# HMIN103 Presentation des données du Web : XML

Federico Ulliana UM, LIRMM, INRIA GraphIK Team

Slides collected from J. Cheney, S. Abiteboul, I. Manolescu, P. Senellart, P. Genevès, D. Florescu, and the W3C

#### Intervenants

Pierre Pompidor

- Federico Ulliana <u>ulliana@lirmm.fr</u>
  - Accueuil étudiants lundi 16h au LIRMM (bat.5 salle 3-130; rdv par mail)

## Programme: XML 360°

Objectif du cours : étudier le langage XML et voir ses applications principales

La famille des langages XML (et relatifs)
 XML, DTD, XPath/XQuery, Updates, JSON

Présentation de données XML
 XSLT, HTML5/Js, WebGL

#### MCC

•50% examen

(2 sessions)

•25% Projet

(P. Pompidor)

•25% TPs à rendre

(F. Ulliana)

(I session)

•Note TP finale: moyenne des notes attention aux zeros! (documents non consignés)

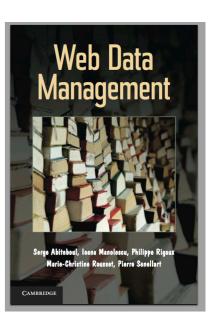
#### Dates Rendus de TP

- Données et Schémas XML (27/09)
- XPath et XQuery (11/10)
- XML <-> Relationnel (18/10)
- XML Updates (25/10)

• Informations sur Moodle, en cas de problème : <u>ulliana@lirmm.fr</u>

### Readings

Web Data Management - Abiteboul & al.



[WDM-XML] Chapter: Data-model

http://webdam.inria.fr//Jorge/files/wdm-datamodel.pdf

[WDM-DTD] Chapter: Schemas (only section 3)

http://webdam.inria.fr//Jorge/files/wdm-typing.pdf

#### What is XML?

eXtensible Markup Language [W3C 1998]

Ask five different people, get five different answers...

- a self-describing data format?
- a generalization of HTML?
- part of the DNA of computer science?
- "best thing since sliced bread?"
- a meta-language?

## If XML is the solution, then what was the problem?

#### Web data = by far the largest information system ever seen

- Billions of textual documents, images, PDF, multimedia files, provided and updated by millions of institutions and individuals.
- An anarchical process which results in highly heterogeneous data organization, steadily growing and extending to meet new requirements.
- New usage and applications of communication appear every day: yesterday P2P file sharing, today social networking, tomorrow?

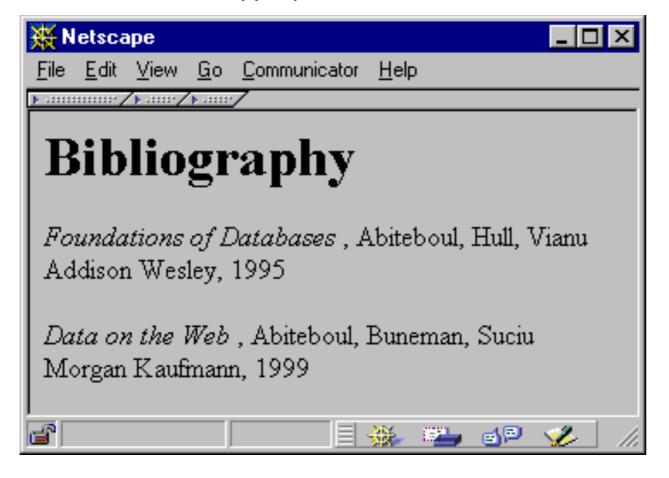
#### Challenges

- Master the size and extreme variability of Web information, and make it usable.
- Ensure long-term access to data. Write documents in 1998 and read in 2087.

#### The role of XML

## Web data management has been primarily based on HTML, which describes presentation

appropriate for humans, but falls short when it comes to software exploitation of data



```
<html>
<html>
<html>
<hl>
<hl>
<hl>
Bibliography </hl>
 <i>> Foundations of Databases </i>
Abiteboul, Hull, Vianu

<br>
<br>
Addison Wesley, 1995</br>
 <i>> Data on the Web </i>
Abiteboul, Buneman, Suciu

<br>
<br>
<br>
Abiteboul, Buneman, 1999</br>
</html>
```

#### The role of XML

XML describes content, and promotes machineto-machine communication and data exchange

```
<br/>
<book>
<title> Foundations... </title>
<author> Abiteboul </author>
<author> Hull </author>
<author> Vianu </author>
<publisher> Addison Wesley
</publisher>
<year> 1995 </year>
</book>...
</bibliography></br/>
```

```
<html>
<html>
<h1> Bibliography </h1>
 <i> Foundations of Databases </i>
    Abiteboul, Hull, Vianu
    <br>
    Addison Wesley, 1995
 <i> Data on the Web </i>
    Abiteboul, Buneman, Suciu
    <br>
    Morgan Kaufmann, 1999
</html>
```

### XML for Data Exchange

### Web data exchange

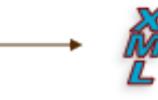
 "Data exchange is the oldest database problem" (Phil Bernstein, 2003)

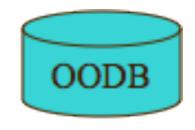
#### Massive demand

- across platforms/DBs
- across enterprises



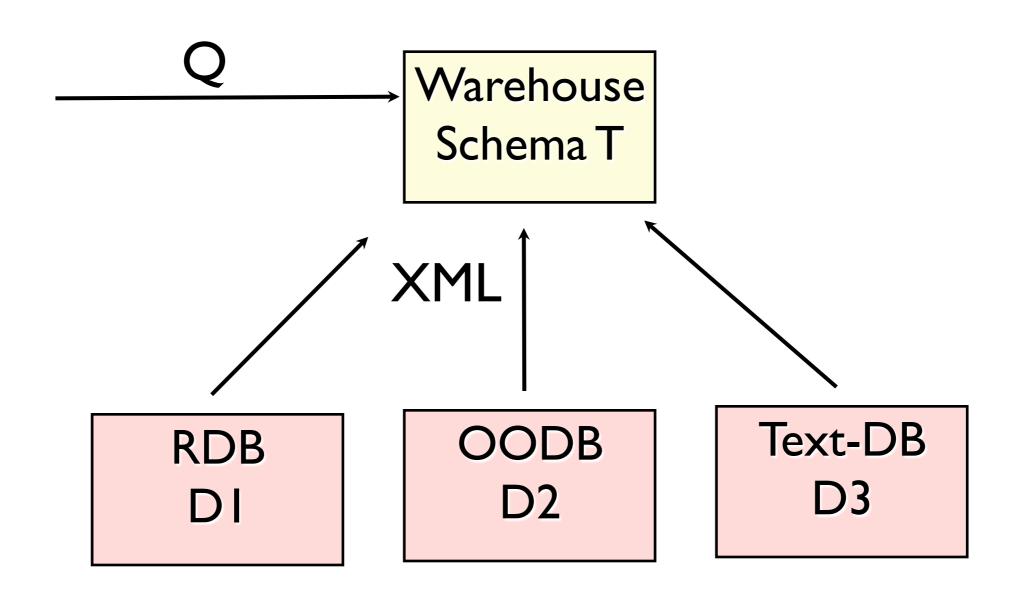






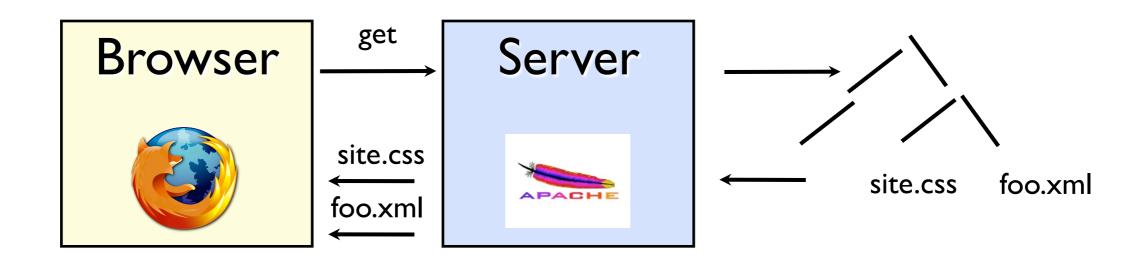
 XML has become the prime standard for data exchange on the Web

### Data integration



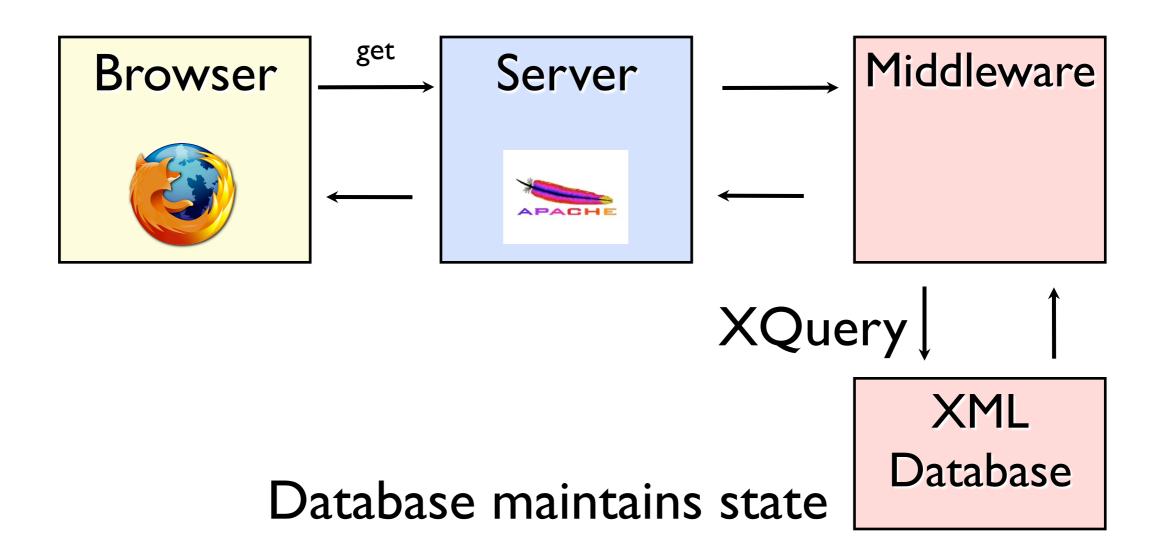
### XML for Web Applications

## Static Web site (XML+ CSS)

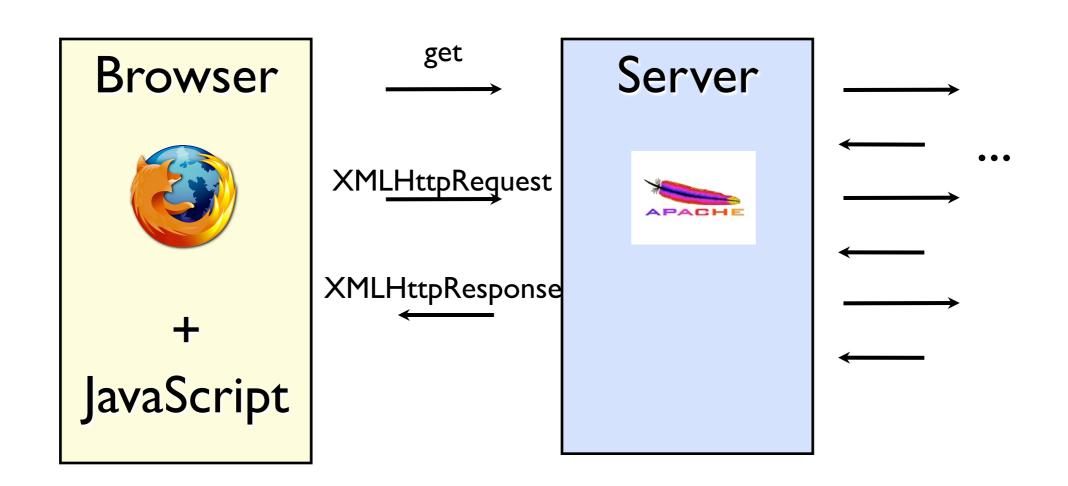


allows better factoring into data + presentation

### Dynamic Web site



## Asynchronous Javascript and XML



## XML support in industry

- Most commercial RDBMSs now provide some XML support
  - Oracle I Ig XML DB



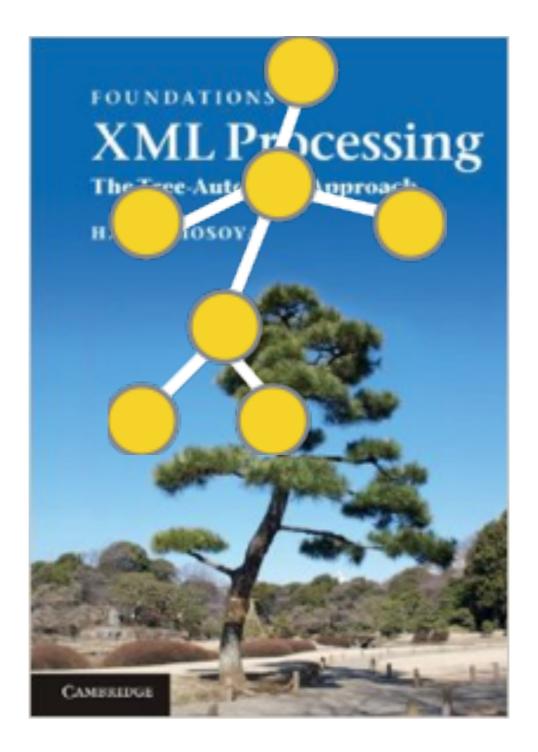
- IBM DB2 pureXML
- Microsoft SQL Server XML support since 2005

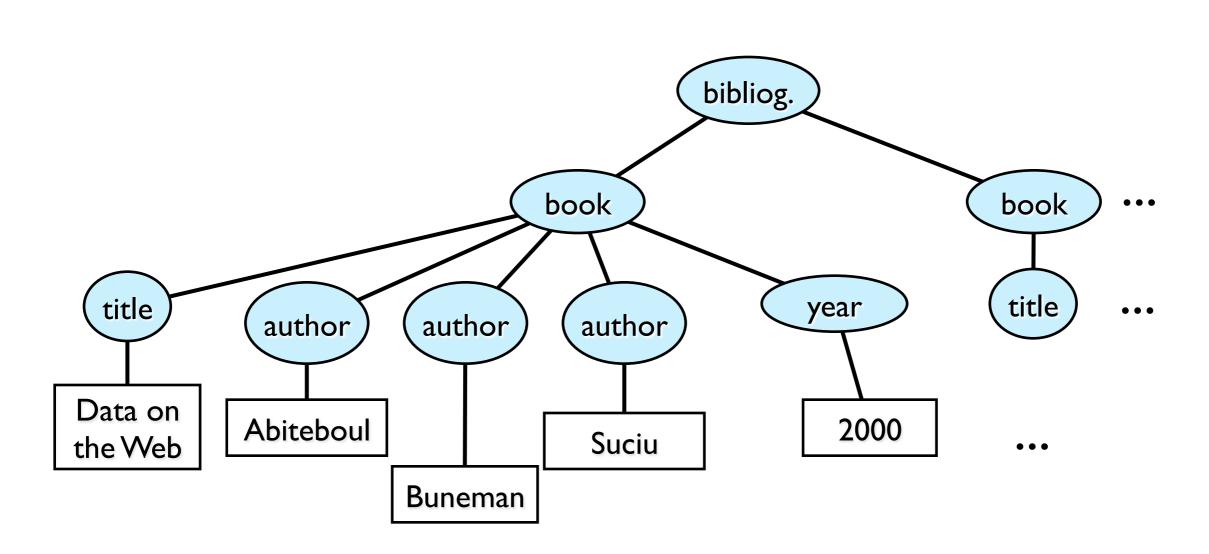


Microsoft

- Language Integrated Query (LINQ) targets SQL & XML in .NET programs
- Data publishing, exchange, integration problems are very important
  - big 3 have products for all of these
  - SQL/XML standard for defining XML views of relational data

#### The essence of XML



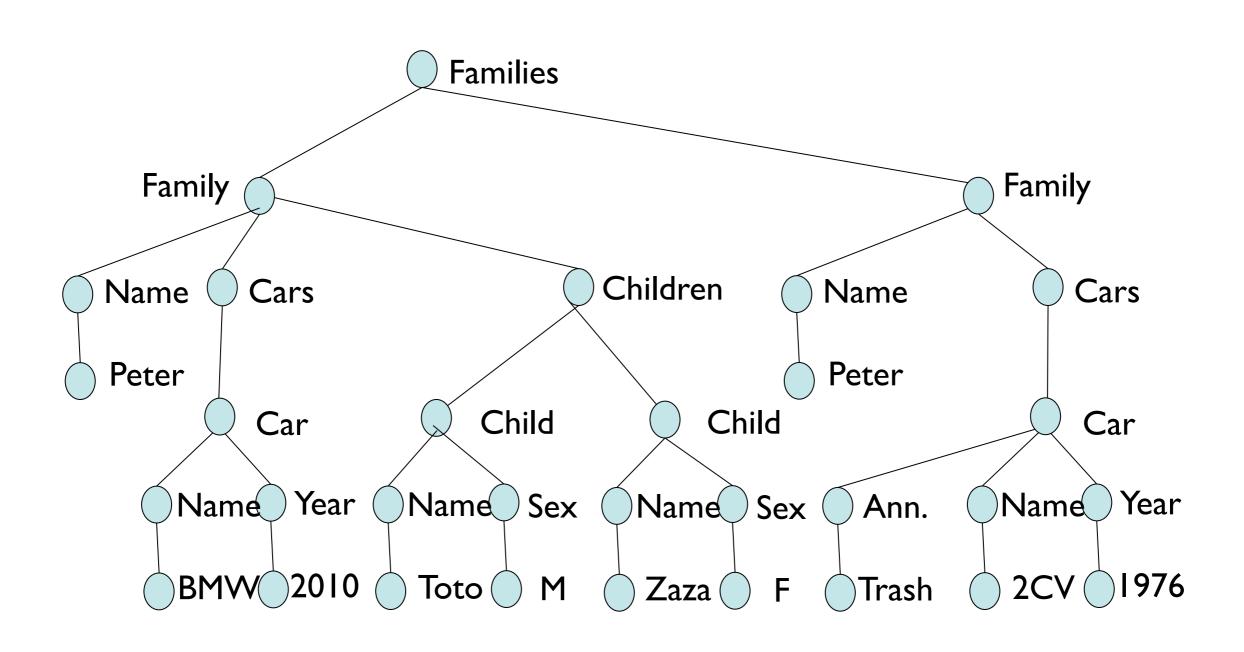


```
<bibliography>
    <book>
         <title>Data on the Web</title>
         <author>Abiteboul</author>
         <author>Buneman</author>
         <author>Suciu</author>
         <year>2000</year>
    </book>
     <book>
         <title>...</title>
Data on
 </bibliography>
```

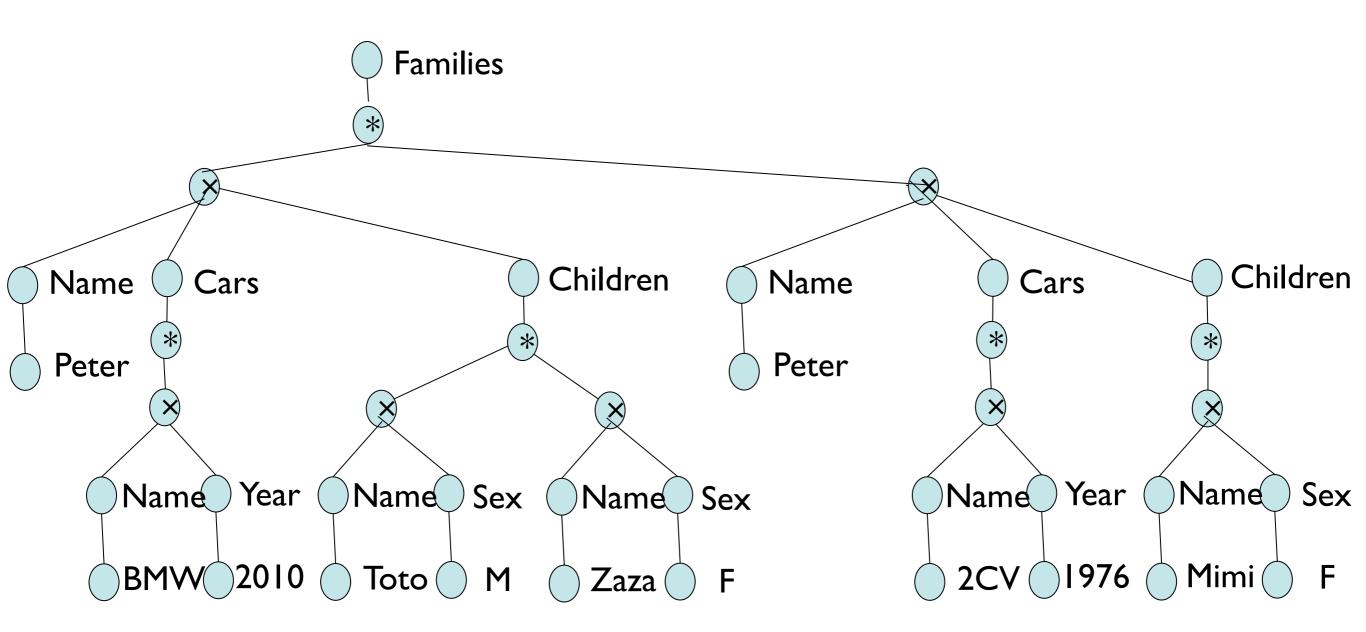
Using trees to represent data: an old idea

- From the 60s and IMS (Hierarchical model, IBM)
- From the 80s and object databases

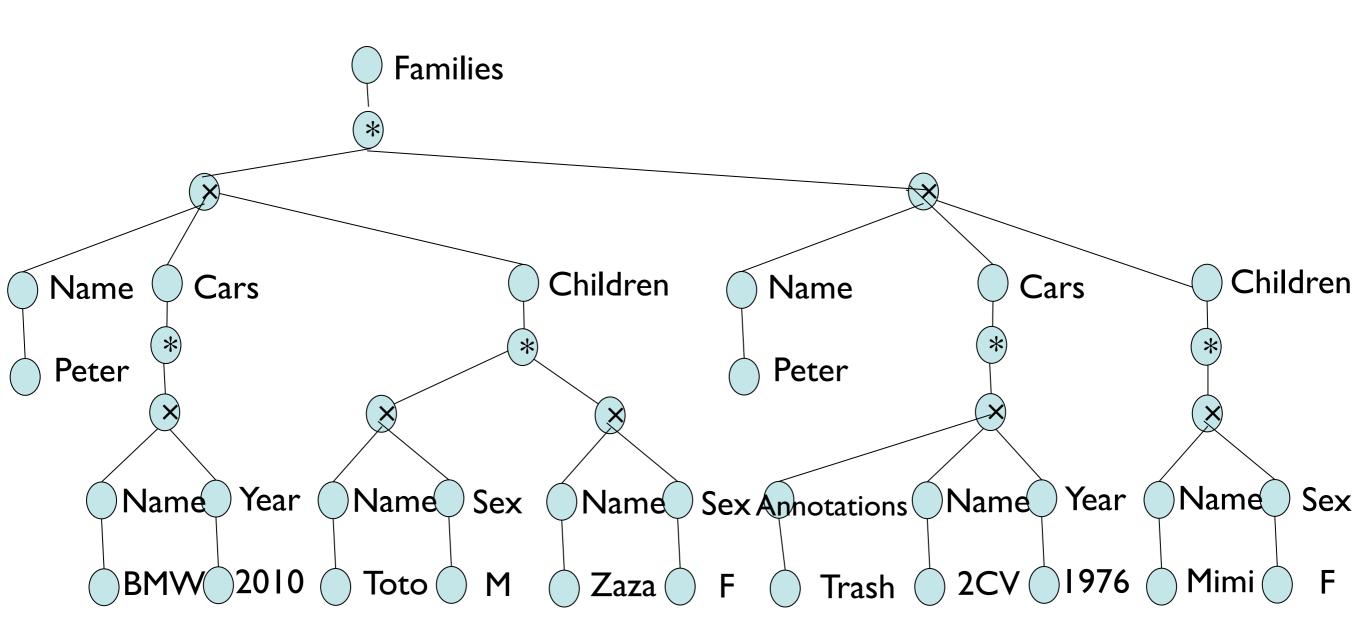
## XML = ordered, labeled, unranked trees



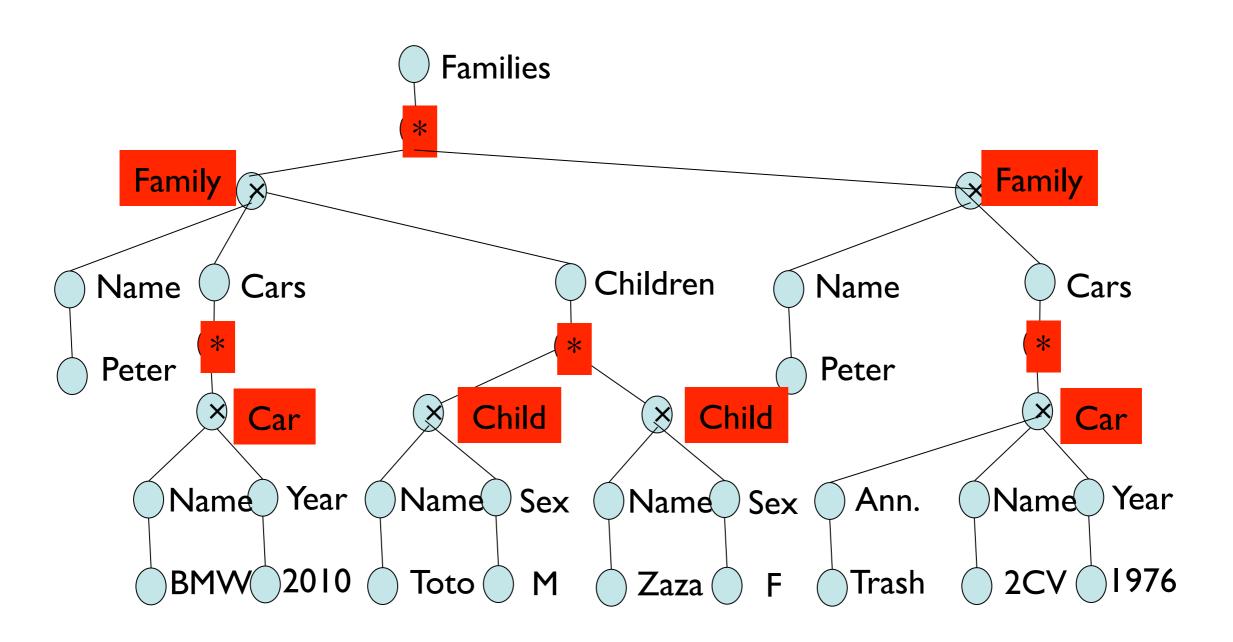
## From database objects to XML trees



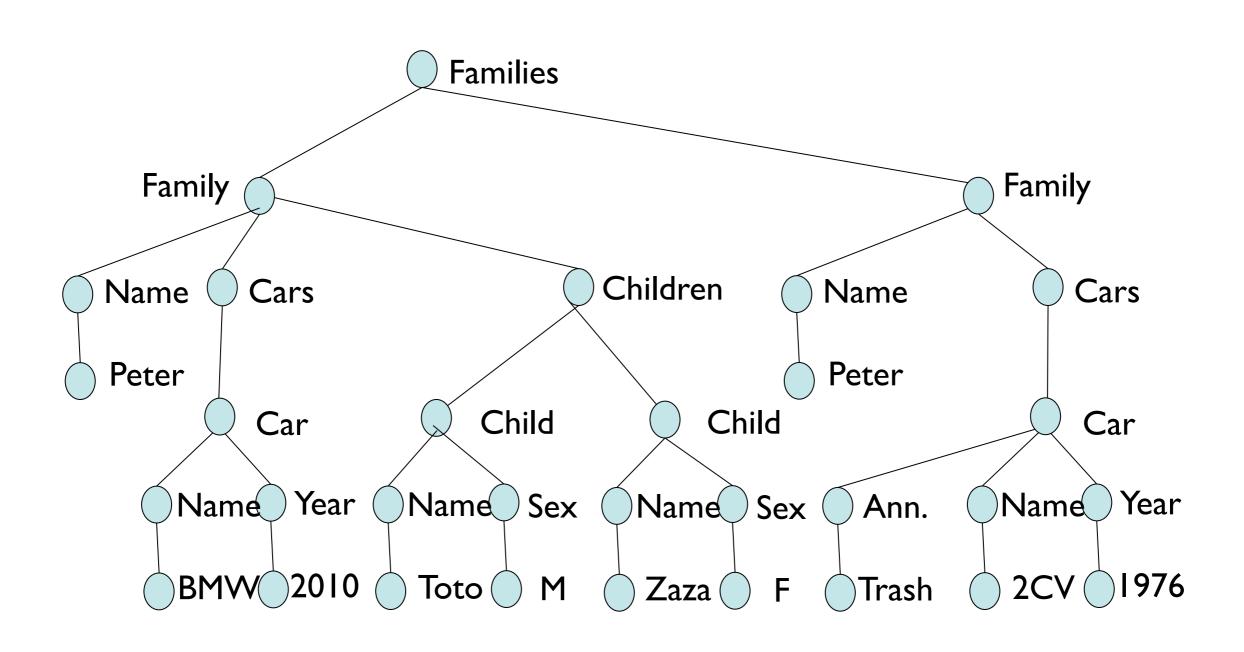
### Revolution I: more flexibility



## Revolution 2: Remove some nodes; name all



## XML = ordered, labeled, unranked trees



## This is better adapted to a Web context

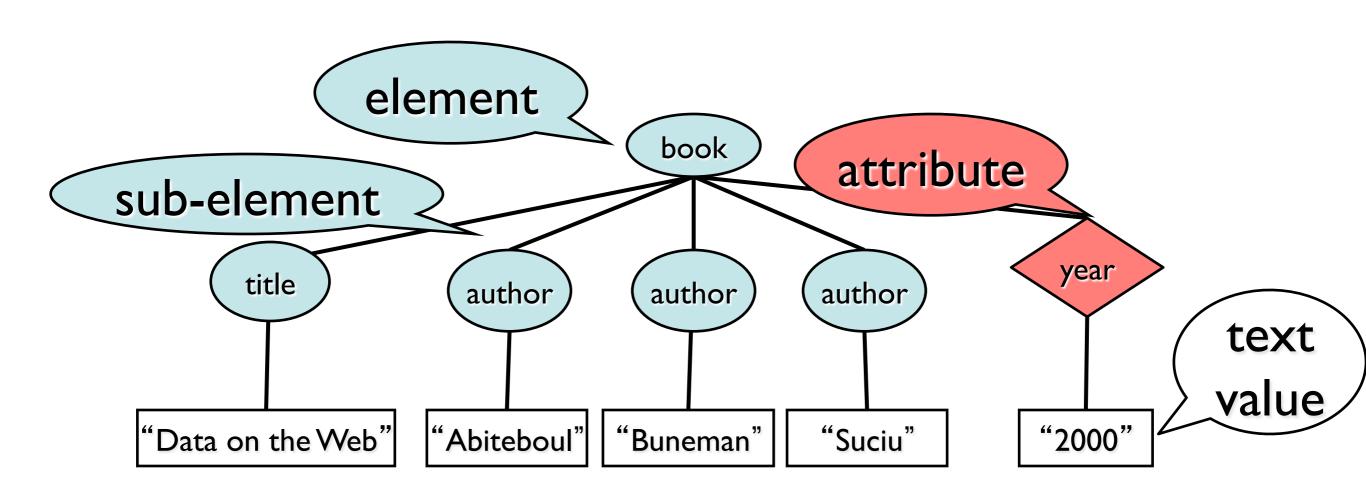
- Self describing data: no separation between schema-vocabulary (tags) and data
  - An XML document can exists without a schema!
  - NB. this was not the case for Object-DBs (and also relationals)

### XML nodes

XML nodes can be of many different kind

- We will focus on the three most important
  - Elements
  - Attributes
  - Text-nodes

## A closer look : elements, attributes, and text-values



### The syntax for elements,

element attribute value value value value

<title>Data on the Web</title>

<author>Abiteboul</author>

sub-element

<author>Bunem text hor>

value

<author>Suciu</author>

</book>

#### Elements

Element: the segment between an start and its corresponding end tag

Unique root element

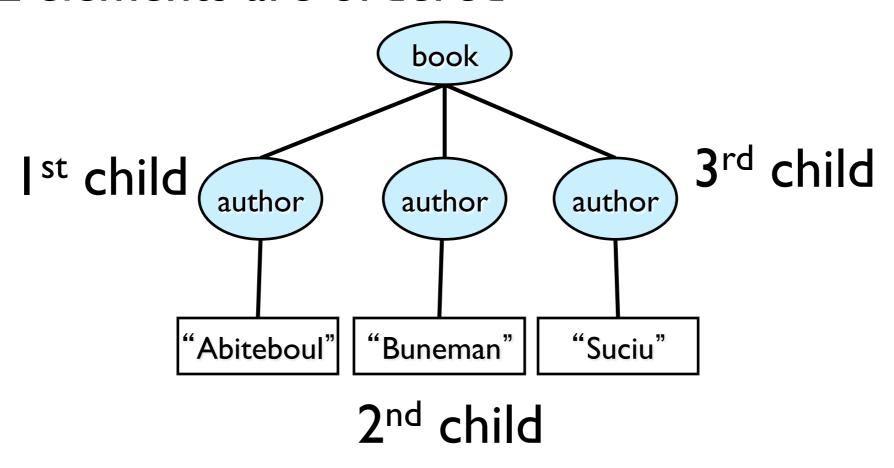
Nested tags can be used to express various "records" and "lists"

#### A document is

 Well-formed: if the opening & closing tags match (condition necessary to be deemed an XML document)

### Ordering

XML elements are ordered



How to represent sets in XML?

## Attributes: Syntax

A start tag may contain attributes describing the element

## Attributes can be used to mimic References

but IDs are not verified without a DTD

## Attribute Structure

XML attributes cannot be nested (they are flat)

The names of XML attributes of an element must be unique.

one can't write

<person pal="Blair" pal="Clinton">

#### Attributes are UNordered

is the same as

#### elements vs attributes

#### How to choose between elements and attributes?

doc I doc 2 doc 3

<note date="12/11/2002"> <note> <note> <to>Tove</to> <date> <date> 12/11/2002 <day> 12 </day> <from>Jani</from> </date> <month>11</month> <year>2002</year> <heading> <to>Tove</to> </date> Reminder </heading> <from>Jani</from> <to>Tove</to> <heading> </note> <from>Jani</from> Reminder </heading> <heading> Reminder </note> </heading> </note>

## elements vs attributes

- attributes cannot contain multiple values (child elements can)
- attributes are not easily expandable (for future changes)
- attributes cannot describe structures (child elements can)
- attributes are more difficult to manipulate by program code
- attribute values are not easy to test against a DTD data
- Use attributes for IDs and Keys.
- Don't end up like this (this is not how XML should be used):

```
<note day="12" month="11" year="2002"
to="Tove" from="Jani" heading="Reminder">
</note>
```

## Quiz: find errors

```
<books>
 <book id="1'>
  <title>Data on the Web</title>
  <authors>
    <author id="a1">Abiteboul
    <author id=a2>Buneman </author>
    <author id='a3'>Suciu</authors>
  </author>
  <year>2000/year>
  <publisher>Addison-Wesley</publisher>
</books>
<foo>bar</foo>
```

# Quiz

```
<books>
<book id="1">
  <title>Data on the Web</title>
  <authors>
    <author id="a1">Abiteboul</author>
    <author id="a2">Buneman </author>
    <author id='a3'>Suciu</authors>
  </author>
  <year>2000
  <publisher>Addison-Wesley</publisher>
</books>
<foo>bar</foo>
```

## Other kinds of nodes

```
entity references: & " >
```

- textual substitution; allows escaping special characters
- you can define your own if you want

```
processing instructions: <? foo : bar ?>
```

can be used to pass information to processors

```
comments: <!-- foo -->
```

CDATA sections: - <!CDATA[[ I <3 XML ]]>

• allows including raw text (<, >, &, etc. uninterpreted)

Luckily, these are mostly irrelevant to use of XML for data

but you need to know about them when writing reading/writing XML as text

# Summing Up

 XML, the standard de-facto for data representation and exchange on the Web.

 Trees are the essence of XML, and there exists a precise syntax to define them.

# DTDs

# Schemas = Types (in XML)

In XML you can define your own markup languages

via, external grammars, aka types, aka schemas

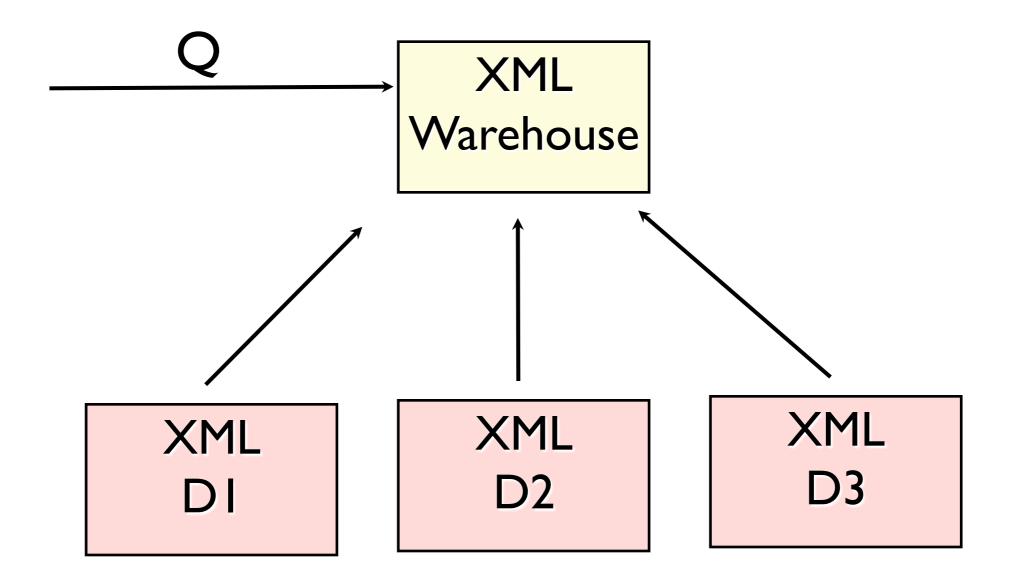
#### A document is

• Valid: means there is a schema and the document matches it

# If they wanted XML to be a new universal language, why again defining schemas?

# Data Exchange/Integration

Easy only when a schema is agreed between peers



# Goals of typing

#### Interoperability/reliability

specify required, optional, default values

#### Consistency

ensure updates or generated output is coherent

#### Efficiency

- use to organize storage
- use as basis for query optimization

## Schemas

Many schema languages/formalisms have been considered

- DTDs (document type definition) (XML I.0)
- XML Schema (W3C)
- Relax/NG (OASIS), DSD, Schematron, ...
- Regular expression types (XDuce, XQuery)

All of these are based on regular expressions

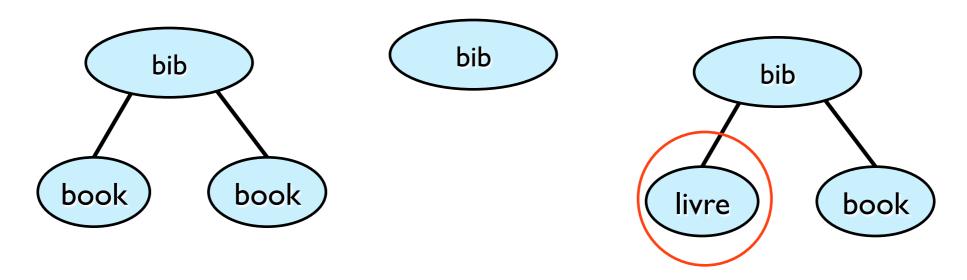
#### DTD

 As for XML, the main components of a DTD are the definitions of elements and attributes

#### Elements

#### Element declarations

- content usually a regular expression over element names
- also allowed: ANY, EMPTY, PCDATA (text)



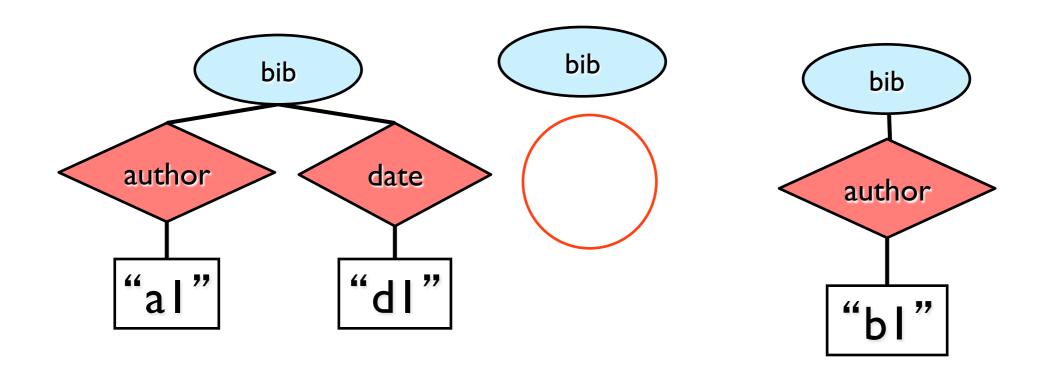
#### Attributes

Attribute declarations make reference to an element (e.g. bib)

<!ATTLIST bib author CDATA #REQUIRED>

<!ATTLIST bib date CDATA #IMPLIED>

•Attributes can also be fixed (#FIXED), and have a specified default value



#### ID/IDRef

ID: Attribute value must be unique within document

<!ATTLIST person pid ID #REQUIRED>

IDREF: Attribute must refer to an ID elsewhere

• <!ATTLIST person pal IDREF #IMPLIED>

## Enumeration

Enumerations: one of a list

<!ATTLIST book type (comic | novel)>

# DTD example: bibliography

```
<!DOCTYPE bib[
<!ELEMENT bib (book* )>
<!ELEMENT book (title, (author+ | editor+ ), publisher, price )>
<!ATTLIST book year CDATA #REQUIRED >
<!ELEMENT author (last, first )>
<!ELEMENT editor (last, first, affiliation )>
<!ELEMENT title (#PCDATA )>
<!ELEMENT last (#PCDATA )>
<!ELEMENT first (#PCDATA )>
<!ELEMENT affiliation (#PCDATA )>
<!ELEMENT publisher (#PCDATA )>
<!ELEMENT price (#PCDATA )>
]>
```

# Quiz: find the error(s)

```
<!ELEMENT root (row*)>
<!ATTLIST root title CDATA #REQUIRED>
<!ELEMENT row (A,(B|C))>
<!ATTLIST row (A | C) >
<!ELEMENT A (#PCDATA)>
<!ELEMENT B (#PCDATA)>
<!ELEMENT C (#PCDATA)>
```

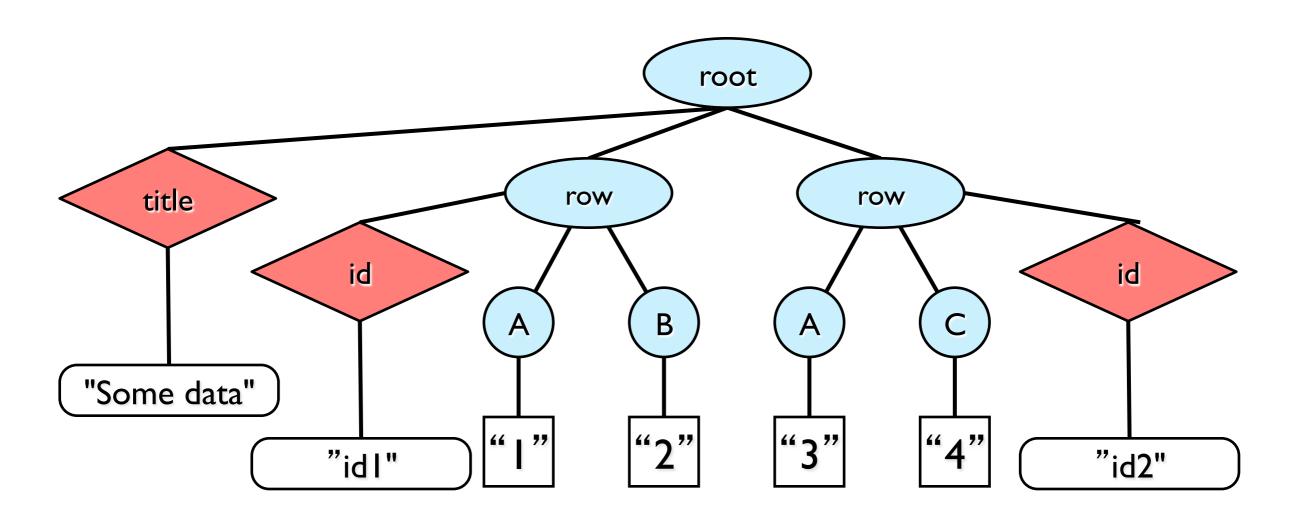
# Quiz: find the error(s)

```
<!ELEMENT root (row*)>
<!ATTLIST root title CDATA #REQUIRED>
<!ELEMENT row (A, (B C))>
<!ELEMENT row (A C) >
                           cannot define twice
                           the same tag row!
<!ELEMENT A (#PCDATA)>
<!ELEMENT B (#PCDATA)>
<!ELEMENT C (#PCDATA)>
```

```
<!ELEMENT root (row*)>
<!ATTLIST root title CDATA #REQUIRED>
```

- <!ELEMENT row (A, (B | C))>
- <!ATTLIST row id ID #REQUIRED>
- <!ELEMENT A (#PCDATA)>
- <!ELEMENT B (#PCDATA)>
- <!ELEMENT C (#PCDATA)>

## Valid



```
<!ELEMENT root (row*)>
<!ATTLIST root title CDATA #REQUIRED>
<!ELEMENT row (A,(B|C))>
```

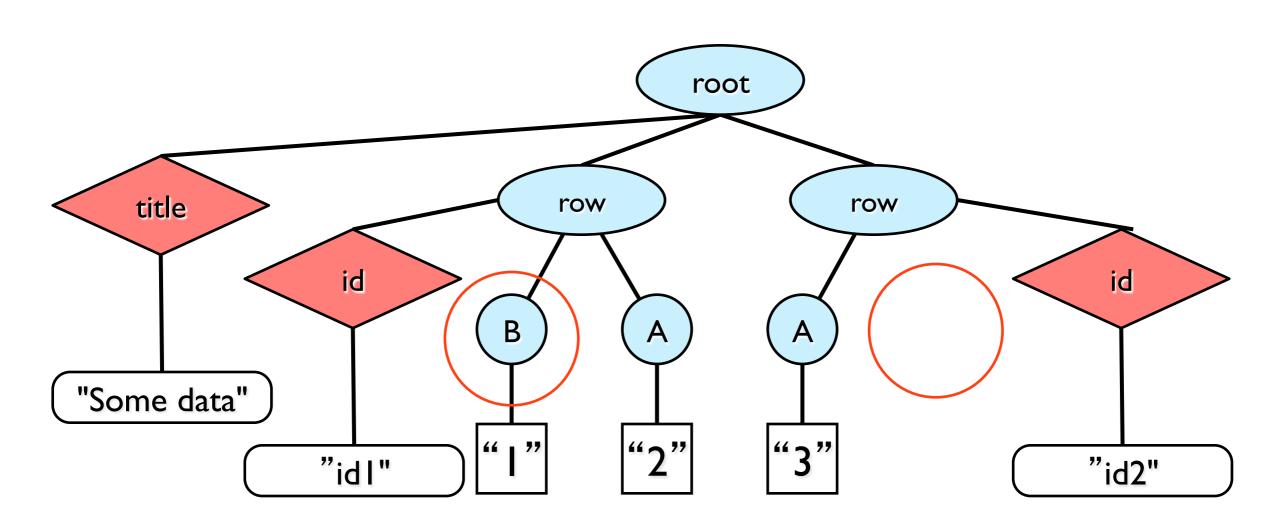
<!ATTLIST row id ID #REQUIRED>

<!ELEMENT A (#PCDATA)>

<!ELEMENT B (#PCDATA)>

<!ELEMENT C (#PCDATA)>

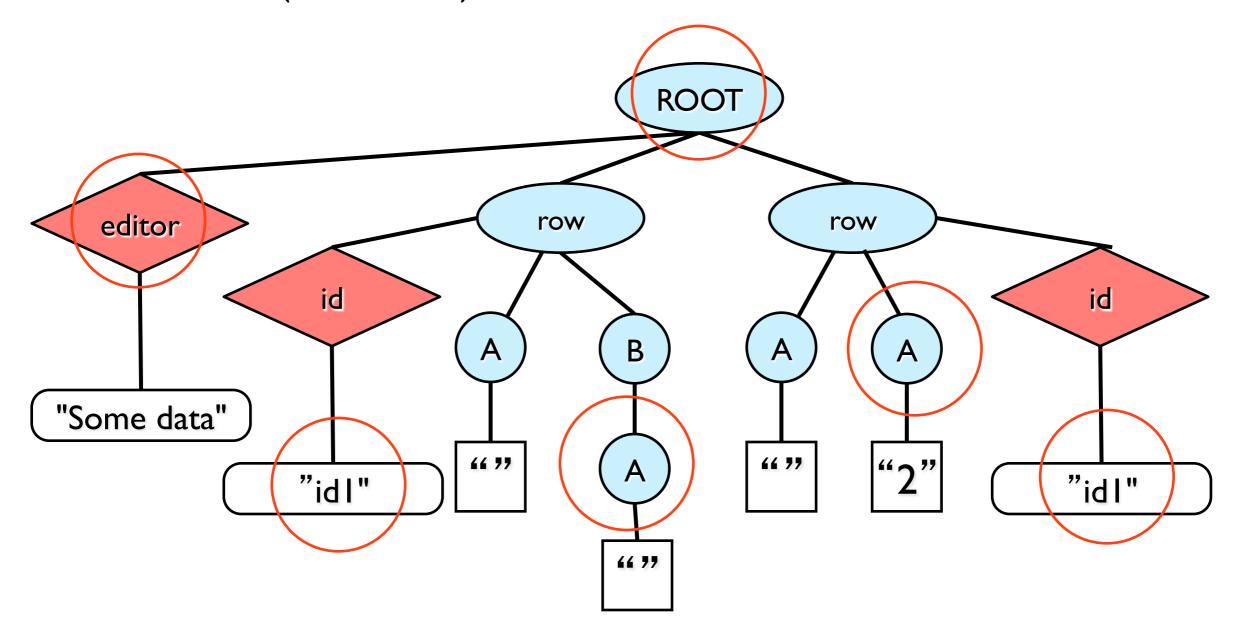
## Invalid



```
<!ELEMENT root (row*)>
```

- <!ATTLIST root title CDATA #REQUIRED>
- <!ELEMENT row (A, (B | C))>
- <!ATTLIST row id ID #REQUIRED>
- <!ELEMENT A (#PCDATA)>
- <!ELEMENT B (#PCDATA)>
- <!ELEMENT C (#PCDATA)>





#### Recursive DTDs

#### DTD rules can be recursive

• node → (node, node)?

#### Recursion increases complexity of DTD

- This leads to documents of unbounded depth
- Some element types might not have any finite matching trees
- but this is easy to detect (look for unguarded cycles)
  - silly → (silly, silly)

## Limitations of DTDs

Can't constrain text / attribute content (except in very limited ways)

Element, attribute content are context insensitive

• can't use same tag, e.g. "name", in different ways

# Quiz

Give a document valid for this DTD, if it exists; otherwise explain why it does not exists.

```
<!ELEMENT X (Y)>
<!ELEMENT Y (A,B,X)>
<!ELEMENT A (#EMPTY)>
<!ELEMENT B (A,B)*>
```

Give a DTD for which only the following XML tree is valid (=no other XML tree is valid!).

Give a document valid for this DTD, if it exists; otherwise explain why it does not exists.

```
<!ELEMENT Y (A)>
<!ELEMENT A (#EMPTY)>
<!ELEMENT A (A,B)*>
```

## XML and DTD together

Coupled Decoupled

```
<?xml version="1.0"?>
<!DOCTYPE bib [
    <!ELEMENT bib book*>
...

J>
<bib> </bib>
```

```
<!DOCTYPE bib [

<!ELEMENT bib book*>
...
]>
```

```
<?xml version="1.0">
<!DOCTYPE bib SYSTEM "bib.dtd">
<bib> </bib>
```

# DTD and Regular Tree Grammars

(fun with regular expressions)

## Plan

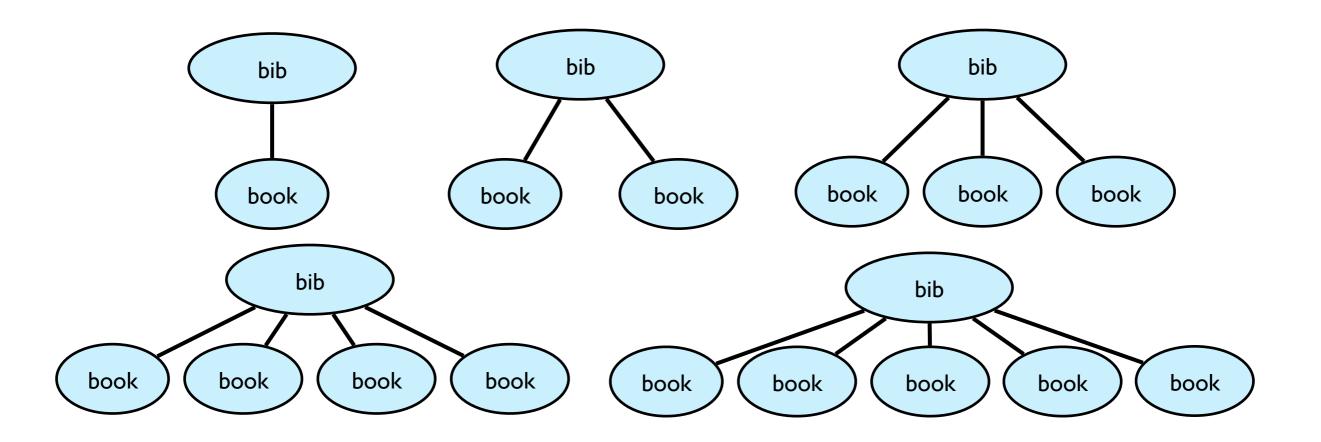
- Regular expressions
- Validation
- Determinism

# Regular Trees

A DTD defines a possibly infinite set of regular XML trees

<!ELEMENT bib book+>

<!ELEMENT book EMPTY>



# Regular Trees

DTD are a subclass of regular tree-grammars called "local"

•This is because any element has at most one definition

<!ELEMENT root child\*>

<!ELEMENT child #PCDATA>

<!ELEMENT child EMPTY>

Regular tree grammars are equivalent to nested regular expressions

# Regular expressions

$$r + = r * | r$$
  $r? = r | \epsilon$ 

## XML Validation

Problem: check if a sequence of children match regular expression

#### **DTD Element**

```
(title, (author+ | editor+ ), publisher, price )
```

#### **XML**

```
<title><author><publisher><price> OK
```

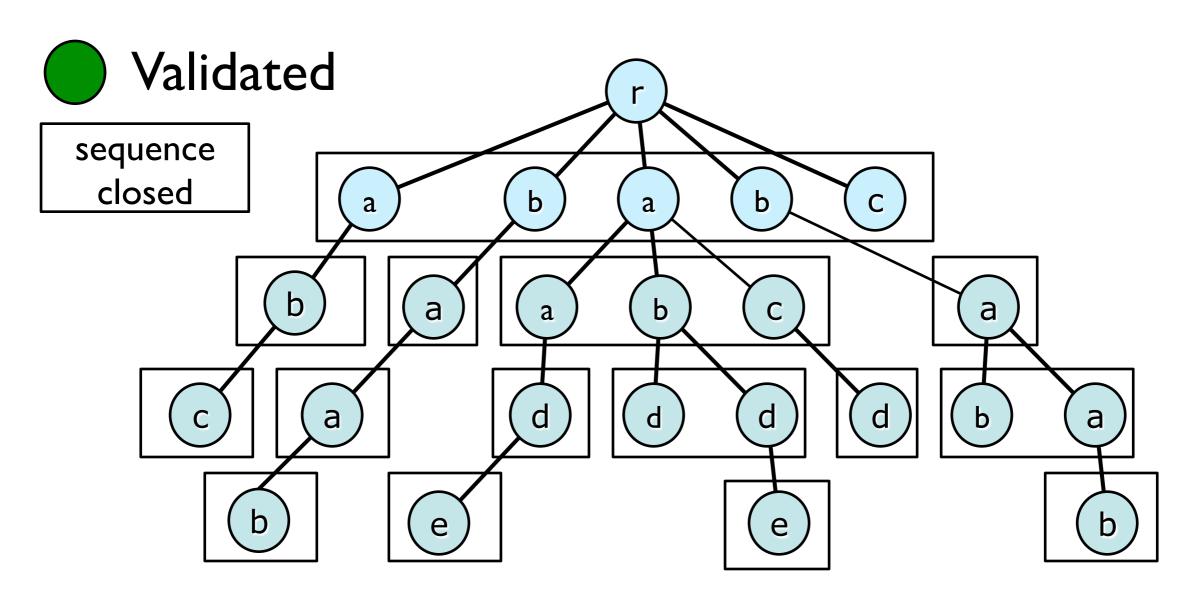
```
<title><author><editor><publisher><price> NO
```

# Validation Algorithm

Traverse XML tree in document-order (=pre-order)

- 1. Check each element's children match regular expression
- 2.Check attribute types
- 3. Check ids are unique and idrefs refer to ids

#### XML Validation



Needs a buffer (worst-case) proportional to the document depth

#### W3C Restriction

Regular expressions in DTDs must be deterministic:

"there must be only one way to match any sequence of tags, no backtrack or look-ahead is required"

```
( title, author* ) | ( title , editor* ) NO
can't decide if <title> matches first or second "title"
Better to write title , ( author* | editor* )
```

#### How to test Determinism?

Simplified algorithm of [Brueggemann-Klein & Wood '98]

Ingredients: three auxiliary functions

FirstTag()

LastTag()

FollowsTag()

# (I/3) FirstTag

What can be the **first** tag of a sequence matching r?

```
r_1 = (title, (author+ | editor+ ), publisher, price )

FirstTag(r_1) ? title

r_2 = (author+ | editor+ )

FirstTag(r_2) ? author, editor
```

# (2/3) LastTag

What can be the **last** tag of a sequence matching r?

```
r_1 = (title, (author+ | editor+ ), publisher, price )

LastTag(r_1) ? price

r_2 = (author+ | editor+ )

LastTag(r_2) ? author, editor
```

## (3/3) Follows Tag

What tag can follow x in r?

```
r_3 = (title, (author+ | editor+ ), publisher, price ) 

FollowsTag(r_3, title) ? author, editor 

r_4 = (author | editor )* 

FollowsTag(r_4, author) ? author, editor
```

## Disambiguation

$$r_5$$
 = (author, title)? , author

We resolve ambiguity by enumerating the tag occurrences

$$r_5^{\#}$$
 = (author<sub>1</sub>, title)? , author<sub>2</sub>

FirstTag(
$$r_5^{\#}$$
) = author<sub>1</sub>, author<sub>2</sub>

LastTag(
$$r_5^\#$$
) = author<sub>2</sub>

FollowsTag(
$$r_5^{\#}$$
, title) = author<sub>2</sub>

# Determinism: Algorithm

- 1) Enumerate all the occurrences of a tag in r
- 2) Build a graph were
  - there is a node x for each tag in  $(r^{\#})$ , plus a root  $x_0$
  - there is a directed edge  $(x_0,y)$  if y belongs to FirstTag $(r^{\#})$
  - there is a directed edge (x,y) if y belongs to FollowsTag $(r^{\#},x)$

3) return **false** if there exists edges  $(x,y_i)$  and  $(x,y_j)$  with  $i\neq j$ 

4) return **true** otherwise

## Testing Determinism

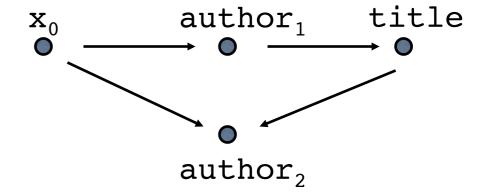
```
r_5 = (author, title)? , author
```

$$r_5^{\#}$$
 = (author<sub>1</sub>, title)? , author<sub>2</sub>

FirstTag( $r_5^\#$ ) = author<sub>1</sub>, author<sub>2</sub>

FollowsTag( $r_5^{\#}$ , author<sub>1</sub>) = title

FollowsTag( $r_5^{\#}$ , title) = author<sub>2</sub>

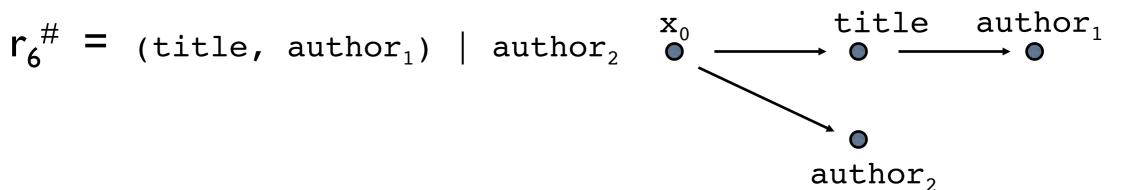


r<sub>5</sub> not deterministic

## Testing Determinism

$$r_6 = (title, author) \mid author$$

$$r_6^{\#}$$
 = (title, author<sub>1</sub>) | author<sub>2</sub>



FirstTag(
$$r_6^{\#}$$
) = title, author<sub>2</sub>

FollowsTag(
$$r_6^{\#}$$
,title) = author<sub>1</sub>

r<sub>6</sub> deterministic

#### Determinism - Quiz

Are the following regular expressions deterministic?

- •(e|cb)\*(c|bed)\*
- •(a,(a|c))|(b,(a|c))

Why did we define **LastTag**(r) afterall?

It is hidden behind the definition of FollowsTag(r,x)

Exercise (pro): define (formally) the 3 auxiliary functions (idea: by induction on the structure of a regular expression)

#### Back to XML Validation

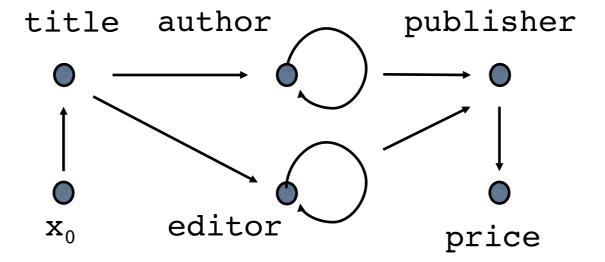
How to check if a sequence of children matches a regular expression?

Comes for free once we computed the graph of r

Simply check if the sequence of children corresponds to a sequence of nodes of a path in the graph ending in a **LastTag**() element

#### Sequence Validation

```
r = (title, (author+ | editor+ ), publisher, price )
```



<title><author><publisher><price> OK
<title><author><editor><publisher><price> NO

# Research Highlights

#### **Checking Determinism**

- Quadratic algorithm based on Glushkov automata [Brueggemann-Klein & Wood, '98]
- (best) Linear algorithm [Groz, Staworko, Maneth 'II]

#### **Checking Validity**

(best) Sublinear space algorithm [Konrad, Magniez 'I I]

# TD/TP (à rendre le 27/09; Moodle)

- Écrire un documents XML
  - Bien formé / Valide par rapport à une DTD
- Écrire une DTD
  - Correcte / Qui permet de valider un document XML
- Vérifier la condition sur le determinisme