Homework XI

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Abstract

This is Daniel's homework of "Numerical Algorithms with Case Studies II".

1 Problems

Q1 We just verify this assumption with n eigenvectors $x_j = \begin{pmatrix} 1 & \omega_n^j & \omega_n^{2j} & \cdots & \omega_n^{(n-1)j} \end{pmatrix}^T$ where $j = 0, 1, 2, \cdots, n-1$ and $\omega_n = e^{-i\frac{2\pi}{n}}$

$$\begin{pmatrix} 0 & 1 & & & \\ & 0 & 1 & & \\ & & \ddots & \ddots & \\ & & & 0 & 1 \\ 1 & & & & 0 \end{pmatrix} \begin{pmatrix} 1 \\ \omega_n^j \\ \omega_n^{2j} \\ \ddots \\ \omega_n^{(n-1)j} \end{pmatrix} = \begin{pmatrix} 1 \\ \omega_n^j \\ \omega_n^{2j} \\ \vdots \\ \omega_n^{(n-1)j} \end{pmatrix} w_n^j$$

$$(1)$$

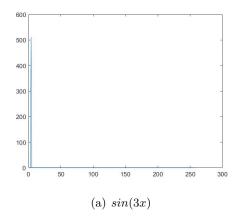
And so we will see that,

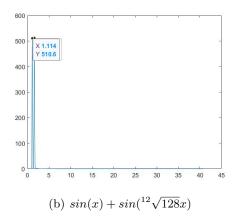
$$F_n^{-1} J_n F_n = diag\{1, \omega_n^1, \omega_n^2, \cdots, \omega_n^{(n-1)}\}$$
 (2)

Q2 First of all, we will explain why two-sided power spectrum appear on the plot. Discrete Fourier transformation can be interpreted as follows.

$$X(f) = \sum_{k=0}^{N} x(k)e^{-i\frac{2\pi fk}{N}} = \sum_{k=N}^{0} x(N-k)e^{-i\frac{2\pi f(N-k)}{N}} = \sum_{k=0}^{N} x(k)e^{i\frac{2\pi fk}{N}} = X(-f) = X(N-f) \quad (3)$$

Eliminate the higher mirrored copy of peaks. And check it with Matlab plots, only one peak for sin(3x), while $sin(x) + sin(\sqrt[12]{128}x)$ has two peaks.





Q3 The convolution theorem based on DFT is describe as follows,

$$\widehat{u * v_k} = \widehat{u}_k \cdot \widehat{v_k} \tag{4}$$

Represent $\widehat{u*v_k}$ using DFT formula,

$$(u * v)_k = \sum_{j=0}^{N-1} u(j) \cdot v(k-j)$$
 (5)

$$= \sum_{j=0}^{N-1} \left(\frac{1}{N} \sum_{n=0}^{N-1} U(n) \cdot e^{i\frac{2\pi nj}{N}}\right) \left(\frac{1}{N} \sum_{m=0}^{N-1} V(m) \cdot e^{i\frac{2\pi m(k-j)}{N}}\right)$$
(6)

$$= \frac{1}{N} \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} U(n)V(m) \cdot e^{i\frac{2\pi mk}{N}} \left(\frac{1}{N} \sum_{j=0}^{N-1} e^{i\frac{2\pi(n-m)j}{N}}\right)$$
 (7)

And we also have the following Lemma that,

$$\sum_{i=0}^{N-1} e^{i\frac{2\pi(n-m)j}{N}} = \sum_{i=0}^{N-1} \left(e^{i\frac{2\pi(n-m)}{N}}\right)^j \tag{8}$$

$$= \frac{1 - \left(e^{i\frac{2\pi(n-m)}{N}}\right)^N}{1 - e^{i\frac{2\pi(n-m)}{N}}} \tag{9}$$

$$= \begin{cases} 0 & \text{if } m \neq n \\ 1 & \text{if } m = n \end{cases} \tag{10}$$

Therefore, we also have,

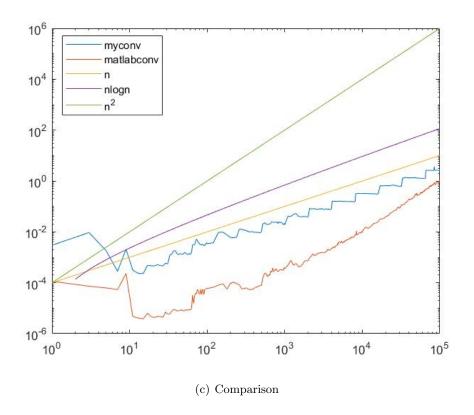
$$(u*v)_k = \frac{1}{N} \sum_{n=0}^{N-1} \sum_{m=0}^{N-1} U(n)V(m) \cdot e^{i\frac{2\pi mk}{N}} \left(\frac{1}{N} \sum_{j=0}^{N-1} e^{i\frac{2\pi(n-m)j}{N}}\right)$$
(11)

$$= \frac{1}{N} \sum_{n=0}^{N-1} U(n)V(n) \cdot e^{i\frac{2\pi nk}{N}}$$
 (12)

$$=\widehat{\widehat{u}_k\cdot\widehat{v}_k}\tag{13}$$

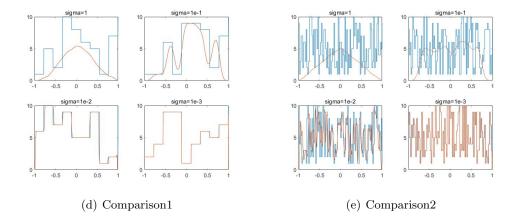
Which means $\widehat{u*v_k} = \widehat{u}_k \cdot \widehat{v}_k$

Q4 The correctness of myconv() is guaranteed by conv(). We use tic and toc in Matlab to examine the time consumption of conv() and myconv().

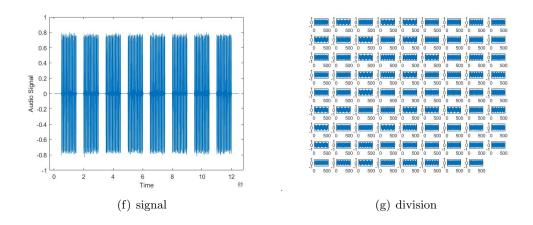


Surprisingly, the Matlab conv() is $O(n^2)$ level algorithm, while myconv() is indeed approximately O(nlogn). It's natural to think that myconv() will definitely perform better than conv() for large problems.

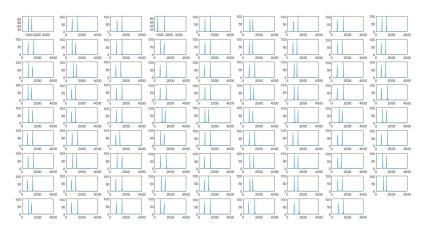
Q5 We construct wierd() to generate a insanely and randomly discontinuous function. Use $\sigma=\{1,10^{-1},10^{-2},10^{-3}\}$ from Gaussian function family $\phi(x)=\frac{1}{\sqrt{2\pi}\sigma}e^{-\frac{x^2}{2\sigma^2}}$



Q6 The audio file and every single tone is divided into small parts using some trivial techniques.

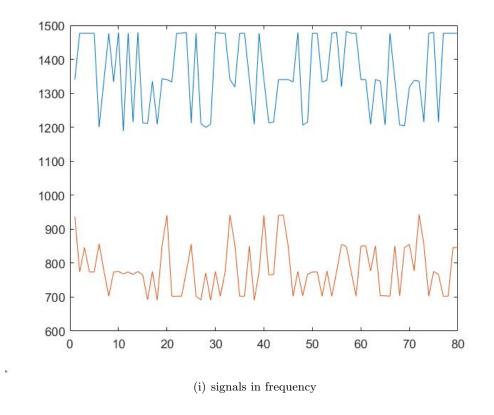


Using fft(). Note that the frequency corresponds to the signal frequency, which is $j \cdot fs$.



(h) signal FFT

The higher part of frequency can be ignored because they are simply mirror copies of spikes. Determine the two peaks for every sound, and look up for approximation in the table, we have,



The recovered keys are " $0696675356\ 4646415180\ 2336731416\ 3608338160\ 4400826146\ 6253689638\ 8482138178\ 5073643399$ ".

Q8 The comparison of myfft3() and myfft5() with fft() in Matlab.

