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RSA Cryptosystem:

a widely used public key cryptographic system that facilitates secure data transmission.

→ Key components:

1) public key : $\{n, e\}$

2) private key : $\{d, n\}$

Math Step:

1) Given, p, q (large prime numbers)

2) $n = p \times q$

3) $\phi(n) = (p-1)(q-1)$

4) choose e such that

a) $1 < e < \phi(n)$

b) $\gcd(e, \phi(n)) = 1$

5) find d , $d \times e \equiv 1 \pmod{\phi(n)}$

6) Encryption, $C = M^e \pmod{n}$ → use public key

7) Decryption, $M = C^d \pmod{n}$ → use private key

Example: Let, $p = 61, q = 53$

$$n = 61 \times 53 = 3233$$

$$\phi(n) = (61-1)(53-1) = 3120$$

Suppose, $e = 17$

$$\therefore \gcd(17, 3120) = 1$$

$$d = 17^{-1} \bmod 3233$$

$$= 2092$$

Let $m = 65$. $C = 65^{17} \bmod 3233 = \text{---} \textcircled{1}$

$$M = (65^{17})^{2092} \bmod 3233 = \text{---} \textcircled{2}$$

$$17 = 8 + 4 + 2 + 2 + 1$$

$$\textcircled{1} \rightarrow \therefore 65^{17} = (65^8 \cdot 65^4 \cdot 65^2 \cdot 65^2 \cdot 65^1) \bmod 3233$$

$$65^1 \bmod 3233$$

$$= 65$$

$$65^2 \bmod n$$

$$= 4225 \bmod n$$

$$= 992$$

$$65^4 \bmod n$$

$$= (992 \times 992) \bmod n$$

$$= 1232$$

$$65^8 \bmod n$$

$$= (2237) \bmod n$$

$$= 1547$$

$$= 2790$$

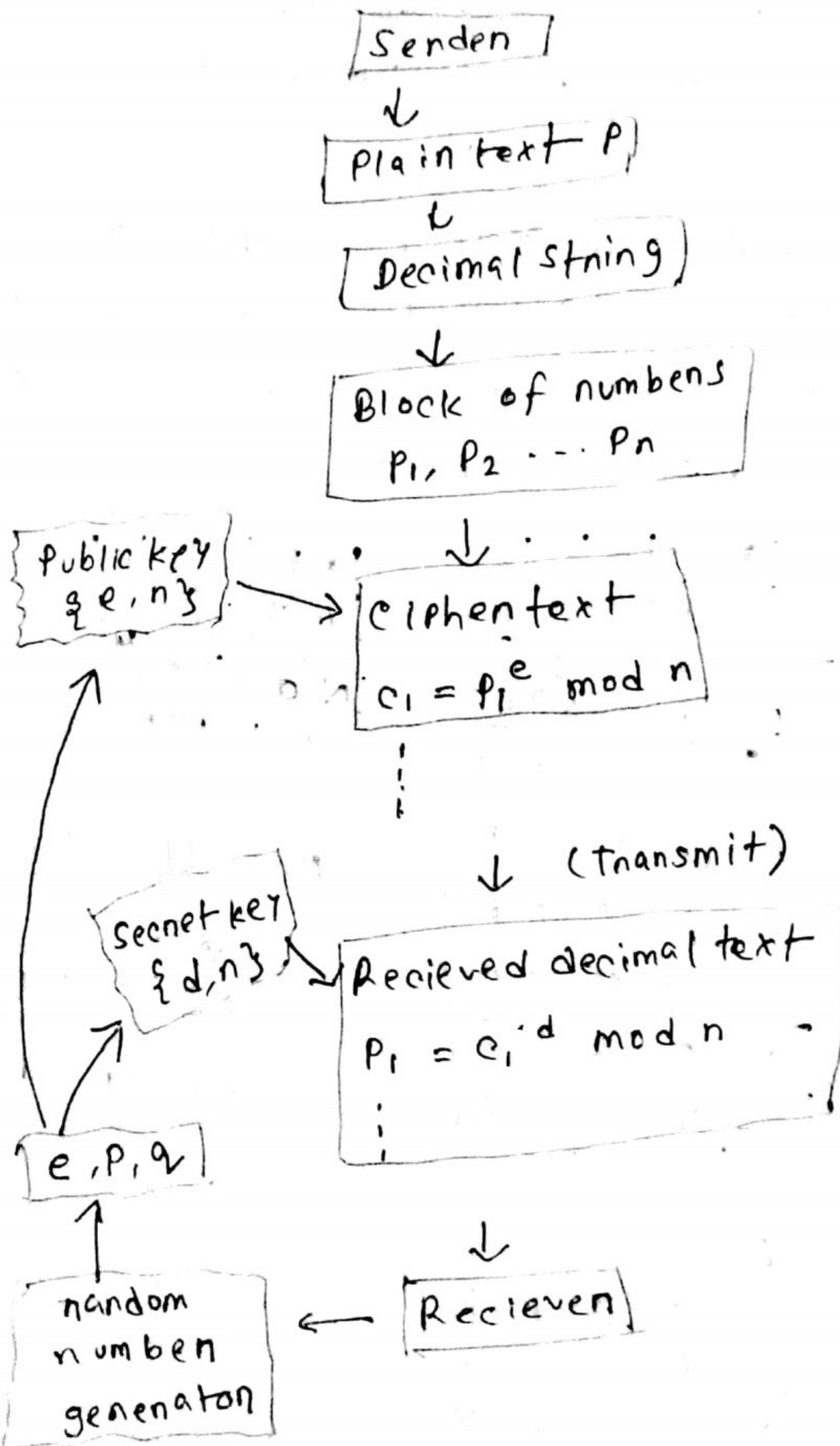
$$\therefore C = 2790$$

$$M' = 2790^{2092} \bmod 3233$$

$$= 2699$$

$$\therefore M \neq M'$$

* General pipeline:



* Playfair cipher:

- symmetric encryption
- digram substitution cipher

Ques
TT * Suppose a group of people agreed upon a key, $K = \text{"PUZZLE"}$. Now find the ciphertext of the plaintext, $M = \text{"LITTLE"}$ using playfair cipher.

→

$K =$

| | | | | |
|---|---|-----|---|---|
| P | U | Z | L | E |
| A | B | C | D | F |
| G | H | S/J | K | M |
| N | O | R | S | |
| T | V | W | X | Y |

$M = \underline{\text{LITTLE}}$

$= \text{LI TX TL EX}$
 ZK VY XP LY

* Diffie - Hellman key exchange protocol.

→ 2 person share a value on public channel $\{p, g\}$

↓
large prime
→ base

→ each person chooses secret number (a/b)

→ compute public keys value:

$$g^a \bmod p ; g^b \bmod p$$

→ exchange public value

→ computed shared secret:

$$(g^a)^b \bmod p = (g^b)^a \bmod p$$

Pos of IPsec (Internet Protocol Security)
→ Internet secure at IP layer

① provides strong encryption

and authentication to protect data

② Ensures data is not altered.

③ works well with existing net/app and protocols without modification

④ flexible and transparent

* Playfair: $K = \text{ASGARD}$

$$K = \begin{array}{c|c|c|c|c} A & S & G & R & D \\ \hline B & C & E & F & H \\ \hline I/J & K & L & M & N \\ \hline O & P & Q & T & U \\ \hline V & W & X & Y & Z \end{array}$$

$M = \text{GROOT}$
 $\underline{GR} \underline{OX} \underline{OT}$
 $RD \quad QV \quad PU$

* Elgamal Cryptosystem:

Given, $p = 7$, $g = 3$

(a) ① select private key. $1 \leq S_k \leq p-2$
 $1 \leq S_k \leq 7-2$
 $\therefore \boxed{S_k = 5}$

② $P_k = g^{S_k} \bmod p$
 $\gamma = 3^5 \bmod 7$
 $= 243 \bmod 7 = 5$

\therefore public key $(p, g, \gamma) = (7, 3, 5)$
 Ans: ✓ private key, $x = 5$

(b) $P_k : (p=7, g=3, y=5), M=13$

random $k : 1 \leq k \leq p-2 \dots k=4$

$$C_1 = g^k \bmod p = 3^4 \bmod 7 = 4$$

$$C_2 = M \cdot y^k \bmod p = 13 \cdot 5^4 \bmod 7 = 5$$

$$\therefore (C_1, C_2) = (4, 5)$$

* Why are RAIDs used?

→ Redundant Array of Independent Disks.

→ data is stored in multiple disks. If 1 disk fails, no data loss.

→ data reliability and availability through redundancy and fault tolerance

→ distribute data across multiple disks