



**Multi-Agent Systems**

# **Voting Theory**

UvA

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# Which Voting Rules?

- We need some general principles to distinguish between voting rules
  - Let's take the point of view of someone who wants to design a voting rule from scratch, and think about what properties, or **axioms**, we would want the voting rule to satisfy

# Recall: A Formal Model

## Definition

- $N = \{1, \dots, n\}$  is a set of agents, or voters
- $A = \{a, b, c, \dots\}, |A| = m$  is a finite set of alternatives, or candidates
- $\succ_i$  is a preference order of voter  $i$  (a linear order on alternatives)
- $L = \{ \succ \mid \succ \text{ is a linear order on } A \}$  is the set of all possible preferences
- $R = (\succ_1, \dots, \succ_n) \in L^n$  is a preference profile
- $F : L^n \rightarrow 2^A \setminus \emptyset$  is a social choice function
- $F : L^n \rightarrow A$  is a resolute social choice function
- $F : L^n \rightarrow L$  is a social welfare function

# Anonymity

- The first axiom we look at is **anonymity**
- It says that the order in which we arrange the voters does not matter for the final result

## Definition (Anonymity Axiom)

A voting rule  $F$  satisfies **anonymity** if, for any permutation  $\sigma$  of the set  $N$  of voters, it holds that:

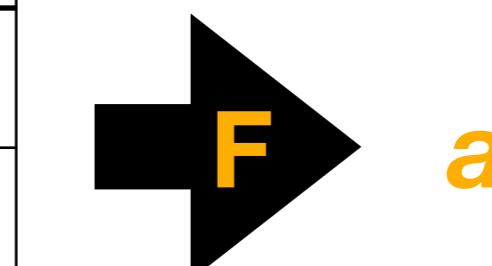
$$F(\succ_1, \dots, \succ_n) = F(\succ_{\sigma(1)}, \dots, \succ_{\sigma(n)})$$

# Anonymity

- **Anonymity** requires invariance under permutations of the voters in the profile

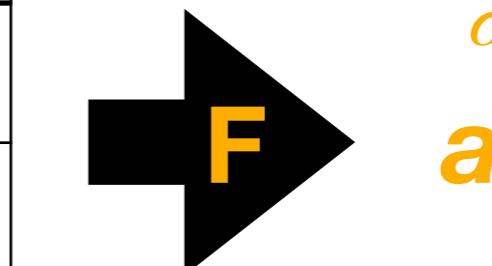
If

| Ann      | Bob      | Cat      |
|----------|----------|----------|
| <i>a</i> | <i>a</i> | <i>b</i> |
| <i>b</i> | <i>c</i> | <i>a</i> |
| <i>c</i> | <i>b</i> | <i>c</i> |



*a*

Permutation here is  
 $\sigma(\text{Ann}) = \text{Bob}$ ,  
 $\sigma(\text{Bob}) = \text{Cat}$ ,  
 $\sigma(\text{Cat}) = \text{Ann}$



*a*

Then

| Bob      | Cat      | Ann      |
|----------|----------|----------|
| <i>a</i> | <i>b</i> | <i>a</i> |
| <i>c</i> | <i>a</i> | <i>b</i> |
| <i>b</i> | <i>c</i> | <i>c</i> |

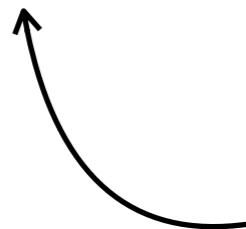
# Neutrality

- The next property is **neutrality**
- It says that the names we give to alternatives do not matter

## Definition (Neutrality Axiom)

A voting rule  $F$  satisfies **neutrality** if, for any permutation  $\sigma$  of the set  $A$  of alternatives, it holds that:

$$\sigma(F(\succ_1, \dots, \succ_n)) = F(\sigma(\succ_1), \dots, \sigma(\succ_n))$$



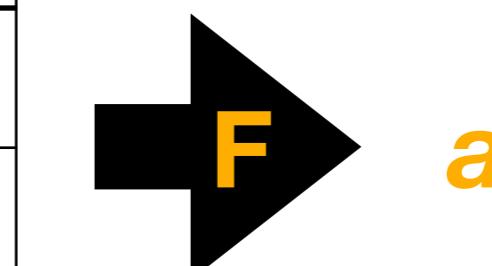
Every alternative is replaced with its  
image under  $\sigma$

# Neutrality

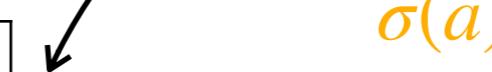
- **Neutrality** requires that permutations of the alternatives in the profile are reflected by permutations of the alternatives in the result

If

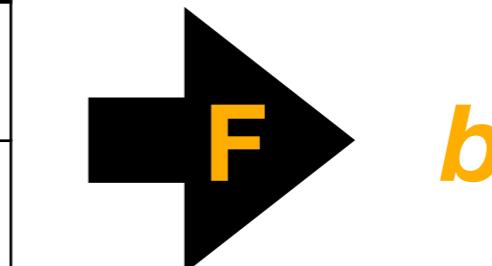
| Ann | Bob | Cat |
|-----|-----|-----|
| a   | a   | b   |
| b   | c   | a   |
| c   | b   | c   |



*a*



$\sigma(a) = b, \sigma(b) = c,$   
 $\sigma(c) = a$



*b*

Then

| Ann | Bob | Cat |
|-----|-----|-----|
| b   | b   | c   |
| c   | a   | b   |
| a   | c   | a   |

# Positive Responsiveness

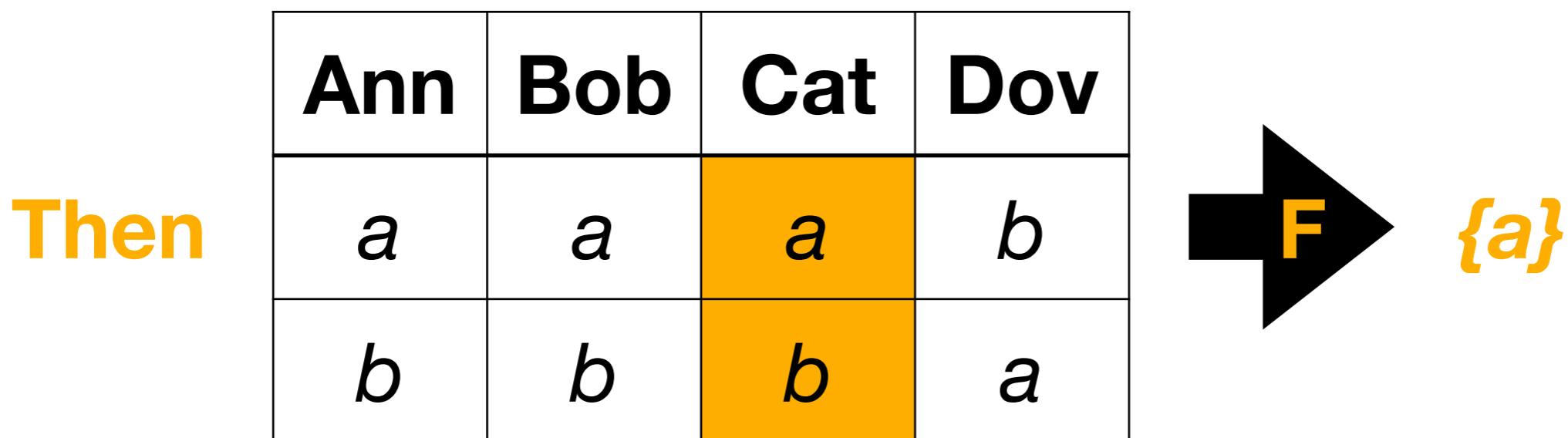
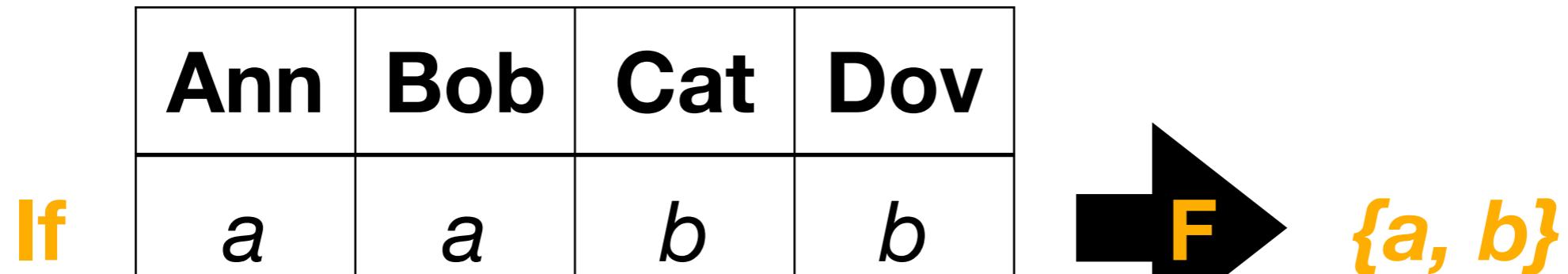
- Now, we look at something a bit more involved: **positive responsiveness**
- It says, roughly, that increased support for some alternative has the power to break a tie in favor of that alternative

## Definition (Positive Responsiveness Axiom)

A social choice function  $F$  satisfies **positive responsiveness** if, for any distinct profiles  $R$  and  $R'$ , and alternative  $x^*$ , we have that  $R$  and  $R'$  are the same except that in  $R'$  some voters move  $x^*$  up some positions in their preference rankings, then it holds that if  $x^* \in F(R)$ , then  $F(R') = \{x^*\}$ .

# Positive Responsiveness

- If in  $R'$  some voters raise  $x^*$ , while leaving everything else untouched, then  $x^*$  goes from being a (possibly tied) winner to the unique winner



# May's Theorem

- For two alternatives, it turns out that these properties are satisfied only by the majority voting rule
  - Note that when there are only two alternatives, the majority rule is well-defined

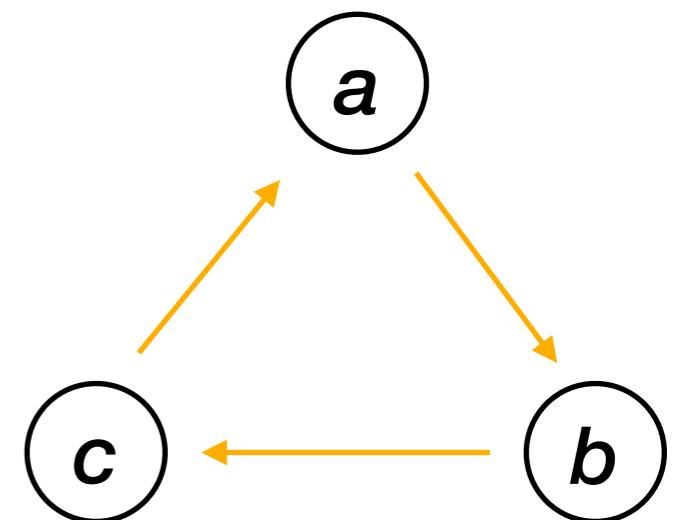
## Theorem (May, 1952)

If there are only two alternatives, then the only social choice function that satisfies **anonymity**, **neutrality**, and **positive responsiveness** is the majority rule.

# More on May's Theorem

- For two alternatives, we cannot do better than using majority
  - And note that when there are only two alternatives, all the voting rules we have looked at so far are equivalent to the majority rule
  - Now for more than two alternatives...
- We know that majority comparisons can get us into troubles with cycles
  - But maybe there is some other clever way to combine preferences into a coherent social ranking?

|   |   |   |
|---|---|---|
| 1 | 1 | 1 |
| a | b | c |
| b | c | a |
| c | a | b |



# Reasonable Properties of Voting Rules

- For the next results we will focus on **social welfare functions**: voting rules that return a ranking of the alternatives
- Let's write down some more reasonable properties...

# Pareto Efficiency

- Pareto efficiency captures the intuition that if everyone thinks some alternative is better than another, then this should be reflected in the result

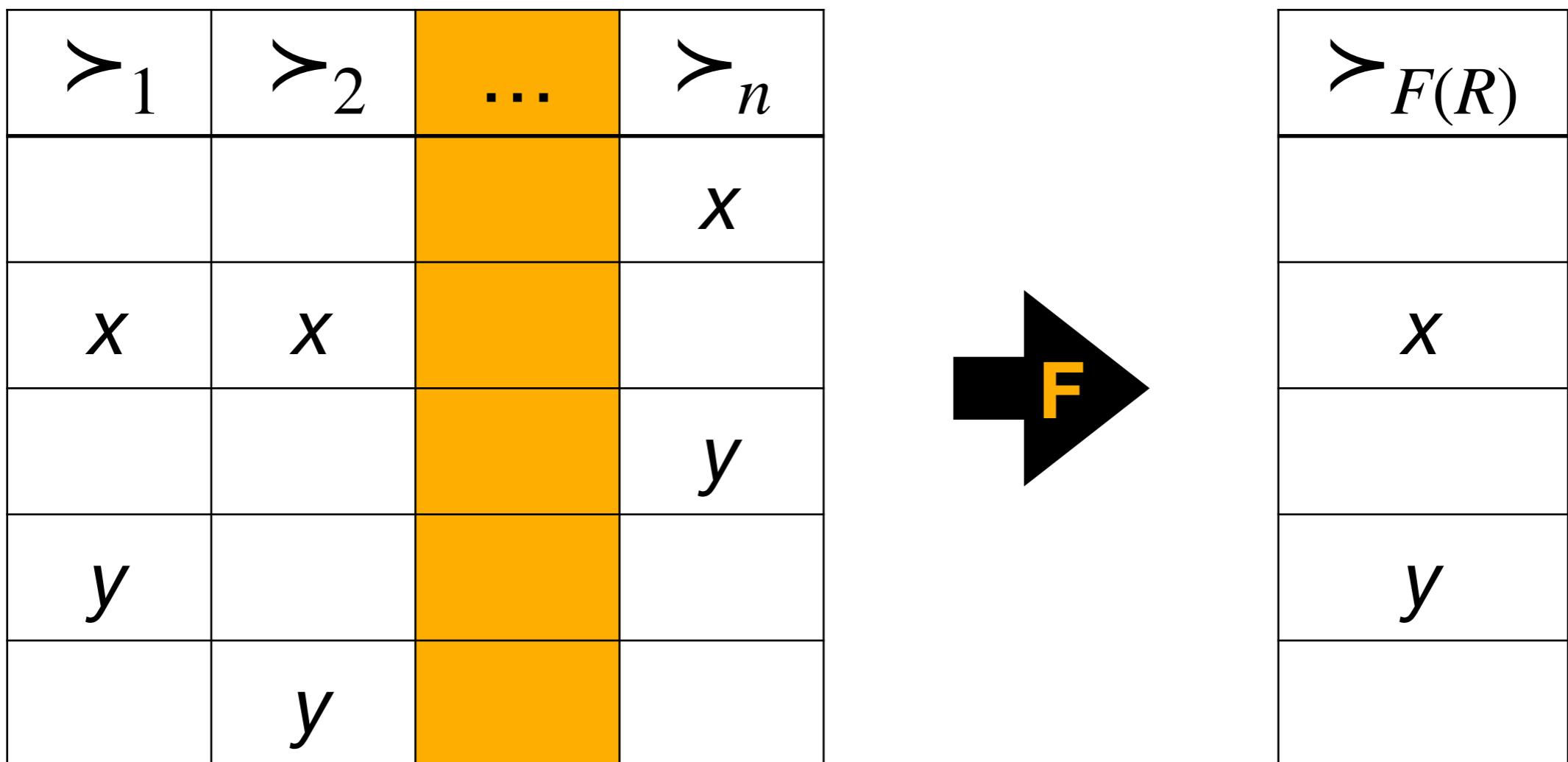
## Definition (Pareto Efficiency Axiom)

A social welfare function  $F$  satisfies **Pareto efficiency** if, for any alternatives  $x$  and  $y$ , it holds that if  $x >_i y$ , for every voter  $i \in N$ , then  $x >_{F(R)} y$ .



# Pareto Efficiency

- If there is unanimous agreement that  $x$  is better than  $y$ , then  $x$  is ranked above  $y$  in the aggregated ranking



# Independence of Irrelevant Alternatives

- Society's ranking between two alternatives  $x$  and  $y$  should depend on how voters in the profile rank  $x$  and  $y$ ... and nothing else

## Definition (Independence of Irrelevant Alternatives, or IIA, Axiom)

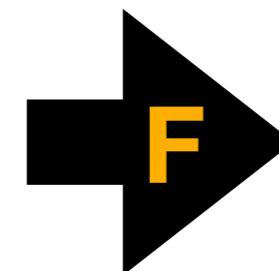
A social welfare function  $F$  satisfies **independence of irrelevant alternatives (IIA)** if, for any alternatives  $x$  and  $y$ , and profiles  $R$ ,  $R'$  such that for any agent  $i \in N$  it holds that  $x >_i y$  if and only if  $x >'_i y$ , then it holds that  $x >_{F(R)} y$  if and only if  $x >_{F(R')} y$ .

# Independence of Irrelevant Alternatives

- If voters rank  $x$  and  $y$  in the same way across the two profiles, then the final ranking between  $x$  and  $y$  is the same for both profiles

If

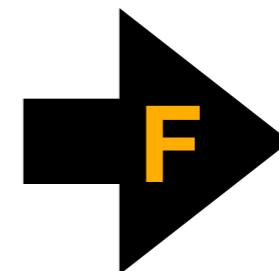
| $\succ_1$ | $\succ_2$ | ... | $\succ_n$ |
|-----------|-----------|-----|-----------|
|           |           |     |           |
| $x$       | $x$       |     | $y$       |
| $y$       |           |     |           |
|           | $y$       |     | $x$       |
|           |           |     |           |



| $\succ_{F(R)}$ |
|----------------|
|                |
|                |
| $y$            |
|                |
| $x$            |
|                |

Then

| $\succ'_1$ | $\succ'_2$ | ... | $\succ'_n$ |
|------------|------------|-----|------------|
|            |            |     |            |
| $x$        | $x$        |     | $y$        |
| $y$        |            |     |            |
|            | $y$        |     | $x$        |
|            |            |     |            |



| $\succ_{F(R')}$ |
|-----------------|
|                 |
|                 |
| $y$             |
|                 |
| $x$             |
|                 |

# Non-Dictatorship

- Non-dictatorship is about making sure that there is no one voter who has a final say, regardless of the preferences of the other voters

## Definition (Dictator)

An agent  $i \in N$  is a **dictator** for a social welfare function  $F$  if, for any alternatives  $x$  and  $y$ , and profile  $R$ , it holds that if  $x >_i y$ , then  $x >_{F(R)} y$ .

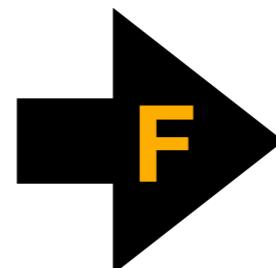
## Definition (Non-Dictatorship Axiom)

A social welfare function  $F$  satisfies **non-dictatorship** if no agent is a dictator.

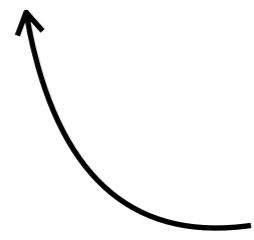
# Dictatorship

- A dictator decides the final ranking of every pair of alternatives, and thus the full final ranking

| Ann      | Bob      | Cat      |
|----------|----------|----------|
| <i>a</i> | <i>a</i> | <i>b</i> |
| <i>b</i> | <i>c</i> | <i>a</i> |
| <i>c</i> | <i>b</i> | <i>c</i> |



| $\succ_{F(R)}$ |
|----------------|
| <i>a</i>       |
| <i>b</i>       |
| <i>c</i>       |



Designated dictator

# Arrow's Theorem

- These properties seem reasonable enough. But it turns out that, together, they spell trouble...

## Theorem (Arrow, 1951)

If there are at least three alternatives, then any social welfare function that satisfies **Pareto efficiency** and **independence of irrelevant alternatives** is a **dictatorship**.

# **What now?**

**Let's look closer at another type  
of strategy: lying...**

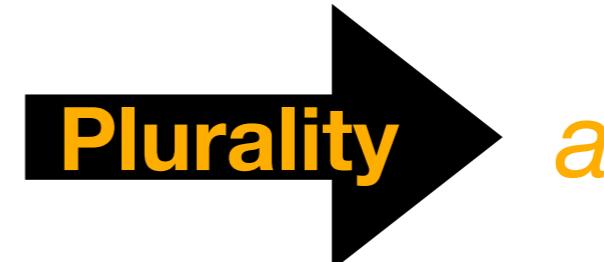
# Lying

- We've seen that many voting rules are afflicted by a common problem: they create incentives for voters to lie about their preferences
- For now, we return to resolute social choice functions  $F : L^n \rightarrow A$ , where  $F$  outputs a single candidate
- Recall...

# Plurality

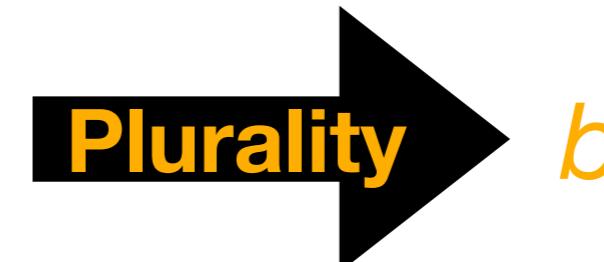
- Under plurality, voters don't want to support a losing candidate

|    |    |   |
|----|----|---|
| 49 | 48 | 3 |
| a  | b  | c |
| b  | c  | b |
| c  | a  | a |



|    |    |   |
|----|----|---|
| 49 | 48 | 3 |
| a  | b  | b |
| b  | c  | c |
| c  | a  | a |

Submitting insincere ballot



# Borda

- Under Borda, voters can manipulate by pushing alternatives they don't like down their list

|          |          |
|----------|----------|
| 2        | 1        |
| <i>b</i> | <i>a</i> |
| <i>a</i> | <i>b</i> |
| <i>c</i> | <i>c</i> |
| <i>d</i> | <i>d</i> |



|          |          |
|----------|----------|
| 2        | 1        |
| <i>b</i> | <i>a</i> |
| <i>a</i> | <i>c</i> |
| <i>c</i> | <i>d</i> |
| <i>d</i> | <i>b</i> |



**It would be great if truthfulness  
turned out to be a dominant  
strategy...**

# Strategyproofness

## Definition (Strategyproofness)

A resolute social choice function  $F$  is **strategyproof** if for all voters  $i \in N$  it holds that there does not exist a profile  $R$  (containing  $i$ 's truthful preference) and some order  $\succ'_i$  (representing some untruthful preference of  $i$ ) such that:

$$F(\succ'_i, R_{-i}) \succ_i F(R)$$



In other words, it is not possible for any voter  $i$  to get a better result by submitting a preference order  $\succ'_i$  instead of its true preference order  $\succ_i$ .

**Can we design strategyproof  
voting rules?**

# Dictatorship

## Definition (Dictatorship)

Choose an agent  $i \in N$ , called **the dictator**. The winner is the top choice of the dictator.

# Dictatorship

- Under dictatorship there is no point in manipulating

Dictator

| Ann | Bob | Cat |
|-----|-----|-----|
| a   | a   | c   |
| b   | c   | a   |
| c   | b   | b   |

Dictatorship → a

Dictator

| Ann | Bob | Cat |
|-----|-----|-----|
| a   | a   | b   |
| b   | c   | a   |
| c   | b   | c   |

Submitting insincere ballot

Dictatorship → a

# Another Impossibility Result!

**Theorem (Gibbard-Satterthwaite, 1973, 1975)**

If a resolute social choice function  $F$  has at least three possible outcomes, then  $F$  is strategyproof if and only if it is a dictatorship.

# What Now?

- What can we take from the impossibility theorems?
  - Certain intuitive desirable properties are incompatible
  - Illustrates that there are trade-offs when designing a voting system
  - Democracy is mathematically limited
  - It does not mean that voting is pointless, but we must decide what kind of unfairness we can live with