

AMRITSAR GROUP OF COLLEGES  
(AUTONOMOUS COLLEGE)  
DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING

B. Tech. (CSE-I) 4<sup>th</sup> SEM  
OPERATING SYSTEM

ACCS - 16402

Assignment-1

Total Marks: 24

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Section-A

(6\*2=12)

Q1.

a) Differentiate between multiprogramming and multiprocessing? [CO5]

Ans: Multiprogramming:- Multiprogramming is more than one process running at a time it increases CPU utilization by organizing job so that the CPU always has one to execute. To motive is keep multiple jobs in main memory.

Multiprocessing:- The availability of more than one processor per system, that can execute several set of instructions in parallel is known as multiprocessing.

b) What is process synchronization? Discuss about critical section problem. [CO9]

Ans: Process synchronization:- Is the task of coordinating the ~~executing~~ execution of process in a way that no two process can have to the same shared data and resources it is specially in a multi-process system when multiple processes are running together.

Critical section problem:- A critical section is a segment of code which can be accessed by a single process at a specific point of time. The section consists of shared data resources that required be accessed by other process.



c) Differentiate between preemptive and non-preemptive scheduling? [CO2]

Ans: Preemptive scheduling: - (i) The CPU allocated to a process for a limited time. (ii) process can be interrupted in between.

(iii) used in time sharing system

(iv) CPU utilization is higher than non-preemptive scheduling

Non-Preemptive scheduling: - The CPU allocated to a process unit it terminate or switches to waiting state. A process can not be interpreted until its terminate or switches to waiting state.

d) What do you mean by process control block? What information does it contain? [CO1]

Ans: A process control block (PCB) is a data structure used by computer operating systems to store all the information about a process. It is also known as a process describe to when a process is created (initialized) or installed. The operating information contains in this process i.e. registers Quantum, priority, etc. In this process control stores many data items that are needed for efficient process management.

e) Describe Round Robin with example. [CO2]

Ans: In this scheduling Time Quantum is given and we have to make ready queue of running queue. It is a preempting scheduling which every process get execute in a cyclic way.

Example:-

Process	AT	BT	CT	TAT	WT
P <sub>1</sub>	0	5	12	12	7
P <sub>2</sub>	1	4	11	10	6
P <sub>3</sub>	2	2	6	4	2
P <sub>4</sub>	4	1	9	3	2

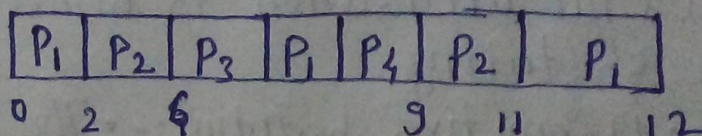
$$Avg = 7.5 \quad wt = 4.7$$



Ready Queue



Runing Queue





f) Define Turnaround time, Waiting time, Response time. [CO2]

Ans:

→ Turnaround time:- {completion time - arrival time}

→ Waiting time:- {Turnaround time - Burst time}

Response time:- {The time at which a process get cpu first time} - (arrival time)

### Section-B

(3\*4=12)

Q2. Explain Producer/Consumer problem with example in process synchronization. [CO9]

Ans: The producer consumer problem is a synchronization problem. There is a fixed size buffer and the producer produces items and enters them into buffer and consumers them. A producer should not produce item into the buffer when the consumer is consuming an item from the buffer and vice versa. So buffer should not be accessed by the producer or consumer at a time. The producer consumer problem can be resolved using semaphores. Semaphore is an atomic action wait/signal and wait/signal that are used for process synchronization.

```
int s = 1;

wait() { s--; }
signal() { s++; }

void producer()
{
    while(s <= 0)
        s = s - 1;
}

void consumer()
{
    while(True)
    {
        s = s + 1;
    }
}
```



```

{
  producer()
  wait(s)
  wait(s)
  append(x)
  signal(s)
  signal(f)
}
}

```

```

void consumer()
{
  while (True)
  {
    wait(f)
    wait(s)
    take(x)
    signal(s)
    signal(s)
    signal(e)
    use()
  }
}

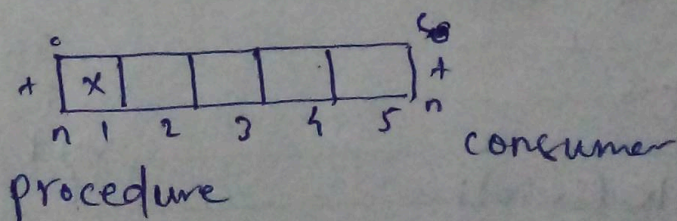
```

$S = 1 \ 0 \ 1 \ 0 \ 1 \ 0 \ 1$

$E = 0 \ 1 \ 0 \ 1 \ 0 \ 1$

$f = 0 \ 1 \ 0 \ 1 \ 0$

(empty)  
(full)



procedure



Q3. What do you mean by SJF? Explain with example (Preemptive and non-preemptive). [CO2]

Ans: Shortest Job first (SJF) is an algorithm in which the process having the smallest execution time is chosen for execution. On a non-preemptive, it significantly reduces the average waiting time for other processes awaiting execution. There are basically two types of SJF method.

- preemptive sjf
- Non-preemptive sjf

Process	Arrival Time	BT	CT	TAT	WT
P <sub>1</sub>	0	6	5	11	5
P <sub>2</sub>	1	8	7	24	15
P <sub>3</sub>	2	7	6	5	3
P <sub>4</sub>	3	3	2	6	3
				$\Sigma = 1.25$	$\Sigma = 5.7$

Grant chart

$P_1$	$P_2$	$P_3$	$P_4$	$P_4$	$P_1$	$P_3$	$P_2$	
0	1	2	3	4	6	11	17	24

Q4. What is deadlock? State necessary conditions for deadlock characterization. [CO3]

Ans: A deadlock is a condition that may happen in a system composed of multiple processes that can access shared resources. A deadlock is said to occur when two or more processes are waiting for each other to release a resource. None of the processes can make any progress. There are four necessary conditions for the occurrence of a deadlock.

Mutual exclusion: - When two people meet in the hallway, they can't just walk through because there is



space only for one person. This condition to allow only one person (or process) to use the step between them or the resource is the first condition necessary for the occurrence of the deadlock.

# Hold and wait: — When the 2 people refuses and hold their grounds it is called holding. This is the next necessary condition for the deadlock.

# No preemption: — for resolving the deadlock, one can simply cancel on of the process for other to continue. But operating system doesn't do so. It all allocates the resources to the processors for as much time added until the task is completed. Hence, there is no temporary reallocation of the resources.

# Circular wait: — When the two people refuses to retreat and for each other to retreat, so that they can complete their task, it is called circular wait. It is the last condition for the deadlock to occur.