside numbers tell the dimensions or the order of the resulting matrix.



The answer will be a 3 x 5 matrix.

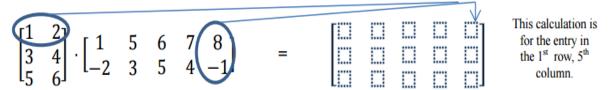
The position of each element (row, column) in the answer is a clue to how to multiply.

$$(1,1)(1,2)(1,3)(1,4)(1,5)$$

$$(2,1)(2,2)(2,3)(2,4)(2,5)$$

$$(3,1)(3,2)(3,3)(3,4)(3,5)$$
This entry is in the 1st row and 5th column so it is labeled (1, 5).

To do the multiplication of the two matrices, a calculation must be completed with the row and columns as follows: To obtain each entry in the solution matrix, we will look at the row in the first matrix and the column in the second matrix that correspond to the solution matrix entry. So, for the entry that belongs in the solution matrix in the location (1, 5) we will use the 1st row in the first matrix and the 5th column in the second matrix.



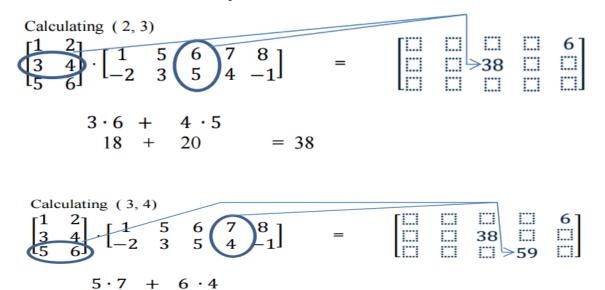
We will multiply the first entry in each and the second entry in each, and then we will add those two results together:

$$1 \cdot 8 + 2 \cdot -1$$

$$8 + -2 = 6$$

$$\begin{bmatrix} \vdots & \vdots & \vdots & \vdots & \vdots \\ \vdots & \vdots & \vdots & \vdots & \vdots \end{bmatrix}$$

This process must be done for each entry in the solution matrix. Below are a few more examples. Then, the final matrix after all calculations are completed.



Calculating (1, 1)
$$\begin{bmatrix}
1 & 2 \\
3 & 4 \\
5 & 6
\end{bmatrix}
\cdot
\begin{bmatrix}
1 \\
-2
\end{bmatrix}
\begin{bmatrix}
5 & 6 & 7 & 8 \\
3 & 5 & 4 & -1
\end{bmatrix}
=
\begin{bmatrix}
-3 & \Box & \Box & \Box & 6 \\
\Box & \Box & 38 & \Box & \Box \\
\Box & \Box & 59 & \Box
\end{bmatrix}$$

$$1 \cdot 1 + 2 \cdot -2$$

$$1 + -4 = -3$$

The final answer for this matrix multiplication:
$$\begin{bmatrix} -3 & 11 & 16 & 15 & 6 \\ -5 & 27 & 38 & 37 & 20 \\ -7 & 43 & 60 & 59 & 34 \end{bmatrix}$$

Input: Enter the data for first matrix and second matrix.

Output: addition, subtraction & multiplication of entered matrix and transpose of matrix.

Algorithm/Pseudocode:

Addition of two matrices:

```
def addition_matrix(M1,M2,M3,r,c) :
  for i in range(r) :
    A = []
  for j in range(c):
    A.append(M1[i][j] + M2[i][j])
    M3.append(A)
```

Subtraction of two matrices:

```
def substraction_matrix(M1,M2,M3,r,c):
  for i in range(r):
    A = []
    for j in range(c):
        A.append(M1[i][j] - M2[i][j])
    M3.append(A)
```

Multiplication of two matrices:

```
def multiplication_matrix(M1,M2,M3,r1,c1,c2) :
   for i in range(r1):
     A = []
     for j in range(c2):
       sum = 0
       for k in range(c1):
          sum = sum + (M1[i][k] * M2[k][j])
       A.append(sum)
     M3.append(A)
          Transpose of matrix:
def find_transpose_matrix(M,r,c,T):
  for i in range(c):
     A = []
     for j in range(r):
      A.append(M[j][i])
    T.append(A)
```

Software required: Open Source Python, Programming tool like Jupyter Notebook, Spyder.

Conclusion: thus, we have studied and implemented the matrix and performed different operations on it.

Assignment No- 04

Aim: To illustrate the various searching techniques.

Problem Statement: a) Write a Python program to store names and mobile numbers of your friends in sorted order on names. Search your friend from list using binary search (recursive and non-recursive). Insert friend if not present in phonebook

b) Write a Python program to store names and mobile numbers of your friends in sorted order on names. Search your friend from list using Fibonacci search. Insert friend if not present in phonebook.

Learning Objectives:

To understand concept of searching.

To compare time complexity of various searching techniques.

Learning Outcome: Students will be able to analyze problems to apply suitable searching and sorting algorithm to various applications.

Theory:

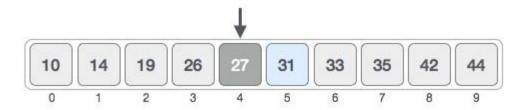
Binary Search:

For a binary search to work, it is mandatory for the target array to be sorted. We shall learn the process of binary search with a pictorial example. The following is our sorted array and let us assume that we need to search the location of value 31 using binary search.

First, we shall determine half of the array by using this formula –

$$mid = low + (high - low) / 2$$

Here it is, 0 + (9 - 0) / 2 = 4 (integer value of 4.5). So, 4 is the mid of the array.



Now we compare the value stored at location 4, with the value being searched, i.e. 31. We find that the value at location 4 is 27, which is not a match. As the value is greater than 27 and we have a sorted array, so we also know that the target value must be in the upper portion of the array.



We change our low to mid + 1 and find the new mid value again.

$$low = mid + 1$$
$$mid = low + (high - low) / 2$$

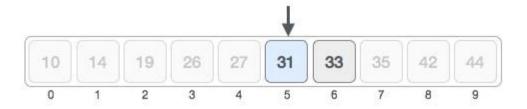
Our new mid is 7 now. We compare the value stored at location 7 with our target value 31.



The value stored at location 7 is not a match, rather it is more than what we are looking for. So, the value must be in the lower part from this location.



Hence, we calculate the mid again. This time it is 5.



We compare the value stored at location 5 with our target value. We find that it is a match.



We conclude that the target value 31 is stored at location 5.

Binary search halves the searchable items and thus reduces the count of comparisons to be made to very less numbers.

Fibonacci Search:

Fibonacci search is an efficient search algorithm based on divide and conquer principle that can find an element in the given sorted array with the help of Fibonacci series in O(log N) time complexity. This is based on Fibonacci series which is an infinite sequence of numbers denoting a pattern which is captured by the following equation:

$$F(n+1)=F(n)+F(n-1)$$

where F(i) is the ith number of the Fibonacci series where F(0) and F(1) are defined as 0 and 1 respectively.

The first few Fibonacci numbers are:

0,1,1,2,3,5,8,13....

$$F(0) = 0$$

$$F(1) = 1$$

$$F(2) = F(1) + F(0) = 1 + 0 = 1$$

$$F(3) = F(2) + F(1) = 1 + 1 = 2$$

$$F(4) = F(3) + F(2) = 1 + 2 = 3$$
 and so continues the series

Other searches like binary search also work for the similar principle on splitting the search space to a smaller space but what makes Fibonacci search different is that it divides the array in unequal parts and operations involved in this search are addition and subtraction only which means light arithmetic operations takes place and hence reducing the work load of the computing machine.

Input: Enter the names and mobile numbers of your friend in sorted order

Output: Searching by Binary search and Fibonacci search

Algorithm/Pseudo code:

```
Binary Search:
Procedure binary_search
 A \leftarrow sorted array
 n \leftarrow size of array
 x \leftarrow value to be searched
 Set lowerBound = 1
 Set upperBound = n
 while x not found
   if upperBound < lowerBound
     EXIT: x does not exists.
   set midPoint = lowerBound + ( upperBound - lowerBound ) / 2
   if A[midPoint] < x
     set lowerBound = midPoint + 1
   if A[midPoint] > x
     set upperBound = midPoint - 1
```

if A[midPoint] = x

EXIT: x found at location midPoint

end while

end procedure

Fibonacci Search:

Let the length of given array be n [0...n-1] and the element to be searched be x.

Then we use the following steps to find the element with minimum steps:

1. Find the smallest Fibonacci number greater than or equal to n. Let this number be fb(M) [m'th Fibonacci number]. Let the two Fibonacci numbers preceding it be fb(M-1) [(m-1)'th Fibonacci number] and fb(M-2) [(m-2)'th Fibonacci number].

- 2. While the array has elements to be checked:
- -> Compare x with the last element of the range covered by fb(M-2)
- -> If x matches, return index value
- -> Else if x is less than the element, move the third Fibonacci variable two Fibonacci down, indicating removal of approximately two-third of the unsearched array.
- -> Else x is greater than the element, move the third Fibonacci variable one Fibonacci down. Reset offset to index. Together this results into removal of approximately front one-third of the unsearched array.
- 3. Since there might be a single element remaining for comparison, check if fbMm1 is '1'. If Yes, compare x with that remaining element. If match, return index value.

Software required: Open Source Python, Programming tool like Jupyter Notebook, PyCharm, Spyder.

Conclusion: Thus, we have studied the implementation of various sorting techniques (Bubble and Selection)

Assignment No- 05

Aim: To illustrate the various sorting techniques.

Problem Statement: Write a Python program to store first year percentage of students in array.

Write function for sorting array of floating-point numbers in ascending order using a) Selection

Sort b) Bubble sort and display top five scores

Learning Objectives:

To understand concept of sorting.

To compare time complexity of various sorting techniques.

Learning Outcome: Students will be able to analyze problems to apply suitable searching and

sorting algorithm to various applications.

Theory:

Bubble Sort

Bubble sort is a simple and well-known sorting algorithm. It is used in practice once in a blue

moon and its main application is to make an introduction to the sorting algorithms. Bubble sort

belongs to O(n2) sorting algorithms, which makes it quite inefficient for sorting large data

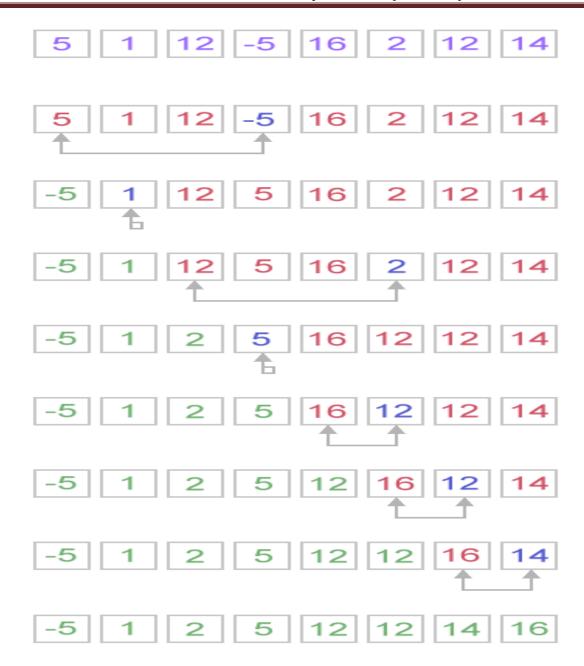
volumes. Bubble sort is **stable** and **adaptive**.

Example:

5 1 12 -5 16	unsorted
5 1 12 -5 16 1 5 12 -5 16 1 5 12 -5 16 1 5 -5 12 16	5 > 1, swap 5 < 12, ok 12 > -5, swap 12 < 16, ok
1 5 -5 12 16 1 5 -5 12 16 1 -5 5 12 16	1 < 5, ok 5 > -5, swap 5 < 12, ok
1 -5 5 12 16 -5 1 5 12 16	1 > -5, swap 1 < 5, ok
-5 1 5 12 16	-5 < 1, ok
-5 1 5 12 16	sorted

Selection Sort

Selection sort is one of the $O(n^2)$ sorting algorithms, which makes it quite inefficient for sorting large data volumes. Selection sort is notable for its programming simplicity and it can over perform other sorts in certain situations (see complexity analysis for more details).



Input: Enter the first percentage of students

Output: sorting by selection and bubble sort and display top 5 students marks

Algorithm for Bubble Sort:

We assume list is an array of n elements. We further assume that swap function, swaps the values of given array elements.

begin BubbleSort(list)

```
for all elements of list

if list[i] > list[i+1]

swap(list[i], list[i+1])

end if

end for

return list

end BubbleSort
```

Time Complexity = $O(n^2)$

Algorithm for Selection Sort:

We assume list is an array of n elements. We further assume that swap function, swaps the values of given array elements.

```
Begin SelectionSort(list)

for all elements of list

for j to all element of list

if list[i] > list[j]

swap(list[i], list[j])

end if

end for

end for

return list

end BubbleSort
```

Software required: Open Source Python, Programming tool like Jupyter Notebook, PyCharm, Spyder.

Conclusion: Thus, we have studied the implementation of various sorting techniques (Bubble and Selection)

Assignment No-06

Aim: To illustrate Quick Sort.

Problem Statement: Write a Python program to store first year percentage of students in array.

Write function for sorting array of floating-point numbers in ascending order using quick sort and

display top five scores.

Learning Objectives:

To understand concept of sorting.

To find time complexity of Quick Sort.

Learning Outcome: Students will be able to analyze problems to apply suitable searching and

sorting algorithm to various applications.

Theory of Quick sort:

Quicksort is a fast sorting algorithm, which is used not only for educational purposes, but

widely applied in practice. On the average, it has O(n log n) complexity, making quicksort

suitable for sorting big data volumes. The idea of the algorithm is quite simple and once you

realize it, you can write quicksort as fast as bubble sort.

Algorithm

The divide-and-conquer strategy is used in quicksort. Below the recursion step is described:

• Choose a pivot value. We take the value of the middle element as pivot value, but it can be

any value, which is in range of sorted values, even if it doesn't present in the array.

Partition. Rearrange elements in such a way, that all elements which are lesser than the

pivot go to the left part of the array and all elements greater than the pivot, go to the right part

of the array. Values equal to the pivot can stay in any part of the array. Notice, that array may

be divided in non-equal parts.

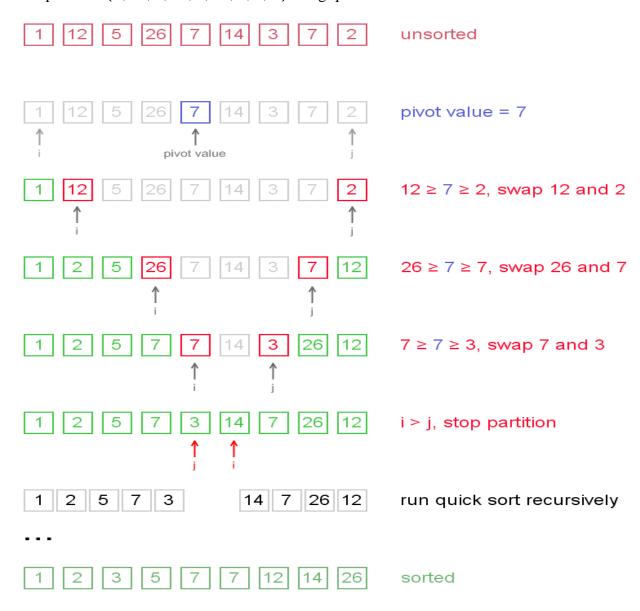
• Sort both parts. Apply quicksort algorithm recursively to the left and the right parts.

Partition algorithm in detail

There are two indices i and j and at the very beginning of the partition algorithm i points to the first element in the array and j points to the last one. Then algorithm moves i forward, until an element with value greater or equal to the pivot is found. Index j is moved backward, until an element with value lesser or equal to the pivot is found. If $i \le j$ then they are swapped and i steps to the next position (i+1), j steps to the previous one (j-1). Algorithm stops, when i becomes greater than j.

After partition, all values before i-th element are less or equal than the pivot and all values after j-th element are greater or equal to the pivot.

Example. Sort {1, 12, 5, 26, 7, 14, 3, 7, 2} using quicksort.



Time complexity: O(n log n) **Input:** Enter the first-year percentage of students Output: Sorting by Quick sort Algorithm /Pseudo Code: void quickSort(int arr[], int left, int right) { int i = left, j = right; int tmp; int pivot = arr[(left + right) / 2];/* partition */ while $(i \le j)$ { while (arr[i] < pivot) i++; while (arr[j] > pivot)j--; if $(i \le j)$ { tmp = arr[i];arr[i] = arr[j];arr[j] = tmp;i++; j--; **}**; /* recursion */ if (left < j)quickSort(arr, left, j); if (i < right)

quickSort(arr, i, right);

Software required: Open Source Python, Programming tool like Jupyter Notebook, PyCharm, Spyder.

Conclusion: Thus, we have studied the implementation of Quick Sort.

Assignment no: 07

Aim: To understand and implement singly linked list

Problem definition: Department of Computer Engineering has student's club named 'Pinnacle

Club'. Students of Second, third and final year of department can be granted membership on

request. Similarly, one may cancel the membership of club. First node is reserved for president of

club and last node is reserved for secretary of club. Write C++ program to maintain club member's

information using singly linked list. Store student PRN and Name. Write functions to

a) Add and delete the members as well as president or even secretary.

b) Compute total number of members of club

c) Display members

d) Display list in reverse order using recursion

e) Two linked lists exist for two divisions. Concatenate two lists.

Learning objectives

To understand concept of sll.

To analyze the working of functions.

Learning outcome

Students will be able to analyze problems to use variants of linked list and solve various real-life

problems.

Theory:

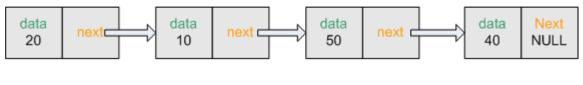
Linked list:

Linked list is one of the fundamental data structures, and can be used to implement other data

structures. In a linked list there are different numbers of nodes. Each node is consists of two fields.

The first field holds the value or data and the second field holds the reference to the next node or

null if the linked list is empty.



Linked list

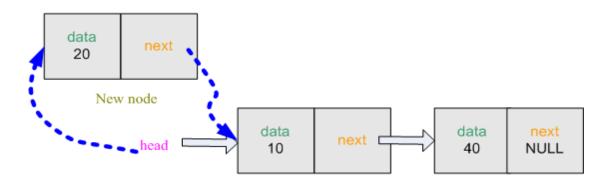
The singly-linked list is the easiest of the linked list, which has one link per node.

Linked list operation:

First, we create a structure "node". It has two members and first is int data which will store the information and second is node *next which will hold the address of the next node. Linked list structure is complete so now we will create linked list. We can insert data in the linked list from 'front' and at the same time from 'back'. Now we will examine how we can insert data from front in the linked list.

1) Insert from front:

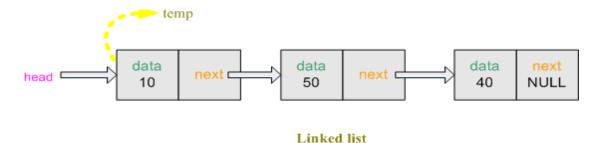
At first initialize node type. Then we take the data input from the user and store in the node info variable. Create a temporary node node *temp and allocate space for it. Then place info to temp->data. So the first field of the node *temp is filled. Now temp->next must become a part of the remaining linked list (although now linked list is empty but imagine that we have a 2 node linked list and head is pointed at the front) So temp->next must copy the address of the *head (Because we want insert at first) and we also want that *head will always point at front. So *head must copy the address of the node *temp.



Linked list

2) Traverse

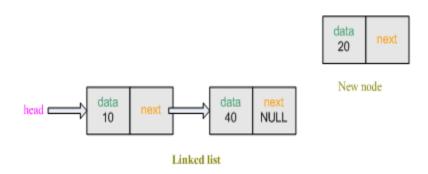
Now we want to see the information stored inside the linked list. We create node*temp1. Transfer the address of *head to *temp1. So *temp1 is also pointed at the front of the linked list. Linked list has 3 nodes. We can get the data from first node using temp1->data. To get data from second node, we shift *temp1 to the second node. Now we can get the data from second node.

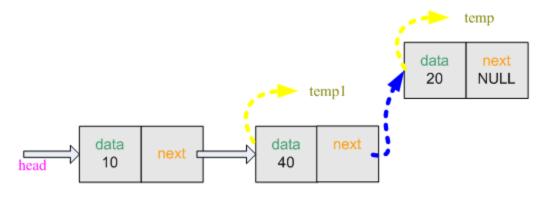


This process will run until the linked list's next is NULL.

3) Insert from back

Insert data from back is very similar to the insert from front in the linked list. Here the extra job is to find the last node of the linked list. Now, Create a temporary node node *temp and allocate space for it. Then place info to temp->data, so the first field of the node node *temp is filled. node *temp will be the last node of the linked list. For this reason,temp->next will be NULL. To create a connection between linked list and the new node, the last node of the existing linked list node *temp1's second field temp1->next is pointed to node *temp.

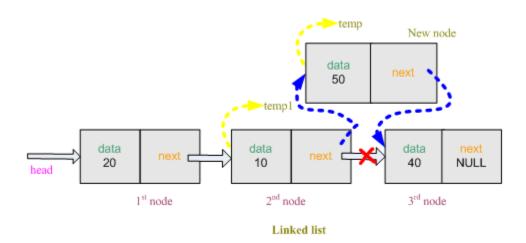




Linked list

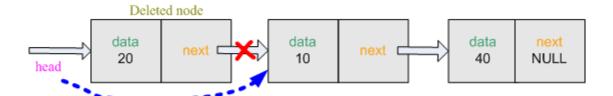
4) Insert after specified number of nodes

Insert data in the linked list after specified number of node is a little bit complicated. But the idea is simple. Suppose, we want to add a node after 2nd position. So, the new node must be in 3rd position. The first step is to go the specified number of node. Let, node *temp1 is pointed to the 2nd node now. Now, Create a temporary node node *temp and allocate space for it. Then place info to temp->next, so the first field of the node node *temp is filled. To establish the connection between new node and the existing linked list, new node's next must pointed to the 2nd node's (temp1) next. The 2nd node's (temp1) next must pointed to the new node(temp).



5) Delete from front:

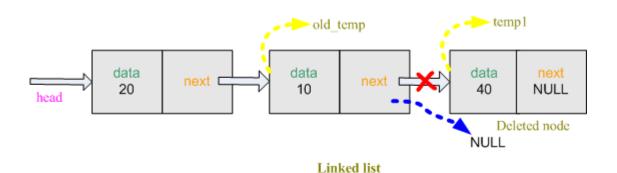
Delete a node from linked list is relatively easy. First, we create node*temp. Transfer the address of *head to*temp . So *temp is pointed at the front of the linked list. We want to delete the first node. So transfer the address of temp->next tohead so that it now pointed to the second node. Now free the space allocated for first node.



Linked list

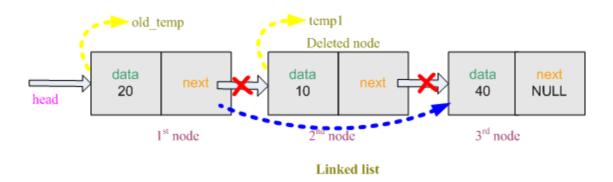
6) Delete from back

The last node's next of the linked list always pointed to NULL. So when we will delete the last node, the previous node of last node is now pointed at NULL. So, we will track last node and previous node of the last node in the linked list. Create temporary node * temp1 and *old_temp . Nownode *temp1 is now pointed at the last node and *old_temp is pointed at the previous node of the last node. Now rest of the work is very simple. Previous node of the last node old_temp will be NULL so it become the last node of the linked list. Free the space allocated for last lode.



7) Delete specified number of node

To delete a specified node in the linked list, we also require to find the specified node and previous node of the specified node. Create temporary node * temp1 , *old_temp and allocate space for it. Take the input from user to know the number of the node. Nownode *temp1 is now pointed at the specified node and *old_temp is pointed at the previous node of the specified node. The previous node of the specified node must connect to the rest of the linked list so we transfer the address of temp1->next to old_temp->next. Now free the space allocated for the specified node.



Input:

Enter members, president and secretary having their PRN and Name to be stored in the singly linked list. Every node would contain 2 fields: data and next pointer.

Output:

Insertion of node at front, at end, at any given position, deletion of node, reverse of linked list, count total number of nodes, merge two linked list.

Algorithm:

• Algorithm to insert President at first node:

Step 1: p - pointer to a linked list

Step 2: e - element to be added

Step 3: Getnode() returns a new node

q = Getnode()

info(q) = e

next(q) = p

p = q

return p

Step 4: end

• Algorithm to insert members at any position in link list:

```
Step 1: p - pointer to a linked list
Step 2: Getnode() returns a new node
Step 3: x - key node after which e is inserted
Step 4: e - element to be added
 k = p
       while k \neq \text{NIL} and info(k) = x \text{ do } // \text{ find the key node}
k = next(k) if k = NIL then
Write "Node not found" return p
q = Getnode()
info(q) = e
next(q) = next(k)
next(k) = q
return p
Step5:Stop
        Algorithm to insert secretary at Last node:
Step 1: p - pointer to a linked list
Step 2: Getnode() returns a new node
Step 3: e - element to be added
if p = NIL then
return p
k = p
while next(k) \neq NIL do // find the last node k = next(k)
q = Getnode()
info(q) = e
next(k) = q
next(q) = NIL
return p
Step 4: end.
```

• Algorithm to Delete any member Node

```
Step 1: p - pointer to linked list
Step 2: x - key node to be deleted
Step 3: k and pred – temporary variables
Step 4: k = p; pred = NIL
  while k \neq \text{NIL} and info(k) = x \text{ do } // \text{ find the key node } pred = k
 k = next(k)
  if k = NIL then
  Write "Node not found" else
 if pred = NIL // only one node in the list then
          p = next(p) else
 next(pred) = next(k)
return p
         end
     DeleteNode.
Step 5: end
```

• Algorithm to Display the members of link list.

Step 1: Create link list

Step 2: Scan and print the entire link list of the members having their PRN and name one by one.

Step 3: end

• Algorithm to Count the total number of members nodes in the link list.

```
Step 1: count = 0
Step 2: Increment count as each node is traversed current = head
while (current != NULL) then
count++
Step 3: end
```

• Algorithm to merge two linked list

Step 1: enter two linked list

Step 2: after first list is traverse merge the second linked list.

Step 3: end

• Algorithm to reverse the link list

Step 1: Iterative list reverse. Iterate through the list left-right. Move/insert each node to the front of the result list like a Push of the node.

```
Step 2: result = NULL;
Step 3: while (current != NULL)

next = current->next ( note the next node current->next = result move the node onto the result)

result = current

current = next

Step 4 : end
```

Software required: g++ / gcc compiler- / 64 bit fedora.

Conclusion: We understand & implement different operations on of linked list.

Assignment no: 08

Aim: To understand and implement set operation using linked list.

Problem definition: Second year Computer Engineering class, set A of students like Vanilla Ice-cream and set B of students like butterscotch ice-cream. Write C/C++ program to store two sets using linked list. compute and display i. Set of students who like either vanilla or butterscotch or both ii. Set of students who like both vanilla and butterscotch iii. Set of students who like only vanilla not butterscotch iv. Set of students who like only butterscotch not vanilla v. Number of students who like neither vanilla nor butterscotch

Learning objectives

To understand concept of set operations and linked list.

Learning outcome

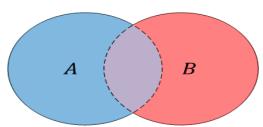
Students will be able to analyze problems to use variants of linked list and solve various reallife problems.

Theory:

Set

Elements are the objects contained in a **set**. A set may be defined by a common property amongst the objects. For example, the set E of positive even integers is the set $E=\{2,4,6,8,10..\}$

Two data sets A and B



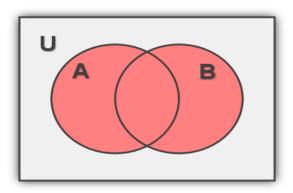
Definition (Union): The union of sets A and B, denoted by $A \cup B$, is the set defined as

$$A \cup B = \{ x \mid x \in A \lor x \in B \}$$

Example 1: If $A = \{1, 2, 3\}$ and $B = \{4, 5\}$, then $A \cup B = \{1, 2, 3, 4, 5\}$.

Example 2: If $A = \{1, 2, 3\}$ and $B = \{1, 2, 4, 5\}$, then $A \cup B = \{1, 2, 3, 4, 5\}$.

Note that elements are not repeated in a set.



Definition (Intersection): The intersection of sets A and B, denoted by $A \cap B$, is the set defined as

$$A \cap B = \{ x \mid x \in A \land x \in B \}$$

Example 3: If $A = \{1, 2, 3\}$ and $B = \{1, 2, 4, 5\}$, then $A \cap B = \{1, 2\}$.

Example 4: If $A = \{1, 2, 3\}$ and $B = \{4, 5\}$, then $A \cap B = \emptyset$.

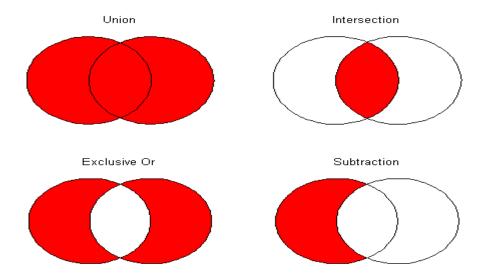
Definition (Difference): The difference of sets A from B, denoted by A - B, is the set defined as

$$A - B = \{ x \mid x \in A \land x \not\in B \}$$

Example 5: If $A = \{1, 2, 3\}$ and $B = \{1, 2, 4, 5\}$, then $A - B = \{3\}$.

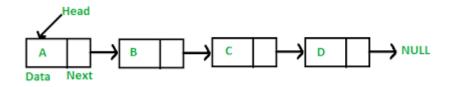
Example 6: If $A = \{1, 2, 3\}$ and $B = \{4, 5\}$, then $A - B = \{1, 2, 3\}$.

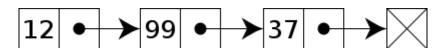
Note that in general $A - B \neq B - A$



Singly linked list

Singly linked lists contain nodes which have a data field as well as a 'next' field, which points to the next node in line of nodes. Operations that can be performed on singly linked lists include insertion, deletion and traversal.





A singly linked list whose nodes contain two fields: an integer value and a link to the next node

Union and Intersection of two Linked Lists

Input:

List1: 10->15->4->20

lsit2: 8->4->2->10

Output:

```
Intersection List: 4->10
```

Union List: 2->8->20->4->15->10

Intersection (list1, list2):

Initialize result list as NULL. Traverse list1 and look for its each element in list2, if the element is present in list2, then add the element to result.

Union (list1, list2):

Initialize result list as NULL. Traverse list1 and add all of its elements to the result. Traverse list2. If an element of list2 is already present in result then do not insert it to result, otherwise insert.

Input: enter set A of students like Vanilla Ice-cream and set B of students like butterscotch ice-cream.

Output: Intersection, Union, Neither nor.

Algorithm:

Set of students who like both vanilla and butterscotch

```
}
                     else
                            cur2=cur2->next;
              if(found==1)
              {
                     cout<<curl->data<<" ";
                     cur1=cur1->next;
                     found=0;
              }
              else
                     cur1=cur1->next;
              cur2=Vanilla;
       }
}
      Set of students who like only vanilla
void onlyV()
{
       Node *cur1=Butter;
       Node *cur2=Vanilla;
       int found=0;
       while(cur2!=NULL)
              while(cur1!=NULL)
                     if(cur2->data==cur1->data)
                            found=1;
                            break;
                     }
                     else
                            cur1=cur1->next;
```

```
}
              }
              if(found==1)
              {
                     cur2=cur2->next;
                     found=0;
              }
              else
              {
                     cout<<cur2->data<<" ";
                     cur2=cur2->next;
              cur1=Butter;
       }
}
      Set of students who like only butterscotch
void onlyB()
{
       Node *cur1=Butter;
       Node *cur2=Vanilla;
       int found=0;
       while(cur1!=NULL)
              while(cur2!=NULL)
                     if(cur1->data==cur2->data)
                            found=1;
                            break;
                     else\{
                            cur2=cur2->next;
```

• Set of students who like either vanilla or butterscotch or both

• Number of students who like neither vanilla nor butterscotch

```
void neither()
{
  cout<<"\nstudents who like neither vanila nor butterscotch\n";
  temp=h1;
  while(temp!=NULL)
  { temp3=head3;
   f=0;
  while(temp3!=NULL)</pre>
```

```
{ if(temp->roll==temp3->roll)
    { f=1;     }
    temp3=temp3->next;
}

if(f==0)
{ cout<<"\n"<<temp->roll; }
    temp=temp->next;
}
```

Software required: g++ / gcc compiler- / 64 bit fedora.

Conclusion: We understand & implement different operations on of linked list.

Assignment No- 09

Aim: To illustrate the concept of stack and string.

Problem Statement: A palindrome is a string of character that's the same forward and backward. Typically, punctuation, capitalization, and spaces are ignored. For example, ||Poor Dan is in a droop is a palindrome, as can be seen by examining the characters —poor danisina droop and observing that they are the same forward and backward. One way to check for a palindrome is to reverse the characters in the string and then compare with them

the original-

in a palindrome, the sequence will be identical. Write C++ program with functions-

1. To check whether given string is palindrome or not that uses a stack to determine

whether a string is a palindrome.

2. To remove spaces and punctuation in string, convert all the Characters to

lowercase, and then call above Palindrome checking function to check for a

palindrome

3. To print string in reverse order using stack

Learning Objectives:

To understand concept of stack and string.

To analyze the string is palindrome or not.

Learning Outcome: Students will be able to implement stack and queue data structures and

algorithms for solving different kinds of problems.

Theory:

Stack

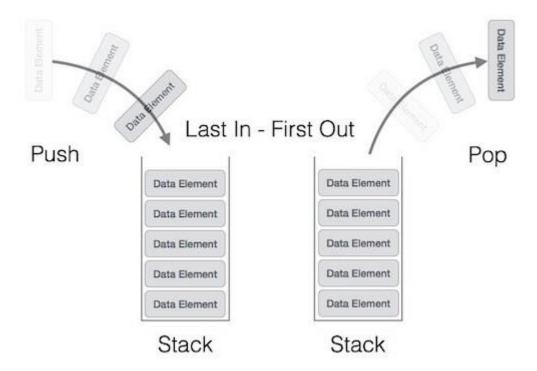
It is named stack as it behaves like a real-world stack, for example – deck of cards or pile of plates etc. A stack is an abstract data type that serves as a collection of elements, with two principal operations: push, which adds an element to the collection, and pop, which removes the most recently added element that was not yet removed. The order in which elements come off a stack gives rise to its alternative name, LIFO (for last in, first out). Additionally, a peek operation may give access to the top without modifying the stack.

A real-world stack allows operations at one end only. For example, we can place or remove a card or plate from top of the stack only. Likewise, Stack ADT allows all data operations at one end only. At any given time, We can only access the top element of a stack.

This feature makes it LIFO data structure. LIFO stands for Last-in-first-out. Here, the element which is placed (inserted or added) last, is accessed first. In stack terminology, insertion operation is called PUSH operation and removal operation is called POP operation.

Stack Representation

Below given diagram tries to depict a stack and its operations –



A stack can be implemented by means of Array, Structure, Pointer and Linked-List. Stack can either be a fixed size one or it may have a sense of dynamic resizing. Here, we are going to implement stack using arrays which makes it a fixed size stack implementation. Basic Operations

Stack operations may involve initializing the stack, using it and then de-initializing it. Apart from these basic stuffs, a stack is used for the following two primary operations –

- **push()** pushing (storing) an element on the stack.
- **pop()** removing (accessing) an element from the stack.

When data is PUSHed onto stack.

To use a stack efficiently we need to check status of stack as well. For the same purpose, the following functionality is added to stacks –

- **peek**() get the top data element of the stack, without removing it.
- **isFull()** check if stack is full.
- **isEmpty**() check if stack is empty.

At all times, we maintain a pointer to the last PUSHed data on the stack. As this pointer always represents the top of the stack, hence named **top**. The **top**pointer provides top value of the stack without actually removing it.

Input:

Enter the string.

Output:

Entered string is palindrome or not

Algorithm:

Step 7: End

```
Step 1: Input S (string)
Step 2: Len = 0, Flag = 0
Step 3: While (S[Len] != NULL)
              Len++
Step 4: I = 0, J = Len-1
Step 5: While (I < (Len/2)+1)
              If (S[I] == S[J])
                    Flag=0
              else
                    Flag=1
              I++ , J-
Step 6:
          If (Flag == 0)
                   Print Key Is a Palindrome
            else
                   Print Key Is Not a Palindrome
```

	Data Structure	Laboratory M	anual (210246)	
Software rec	uired: g++/gcc compi	ler- / 64 bit fedora	ì.	
Conclusion :	Thus we have studied the	he implementation	n of string is palindro	me or not

Assignment No- 10

Aim: To illustrate the various concept of stack.

Problem Statement: In any language program mostly syntax error occurs due to unbalancing

delimiter such as (),{},[]. Write C++ program using stack to check whether given expression

is well parenthesized or not.

Learning Objectives:

To understand concept of stack.

To analyze the working of various functions of stack.

Learning Outcome: Students will be able to implement stack and queue data structures and

algorithms for solving different kinds of problems.

Theory:

Stack

A stack is a container of objects that are inserted and removed according to the last-in

first-out (LIFO) principle.

In the pushdown stacks only two operations are allowed:

push the item into the stack, and

pop the item out of the stack.

A stack is a limited access data structure - elements can be added and removed from

the stack only at the top. push adds an item to the top of the stack, pop removes the item

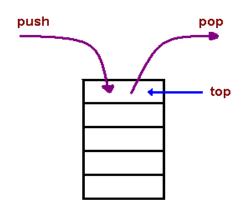
from the top. A helpful analogy is to think of a stack of books; you can remove only the top

book, also you can add a new book on the top.

A stack is a **recursive** data structure. Here is a structural definition of a Stack:

a stack is either empty or

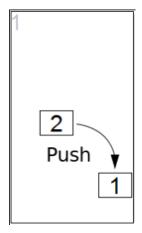
it consists of a top and the rest which is a stack;

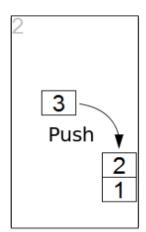


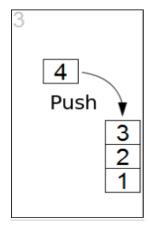
A **stack** is an abstract data type that serves as a collection of elements, with two principal operations: push, which adds an element to the collection, and pop, which removes the most recently added element that was not yet removed. The order in which elements come off a stack gives rise to its alternative name, **LIFO** (for **last in, first out**). Additionally, a peek operation may give access to the top without modifying the stack.

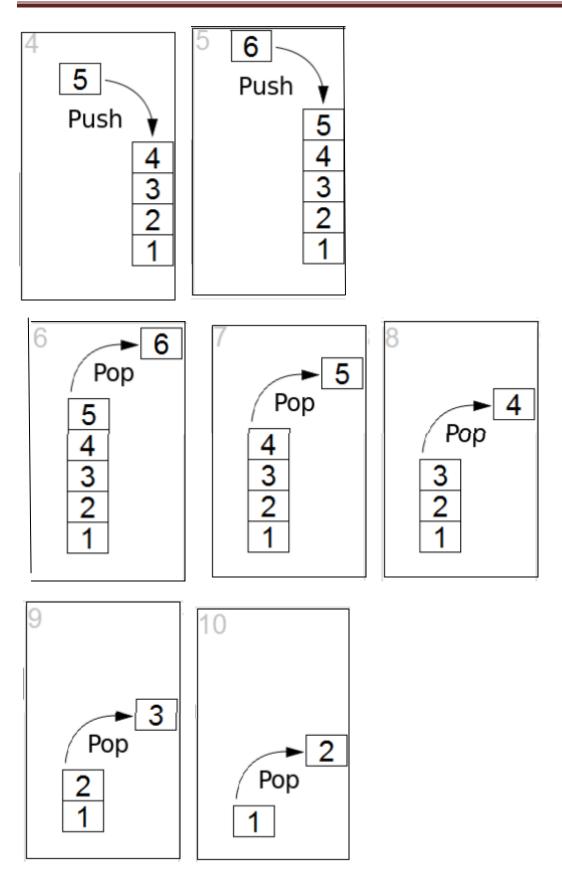
The name "stack" for this type of structure comes from the analogy to a set of physical items stacked on top of each other, which makes it easy to take an item off the top of the stack, while getting to an item deeper in the stack may require taking off multiple other items first.^[1]

Considered as a linear data structure, or more abstractly a sequential collection, the push and pop operations occur only at one end of the structure, referred to as the top of the stack. A stack may be implemented to have a bounded capacity. If the stack is full and does not contain enough space to accept an entity to be pushed, the stack is then considered to be in an overflow state. The pop operation removes an item from the top of the stack.









Balanced parentheses

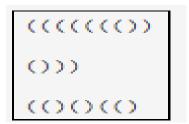
We now turn our attention to using stacks to solve real computer science problems. You have no doubt written arithmetic expressions such as

$$(5+6)*(7+8)/(4+3)(5+6)*(7+8)/(4+3)$$

where parentheses are used to order the performance of operations.

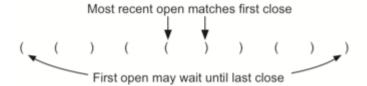
Balanced parentheses means that each opening symbol has a corresponding closing symbol and the pairs of parentheses are properly nested. Consider the following correctly balanced strings of parentheses:

Compare those with the following, which are not balanced:



The ability to differentiate between parentheses that are correctly balanced and those that are unbalanced is an important part of recognizing many programming language structures.

The challenge then is to write an algorithm that will read a string of parentheses from left to right and decide whether the symbols are balanced. To solve this problem we need to make an important observation. As you process symbols from left to right, the most recent opening parenthesis must match the next closing symbol . Also, the first opening symbol processed may have to wait until the very last symbol for its match. Closing symbols match opening symbols in the reverse order of their appearance; they match from the inside out. This is a clue that stacks can be used to solve the problem.



Input:

Enter any expression having number of opening & closing brackets.

Output:

Entered expression is well parenthesized or not.

Algorithm:

```
Step 1: Declare a character stack S.

Step 2: Now traverse the expression string exp.

If the current character is a starting bracket ('(' or '{' or '[')}) then

push it to stack.

If the current character is a closing bracket (')' or '}' or ']')

then

pop from stack and

if the popped character is the matching starting bracket then

fine

else
```

Step 4: After complete traversal, if there is some starting bracket left in stack then "not balanced"

parenthesis are not balanced

Step 5:Exit

Software required: g++ / gcc compiler- / 64 bit fedora.

Conclusion: Thus we have studied the implementation of expression is well parenthesized or not using stack.

Assignment No - 11

Aim: To illustrate the concept of queue.

Problem Statement: Queues are frequently used in computer programming, and a typical example is the creation of a job queue by an operating system. If the operating system does not use priorities, then the jobs are processed in the order they enter the system. Write C++ program for simulating job queue. Write functions to add job and delete job from queue.

Learning Objectives:

To understand concept of queue.

To analyze the various functions of queue.

Learning Outcome: Students will be able to implement stack and queue data structures and algorithms for solving different kinds of problems.

Theory:

A queue is logically a first in first out (FIFO or first come first serve) linear data structure. The concept of queue can be understood by our real life problems. For example a customer come and join in a queue to take the train ticket at the end (rear) and the ticket is issued from the front end of queue. That is, the customer who arrived first will receive the ticket first. It means the customers are serviced in the order in which they arrive at the service centre. It is a homogeneous collection of elements in which new elements are added at one end called rear, and the existing elements are deleted from other end called front. The basic operations that can be performed on queue are A minimal set of operations on a queue is as follows:

1. create()—creates an empty queue, Q

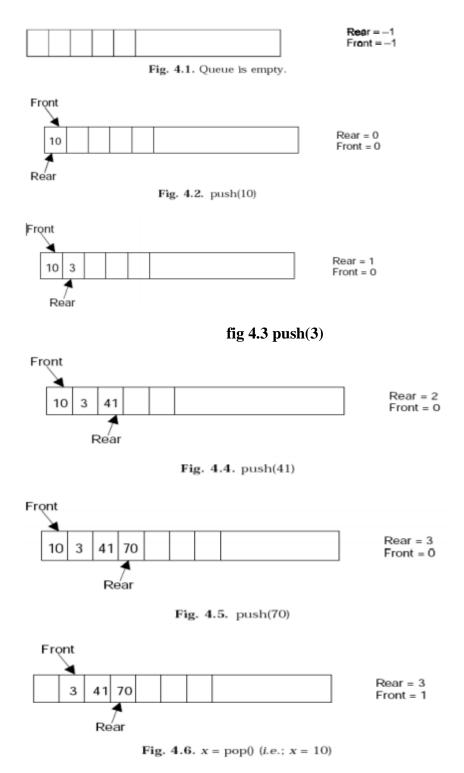
2. add(i,Q)—adds the element i to the rear end of the queue, Q and returns the new queue

3. delete(Q)—takes out an element from the front end of the queue and returns the resulting queue

4. getFront(Q)—returns the element that is at the front position of the queue

5. Is_Empty(Q)—returns true if the queue is empty; otherwise returns false

Push operation will insert (or add) an element to queue, at the rear end, by incrementing the array index. Pop operation will delete (or remove) from the front end by decrementing the array index and will assign the deleted value to a variable. Total number of elements present in the queue is front-rear+1, when implemented using arrays. Following figure will illustrate the basic operations on queue.



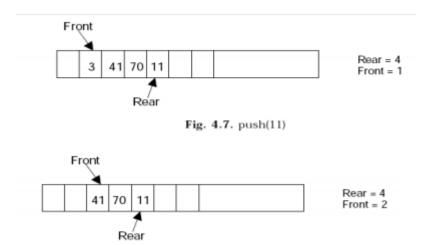


Fig. 4.8. x = pop() (i.e.; x = 3)

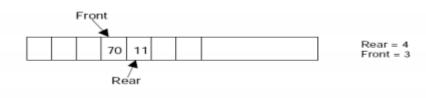
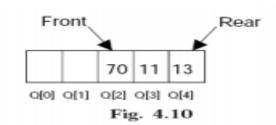


Fig. 4.9. x = pop() (i.e., x = 41)



Input: Enter the jobs to an operating system

Output: Add jobs and delete jobs from queue.

Algorithm:

• Algorithm to define class

Step 1: class queue(element)

Step 2: declare create() A queue

Step 3: add(element, queue) A queue

Step 4: delete(queue) A queue

Step 6: getFront(queue) A queue

Step 7: Is_Empty(queue) A Boolean;

Step 8:For all Q belongs to queue, i belongs to element let

```
Step 9:Is_Empty(create()) = true
Step 10:Is_Empty(add(i,Q)) = false
Step 11:delete(create()) = error
Step 12:delete(add(i,Q)) =
              if Is_Empty(Q)
              then
                      create
              else
                      add(i, delete(Q))
Step 13:getFront(create) = error
Step 14:getFront(add(i, Q)) =
              if Is_Empty(Q)
              then
                      i
              else
                      getFront(Q)
                  end
          Step 15:End
```

• Algorithm to create an empty queue

```
Step 1: Initialize queue

Front = Rear = -1;

Step 2: End
```

• Algorithm to checks whether the queue is empty or not

```
Step 1: Is_Empty()
Step 2: if(Front == Rear)
return 1;
else
return 0;
End
Step 3: End
```

• Algorithm to checks whether queue is full or not.

```
Step 1:Is_Full()
tep 2: if(Rear == max - 1)
              return 1;
         else
              return 0;
     End
Step 3: End
       Algorithm to adds an element in the queue
Step 1: Add(int Element)
Step 2:if(Is_Full())
      then
              print Error, Queue is full
         else
              Queue[++Rear] = Element;
         End
Step 3: End
       Algorithm to deletes an element from the front of the queue
Step 1: Delete()
Step 2:if(Is_Empty())
         then
              print Sorry, queue is Empty
      else
              return(Queue[++Front]);
         End
Step 3: End
       Algorithm to returns the element at the front
Step 1:getFront()
Step 2:if(Is_Empty())
          then
              print Sorry, queue is Empty
          else
              return(Queue[Front + 1])
```

End

Step 3: End

Software required: g++ / gcc compiler- / 64 bit fedora.

Conclusion : Thus we have studied the implementation of queue.

Assignment No-12

Aim: To illustrate the concept of double-ended queue (deque).

Problem Statement: A double-ended queue (deque) is a linear list in which additions and deletions may be made at either end. Obtain a data representation mapping a deque into a one- dimensional array. Write C++ program to simulate deque with functions to add and delete elements from either end of the deque.

Learning Objectives:

To understand concept of double-ended queue (deque).

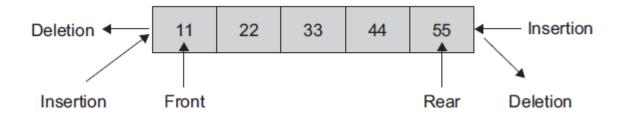
To analyze the various functions of double-ended queue (deque).

Learning Outcome: Students will be able to implement stack and queue data structures and algorithms for solving different kinds of problems.

Theory:

DeQueue:

Dequeue is a data structure in which elements may be added to or deleted from the front or the rear. Like an ordinary queue, a double-ended queue is a data structure it supports the following operations: enq_front, enq_back, deq_front, deq_back, and empty. Dequeue can be behave like a queue by using only enq_front and deq_front, and behaves like a stack by using only enq_front and deq_rear. The DeQueue is represented as follows.



The output restricted Dequeue allows deletions from only one end and input restricted Dequeue allow insertions at only one end.

Input : Enter the elements in dequeue.

Output: Add elements from either end, delete elements from both ends.

Algorithm:

Algorithm to add an element into DeQueue:

Assumptions: pointer f,r and initial values are -1,-1 Q[] is an array max represent the size of a queue

Algorithm to add an element from front :

```
step1. Start step2. Check the queue is full or not as if (f <> step3. If false update the pointer f as f= f-1 step4. Insert the element at pointer f as Q[f] = element step5. Stop
```

Algorithm to add an element from back :

```
step1. Start step2. Check the queue is full or not as if (r == max-1) if yes queue is full step3. If false update the pointer r as r = r+1 step4. Insert the element at pointer r as Q[r] = element step5. Stop
```

Algorithm to delete an element from the DeQueue

• Algorithm to delete an element from front :

```
step1. Start step2. Check the queue is empty or not as if (f == r) if yes queue is empty step3. If false update pointer f as f = f+1 and delete element at position f as element = Q[f] step4. If (f == r) reset pointer f and r as f == r-1 step5. Stop
```

Algorithm to delete an element from back :

```
step1. Start step2. Check the queue is empty or not as if (f == r) if yes queue is empty step3. If false delete element at position r as element = Q[r] step4. Update pointer r as r = r-1 step5. If (f == r) reset pointer f and r as f = r= -1 step6. Stop
```

	Data Structure Laboratory I	- Tantaar (210210)
Software	required: g++ / gcc compiler.	
Conclusio	1: Thus, we have studied the implementation	on of double-ended queue(deque).

Assignment No - 13

Aim: To illustrate the concept of circular queue.

Problem Statement: Pizza parlor accepting maximum M orders. Orders are served in first come first served basis. Order once placed cannot be cancelled. Write C++ program to simulate the system using circular queue using array.

Learning Objectives:

To understand concept of circular queue.

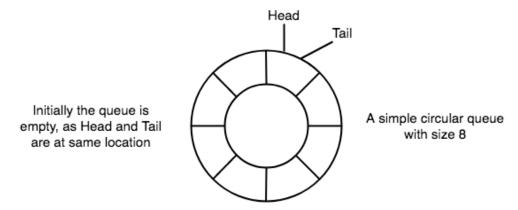
To analyze the various functions of circular queue.

Learning Outcome: Students will be able to implement stack and queue data structures and algorithms for solving different kinds of problems.

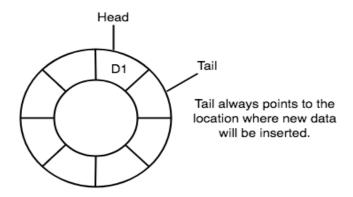
Theory:

Circular Queue:

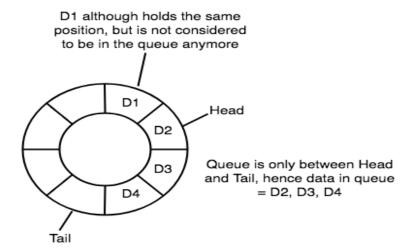
- 1. In case of a circular queue, head pointer will always point to the front of the queue, and tail pointer will always point to the end of the queue.
- 2. Initially, the head and the tail pointers will be pointing to the same location, this would mean that the queue is empty.



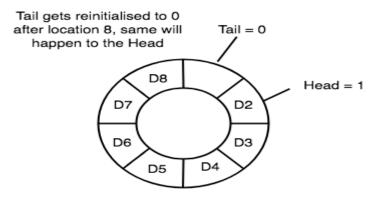
3. New data is always added to the location pointed by the tail pointer, and once the data is added, tail pointer is incremented to point to the next available location.



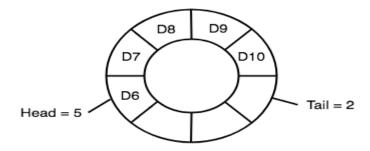
4. In a circular queue, data is not actually removed from the queue. Only the head pointer is incremented by one position when dequeue is executed. As the queue data is only the data between head and tail, hence the data left outside is not a part of the queue anymore, hence removed.



5. The head and the tail pointer will get reinitialized to 0 every time they reach the end of the queue



6. Also, the head and the tail pointers can cross each other. In other words, head pointer can be greater than the tail. Sounds odd? This will happen when we dequeue the queue a couple of times and the tail pointer gets reinitialized upon reaching the end of the queue.



In such a situation the value of the Head pointer will be greater than the Tail pointer

Input: Enter the orders of Pizza Parlor.

Output: Add orders and serve orders of pizza.

Algorithm: Implementation of circular queue

1. Initialize the queue, with size of the queue defined (maxSize), and head and tail pointers.

- 2. enqueue: Check if the number of elements is equal to maxSize 1:
 - o If Yes, then return Queue is full.
- o If No, then add the new data element to the location of tail pointer and increment the tail pointer.
- 3. dequeue: Check if the number of elements in the queue is zero:
 - o If Yes, then return Queue is empty.
 - o If No, then increment the head pointer.
- 4. Finding the size:
 - o If, tail >= head, size = (tail head) + 1
 - o But if, head > tail, then size = maxSize (head tail) + 1

Software required: g++ / gcc compiler- / 64 bit fedora.

Conclusion : Thus we have studied the implementation of circular queue.