



UvA

Multi-Agent Systems

Game Theory

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Decision Theory

- **Decision theory** is about deciding what is best when one agent acts against nature (not other agents)
- Example: Two lotteries, you can enter one.
 - Lottery 1: 50% chance to win 150€
 - Lottery 2: 90% chance to win 100€
 - Which one would you choose?
- You want to **maximize expected utility**



$$\underset{\substack{\uparrow \\ \text{expected utility}}}{\mathbb{E}[u(a)]} = \sum_o (\overset{\substack{\text{utility of acting } a \\ \text{when the state is } o}}{u(a, o)} \cdot \overset{\substack{\text{probability that } o \\ \text{occurs}}}{\mathbb{P}[o]})$$

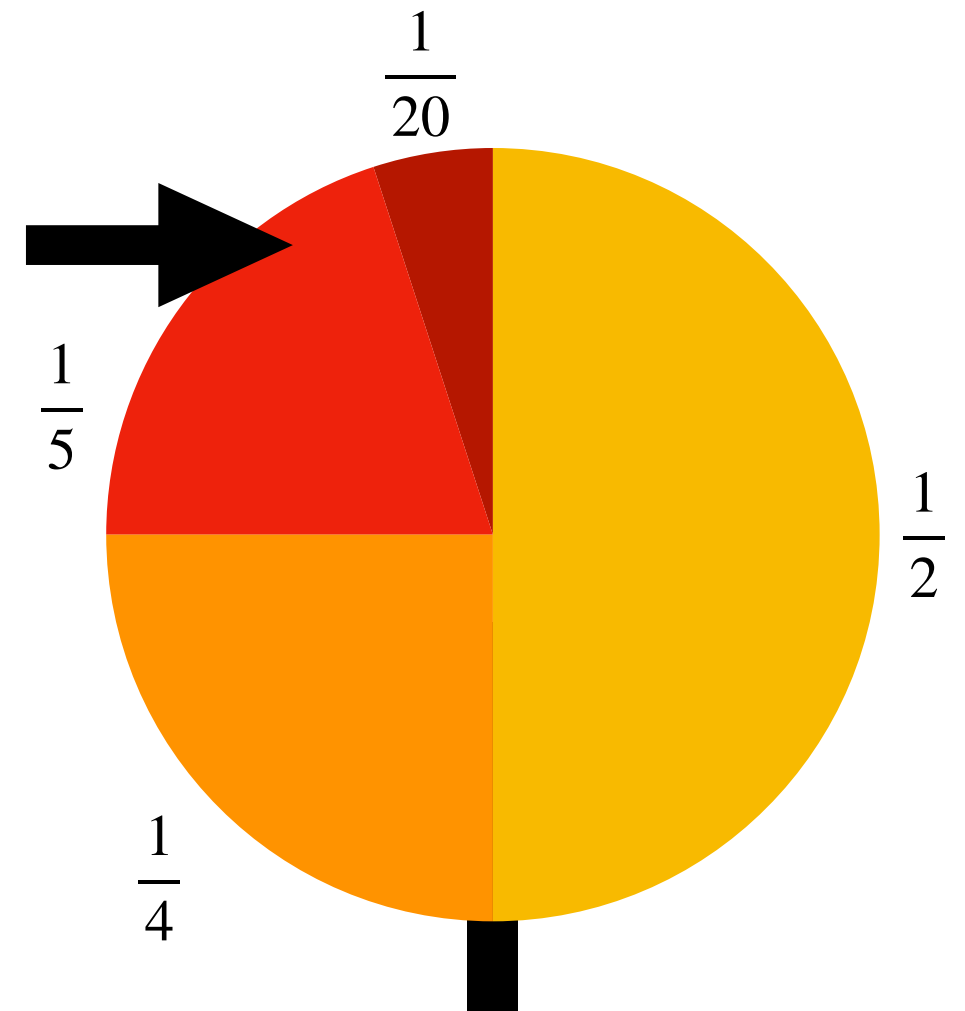
$o \leftarrow$ possible state of nature

$$\mathbb{E}[u(L1)] = 150 \cdot \frac{1}{2} = 75$$

$$\mathbb{E}[u(L2)] = 100 \cdot \frac{9}{10} = 90$$

Decision Theory: Another Example

$$\mathbb{E}[u(a)] = \sum_o (u(a, o) \cdot \mathbb{P}[o])$$



$$\mathbb{E}[u(\text{spinning the wheel})] =$$

$$u(\text{spinning, win } 0) \cdot \mathbb{P}[\text{win } 0] + u(\text{spinning, win } 5) \cdot \mathbb{P}[\text{win } 5] + (\dots) =$$

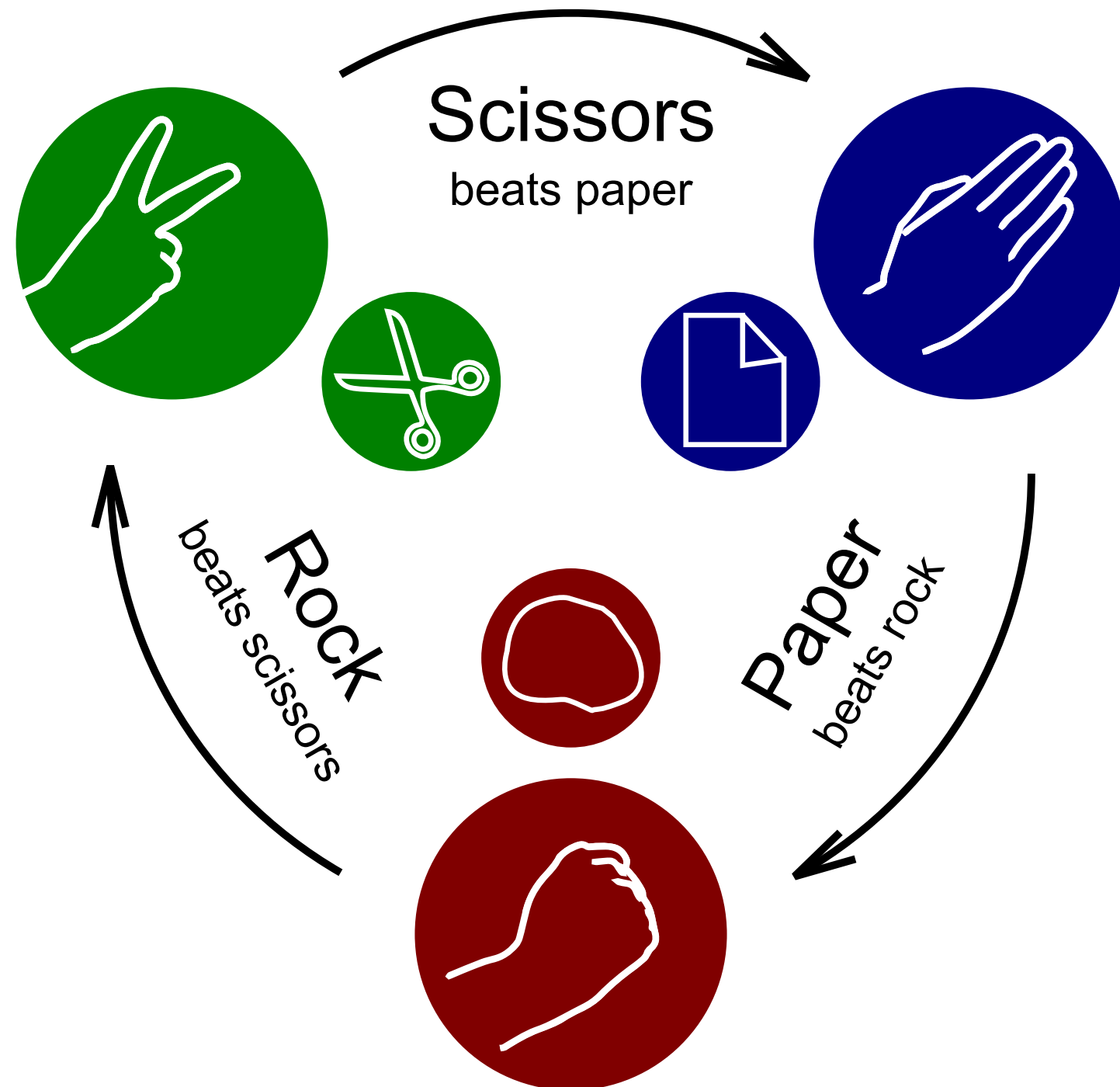
$$0 \cdot \frac{1}{2} + 5 \cdot \frac{1}{4} + 10 \cdot \frac{1}{5} + 100 \cdot \frac{1}{20} = \frac{33}{4} = 8,25$$

Coming up: Playing with others...

What is Game Theory?

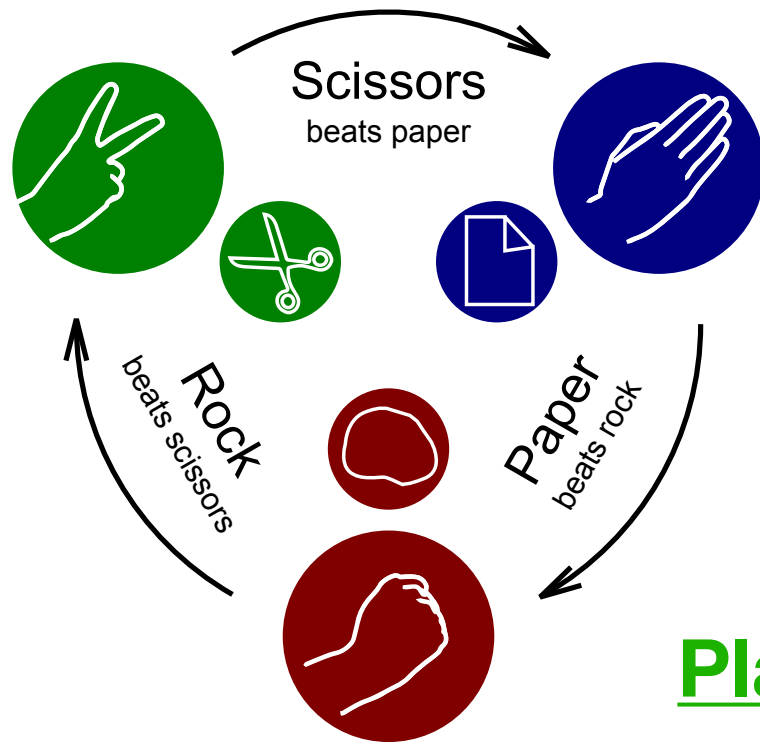
- **Game theory** is about interactions among independent, self-interested agents
- In a way: thinking about conflict, natural phenomena or daily-life scenarios as games
- **Cooperative** and **non-cooperative** game theory
 - In this course: non-cooperative
 - Players are:
 1. **Intelligent** (reason perfectly and quickly)
 2. **Rational** (always want to maximize their payoff)
 3. **Selfish** (only care for their own payoff)

Example: Rock Paper Scissors



Example: Rock Paper Scissors

You get 1 point for winning, and -1 for losing.



Player 1

Rock

Paper

Scissors

Player 2

Rock

Paper

Scissors

	Rock	Paper	Scissors
Rock	(0,0)	(-1,1)	(1,-1)
Paper	(1,-1)	(0,0)	(-1,1)
Scissors	(-1,1)	(1,-1)	(0,0)

This way of presenting a game is called a **normal form game**.

Generally, we assume Player 1 is the row player, and Player 2 is the column player.

Another Example: Stag Hunt

- Two hunters
- Have to decide what to hunt: one stag or two hares
- A stag takes two to catch: hunting a stag solo is certain failure
- One hunter alone can catch both hares, but if both hunters go for the hares they end up with one each
- Even combined, the hares are worth less than the stag



Another Example: Stag Hunt



Hunter 1

Hunter 2

Stag

Hare

Stag

(10,10)

(0,6)

Hare

(6,0)

(3,3)



Games in Normal Form

- A game in normal form consists of **players** who can take **actions**, which lead to **payoffs**

Definition (Normal-Form Game)

A normal-form game is a tuple (N, A, u) , where:

- $N = \{1, \dots, n\}$ is a finite set of agents, or players
- A_i is a finite set of actions available to player i
- $a = (a_1, \dots, a_n)$ is an action profile
- $A = A_1 \times \dots \times A_n$ is the set of all action profiles
- $u_i : A \rightarrow \mathbb{R}$ is a utility (or payoff) function for player i
- $u = (u_1, \dots, u_n)$ is a utility profile

Back to the Stag Hunt

Hunter 2

Hunter 1

Stag

Hare

Stag

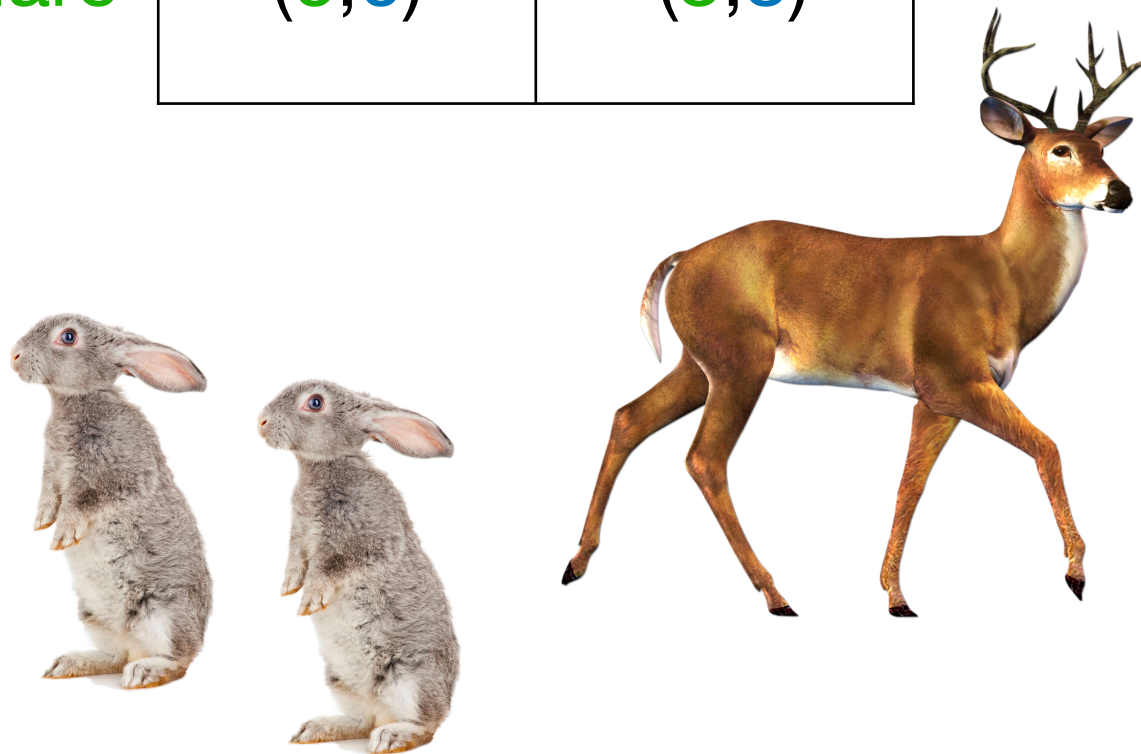
(10,10)

(0,6)

Hare

(6,0)

(3,3)



- $N = \{1,2\}$
- $A_1 = \{Stag, Hare\}$
- $A_2 = \{Stag, Hare\}$
- $A = \{(Stag, Stag), (Stag, Hare), (Hare, Stag), (Hare, Hare)\}$
- $u_1(Stag, Stag) = 10$
- $u_2(Stag, Stag) = 10$
- $u_1(Stag, Hare) = 0$
- $u_2(Stag, Hare) = 6$
- (...)

What Should Players Do: Strategies

- A **pure strategy** for an agent is to select a single action and play it

- $N = \{1, \dots, n\}$ is a finite set of agents, or players
- A_i is a finite set of actions available to player i
- $a = (a_1, \dots, a_n)$ is an action profile
- $A = A_1 \times \dots \times A_n$ is the set of all action profiles
- $u_i : A \rightarrow \mathbb{R}$ is a utility (or payoff) function for player i
- $u = (u_1, \dots, u_n)$ is a utility profile

Definition (Pure Strategy)

- $S_i = A_i$ are the pure strategies of player i
- $s = (s_1, \dots, s_n)$ is a strategy profile
- $S = S_1 \times \dots \times S_n$ is the set of all strategy profiles
- $u_i(s) = u_i(a)$ is the utility of i wrt. strategy profile s
- $s_{-i} = (s_1, \dots, s_{i-1}, s_{i+1}, \dots, s_n)$ is s without s_i
- $s = (s_i, s_{-i})$ is another way to write s

What Should Players Do: Solution Concepts

- A **solution concept** describes what strategies we might expect the players will adopt
 - And therefore, the result of the game

Dominance Among Strategies

- A player has a **dominated strategy** if the player would always be as well (or even better) off playing something else

Strict Dominance

Definition (Strict Dominance Among Strategies)

- Strategy s_i **strictly dominates** s'_i if $u_i(s_i, s_{-i}) > u_i(s'_i, s_{-i})$, for any profile s_{-i} of other agents' strategies
- Strategy s_i is **strictly dominant** if it strictly dominates every other strategy s'_i

Example

	L	C	R
T	(3,0)	(2,1)	(0,0)
M	(1,1)	(1,1)	(5,0)
B	(0,1)	(4,2)	(0,1)

- For Player 1, does T strictly dominate M?

- If P2 plays L: $3 > 1$ ✓
- If P2 plays C: $2 > 1$ ✓
- If P2 plays R: $0 < 5$ ✗
- No

- For Player 2, does C strictly dominate L?

- If P1 plays T: $1 > 0$ ✓
- If P1 plays M: $1 = 1$ ✗
- If P1 plays B: $2 > 1$ ✓
- No

Example

	L	C	R
T	(3,0)	(2,1)	(0,0)
M	(1,1)	(1,1)	(5,0)
B	(0,1)	(4,2)	(0,1)

- For Player 2, does C strictly dominate R?
 - If P1 plays T: $1 > 0$ ✓
 - If P1 plays M: $1 > 0$ ✓
 - If P1 plays B: $2 > 1$ ✓
 - Yes

Player 2 should always play C over R...

...this suggests a way of filtering out strategies

Eliminating Strictly Dominated Strategies

Definition (Iterated Elimination of Strictly Dominated Strategies, or IESDS)

- For each player, go through its strategies. If any are strictly dominated, eliminate them. Repeat until there are no further strategies to eliminate.

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R

	L	C	R
T	(3,0)	(2,1)	(0,0)
M	(1,1)	(1,1)	(5,0)
B	(0,1)	(4,2)	(0,1)

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R
 - T strictly dominates M:
 - Eliminate M

	L	C
T	(3,0)	(2,1)
M	(1,1)	(1,1)
B	(0,1)	(4,2)

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R
 - T strictly dominates M:
 - Eliminate M
 - C strictly dominates L:
 - Eliminate L

	L	C
T	(3,0)	(2,1)
B	(0,1)	(4,2)

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R
 - T strictly dominates M:
 - Eliminate M
 - C strictly dominates L:
 - Eliminate L
 - B strictly dominates T:
 - Eliminate T

	C
T	(2,1)
B	(4,2)

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R
 - T strictly dominates M:
 - Eliminate M
 - C strictly dominates L:
 - Eliminate L
 - B strictly dominates T:
 - Eliminate T
 - Nothing left:
 - Stop

	C
B	(4,2)

Back to the Example

- What can we eliminate here?
 - C strictly dominates R:
 - Eliminate R
 - T strictly dominates M:
 - Eliminate M
 - C strictly dominates L:
 - Eliminate L
 - B strictly dominates T:
 - Eliminate T
 - Nothing left:
 - Stop
- IESDS solution: (B, C)

	L	C	R
T	(3,0)	(2,1)	(0,0)
M	(1,1)	(1,1)	(5,0)
B	(0,1)	(4,2)	(0,1)

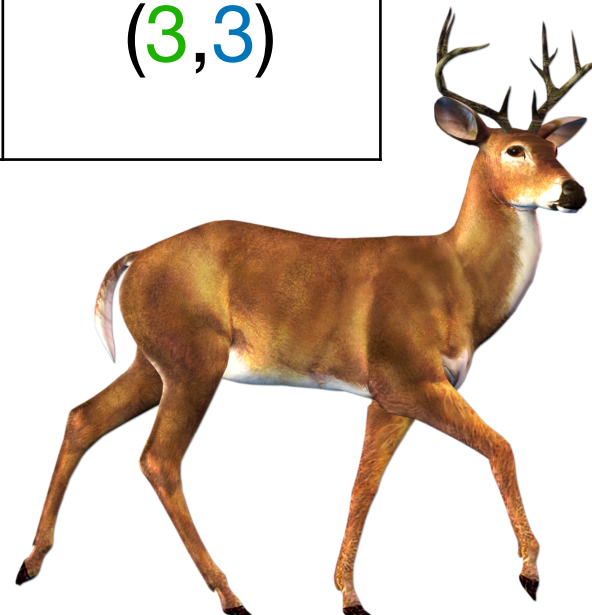
More on IESDS

- The order in which strictly dominated strategies are eliminated does not matter
- We say that a game is **solvable** when IESDS results in a unique strategy profile

Return to the Stag Hunt

		<u>Hunter 2</u>	
<u>Hunter 1</u>		Stag	Hare
Stag		(10,10)	(0,6)
Hare		(6,0)	(3,3)

- Are there any strictly dominated strategies to eliminate here?
 - No



Not all games are solvable...

Weak Dominance

Definition (Weak Dominance Among Strategies)

- Strategy s_i weakly dominates s'_i if:
 - $u_i(s_i, s_{-i}) \geq u_i(s'_i, s_{-i})$, for any profile s_{-i} of other agents' strategies, and
 - there is some strategy profile \tilde{s}_{-i} such that
$$u_i(s_i, \tilde{s}_{-i}) > u_i(s'_i, \tilde{s}_{-i})$$
- Strategy s_i is weakly dominant if it weakly dominates every other strategy s'_i

Example

	L	C	R
T	(3,0)	(2,1)	(0,0)
M	(1,1)	(1,1)	(5,0)
B	(0,1)	(4,2)	(0,1)

- For Player 2, does C weakly dominate L?
 - If P1 plays T: $1 > 0$ ✓
 - If P1 plays M: $1 = 1$ ✓
 - If P1 plays B: $2 > 1$ ✓
 - Yes

More on Weakly Dominated Strategies

- We can eliminate weakly dominated strategies, just like we eliminate strictly dominated strategies
- This time, though, the order in which we eliminate matters

This will be further explored in the homework...

Pareto Optimality

- In a **Pareto optimal outcome** no one can be made better off without making someone else worse off

Definition (Pareto Optimality)

- Strategy profile s **Pareto dominates** strategy profile s' if:
 - $u_i(s) \geq u_i(s')$ for every player $i \in N$, and
 - there exists an agent j such that $u_j(s) > u_j(s')$
- Strategy profile s is **Pareto optimal** if there is no profile s' that Pareto dominates s

The Coordination Game

- A country with no traffic rules
- Two cars on the road, driving towards each other
- They have to decide what side of the road to take
- If they choose the same side, all is well
- If they choose different sides, they crash into each other



The Coordination Game

What are the Pareto optimal outcomes?

	Left	Right
Left	(1,1)	(0,0)
Right	(0,0)	(1,1)

Pareto optimal

Pareto dominated
by (Left, Left) and
(Right, Right)



More on Pareto Optimality

- Pareto optimal outcomes always exist
- They are not always unique (there can be more than one)
- Pareto optimality does not necessarily imply that outcomes are fair

		Farmers		
		Feudalism	Capitalism	Communism
Land-owners	Feudalism	(90,10)	(5,5)	(5,5)
	Capitalism	(5,5)	(70,30)	(5,5)
	Communism	(5,5)	(5,5)	(50,50)

Pareto optimal