

Name:	Sunit Sunil Khaire
Roll No:	19
Class/Sem:	TE/V
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**Aim:** Implementation of Bidirectional search for problem solving.

**Objective:** To study the Bidirectional searching techniques and its implementation for problem solving.

#### Theory:

Bidirectional search is a graph search algorithm which find smallest path from source to goal vertex. It runs two simultaneous search –

- 1. Forward search from source/initial vertex toward goal vertex
- 2. Backward search from goal/target vertex toward source vertex

Bidirectional search replaces single search graph(which is likely to grow exponentially) with two smaller sub graphs – one starting from initial vertex and other starting from goal vertex. The search terminates when two graphs intersect.

#### Algorithm:

#### Steps for Bidirectional Search Algorithm

#### 1. Initialization:

- Create two frontiers:
  - One for the forward search starting from the initial node (start).
  - One for the backward search starting from the goal node (goal).
- o Create two sets to keep track of visited nodes for each search direction:
  - visited\_start for the forward search.
  - visited\_goal for the backward search.
- o Initialize the frontiers by adding the start node to frontier\_start and the goal node to frontier\_goal.
- Initialize the visited\_start and visited\_goal sets with their respective starting nodes.

#### 2. Search Expansion:

- o Repeat the following steps until a meeting point is found or one of the frontiers is empty:
  - Expand Forward Search:
    - Remove the current node from frontier\_start.
    - Expand all neighboring nodes of the current node.



- For each neighbor:
  - If the neighbor is not in visited\_start:
    - Add the neighbor to frontier\_start.
    - Mark the neighbor as visited in visited\_start.
    - Check if the neighbor is already in visited\_goal:
      - If yes, a meeting point is found. Proceed to reconstruct the path.

#### Expand Backward Search:

- Remove the current node from frontier\_goal.
- Expand all neighboring nodes of the current node.
- For each neighbor:
  - If the neighbor is not in visited\_goal:
    - Add the neighbor to frontier\_goal.
    - Mark the neighbor as visited in visited\_goal.
    - Check if the neighbor is already in visited\_start:
      - If yes, a meeting point is found. Proceed to reconstruct the path.

#### 3. Meeting Point:

The search stops when a node from frontier\_start is found in visited\_goal or a node from frontier\_goal is found in visited\_start. This node is the meeting point.

#### 4. Path Reconstruction:

o Reconstruct the path from the start node to the goal node by combining the paths from both



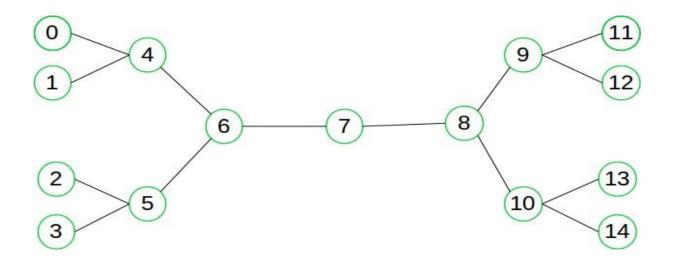
searches at the meeting point.

- Start from the meeting point:
  - Trace back to the start node using the information in visited\_start.
  - Trace back to the goal node using the information in visited\_goal.
- o Concatenate the two paths to form the complete path from start to goal.

#### 5. Termination:

 If one of the frontiers is empty and no meeting point is found, it means there is no path from the start node to the goal node.

#### Example:



Suppose we want to find if there exists a path from vertex 0 to vertex 14. Here we can execute two searches, one



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from vertex 0 and other from vertex 14. When both forward and backward search meet at vertex 7, we know that we have found a path from node 0 to 14 and search can be terminated now. We can clearly see that we have successfully avoided unnecessary exploration.

#### When to use bidirectional approach?

We can consider bidirectional approach when-

- 1. Both initial and goal states are unique and completely defined.
- 2. The branching factor is exactly the same in both directions.

#### **Performance measures**

**Completeness:** Bidirectional Search is complete if we use BFS in both searches.

**Time Complexity:** Time complexity of bidirectional search using BFS is  $O(b^d)$ .

**Space Complexity:** Space complexity of bidirectional search is  $O(b^d)$ .

**Optimal:** Bidirectional search is Optimal.

#### Advantages:

o Bidirectional search is fast.

Bidirectional search requires less memory

#### **Disadvantages:**

- Implementation of the bidirectional search tree is difficult.
- o In bidirectional search, one should know the goal state in advance.

#### Code:

from collections import deque

```
# Bidirectional search function
def bidirectional_search(graph, start, goal):
  # Initialize visited sets for both directions
  visited_from_start = {start}
  visited_from_goal = {goal}
  # Queues for BFS from both directions
  queue_start = deque([start])
  queue_goal = deque([goal])
  # Parent dictionaries to reconstruct the path
  parent_start = {start: None}
  parent_goal = {goal: None}
  # Perform BFS from both directions
  while queue_start and queue_goal:
    # Expand from start side
    if queue_start:
       node_start = queue_start.popleft()
       for neighbor in graph[node_start]:
          if neighbor not in visited from start:
```

visited\_from\_start.add(neighbor)



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```
parent_start[neighbor] = node_start
            queue_start.append(neighbor)
            # If a connection is found, reconstruct the path
            if neighbor in visited from goal:
               return reconstruct_path(parent_start, parent_goal, neighbor)
    # Expand from goal side
    if queue_goal:
       node_goal = queue_goal.popleft()
       for neighbor in graph[node_goal]:
          if neighbor not in visited_from_goal:
            visited from goal.add(neighbor)
            parent_goal[neighbor] = node_goal
            queue_goal.append(neighbor)
            # If a connection is found, reconstruct the path
            if neighbor in visited_from_start:
               return reconstruct_path(parent_start, parent_goal, neighbor)
  return None # No path found
# Function to reconstruct the path from the meeting point
def reconstruct_path(parent_start, parent_goal, meeting_point):
  # Reconstruct the path from start to meeting point
  path_start = []
  node = meeting_point
  while node is not None:
    path_start.append(node)
    node = parent_start[node]
  path start.reverse()
  # Reconstruct the path from meeting point to goal
  path_goal = []
  node = parent_goal[meeting_point]
  while node is not None:
    path goal.append(node)
    node = parent_goal[node]
  # Combine both parts of the path
  return path_start + path_goal
# Example usage
if __name__ == "__main__":
  # Example graph as an adjacency list
  graph = \{
    'A': ['B', 'C'],
    'B': ['A', 'D', 'E'],
    'C': ['A', 'F'],
    'D': ['B'],
    'E': ['B', 'F'],
    'F': ['C', 'E']
  start_node = 'A'
  goal\_node = 'F'
```

result = bidirectional\_search(graph, start\_node, goal\_node)



if result:
 print("Path found:", result)
else:
 print("No path found")

#### Output

Path found: ['A', 'C', 'F']

#### **Conclusion:**

Bidirectional search is used in real-world applications like GPS navigation for finding the shortest routes, social networks for discovering connections, robot path planning in complex environments, network routing for efficient data transmission, and puzzle solving to find solutions faster. It reduces the search space by exploring from both the start and goal simultaneously, improving efficiency irectional search techniques.



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