Binary Search:

- Operates on a contiguous sequence with a specified left and right index (Search Space).
- Maintains the left, right, and middle indices of the search space.
- Compares the search target or applies the search condition to the middle value of the collection:
 - if the condition is unsatisfied or values unequal, keep searching within satisfied range.
 - If the search ends with an empty half, the condition cannot be fulfilled and target is not found.

Procedure:

- 1.Pre-processing Sort if collection is unsorted.
- 2.Binary Search Using a loop or recursion to divide search space in half after each comparison.
- 3. Post-processing Determine viable candidates in the remaining space.

Notice:

- Instead of searching a specific value, Binary Search can take many alternate forms.
- Sometimes you will have to apply a specific condition or rule to determine which side (left or right) to search next.

When to Use?

• Should be considered every time you need to search for an index or element in a collection.

Q1:

Given a **sorted** (in ascending order) integer array **nums** of **n** elements and a **target** value, write a function to search **target** in **nums**. If **target** exists, then return its index, otherwise return -1

Example 1:

Input: nums = [-1,0,3,5,9,12], target = 9 **Output:** 4

Explanation: 9 exists in nums and its index is 4

Example 2:

Input: nums = [-1,0,3,5,9,12], target = 2 **Output:** -1

Explanation: 2 does not exist in nums so return -1

```
def binarySearch(nums, target):
    :type nums: List[int]
    :type target: int
    :rtype: int
    if len(nums) == 0:
        return -1
    left, right = 0, len(nums) - 1
    while left <= right:
        mid = (left + right) // 2
        if nums[mid] == target:
            return mid
        elif nums[mid] < target:
            left = mid + 1
        else:
            right = mid - 1
    # End Condition: left > right
    return -1
```

Key Syntax:

- Initial Condition: left = 0, right = length-1
- Termination: left > right
- Searching Left: right = mid-1
- Searching Right: left = mid+1

Most basic and elementary form of Binary Search !!!

```
class Solution(object):
    def search(self, nums, target):
        :type nums: List[int]
        :type target: int
        :rtype: int
        if len(nums) == 0:
            return -1
        left, right = 0, len(nums)
        while left < right:
            mid = (left + right) // 2
            if nums[mid] == target:
                return mid
            elif nums[mid] < target:</pre>
                left = mid + 1
            else:
                right = mid
        # Post-processing:
        # End Condition: left == right
        if left != len(nums) and nums[left] == target:
            return left
        return -1
```

Key Syntax:

- Initial Condition: left = 0, right = length
- Termination: left == right
- Searching Left: right = mid
- Searching Right: left = mid+1

Attributes:

- An advanced way to implement Binary Search.
- Search Condition needs to access element's immediate right neighbor
- Use element's right neighbor to determine if condition is met and decide whether to go left or right
- Gurantee Search Space is at least 2 in size at each step
- Post-processing required. Loop/Recursion ends when you have 1 element left. Need to assess if the remaining element meets the condition.

```
class Solution(object):
    def search(self, nums, target):
        :type nums: List[int]
        :type target: int
        :rtype: int
        if len(nums) == 0:
            return -1
        left, right = 0, len(nums) - 1
        while left + 1 < right:
            mid = (left + right) // 2
            if nums[mid] == target:
                return mid
            elif nums[mid] < target:
                left = mid
            else:
                right = mid
        # Post-processing:
       # End Condition: left + 1 == right
        if nums[left] == target: return left
        if nums[right] == target: return right
        return -1
```

Key Syntax:

- Initial Condition: left = 0, right = length 1
- Termination: left + 1== right
- Searching Left: right = mid
- Searching Right: left = mid

Attributes:

- An alternative way to implement Binary Search
- Search Condition needs to access element's immediate left and right neighbors
- Use element's neighbors to determine if condition is met and decide whether to go left or right
- Gurantee Search Space is at least 3 in size at each step
- Post-processing required. Loop/Recursion ends when you have 2 elements left. Need to assess if the remaining elements meet the condition.

Summary:

```
Template #1:
                                  Template #2:
                                                                     Template #3:
// Pre-processing
                                  // Pre-processing
                                                                    // Pre-processing
left = 0; right = length - 1;
                                  left = 0; right = length;
                                                                    left = 0; right = length - 1;
while (left <= right) {
                                  while (left < right) {
                                                                    while (left + 1 < right) {
mid = left + (right - left) / 2;
                                  mid = left + (right - left) / 2;
                                                                    mid = left + (right - left) / 2;
if (nums[mid] == target) {
                                   if(nums[mid] < target) {
                                                                    if (num[mid] < target) {
  return mid:
                                     left = mid + 1;
                                                                       left = mid:
 } else if(nums[mid] < target) {
                                   } else {
                                                                     } else {
  left = mid + 1:
                                                                       right = mid;
                                     right = mid;
} else
  right = mid - 1;
                                  // left == right
                                                                    // left + 1 == right
// right + 1 == left
                                  // 1 more candidate
                                                                    // 2 more candidates
// No more candidate
                                  // Post-Processing
                                                                    // Post-Processing
```

- 99% of binary search problems that you see online will fall into 1 of these 3 templates.
- Template 1 and 3 are the most commonly used and almost all binary search problems can be easily implemented in one of them.

Q2:

Implement int sqrt(int x).

Compute and return the square root of x, where x is guaranteed to be a non-negative integer.

Since the return type is an integer, the decimal digits are truncated and only the integer part of the result is returned.

Example 1:

Input: 4

Output: 2

Example 2:

Input: 8

Output: 2

Explanation: The square root of 8 is 2.82842..., and since

the decimal part is truncated, 2 is returned.

```
class Solution(object):
    def mySqrt(self, x):
        :type x: int
        :rtype: int
        1, r = 1, x
        while 1 <= r:
            mid = (1 + r)/2
            if mid*mid == x:
               return mid
            elif mid*mid < x:
                l = mid + 1
            else:
                r = mid - 1
        return r
```

Q3:

First Bad Version:

You are a product manager and currently leading a team to develop a new product. Unfortunately, the latest version of your product fails the quality check. Since each version is developed based on the previous version, all the versions after a bad version are also bad.

Suppose you have n versions [1, 2, ..., n] and you want to find out the first bad one, which causes all the following ones to be bad.

You are given an API bool isBadVersion(version) which will return whether version is bad

Example 1:

Input: 8

1	2	3	4	5	6	7	8
good	good	good	good	bad	bad	bad	bad

Output: 5

```
class Solution(object):
    def firstBadVersion(self, n):
        :type n: int
        :rtype: int
        11 11 11
        1, r = 1, n
        while 1 < r:
            mid = (1 + r)/2
            if not isBadVersion(mid): #Good Version
                l = mid + 1
            else:
                r = mid
        return 1
```

Q4: Find K Closest Elements

Given a sorted array, two integers k and x, find the k closest elements to x in the array. The result should also be sorted in ascending order. If there is a tie, the smaller elements are always preferred.

Example 1:

Input: [1,2,3,4,5], k=4, x=3

Output: [1,2,3,4]

Example 2:

Input: [1,2,3,4,5], k=4, x=-1

Output: [1,2,3,4]

Analysis:

The final result should be a list of *k* consecutive elements. i.e. *arr[start: start + k]*How to find this *start* point?
Using binary search by comparing a and b:

- a = x arr[start]
- b = arr[start + k] x
- If a > b: start point is too left, left->mid + 1
- If a < b: start point is too right, righr ->mid

```
class Solution(object):
    def findClosestElements(self, arr, k, x):
        :type arr: List[int]
        :type k: int
        :type x: int
        :rtype: List[int]
        left, right = 0, len(arr) - k
        while left < right:
            mid = (left + right) / 2
            if x - arr[mid] > arr[mid + k] - x:
                left = mid + 1
            else:
                right = mid
        return arr[left:left + k]
```

Q5: Search in Rotated Sorted Array

Suppose an array sorted in ascending order is rotated at some pivot unknown to you beforehand. (i.e., [0,1,2,4,5,6,7] might become [4,5,6,7,0,1,2]).

You are given a target value to search. If found in the array return its index, otherwise return -1.

You may assume no duplicate exists in the array.

Example 1:

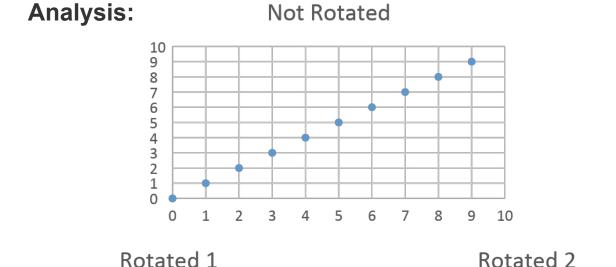
Input: nums = [4,5,6,7,0,1,2], target = 0

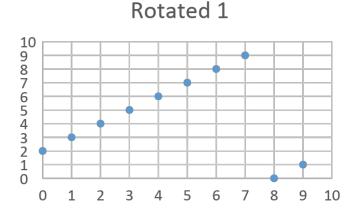
Output: 4

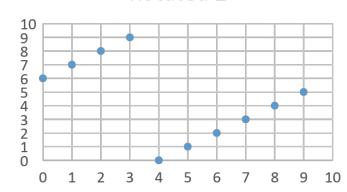
Example 2:

Input: nums = [4,5,6,7,0,1,2], target = 3

Output: -1







```
class Solution(object):
    def search(self, nums, target):
        :type nums: List[int]
        :type target: int
        :rtype: int
        low = 0
        high = len(nums) - 1
        while low <= high:
            mid = (low + high)/2
            if nums[mid] == target:
                 return mid
            if nums[low] <= nums[mid]:</pre>
                 if nums[low] <= target <= nums[mid]:</pre>
                     high = mid - 1
                 else:
                     low = mid + 1
            else:
                 if nums[mid] <= target <= nums[high]:</pre>
                     low = mid + 1
                 else:
                     high = mid - 1
        return -1
```

In rotated cases: find monotone increasing range and JUDGE!!!