

## Session Outcome

After completion of this session, you will be able to

- implement Linear data structures (Apply)



implement Array data structure (Apply)

# Array

**Need:** Storing marks of 20 students need 20 variables

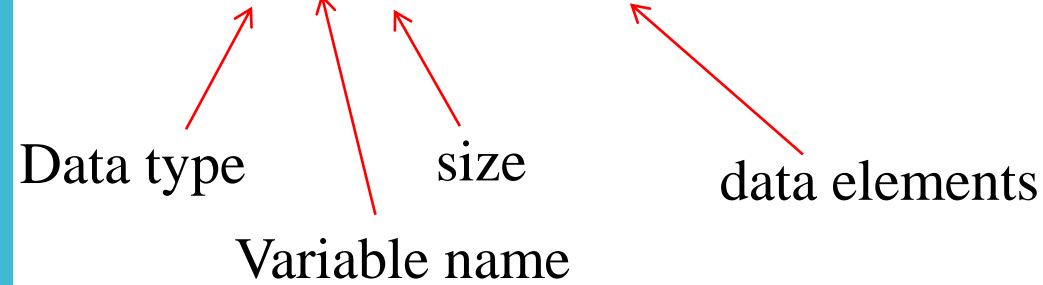
**Definition:** *An array is a fixed size, sequence collection of elements of same data type.*

# Declaration and storing Array Elements

1. `int a[3] = { 1, 2, 5 }`

[Compile Time]

Data type      Variable name      size      data elements



2. `int i, a[3];`  
`for (i=0 ; i < 3; i++)`  
`a[i] = i;`

[Run time]

3. `int i, a[3];`  
`for (i=0 ; i < 3 ; i++)`  
`scanf("%d",&a[i]);`

[Run time]

# Address of Array Elements

$$a[k] = \text{BaseAddress } [a] + w (k - \text{lowerbound})$$

Where  $k$ , is the index of element

$a$ , is the array variable name

$w$ , is the number of words per memory location

$\text{lowerbound}$ , is starting index

$\text{BaseAddress}$ , is the starting address of array or address of first element of array

e.g.

| index   | 0    | 1    | 2    | 3    | 4    | 5    | 6    | 7    |
|---------|------|------|------|------|------|------|------|------|
| a       | 10   | 20   | 30   | 40   | 50   | 60   | 70   | 80   |
| Address | 1000 | 1002 | 1004 | 1006 | 1008 | 1010 | 1012 | 1014 |

$$\begin{aligned} a[3] &= 1000 + 2 (3 - 0) \\ &= 1006. \end{aligned}$$

# Length of Array

$$\text{Length} = \text{upperbound} - \text{lowerbound} + 1$$

Where, upperbound is the last index of array  
lowerbound is the first index of array

e.g.

|          |      |      |      |      |      |      |      |      |
|----------|------|------|------|------|------|------|------|------|
| index    | 12   | 13   | 14   | 15   | 16   | 17   | 18   | 19   |
| <b>a</b> | 10   | 20   | 30   | 40   | 50   | 60   | 70   | 80   |
| Address  | 1000 | 1002 | 1004 | 1006 | 1008 | 1010 | 1012 | 1014 |

$$\text{Length} = 19 - 12 + 1 = 8.$$

Operations of  
Array  
{ **traversal**,  
insertion,  
deletion,  
merging,  
searching }

## 1. **Traversal**

A is the array

lowerbound is starting index

upperbound is last index

step 1: set  $i = \text{lowerbound}$

step 2: Repeat step 2 to step 4 while  $i \leq \text{upperbound}$

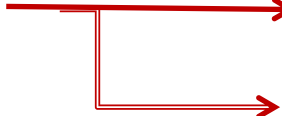
step 3: print (  $A[ i ]$  )

step 4: set  $i = i + 1$

step 5: Exit.

Operations of  
Array  
{ traversal,  
**insertion**,  
deletion,  
merging,  
searching }

## 2. Insertion

 I ) at the end of Array  
II ) middle of Array

### (I) Insertion at the end of Array

➡ **Simply add new element at end of array**

step 1:  $\text{upperbound} = \text{upperbound} + 1$

step 2:  $a[\text{upperbound}] = \text{new\_value}$

step 3: Exit.

### (II) Insertion at Middle of Array

Operations of  
Array  
{ traversal,  
**insertion**,  
deletion,  
merging,  
searching }

## (II) Insertion at Middle of Array

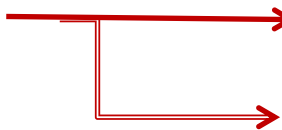
|                    |    |    |    |    |       |    |  |
|--------------------|----|----|----|----|-------|----|--|
| index              | 0  | 1  | 2  | 3  | 4     |    |  |
| a                  | 10 | 20 | 30 | 40 | 50    |    |  |
| Insert at a[1], 45 |    |    |    |    | POS=1 |    |  |
| index              | 0  | 1  | 2  | 3  | 4     |    |  |
| a                  | 10 | 20 | 30 | 40 | 50    |    |  |
| index              | 0  | 1  | 2  | 3  | 4     | 5  |  |
| a                  | 10 | 20 | 30 | 40 | 50    | 50 |  |
| index              | 0  | 1  | 2  | 3  | 4     | 5  |  |
| a                  | 10 | 20 | 30 | 40 | 40    | 50 |  |
| index              | 0  | 1  | 2  | 3  | 4     | 5  |  |
| a                  | 10 | 20 | 30 | 30 | 40    | 50 |  |
| index              | 0  | 1  | 2  | 3  | 4     | 5  |  |
| a                  | 10 | 20 | 20 | 30 | 40    | 50 |  |
| index              | 0  | 1  | 2  | 3  | 4     | 5  |  |
| a                  | 10 | 45 | 20 | 30 | 40    | 50 |  |



# Operations of Array

- { traversal,
- insertion,**
- deletion,
- merging,
- searching }

## 2. Insertion



- I ) at the end of Array
- II ) middle of Array

### (II) Insertion at Middle of Array

#### **insert (A, N, POS, new\_value)**

A, the array in which the element has to be inserted

N, number of elements in the array

POS, the position at which the element has to be inserted, and  
new\_value, the value has to be inserted

step 1: set  $i = N - 1$

step 2: Repeat step 2 to step 4 while  $i \geq \text{POS}$

step 3: set  $a[i+1] = a[i]$

step 4: set  $i = i - 1$

step 5: set  $N = N + 1$

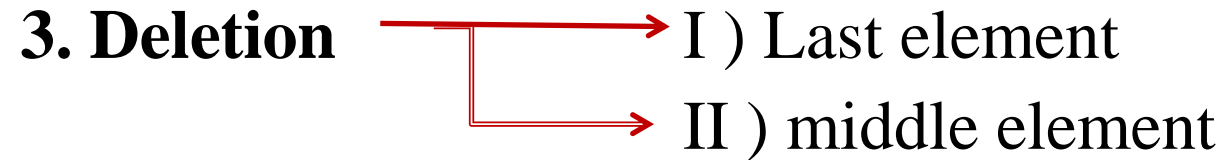
step 6: set  $a[\text{POS}] = \text{new\_value}$

step 7: Exit.

# Operations of Array

- { traversal,
- insertion,
- deletion,**
- merging,
- searching }

## 3. Deletion



### (I) Delete last element

➡ **Simply delete last**

step 1:  $\text{upperbound} = \text{upperbound} - 1$

step 2: Exit.

### (II) Delete middle element

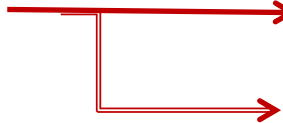
Operations of  
Array  
{ traversal,  
insertion,  
**deletion**,  
merging,  
searching }

## (II) Delete middle element

|             |    |    |    |    |    |       |       |
|-------------|----|----|----|----|----|-------|-------|
| index       | 0  | 1  | 2  | 3  | 4  |       |       |
| a           | 10 | 20 | 30 | 40 | 50 |       |       |
| Delete, POS |    |    |    |    |    | POS=1 |       |
| index       | 0  | 1  | 2  | 3  | 4  |       | i=POS |
| a           | 10 | 20 | 30 | 40 | 50 |       | i=1   |
| index       | 0  | 1  | 2  | 3  | 4  |       |       |
| a           | 10 | 30 | 30 | 40 | 50 |       | i=2   |
| index       | 0  | 1  | 2  | 3  | 4  |       |       |
| a           | 10 | 30 | 40 | 40 | 50 |       | i=3   |
| index       | 0  | 1  | 2  | 3  | 4  |       |       |
| a           | 10 | 30 | 40 | 50 | 50 |       | i=4   |
| index       | 0  | 1  | 2  | 3  | 4  |       |       |
| a           | 10 | 30 | 40 | 50 | 50 |       | N=N-1 |

Operations of  
Array  
{ traversal,  
insertion,  
**deletion**,  
merging,  
searching }

### 3. Deletion

- 
- I ) Last Element
  - II ) Middle Element

#### (II) Delete Middle Element

**delete (A, N, POS)**

A, the array in which the element has to be inserted

N, number of elements in the array

POS, the position at which the element has to be deleted

step 1: set  $i = \text{POS}$

step 2: Repeat step 2 to step 4 while  $i < N-1$

step 3: set  $a[i] = a[i+1]$

step 4: set  $i = i + 1$

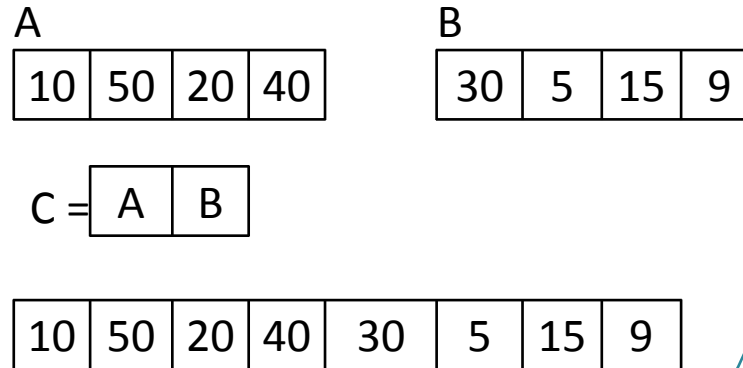
step 5: set  $N = N - 1$

step 6: Exit.

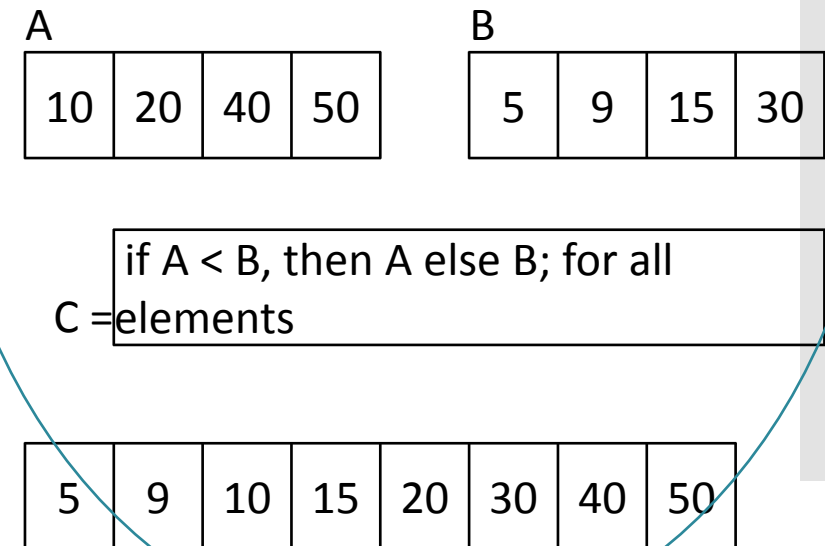
Operations of  
Array  
{ traversal,  
insertion,  
deletion,  
**merging**,  
searching }

#### 4. Merging

**two unsorted array**

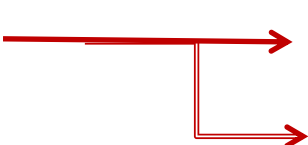


**two sorted array**



# Operations of Array

- { traversal, insertion, deletion, merging, searching }

- 5. Searching 
  - (I) Linear Search
  - (II) Binary Search

## (I) Linear Search

### **linear\_search (A, N, Value, POS)**

A, is the array

N, number of elements in the array

Value, is the data which we want to search

POS, is the index of searched data

step 1: set  $POS = -1$  ,  $i = 0$

step 2: Repeat steps 3 while  $i < N$

step 3: if  $A[i] = Value$  then

    set  $POS = i$

    print POS

    goto step 5

step 4: print “Value is not present in Array”

step 5: Exit.

Operations of  
Array  
{ traversal,  
insertion,  
deletion,  
merging,  
searching }

## (II) Binary Search

**binary\_search (A, lb, ub, Value, POS)**

A, is the array

lb, is the starting index of array

ub, is the last index of array

Value, is the data which we want to search

POS, is the index of searched data

step 1: set  $POS = -1$ ,  $beg = lb$ ,  $end = ub$

step 2: Repeat steps 2 to step 4 while  $beg \leq end$

step 3: set  $mid = (beg + end) / 2$

step 4: if  $A[mid] = Value$

$POS = mid$

    print POS

    goto step 6

    if  $Value < A[mid]$

$end = mid - 1$

    else  $beg = mid + 1$

step 5: if  $POS = -1$ , then

    print "Value is not present in Array"

step 6: Exit.

# 1D Array for Inter-Function Communication

1. Passing Individual Element
  - (i) Passing Data Value
  - (ii) Passing Address
2. Passing the Entire Array



# 1D Array for Inter-Function Communication

## 1. Passing Individual Element

### (i) Passing Data Value

```
void fun(int);  
void main( )  
{  
    int a[5] = {1,2,3,4,5};  
    fun(a[3]);  
}  
*****  
void fun(int num)  
{  
    printf("%d",num);  
}
```

Function declaration

Called function

Calling function

### (ii) Passing Address

```
void fun(int*);  
void main( )  
{  
    int a[5] = {1,2,3,4,5};  
    fun(&a[3]);  
}  
void fun(int *num)  
{  
    printf("%d", *num);  
}
```

Function declaration

Called function

Calling function

# 1D Array for Inter-Function Communication

## 2. Passing the entire Array

```
void fun(int [ ]);  
void main( )  
{  
    int a[5] = {1,2,3,4,5};  
    fun(a);  
}  
  
void fun(int b[ ])  
{  
    int i;  
    for(i=0;i<5;i++)  
        printf("%d,",b[i]);  
}
```

Function declaration

Called function

Calling function

# Pointers and Arrays

```
int a[5] = {1,2,3,4,5};
```

```
int *ptr;
```

```
ptr = &a[0];
```

```
ptr = &a[2];
```

```
ptr = a;
```

Note that  $a[i] = *(a+i)$

# Pointers of Arrays

```
int *ptr[3];  
int a=5, b=7, c=10;  
ptr [0]= &a;  
ptr [1]= &b;  
ptr [2]= &c;  
printf("%d", *ptr[1]);
```

# 2D Arrays

- **Declaration of 2D Array:**

```
data_type array_name [rowsize] [colsize];
```

```
int marks[3][3];
```

|       |    |    |    |
|-------|----|----|----|
| marks | 00 | 01 | 02 |
|       | 10 | 11 | 12 |
|       | 20 | 21 | 22 |

These elements are stored sequentially; Two ways:

- (i) Row Major Order (RMO)
- (ii) Column Major Order (CMO)

## 2D Arrays

### (i) Row Major Order (RMO)

|       |       |       |       |       |       |       |       |       |
|-------|-------|-------|-------|-------|-------|-------|-------|-------|
|       |       |       |       |       |       |       |       |       |
| (0,0) | (0,1) | (0,2) | (1,0) | (1,1) | (1,2) | (2,0) | (2,1) | (2,2) |

**A[M][N];**

**Address  $A[i, j] = \text{BaseAddress} + w * \{ N(i-1) + (j-1) \}$**

Where,  $w$  is the # of words stored per memory location

$N$  is the number of elements in one Row (num of Columns)

$i$  and  $j$  are the subscripts of array element

Example 1:  $A[20][5]$ , Base Address=1000, number of words per memory location = 2, Compute the address of element,  $A[18,4]$ . Assume, the elements are stored in RMO.

## 2D Arrays

### (I) Row Major Order (RMO)

**Example 1:** A[20][5], Base Address=1000, number of words per memory location = 2, Compute the address of element, A[18,4]. Assume, the elements are stored in RMO.

**Solution:**

$$\begin{aligned}\text{Address of } A[i, j] &= \text{BaseAddress} + w * [ N*(i-1) + (j-1) ] \\ \text{Address of } A[18,4] &= 1000 + 2 * [ 5 * (18-1) + (4-1) ] \\ &= 1000 + 2 * [85 + 3] \\ &= 1176\end{aligned}$$

**Example 2:** A[10][5], Base Address=2000, number of words per memory location = 2, Compute the address of element, A[8,5]. Assume, the elements are stored in RMO.

**Ans:**

$$\text{Address of } A[8,5] = 2078.$$

**Example 3:** A[10][10], Base Address=1000, number of words per memory location = 2, Compute the address of element, A[8,5]. Assume, the elements are stored in RMO.

**Ans:**

$$\text{Address of } A[8,5] = 1148.$$

## 2D Arrays

### (II) Column Major Order (CMO)

|   |   |   |
|---|---|---|
| 2 | 4 | 6 |
| 1 | 3 | 5 |
| 3 | 6 | 9 |

|   |   |   |   |   |   |   |   |   |
|---|---|---|---|---|---|---|---|---|
| 2 | 1 | 3 | 4 | 3 | 6 | 6 | 5 | 9 |
|---|---|---|---|---|---|---|---|---|

For MxN Matrix,

Address of  $A[i, j] = \text{BaseAddress} + w * [ M*(j-1) + (i-1) ]$

**Example 1:**  $A[10][10]$ , Base Address=1000, number of words per memory location = 2, Compute the address of element,  $A[8,5]$ . Assume, the elements are stored in CMO.

**Solution:**

Address of  $A[i, j] = \text{BaseAddress} + w * [ M*(j-1) + (i-1) ]$

Address of  $A[8,5] = 1000 + 2 * [ 10 * (5-1) + (8-1) ]$   
 $= 1000 + 2 * [ 40 + 7 ]$   
 $= 1094$



# 2D Arrays

- **Initialization of 2D Array**

```
int A[2][3] = { 1,2,3,4,5,6};
```

```
int A[2][3] = { { 1,2,3}, {4,5,6}};
```

Only the size of first dimension can be omitted.

```
int A[ ] [3] = { { 1,2,3}, {4,5,6}};
```

```
int A[2][3] = {0};
```

- **Accessing the Elements**

- Using two FOR loops

## 2D Arrays

- **Operations on 2D Array**

1) Transpose: A is  $M \times N$  then B is  $N \times M$  where,  $B_{i,j} = A_{j,i}$

2) Sum:  $C_{i,j} = A_{i,j} + B_{i,j}$

3) Difference:  $C_{i,j} = A_{i,j} - B_{i,j}$

4) Product: A is  $M \times N$

B is  $P \times Q$

if  $N = P$  then  $C_{i,j} = \sum A_{i,k} B_{k,j}$  , for  $k = 1$  to  $N$

# Quiz

Which of the following C expressions access the  $(i, j)^{\text{th}}$  entry of an  $(m \times n)$  matrix stored in column major order both row and column starting index is 1?

- a).  $n \times (i - 1) + j$
- b).  $m \times (j - 1) + i$
- c).  $m \times (n - j) + j$
- d).  $n \times (m - i) + j$