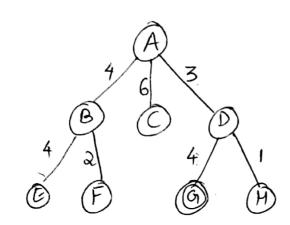
# Artificial Intelligence Assignment-2

SUNDESH RAJ 1001633297

(1)

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· Breadth first search:

$$(A \rightarrow (B \rightarrow C) \rightarrow (D \rightarrow (E) \rightarrow (E) \rightarrow (G)$$

1 Depth first search:

$$(\widehat{A} \rightarrow (\widehat{B} \rightarrow \widehat{E} \rightarrow \widehat{E} \rightarrow \widehat{E} \rightarrow \widehat{C} \rightarrow \widehat{D} \rightarrow \widehat{G})$$

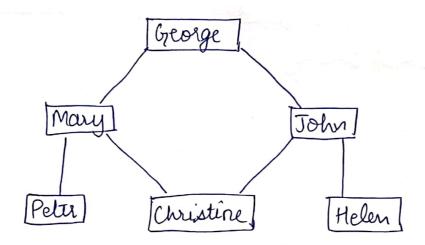
IDS: **(** 

Iteration 1: A Eteration 2: A,B,C,D Iteration 3: A,B,C,D,O,D

1 Uniform cost search:

iteration	Vinited Node	Fringe	Sorted List
0	None	Ao	A
1	A	B4 C6 D3	D3B4C6
2	A-> D	B4 C6 67 H4	B4 M4 C6 G17
3	A>D>B	H4C6G7E8 FG	H496 F6 G7 F8
4	A>D>B>H	G FG G17 158	C6 F6 G7 E8
5	A>D>B>H>X		F6 G7 E8
6	A-> D-> B-> H-> C->	F   G1 68	GT ES
7	A-> D->B>H->B->F->	G E8	Eg
			1

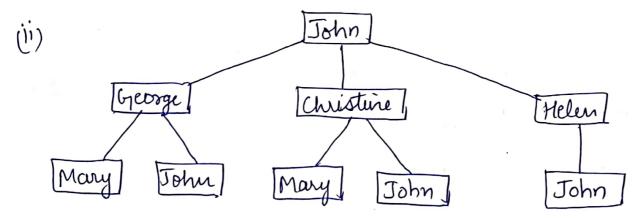
Result: A>D>B>H>C>F>G



(i) BFS and DFS will generate same results. Moreover, they will also provide the same degrees of freedom.

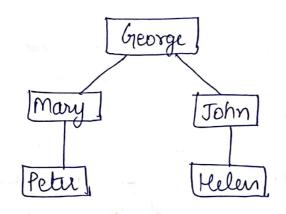
DFS can possibly not terminate even if it finds a solution, there is no confirmation that it will be the right answer.

IDS will provide the correct answer if nodes are just expanded in a DFS method and iterated.



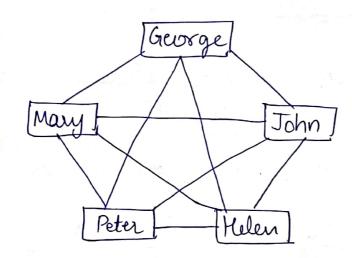
The vertex John in the graph corresponds to multiple vertices in the tree, so there is no one to one correspondence between the 10 2 nodes.

(iii)



Peter and Helen have 4 degrees of seperation

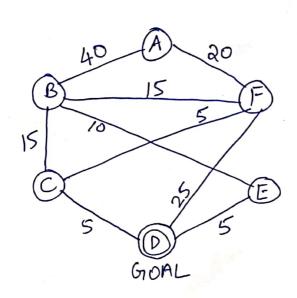
(iv)



Each puson having of l degree of seperation between them.

We we can achieve this by converting the tree search into a graph search or maintain a list of visited states. Do not generate successor nodes for nodes which are present in the visited states list.

Q3



#### Herristic 1

h(A)	=	50
h(B)	-	35
h(c)		
h(D)	=	0

h(F) = 10

## Henristic 2

h(F)= 70

There hewristic is not admissible because the actual cost for travelling from E to D is 5 but the heuristic is h(E) = 45 which is greater.

For heuristic 1 to be admissible the value for h(E) must be 5 or less than 5. New heuristic: h(A) = 50 h(d) = 0 h(B) = 15 h(C) = 5 h(B) = 10

This heuristic is not admissible because the values of h(D) is not equal

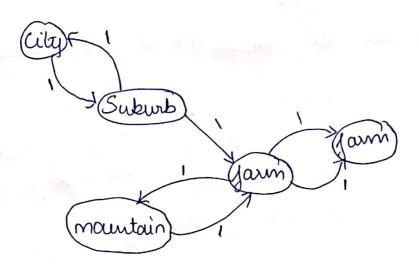
to 0 and h(E), H(F), h(c), h(B), h(A) are all greater than the actual value.

New heuristic: 
$$h(A) = 50$$
  $h(d) = 0$   
 $h(B) = 15$   $h(C) = 5$   
 $h(C) = 5$   $h(A) = 35$ 

### Henristic 3

Admissible heuristic since all the values are less than or equal to the actual values.

Smallest possible sequence of status that is consistent wither the rules is



From the above state diagram ve can obtain the below mentioned heuristic values.

$$h(m) = 0$$
  
 $h(1) = 1$   
 $h(3) = 2$   
 $h(c) = 3$ 

Guiren: b = 4 d = 101 to 208  $C^{*} = 101 \text{ to } 208$  E = 1  $m = \infty$ 

Space complexity:

BFS = [4 101+1 to 4208+1] × 1KB memory

DFS = 4 × 00 × 1 KB memory

UCS = \[ \( \frac{4}{4} \) \( \tau \) \( \frac{208}{1} \) \( \tau \) \( \tau

2000 x 0000 Place do 10 x ept e par monery

.. DFS = 00 KB of memory

BFS =  $2.57 \times 10^{61}$  KB to  $6.77 \times 10^{125}$  KB of memory UCS =  $6.43 \times 10^{60}$  KB to  $1.69 \times 10^{125}$  KB of memory

IDS = 404 to 832 KB of memory.

- a) None of above methods can run in less than 50 kB of memory
- b) IDS will run in less than 1200 kB of memory

## For figure 5:

- -> In some cases Greedy search will perform the same as At , example from 0,0 to 0,8
- -) However in some cases Greedy search will perform better than A\*, example from 0,0 to 8,8, A\* will visit both 0,1 and 1,0 but greedy will only visit one of them.
- -> Minu, Greedy will peyorm better than or the same as A\* depending on the start and end states.

- > Sometimes Gredy performs better than A\*, example 2,6 to 5,8.
- -> Sometimes Greedy performs worse than A\* example 2,0 to 2,2
- -> Sometimes they peyom the same example 2,2 to 2,7
- -> Hence, Greedy may perform better than or the same or worse than A\* despending on the start and end states.