

Antivirus developers continuously incorporate increasingly sophisticated functions and algorithms to detect the latest malware and its evolving variants.

However, there are still straightforward methods that can bypass many of these systems, particularly those that rely solely on signature-based detection and lack advanced heuristic analysis or behavioral scanning."

Created: Sunny thakur

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Chapter 1

Introduction

Over the past decade, antivirus technologies have significantly advanced, evolving from simple signature-based scanners to incorporating more sophisticated heuristic algorithms. Modern antivirus software can scan not only files stored on the hard drive but also opcodes in memory. Opcodes, which are essentially assembly-level commands, represent the lowest level of instructions executed by the CPU for any running application. Typically, programs are written in high-level languages like C or C++, where opcodes are abstracted away. However, during the compilation process, the high-level code is translated into opcodes tailored to the system's architecture.

When an antivirus program scans a file, it reads the memory offsets and their corresponding values. An offset refers to a specific memory address, while the value is an opcode, which can be examined by the antivirus through a binary hex viewer. This process allows the antivirus to search for known malware signatures. If a file successfully passes this scan without being subjected to heuristic sandboxing, it is either deemed safe to execute, or the antivirus software has been effectively bypassed.

This paper will explore several methods and techniques that can be employed to achieve such bypassing

This is for purposes



educational only.

Chapter 2

PE File Structure

A typical PE aka Portable Executable which is the default file format for Windows binaries looks like the picture below. It should be mentioned that not all binaries has all these 5 sections. Sometimes it's only 4 or perhaps 6-7 sections, depending on how the binary is built / made.

The signature which triggers the Anti-Virus application can be located anywhere, though usually it is within one of the actual sections and not section table headers, DOS header, DOS stub etc.

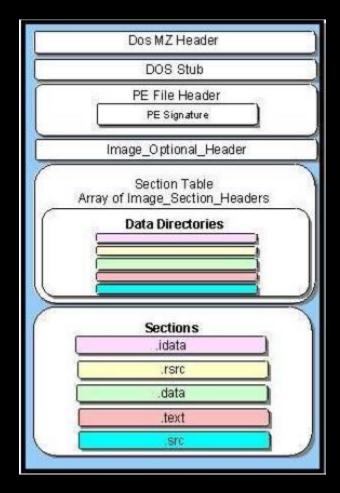


Figure 2.1 – PE File Visualization

2.1 – AV Signatures and the PE file format

Identifying the specific signature that an antivirus application detects is relatively straightforward when using an older technique, which involves splitting the file into multiple parts and scanning each one individually to determine where the signature resides. In some cases, the signature is easy to locate. For instance, if the executable file **ncx99.exe** is used—a basic netcat listener that binds **cmd.exe** to port 99 on the global network interface—the signature becomes apparent. As shown in the example, the main signature can be found between offset **E77E** and **E78F**.

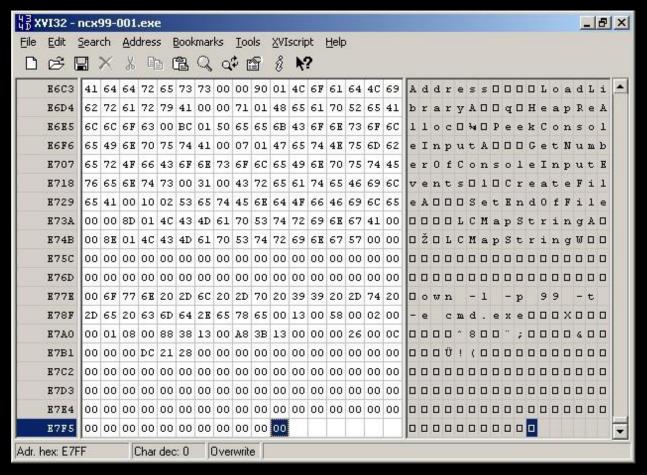


Figure 2.1.1 – Hexadecimal View of a Binary File

In this instance, the signature is located within the .idata section of the executable. This presents a challenge, as encoding the entire .idata section could render the executable non-functional. Instead, a more effective approach would be to modify a portion of the section or encode only the specific signature to bypass antivirus detection.

It's also important to remember that antivirus applications inspect the PE (Portable Executable) headers to determine whether a file is malicious. Even subtle details, such as the time and date stamp within the file, can form part of the signature, making it advisable to modify or replace these values with null. Furthermore, if any section tables, headers, or flags appear invalid to the antivirus scanner, the file might be flagged as malicious or potentially harmful, as the scanner may interpret this as an indication that it cannot properly read the executable

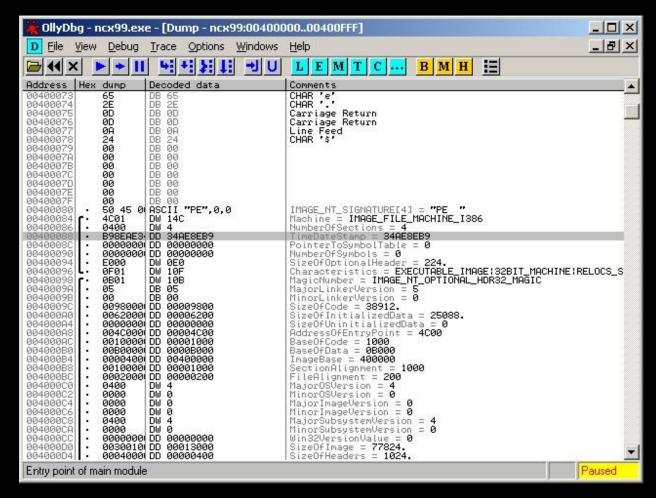


Figure 2.1.2 – Partial View of the PE Header in Ollydbg

2.2 – Modifying AV Signatures in PE Files

Once a signature is identified within one of the sections, it is typically possible to modify it either by directly editing the file with a hex editor, altering the opcodes with a disassembler, or even using a simpler tool such as a debugger (e.g., Ollydbg). In the case of ncx99.exe, as previously mentioned, it is possible to change both the listening port and the program it executes. While altering it to run a harmless application like calc.exe would be ineffective for an attacker, changing the listening port from 99 to a less obvious one, such as 81, could evade detection. Although this technique may not bypass most modern antivirus solutions, a few might still be fooled by it."



Figure 2.2.1 – ncx99.exe – Original (Binds cmd.exe to port 99)



Figure 2.2.2 – ncx99.exe – Modified (Binds cmd.exe to port 81)

As you can see, Avast and Ikarus were bypassed. If we were to attack a computer which used one of these, then we would've succeeded now just by changing the listening port.

2.3 – Polymorphic Techniques and Hijacks

Polymorphic Techniques

Polymorphic viruses maintain the same functionality but vary in their opcode sequences. This technique, often employed by more advanced attackers, makes detection more difficult. For instance, instead of using a simple PUSH -1 instruction, the hacker might implement a sequence like DEC ESI, PUSH ESI, INC ESI—assuming the ESI register is initially set to zero. If ESI is not zero, the value can be temporarily saved by pushing it onto the stack, then zeroed out using XOR ESI, ESI, enabling the hacker to push the value -1 onto the stack. The original value of ESI would

then be restored by popping it from the stack.

This example illustrates how variations in opcode sequences can be used to achieve the same result while evading detection. Although a lone PUSH -1 instruction may not necessarily trigger an AV detection, encountering a signature in executable code that cannot simply be replaced with NOPs (no operation instructions) may require the use of encoding techniques or "polymorphic methods" to bypass detection.

Hijacking

When encoding is necessary, one technique involves hijacking the entry point of the binary file. This can be done by modifying the PE headers or overwriting the initial instruction with a jump (JMP) to a "code cave"—an unused section of the binary where custom code can be inserted without altering the file size. However, it's important to note that some antivirus programs check file sizes as part of their scanning process.

Alternatively, a hacker can modify instructions deeper within the binary. As long as the original functionality of the program is preserved—ensuring it executes without crashing or encountering

errors—there is flexibility in where edits can be applied within the executable •

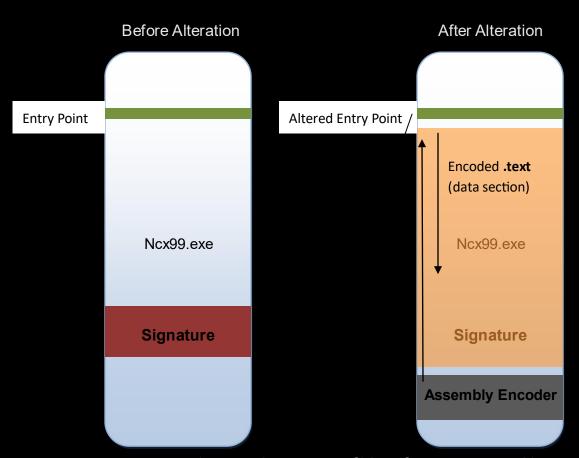
Chapter 3

Encoding Binary Files

Hijacking the entry point of an executable is a frequently employed technique, but it does not inherently bypass antivirus (AV) applications. Rather, it enables an attacker to manipulate the

execution flow of the program as desired. For instance, this could be used to divert the program's execution to a "black hole" intended to deceive heuristic detection systems, or to an encoder designed to evade signature-based AV scanning mechanisms.

The figures below illustrate a comparison between a typical PE (Portable Executable) file and an encoded PE file, highlighting how such modifications can alter the structure and potentially bypass AV signatures.



3.1 – Preparing the PE file for Encoding

First we open our chosen PE file in our favorite disassembler and debugger. In this case we will use Ollydbg to alter the previously mentioned ncx99.exe backdoor. Keep in mind that you don't have to be an Assembly programmer in order to do nor understand this.

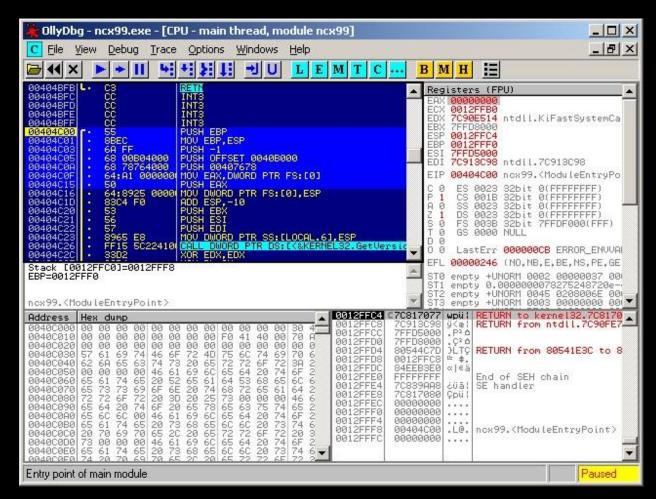


Figure 3.1.1 – Initial Overview of ncx99.exe

First we select the first couple of instructions (opcodes) and copy them to notepad or whatever program we prefer for taking notes. The reason why we're doing this is because we need to reintroduce some of the first overwritten opcodes later on, before we re-route the execution flow back to its original place.

We now have the following opcodes saved which we will need later on:

Address	Нех	dump	Command
00404C00			PUSH EBP
00404C01			MOV EBP, ESP
00404C03			PUSH -1
00404C05			PUSH OFFSET 0040B000
00404C0A			PUSH 00407678
			MOV EAX, DWORD PTR FS:[0]
00404C15		50	PUSH EAX

Figure 3.1.2 – First couple of opcodes inside ncx99.exe

Then we browse through the binary, for a convenient place to implement our custom encoder. After searching for a while inside the .text data section, we may find the following place.

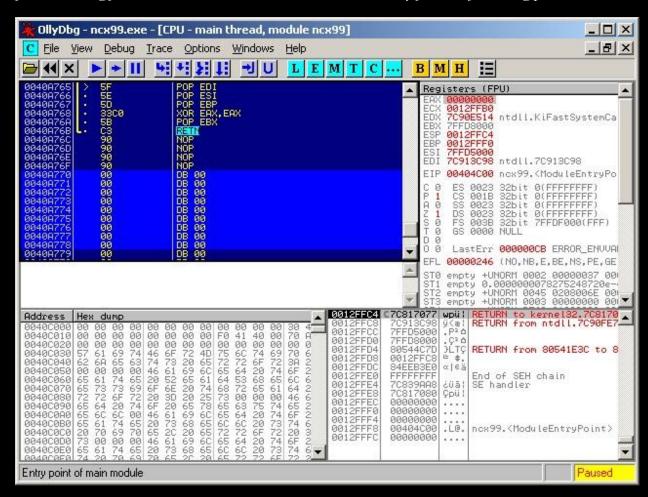


Figure 3.1.3 – A convenient place for a "code cave" in ncx99.exe

After we've noted down the new Entry Point address at 0040A770, we browse to the memory overview by clicking the "M" icon. Then we double-click on the PE Header section and open it.

Simply, because we need to prepare the .text data section by making it writeable and of course, change the old Entry Point to our new one, which points to our "code cave". Adding a few bytes extra to the last section doesn't hurt either, as this may bypass some AV-scanners.

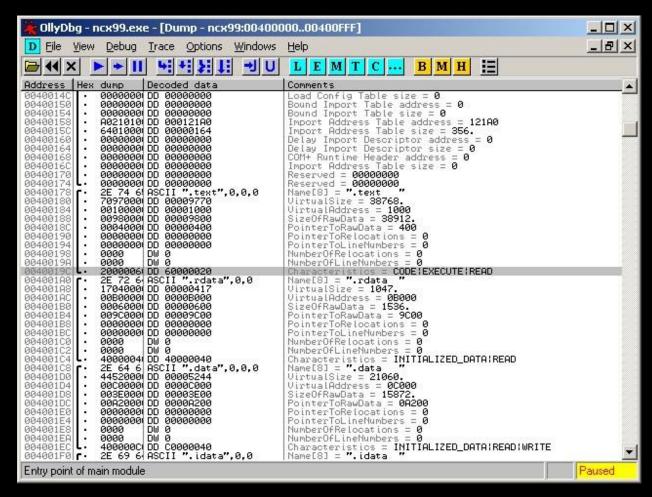


Figure 3.1.4 – PE Header Overview of ncx99.exe

This is the value we need to edit in order to be able to make the .text section writeable. We could use LordPE for this, but knowing the common values by mind, makes us able to do this without.

We will therefore change 60000020 to E0000020 as shown in the next picture, making the .text section writeable, allowing us to encode this while the PE file is executing. If we didn't do this we would get a "permission error" and most likely crash.

Adding a few bytes to one of the sections is a good idea too, if you don't need to be very strict on keeping exactly the same file-size. This is done by taking the hexadecimal value of e.g. the .idata section and then add the number of bytes wanted in hex.

If you're going to do this, then make sure you're calculating in hex and not decimal.

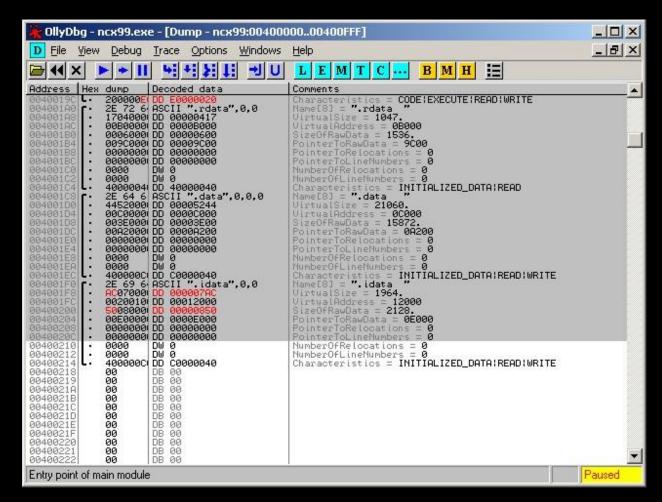


Figure 3.1.5 – Altered PE Header in ncx99.exe

After we've made our modifications, we select the entire section which contains our changes, right click and browse to "Edit", and then "Copy to Executable". Then right click on the new window and choose "Save File".

Keep in mind that this is a bit different in the older version of Ollydbg.

Because we've added a few bytes to the .idata section, the program won't execute. Therefore we need to add the amount of bytes we "added", by using a hex-editor to add the actual amount of bytes that was added to the .idata section.



Figure 3.1.6 – Altered nex99.exe unable to execute

In this case XVI32 (a hex-editor) is sufficient to use. Simply browse to the end of the PE file, open the "Edit" menu and choose "Insert string". Then make sure you either add 00, 90 or CC.

Preferably just add 00 as this will do nothing at all. Under "Insert <n> times" you choose hexadecimal and choose the amount of bytes you added to the .idata section. When you're done click the "save" icon and your executable PE file, should be working again.

3.2 – Implementing the Custom Encoder

With all the preparations made, we're ready to implement the encoder. First we open our modified PE file in Ollydbg and see that we've landed at 0040A770. After taking a closer look on where the base (beginning) address is and where our code cave begins, we note down that from offset 00401000 to 0040A76F, is what we'll encode.

There are many ways to implement an encoder, but the easiest way is the one Mati Aharoni from Offensive Security did in his public video presentation about AV's. The encoder we're going to implement is slightly different in order to hopefully confuse a few more Anti-Virus scanners.

We'll basically encode almost the entire .text section, with an encoder which loops through each byte of the selected code that we want to encode. The encoding mechanism itself will just change the byte to whatever we tell it to become.

Address	Нех	dump	Command
0040A770			
0040A775			
0040A778			
0040А77В			

Figure 3.2.1 – Custom Assembly Encoder

Explanation of the Custom Encoder

First the base (beginning) address is moved into the EAX register.

Then it adds 13 to the byte which EAX is pointing to.

XOR (Exclusive OR) the byte with 0F which EAX is pointing to.

Add 37 to the byte which EAX is pointing to.

Increase EAX to point to the next byte.

Compare EAX with our ending address.

If our ending address hasn't been reached, jump to (2).

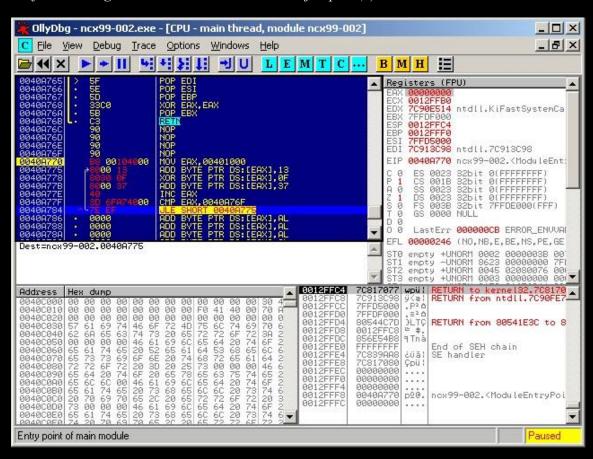


Figure 3.2.2 – Custom Encoder inside Ollydbg

With our custom encoder implemented we're almost done. It should be noted, that we could also use other opcodes too, to encode our .text section. Such opcodes could be: sub, or, and, etc.

But for now we're going to re-introduce some of the first few opcodes that was originally run by the executable in the start. Now we could be simple and just place a jump to 00404C00, but we'll add the first couple of instructions ourselves as shown in the picture below.

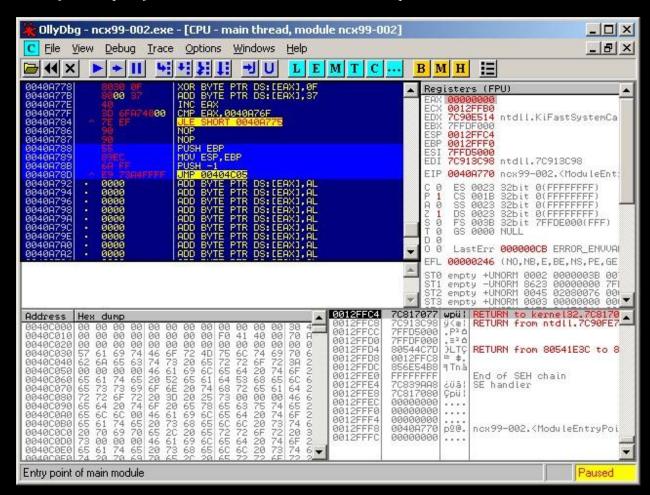


Figure 3.2.3 – Re-Introduced Opcodes in ncx99.exe

With that done you may wonder what the **JMP 00404C05** opcode is. That is a jump to the offset aka memory address where the next instruction after the original **PUSH -1** is located.

It should be noted that if we were to execute the file now, it would simply fail because the PE file as it is right now, will encode the .text section and try to execute it. Since it becomes encoded, then it will most likely fail and crash. But save the changes anyway and re-open it.

This is because we first need to use our encoder, to encode the file and afterwards change it to a decoder, so the execution flow will seem completely normal even without a debugger.

Chapter 4

Decoding Binary Files

After we've successfully implemented our encoder we need to save the encoded contents and then change our encoder to a decoder as previously mentioned. This is relatively simple as you will experience yourself.

When we hit the first instruction which moves 00401000 into the EAX register, we can rightclick this and select "Follow in Dump". By pressing "F7" on your keyboard we can single-step through the custom encoder, and watch our .text section become encoded.

To speed up this process, select the instruction right after the **JLE SHORT 0040A775** opcode, place a breakpoint by pressing "F2" on your keyboard and then press "F9" to execute until the breakpoint stops the execution flow.

Some of the code and even your encoder may seem completely different now. This is because the encoder altered all the opcodes in the .text data section, except your encoder even though it may seem so. Copy the changes to an "executable" and save it as a new file.

4.1 – Altering the Encoder to a Decoder

Now open the newly created file and look at the custom encoder we implemented earlier.

As you can see almost all of the opcodes are "gibberish" and therefore we might have to press CTRL+A to do a quick analysis in order to get our custom encoder visible. When it appears to look like it should, we change two of the opcodes in order to make it a decoder.

Our initial encoder added 13 to the byte which EAX pointed to, then it XOR'd it with 0F and added 37 to end with. Now we need to reverse this, by deducting -37 to start with and then -13.

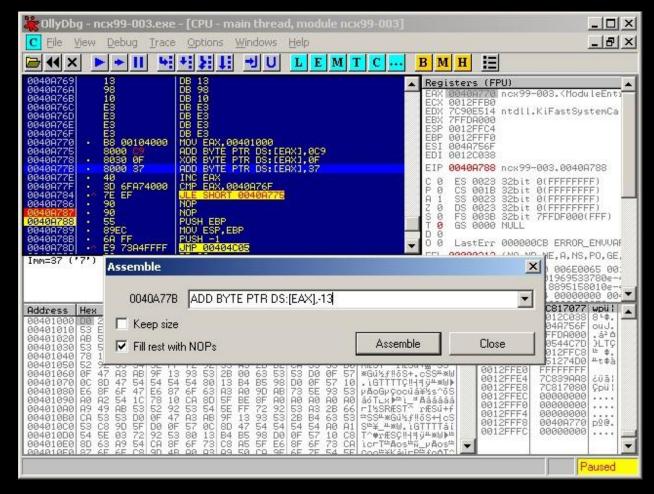


Figure 4.1.1 – Changing the Encoder to a Decoder

We will still use the **add** opcode and of course **xor**, but we'll only need to change the values as mentioned previously, which you can also see in the picture above in figure 4.1.1.

When we've done that we copy our changes to an executable, and save it as a new file.

In theory the PE file should work now just as it did to start with, but it is also encoded too making it able to bypass some AV-scanners, which makes it more interesting.

4.2 – Testing the Custom Decoder

Now it's time to test if our encoded file will decode properly and execute gracefully.

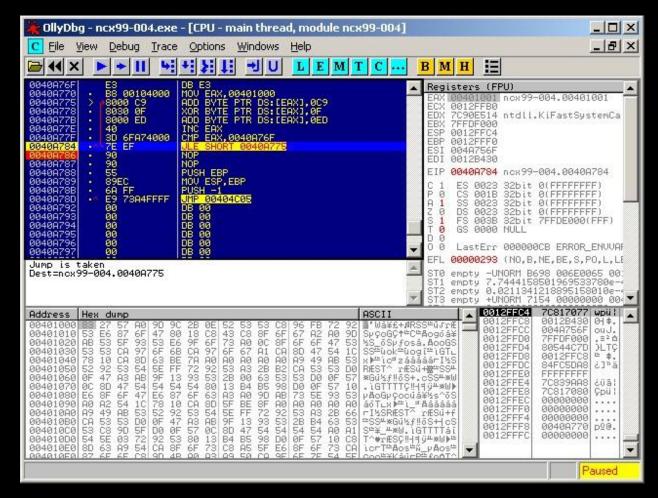


Figure 4.2.1 – Encoded Overview of ncx99.exe

In the picture above, only the first hex character has been decoded back to its original "state". (Please note that a character in this case, is a byte which is the same as 8 bits made of binary.)

By placing a breakpoint right after **JLE SHORT 0040A775**, and then running the PE file until it stops executing, we'll see that the entire .text section has changed back to its original state.

If we execute the first couple of re-introduced opcodes including the long jump back to where the real start of the program is, we'll see that it may still look obfuscated or encoded.

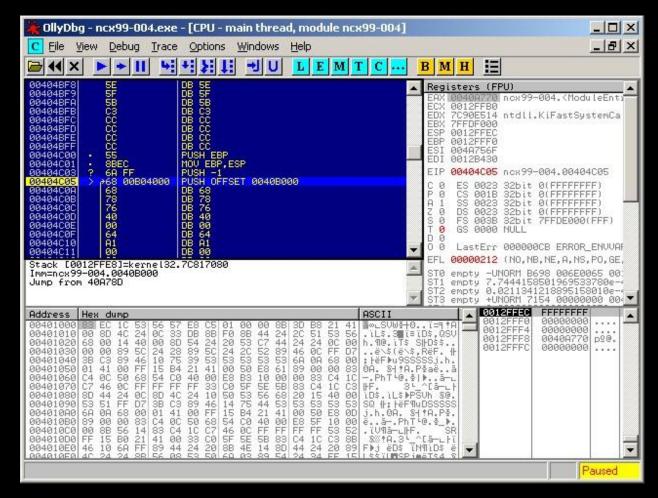


Figure 4.2.2 – The Beginning of the Decoded PE File

This does not really matter, as we can hit CTRL+A and do a quick analysis of the code. When we've done that we may see that our executable PE file is back to its original state again and if we press "F9" we may see that our program is executing without any errors.

If that is the case then we've encoded the binary file and even decoded it successfully.

Even scanning the file now with a lot of different AV-scanners will reveal different results, if the file is just scanned and not executed since we haven't implemented any bypassing techniques for heuristic (malicious) opcode detection.

Figure 4.2.3 – Netstat Overview of ncx99.exe running

If we scan the file with AVG it may be undetectable now, or as script kiddies tends to say: FUD.

This expression means "Fully Undetectable" and is widely used with tools such as cryptors. Most of these use another way of making the PE files able to bypass the AV-scanners, which is e.g. by encrypting the entire file with RC4 and then pack a stub decoder into the file.

This will of course alter the size of the file in many cases. In the picture below you'll see that many of the AV's either didn't know what was scanned or they didn't flag it as malicious at all.



Chapter 5

Conclusion

The purpose of this paper was to demonstrate how easy it is to bypass signature-based AntiVirus scanners which do not use heuristics at all or perhaps only in an insufficient way which makes a hacker able to outsmart the AV-system in use.

Antivirus	Database	Engine	Result
a-squared	18/06/2010	5.0.0.7	Backdoor.Win32.Ncx!IK
Avast	100617-0	5.0	Win32:Nex [Trj]
AVG	271.1.1/2946	9.0.0.725	BackDoor, Generic12, BNOS
Avira AntiVir	7.10.8.127	7.6.0.59	
BitDefender	18/06/2010	7.0.0.2555	Backdoor.NCX_99
ClamAV	18/06/2010	0.96.1	
Comodo	3468	3.13.579	Backdoor.IRC.SdBot.NP
Dr.Web	18/06/2010	5.0	BackDoor.Angel
F-PROT6	20100618	4.5.1.85	W32/Malwarel8370
G-Data	21.371	2.0.7309.847	Backdoor.Win32.Ncx.b A
Ikarus T3	18/06/2010	1.1.84.0	Backdoor.Win32.Ncx
Kaspersky	18/06/2010	9.0.0.736	Backdoor.Win32.Ncx.b
NOD32	5208	4.0.474	Win32/NCX.99
Panda	18/06/2010	10.0.3.0	Bck∕Vonetent
TrendMicro	251	9.120-1004	TROJ_NCX99.A
VBA32	18/06/2010	3.12.12.2	Backdoor.Win32.Ncx.b

Figure 5.1 – ncx99.exe before any modifications

As you can see almost all of the AV-scanners detects the original ncx99.exe by default.

If we alter this file heavily then the results are quite amazing. The amount of code added and changed can be as little as below 30 bytes which is the equivalent of 30 characters you can type on your keyboard. Even the size of the file may have been unaltered too, though the contents of the file may have been encoded with a custom encoder.

Antivirus	Database	Engine	Result	
a-squared	18/06/2010	5.0.0.7	Riskware. Remote Admin. Win 32. Net Cat! IK	
Avast	100617-0	5.0		
AVG	271.1.1/2946	9.0.0.725		
Avira AntiVir	7.10.8.127	7.6.0.59		
BitDefender	18/06/2010	7.0.0.2555	Trojan.Peed.Gen	
ClamAV	18/06/2010	0.96.1		
Comodo	3468	3.13.579		
Dr.Web	18/06/2010	5.0		
F-PROT6	20100618	4.5.1.85		
G-Data	21.371	2.0.7309.847	0086e790991d0c1f376a0a366c1eb7b Possibly infected	
lkarus T3	18/06/2010	1.1.84.0		
Kaspersky	18/06/2010	9.0.0.736	Type_Win32	
NOD32	5208	4.0.474		
Panda	18/06/2010	10.0.3.0		
TrendMicro	251	9.120-1004	PAK_Generic.001	
VBA32	18/06/2010	3.12.12.2		

Figure 5.2 – ncx99.exe after heavy modifications

As you can see for yourself in the picture above, a lot of the AV-scanners were bypassed. Kaspersky detected this file as potentially malicious with its heuristics system, hence the reason Type Win32 is stated which means it is probably a new variant of a virus, trojan, etc.

In any case it is always a good idea to encode the primary signature of the file, though usually there is more than one signature so it isn't piece of cake for any hacker to bypass AV-detection.

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http://debugger.immunityinc.com/

https://forum.immunityinc.com/board/show/0/

http://free.avg.com/ww-en/homepage

http://www.uninformed.org/?v=5&a=3&t=pdf

http://www.offensive-security.com