

Normalization

⊗ What is Normalization?

→ is a database design technique that reduce data redundancy and eliminates undesirable characteristics like (insertion, update delete Anomalies)

1NF (First Normal form)

↙ cant have multiple value in a column attribute

① a table cell should contain a single value

X

101, 102

✓

101
102

↙ Column name unique

② Each Records need to be unique

X

ID	ID

✓

ID	Roll

ID	Name	course
1	M	NET
2	G	JAVA
3	S	C, JAVA

ID	Name	course
1	M	NET
2	G	JAVA
3	S	C
3	S	JAVA

2NF → Second Normal Form

① Table should be in 1NF form

② Partial dependency is not allowed

cannot identify uniquely

PK	NON key	
	ID	address
	1	Dhaka
	2	Banishal
	3	Dhaka

↓
key

Functional dependency
Dependency

all non key column
dependent on key
column (Primary)

↓
Uniquely
Identify
able

... Hence Name, address
dependent on ID
to uniquely Identify
them

Dependent on ID
not on course ID

Dependent on course ID
not on ID

ID	Name	Address	course ID	course Name	credit	Grade
101	A	M	1	C	3	A+
101	A	M	2	JAVA	3	A
102	B	S	1	C	3	B
102	B	S	2	JAVA	3	B+

PK (ID) → PK (course ID)
 Key: ID, course ID; Non Key: Name, Address, credit, Grade

① Candidate key

(multiple primary key
in a table to
uniquely identify)

Here table
need 2 PK
for identifying
a record
uniquely

Partial Dependency

if a table
has candidate key
(multiple primary key)

and any non key value is dependent only
on one of the any primary key

∴ Hence

course_name and credit is non key
and dependent only on course ID

To apply, 2NF Remove partial Dependency
(can have fully on functional dependency)

course ID	course Name	credit
1	C	3
2	Java	3

PK (under course ID)
key (under course ID)
NON key (under course Name)

Functional dependency

2NF → ✓

ID	name	address
1	A	m
2	B	s

PK (under ID)

Functional dependency

2NF → ✓

ID	course ID	Grade
101	1	A+
102	1	B
101	2	B
102	2	A

PK (under ID)
PK (under course ID)

Non key / Grade
dependent on both
primary key

partial dependency → X

2NF → ✓

3NF

- ① Table must follow 2NF
- ② Transitive Dependency not allowed

PK

Transitive

Non key

Course Name	Teacher ID	Teacher Name	Credit
C	1	Fahim	3
OS	2	Sayma	3
DBMS	3	Ahsan	3
C-LAB	1	fahim	1.5

key

Transitive dependency

if a non key column of table is dependent on another non key column of the table

∴ Here Teacher Name is dependent on Teacher i.d which is also a non-key column

~~PK~~

COURSE NAME	Teacher ID	credit
C	1	3
OS	2	3
DBMS	3	3
C-LAB	1	1.5

Transitive dependency $\rightarrow X$

3NF $\rightarrow \checkmark$

~~PK~~

Teacher ID	Teacher Name
1	Fahim
2	Sajma
3	Ahsan

Transitive dependency $\rightarrow X$

3NF $\rightarrow \checkmark$

Boyer-codd

3SNF

\rightarrow

BCNF

\rightarrow

Key column, nonkey
trans dependent
2nd 71

4NF

\rightarrow

no multivalued dependency

Q Fix 3NF

Transitive dependency

member (membershipID, fullname, physical address, saturationID)
 movienented (membershipID, movies nented)

membershipID → saturationID
 NON key → NON key

①

membershipID	Fullname	physical address	saturationID

②

saturationID	Saturation

Q Fix 2NF

member (Fullname, physical address, movie rented, saturation)

①

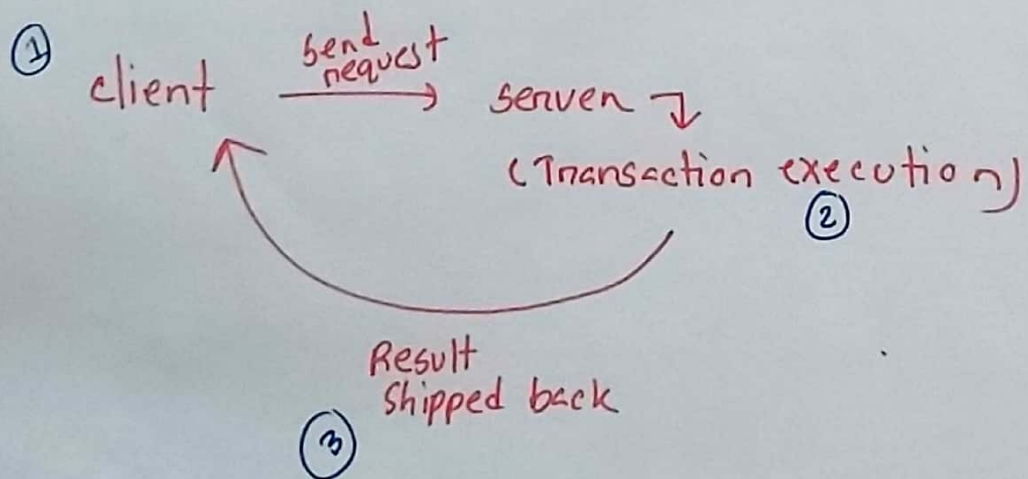
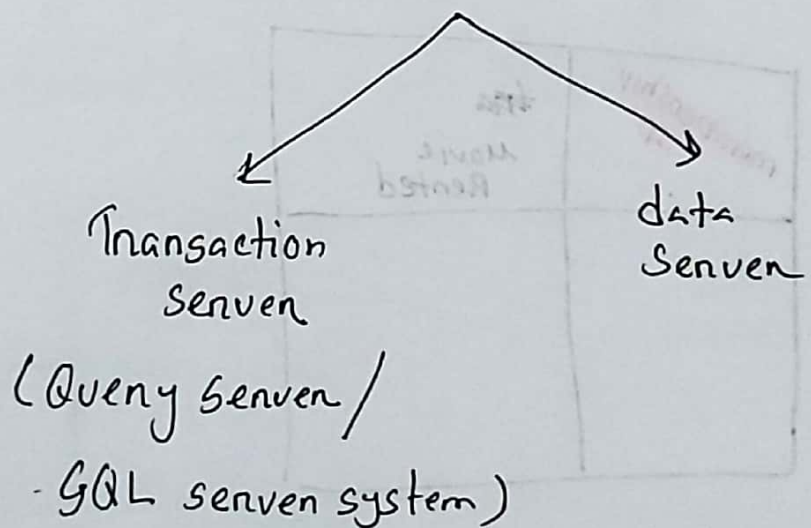
membership ID	Full name	Physical address	saturation

②

membership ID	Full name Movie Rented

Database System Architecture

Centralized Data system



* A typical transaction server consists of multiple process accessing data in shared memory

Transaction process

1. Database writer process
2. Log writer process
3. Check point process
4. Process Monitor Process
5. Lock Manager Process

Shared data

1. Buffer pool
2. Lock table
3. Log buffer
4. Cached Query plan

Two main performance monitors?

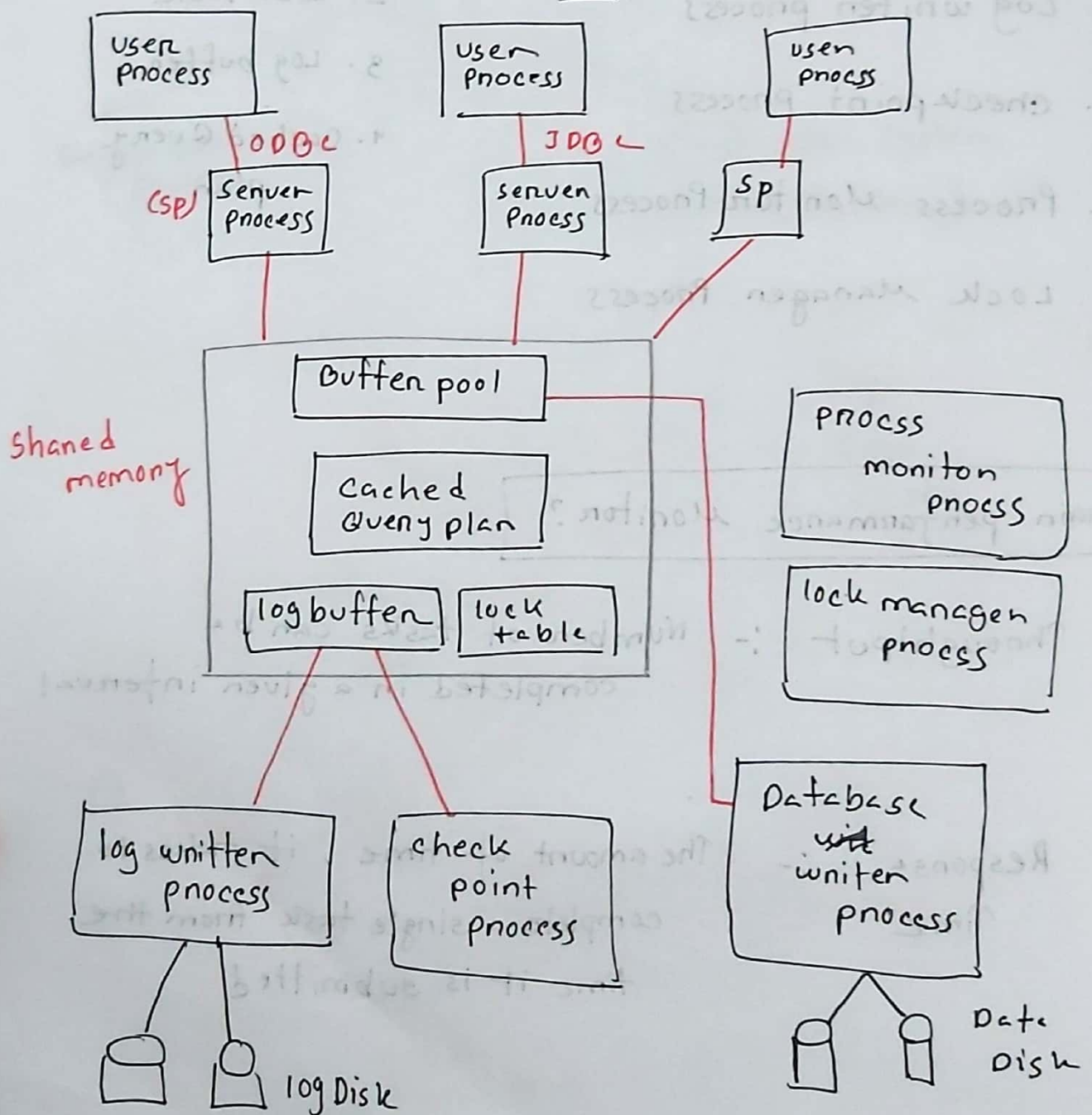
① Throughput :- Number of tasks can be completed in a given interval

② Response Time :- The amount of time, it takes to complete a single task from the time it is submitted

What is transaction?

→ Transaction is a logical unit of work that consists of one or more database operations (reading, writing or modifying data) stored in database.

State Diagram



4 properties of transaction

- ① Atomicity → Test and set
→ compare and swap
- ② Consistency
- ③ Isolation
- ④ Durability

Data Server system

→ is a software application on hardware infrastructure designed to manage and store data efficiently

Data caching

→ is a technique to improve performance of data retrieval operations by storing frequently accessed data in cache

Lock caching

→ is a technique to improve efficiency of Lock operations

① what is parallel system? (PS)

→ ps refers to database architecture that utilizes multiple processor / computing resources to process data operations co-currently

Motivation

→ improve performance by dividing the workload into smaller tasks that can be executed simultaneously on multiple processor

Benefits

① High performance

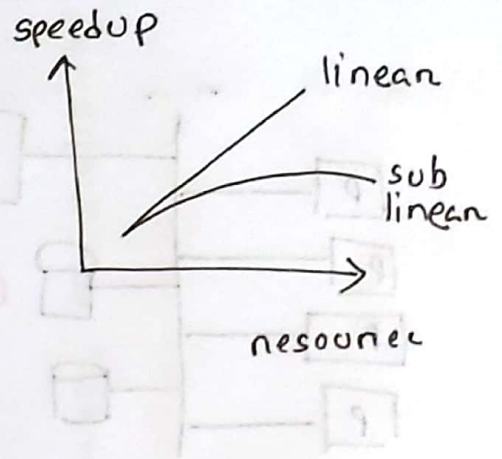
② Decision support

Types

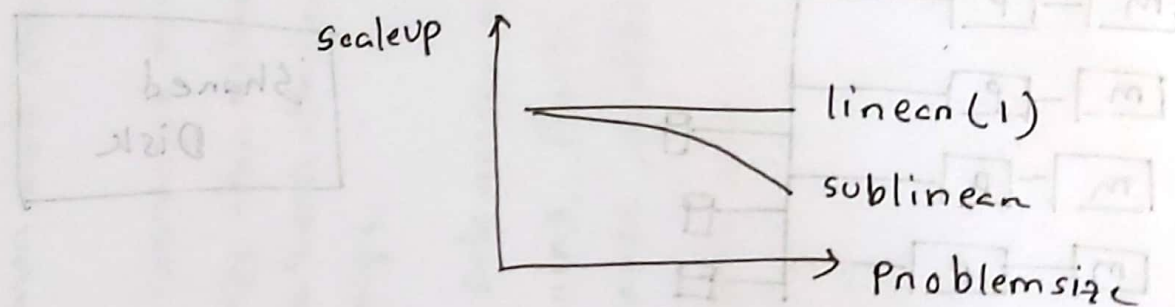
① Coarse grain parallel

② ~~fine~~ gain parallel
fine

$$\text{Speedup} = \frac{\text{small system elapsed time}}{\text{large system elapsed time}}$$



$$\text{Scale Up} = \frac{T_s \text{ Small system small problem elapsed time}}{T_L \text{ large system large problem elapsed time}}$$



Factor Limiting Speedup and ~~scale~~ Scale Up

① Startup/Sequential costs

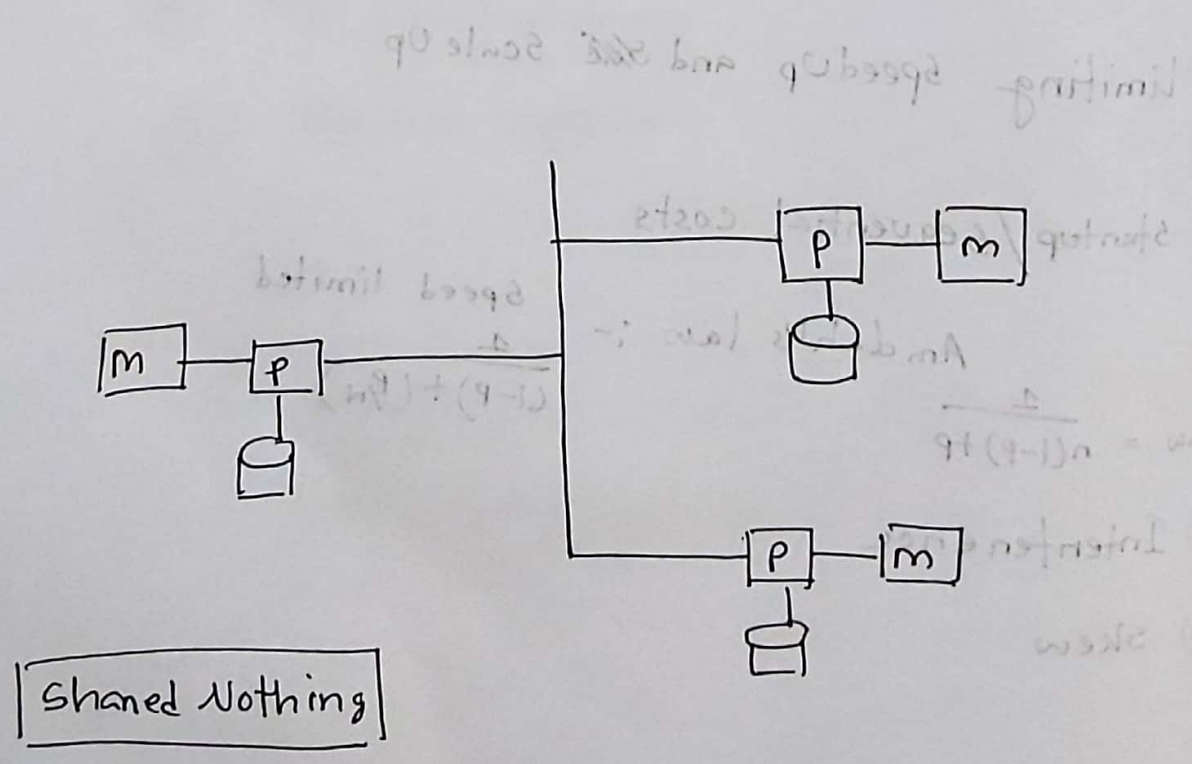
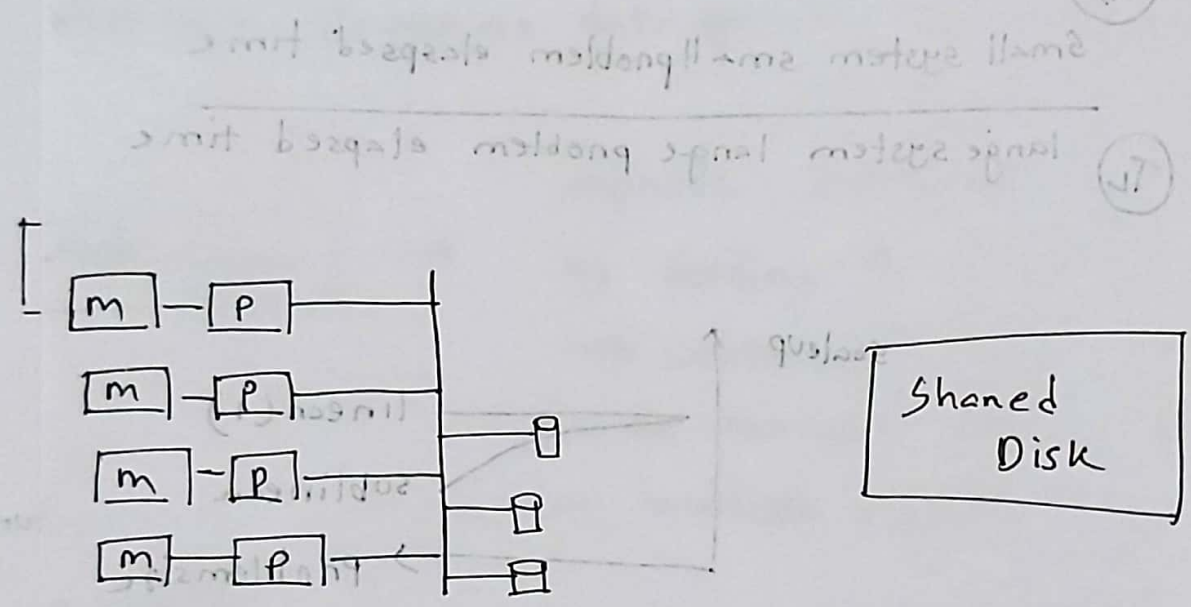
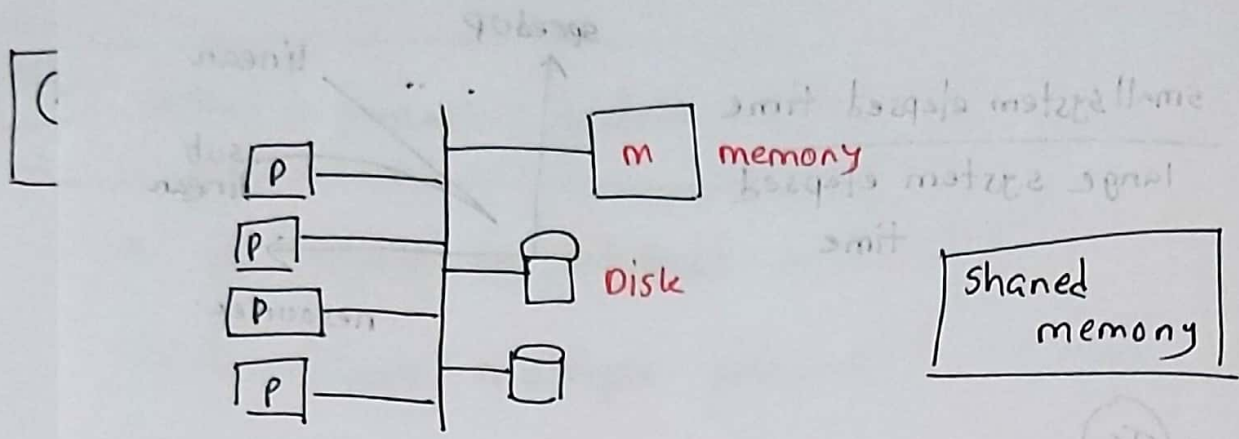
Amdahl's law :-

$$\text{Speed limited} = \frac{1}{(1-P) + (P/N)}$$

$$\text{Gustafson's law} = \frac{1}{n(1-P) + P}$$

② Intenfence

③ skew



shared memory

① multiple processors on nodes access a common physical memory space

② connected by bus or cross bar switch

③ key characteristics :-

communication and data consistency

— (not scalable)

— (Bottle neck issue)

shared Disk

① multiple processors on node have their own private memory but share access to common disk storage

② connected via storage area network or network attached storage

③ data access and sharing, fault tolerance, scalability

(Bottle neck)

shared Nothing

① each processor on node has its own private disk and memory

② do not share memory or storage directly

③ data partitioning, distribution and parallelism

fault tolerance, scalability
(costing high)